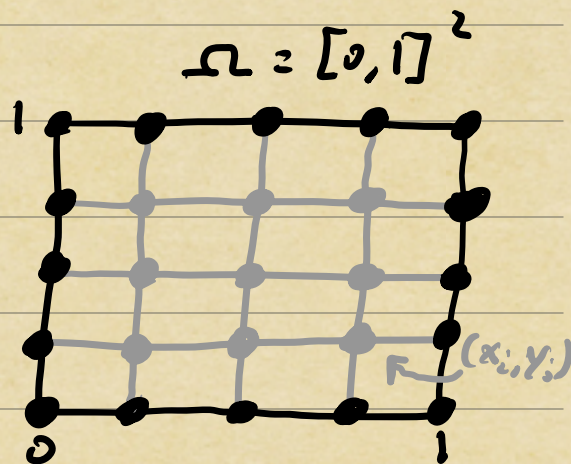


Fast FD Solvers in Mult. Dim.'s

2D Poisson w/ Dirichlet B.C.'s

$$-u_{xx} - u_{yy} = f(x, y)$$

$$u|_{\partial\Omega} = g \quad (\text{Dirichlet B.C.'s})$$



FD discretization on $(n-1) \times (n-1)$ interior grid

$\Rightarrow \mathcal{O}(n^2)$ degrees of freedom

\Rightarrow Aim for $\mathcal{O}(n^2 \log n)$ solver?

$(n-1) \times (n-1)$ matrix $U_{ij} = \text{unknown value of } u(x, y) \text{ at } (i, j)^{\text{th}} \text{ grid point}$

$$-u_{xx}(x_i, y_j) \approx \frac{1}{h^2} (KU)_{ij}, \quad -u_{yy}(x_i, y_j) \approx \frac{1}{h^2} (UK)_{ij}$$

action
cols of U

action
rows of U

$$\frac{1}{h^2} (KU + UK^T) = F_b$$

"Sylvester Equation"

$\hookrightarrow f(x, y)$ modified
on first and last
col/row for D.B.C.'s

Sylvester Matrix Eqn's

\Rightarrow Solved in $O(n^3)$ using Bartels-Stewart, relies on dense matrix factorizations

\Rightarrow Iterative Methods like ADI use fast 1D solvers to build up approx 2D soln.

\Rightarrow For us, Poisson Eq. discretization inherits Fourier eigenseries (fast diagonalization) due to

- (a) separable domain $[0,1] \times [0,1]$
- (b) translation invariant in x and y (up to boundary conditions)

To exploit (a) & (b) in a way that generalizes well to higher dimensions, we'll use

Block Matrices, Vectorization, & Kronecker Products

Caution! Many Eq's below depend on \rightarrow slight mods needed for col-major row-major "vec"

2×2 matrix $X = \begin{bmatrix} x_{11} & x_{12} \\ x_{21} & x_{22} \end{bmatrix}$ $\xrightarrow[\text{ordering}]{\text{row-major}}$ $\text{vec}(X) = \begin{bmatrix} x_{11} \\ x_{12} \\ x_{21} \\ x_{22} \end{bmatrix}$

Apply $A = \begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{bmatrix}$ to rows/cols of X :

XA^T
 \downarrow
 action on rows

$\rightarrow \begin{bmatrix} A & \\ & A \end{bmatrix} \text{vec}(X)$
 $\underbrace{\hspace{1cm}}_{I_2 \otimes A}$

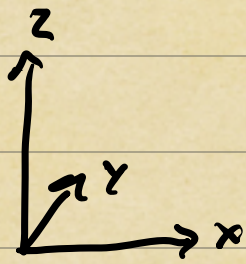
AX
 \downarrow
 action on cols

$\rightarrow \begin{bmatrix} a_{11} & a_{12} & a_{21} & a_{22} \\ a_{11} & a_{12} & a_{21} & a_{22} \\ a_{11} & a_{12} & a_{21} & a_{22} \\ a_{11} & a_{12} & a_{21} & a_{22} \end{bmatrix} \text{vec}(X)$
 $\underbrace{\hspace{1cm}}_{A \otimes I_2}$

Kronecker Products transform linear operators acting along separate array axes into a single linear operator acting on vectorizations.

