

Private-Key Cryptography

1. **Single Key Usage:** Uses one shared key between sender and receiver.
2. **Symmetric Encryption:** Both parties use the same key, making them equals.
3. **Security Issue:** If the key is leaked, communication can be compromised.
4. **Forging Risk:** The receiver could forge a message and claim it was sent by the sender.

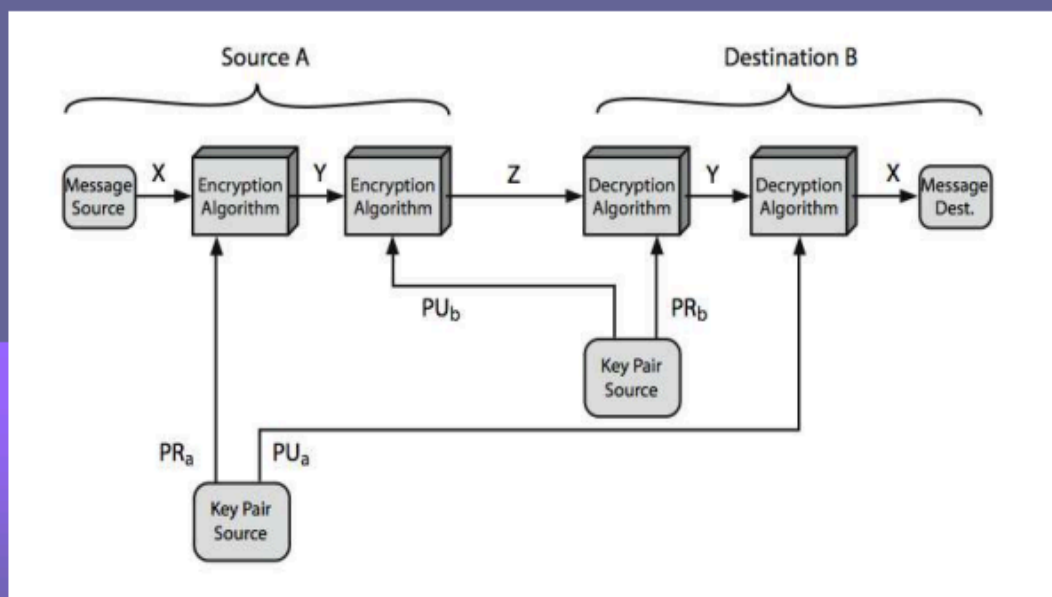
Public-Key Cryptography

1. **Two Keys:** Utilizes a public key (known to everyone) and a private key (kept secret).
2. **Asymmetric Encryption:** The two parties are not equals since they use different keys.
3. **Mathematical Foundation:** Based on complex number theory, making it secure and functional.
4. **Enhances Security:** Complements, rather than replaces, private-key cryptography.

Symmetric vs Public-Key

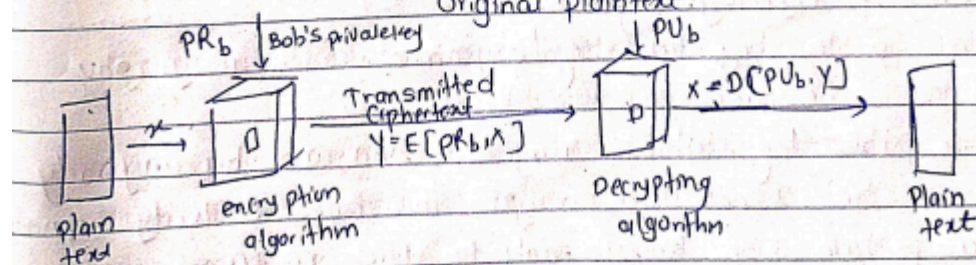
Conventional Encryption	Public-Key Encryption
<p><i>Needed to Work:</i></p> <ol style="list-style-type: none"> 1. The same algorithm with the same key is used for encryption and decryption. 2. The sender and receiver must share the algorithm and the key. <p><i>Needed for Security:</i></p> <ol style="list-style-type: none"> 1. The key must be kept secret. 2. It must be impossible or at least impractical to decipher a message if no other information is available. 3. Knowledge of the algorithm plus samples of ciphertext must be insufficient to determine the key. 	<p><i>Needed to Work:</i></p> <ol style="list-style-type: none"> 1. One algorithm is used for encryption and decryption with a pair of keys, one for encryption and one for decryption. 2. The sender and receiver must each have one of the matched pair of keys (not the same one). <p><i>Needed for Security:</i></p> <ol style="list-style-type: none"> 1. One of the two keys must be kept secret. 2. It must be impossible or at least impractical to decipher a message if no other information is available. 3. Knowledge of the algorithm plus one of the keys plus samples of ciphertext must be insufficient to determine the other key.