$n \times m$ $n' \times m'$ $nn' + mm'$
 $A \otimes B = \begin{bmatrix} a_{11}B & a_{12}B & \dots & a_{1m}B \\ a_{21}B & a_{22}B & & a_{2m}B \\ \vdots & & \ddots & \vdots \\ a_{n1}B & a_{n2}B & \dots & a_{nm}B \end{bmatrix}$

Easily extends to higher dimensions



$\text{vec}(X)$ stacks rows into vector n^3

Apply A to x -axis: $(A \otimes I \otimes I) \text{vec}(X)$

y -axis: $(I \otimes A \otimes I) \text{vec}(X)$

z -axis: $(I \otimes I \otimes A) \text{vec}(X)$

Property 1

$$(A \times B) = (A \otimes B^T) \text{vec}(X)$$

Property 2

$$(A \otimes B) \cdot (C \otimes D) = (AC) \otimes (BD)$$

\Rightarrow Systematic manipulation and derivations of highly structured matrices derived from multi-dimensional problems on separable dom.'s.

\Rightarrow Usually avoid explicitly forming $A \otimes B$ for computation, it's expensive!

Structured Block Matrices

Block matrices from Kronecker products of PDE discretizations often retain important structure from 1D problems.

Example: Circulant Block Circulant

\Rightarrow FD discretization of 2D Poisson w/periodic B.C.s

$$C = \begin{bmatrix} C_{11} & C_{31} & C_{21} \\ C_{21} & C_{11} & C_{31} \\ C_{31} & C_{21} & C_{11} \end{bmatrix}$$

C_{11}, C_{21}, C_{31} circulant

$$\Rightarrow C_{ij} = F_n^{-1} \Lambda_{ij} F_n$$

\uparrow DFT matrix

$$N = m \cdot n$$

$$- / m = 3$$

Each block is circulant:

$$\underbrace{\begin{bmatrix} F_n & & \\ & F_n & \\ & & F_n \end{bmatrix}}_{I_m \otimes F_n} C \underbrace{\begin{bmatrix} F_n^{-1} & & \\ & F_n^{-1} & \\ & & F_n^{-1} \end{bmatrix}}_{I_m \otimes F_n^{-1}} = \begin{bmatrix} F_n C_{11} F_n^{-1} & & \\ F_n C_{21} F_n^{-1} & \dots & \\ F_n C_{31} F_n^{-1} & & \end{bmatrix}$$

$$= \begin{bmatrix} \Lambda_{11} & \Lambda_{31} & \Lambda_{21} \\ \Lambda_{21} & \Lambda_{11} & \Lambda_{31} \\ \Lambda_{31} & \Lambda_{21} & \Lambda_{11} \end{bmatrix}$$

Circulant matrices of diagonal blocks!

$\uparrow \Lambda_R$

$$(F_m \otimes I_n) \Lambda_R (F_m^{-1} \otimes I_n) = \begin{bmatrix} \diagdown & & \\ & \diagup & \\ & & \diagdown \end{bmatrix} = \Lambda_{RC} \text{ diagonal}$$

$$(F_m \otimes I_n)(I_m \otimes F_n) C (I_m \otimes F_n^{-1})(F_m^{-1} \otimes I_n)$$

Apply Property 2

$$= (F_m \otimes F_n) C (F_m^{-1} \otimes F_n^{-1})$$

↖ 2D FFT / IFFT ↗

By Property (3),

$$(F_m \otimes F_n) \text{vec}(X) = F_m X F_n$$

$$(F_m^{-1} \otimes F_n^{-1}) \text{vec}(X) = F_m^{-1} X F_n^{-1}$$

So we can interpret these as transforming separately along each array axis.

What is Λ_{RC} ? From theory of cyclic permutations,

$$\Lambda_{RC} = (I_n \otimes \Lambda_n) + (P \otimes \Lambda_{2n}) + (P^2 \otimes \Lambda_{3n}) + \dots + (P^{m-1} \otimes \Lambda_{mn})$$

where $P = \begin{bmatrix} 0 & & 1 \\ & \ddots & \\ 1 & & 0 \end{bmatrix}^{m \times m}$ so that $P^m = I$

Fast solver for $Cy = x$

matrix
maxim

1) $x' = (I_n \otimes F_n) x$ row FFT $\mathcal{O}(mn \log n)$

2) $x'' = (F_m \otimes I_n) x'$ col FFT $\mathcal{O}(nm \log m)$

3) Compute $y'' = L_{RC}^{-1} x''$ $\mathcal{O}(mn)$

4) $y' = (F_n^{-1} \otimes I_m) y''$ col IFFT $\mathcal{O}(nm \log m)$

5) $y = (I_m \otimes F_n^{-1}) y'$ row IFFT $\mathcal{O}(mn \log n)$

Total $\Rightarrow \mathcal{O}(mn \log n) + \mathcal{O}(nm \log m)$

$= \mathcal{O}(mn \log mn)$

$= \mathcal{O}(N \log N)$

$N = mn = \text{total deg. of freedom}$

Can you design a fast Poisson solver for

$$KU + UK^T = F_b$$

based on the DST - Type 1?