Public-Key Cryptosystems

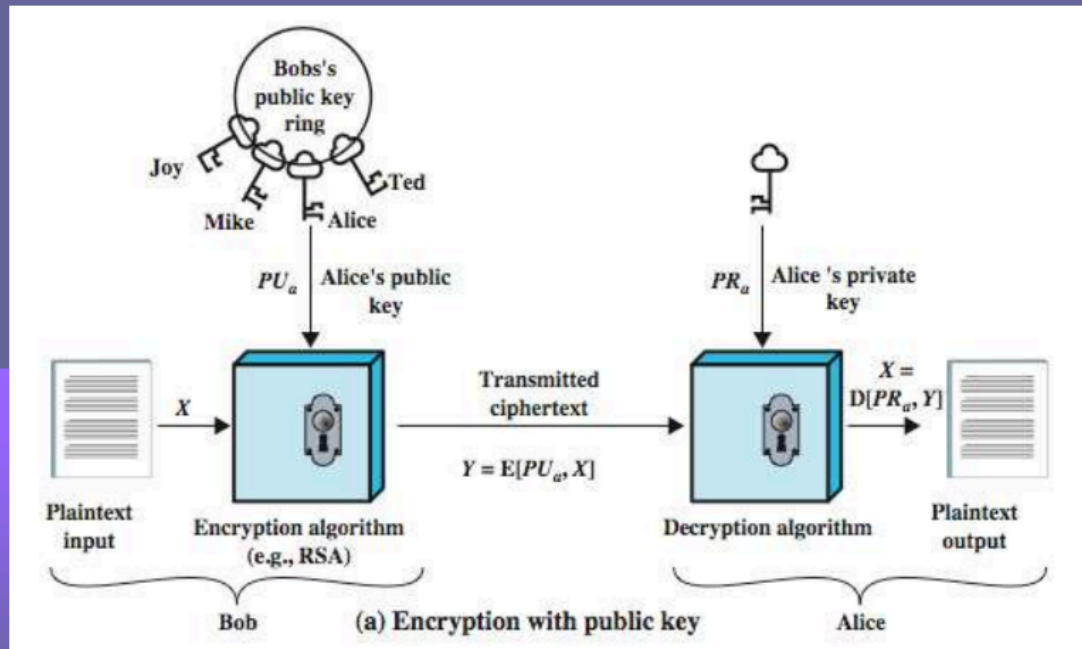


Publickey encryption has 6 ingredients

- ① Plain Text - readable message or data that is fed into the algorithm as input.
- ② Encryption algorithm - performs various transformations on the plain text
- ③ public & private keys - pair of keys that have been selected so that if one is used for encryption, the other is used for decryption.
- ④ Cipher text - encrypted message produced as output depends on plaintext & key. for a given message two keys produce two different ciphertext
- ⑤ Decryption algorithm - accepts the ciphertext & the matching key, produces the original plaintext.



Public-Key Cryptography



Why Public-Key Cryptography?

1. **Key Distribution:** Solves the problem of securely distributing keys without needing a trusted third party.
2. **Digital Signatures:** Verifies the integrity and authenticity of a message.
3. **Invention:** Publicly introduced by Whitfield Diffie & Martin Hellman in 1976.

Public-Key Cryptography Mechanism

1. **Two Keys:**
 - **Public Key:** Used to encrypt messages and verify signatures.
 - **Private Key:** Used to decrypt messages and create signatures.
2. **Asymmetry:** Encryption and decryption keys are distinct, making it impossible to deduce one from the other.

Public-Key Cryptosystems Applications

1. **Encryption/Decryption:** Ensures message secrecy.
2. **Digital Signatures:** Provides authentication.
3. **Key Exchange:** Secures the exchange of session keys.

Public-Key Requirements

1. **Infeasibility:** It should be nearly impossible to determine the private key from the public key.
2. **Efficient Operations:** Encryption and decryption must be computationally easy with the correct key.
3. **Trap-Door One-Way Function:** Allows easy computation in one direction but makes the reverse infeasible without a special piece of information (trapdoor).

Security of Public Key Schemes

1. **Brute Force Attack:** Theoretically possible but impractical due to the large key size (e.g., >512 bits).
2. **Relies on Mathematical Complexity:** Difficult problems like factoring large numbers ensure security.
3. **Large Numbers Requirement:** Security depends on using large primes and numbers.

RSA Algorithm

1. **Inventors:** Developed by Rivest, Shamir, and Adleman in 1977.
2. **Basis:** Works on exponentiation over integers modulo a prime, utilizing large integers (e.g., 1024 bits).
3. **Security:** Relies on the difficulty of factoring large numbers.

RSA Encryption/Decryption Process

1. **Encryption:**
 - **Public Key:** Consists of $PU = \{e, n\}$.
 - **Ciphertext:** $C = M^e \bmod n$, where M is the message, where $0 \leq M < n$
2. **Decryption:**
 - **Private Key:** Consists of $PR = \{d, n\}$.

- **Message Recovery:** $M = C^d \bmod n$.

RSA Key Setup

1. **Choose Primes:** Randomly select two large primes p and q .
2. **Compute Modulus:** $n = p * q$.
3. **Euler's Totient Function:** Calculate $\phi(n) = (p-1)(q-1)$.
4. **Choose Encryption Key:** e , such that $1 < e < \phi(n)$ and $\gcd(e, \phi(n)) = 1$.
5. **Compute Decryption Key:** d such that $e * d = 1 \bmod \phi(n)$.
6. **Publish their public encryption key:** $PU = \{e, n\}$
7. **Keep secret private decryption key:** $PR = \{d, n\}$

Why RSA Works

1. **Based on Euler's Theorem:** Ensures that the encryption and decryption processes correctly recover the original message.

Example RSA Key Setup and Operation

1. **Select Primes:** $p = 17, q = 11$.
2. **Modulus Calculation:** $n = 17 * 11 = 187$.
3. **Totient Calculation:** $\phi(n) = 160$.
4. **Choose e :** $e = 7$.
5. **Compute d :** $d = 23$ (since $7 * 23 = 161 \bmod 160$).
6. **Public Key:** $\{7, 187\}$, **Private Key:** $\{23, 187\}$.
7. **Encrypt/Decrypt Example:**
 - Message $M = 88$.
 - Ciphertext $C = 88^7 \bmod 187 = 11$.
 - Decrypt to get $M = 11^{23} \bmod 187 = 88$.

Efficient Encryption and Decryption

1. **Small Encryption Exponent:** Choosing a small e (e.g., 65537) makes encryption faster.
2. **Chinese Remainder Theorem (CRT):** Used for faster decryption by breaking down calculations.

RSA Security

1. **Factorization Problem:** The difficulty of factoring the modulus n is central to RSA security.
2. **Countermeasures:**
 - **Timing Attacks:** Avoid by constant-time operations.
 - **Chosen Ciphertext Attacks:** Use padding schemes like OAEP for protection.