

# Multidimensional Disjunctive Inequalities for Chance-Constrained Stochastic Problems with Finite Support

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# (Mixed-Integer) Linear Chance-Constrained Problem (LCCP)

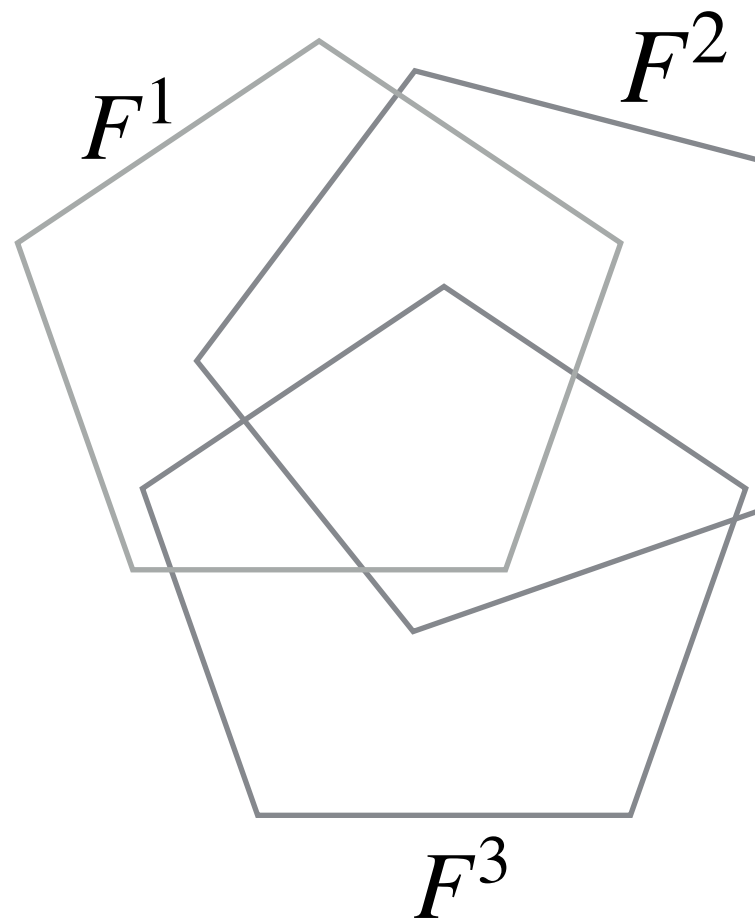
$$\begin{array}{ll}\min_{x} & c^{\top} x \\ \text{s.t.} & x \in X, \\ & \mathbb{P}(A^{\omega}x - b^{\omega} \geq 0) \geq 1 - \tau\end{array}$$

$\omega$  : random variable with discrete support

# LCCP: applications

- Energy: optimal power flow
- Finance: portfolio optimisation
- Communication: reliable network design
- Chemical processing
- Production planning
- Water resources management

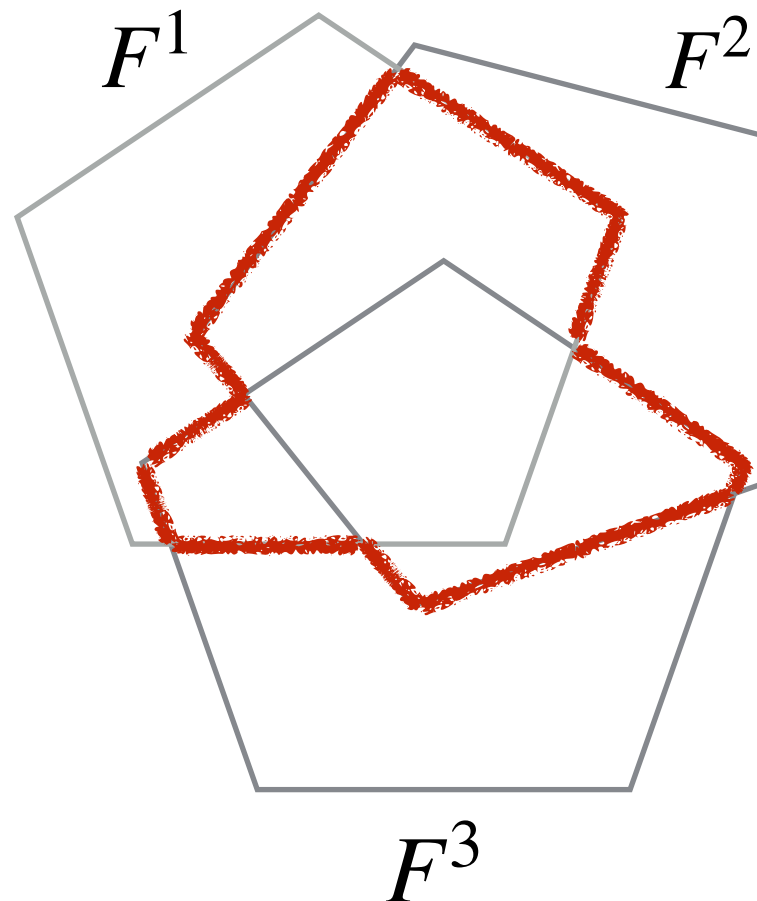
# LCCP



$$\mathbb{P}(F^i) = \frac{1}{3}, i = 1, 2, 3$$

$$1 - \tau = \frac{2}{3}$$

# LCCP



Feasible domain

=

union of exponential number of convex sets

# LCCP: Complexity

LCCP is **strongly NP-hard** even if  
all scenarios have equal probability and

there is only one constraint per scenario  
(Amaldi & Kann 1995)

only the RHS vector is random (Luedtke et al.  
2010)

# LCCP: MIP formulation

$$\begin{array}{ll}\min_{x,y} & c^\top x \\ \text{s.t.} & x \in X, \\ & A^s x \geq \bar{b}^s + (b^s - \bar{b}^s)y^s, \quad s \in S, \\ & \sum_{s \in S} p^s y^s \geq 1 - \tau, \\ & y^s \in \{0, 1\}, \quad s \in S.\end{array}$$

# Performance of CPLEX

(from Qiu et al. 2014)

20 variables, 200 scenarios, 1 constraint per scenario

**Table 1** Performance of CPLEX on CKVLP

$k$	Dense data			Sparse data			Random objective		
	Gap closed (%)	Nodes	Time	Gap closed (%)	Nodes	Time	Gap closed (%)	Nodes	Time
15	2	3,537,864	2,454	7	158,039	83	17	1,777	1
20	2	43,296,679	25,948	6	1,769,574	917	21	6,227	2



# Some related papers

Qiu, F., Ahmed, S., Dey, S. S., & Wolsey, L. A. (2014). Covering linear programming with violations. *INFORMS Journal on Computing*, 26(3), 531-546.

Luedtke, J. (2014). A branch-and-cut decomposition algorithm for solving chance-constrained mathematical programs with finite support. *Mathematical Programming*, 146(1-2), 219-244.

Ahmed, S., Luedtke, J., Song, Y., & Xie, W. (2017). Nonanticipative duality, relaxations, and formulations for chance-constrained stochastic programs. *Mathematical Programming*, 162, 51-81.

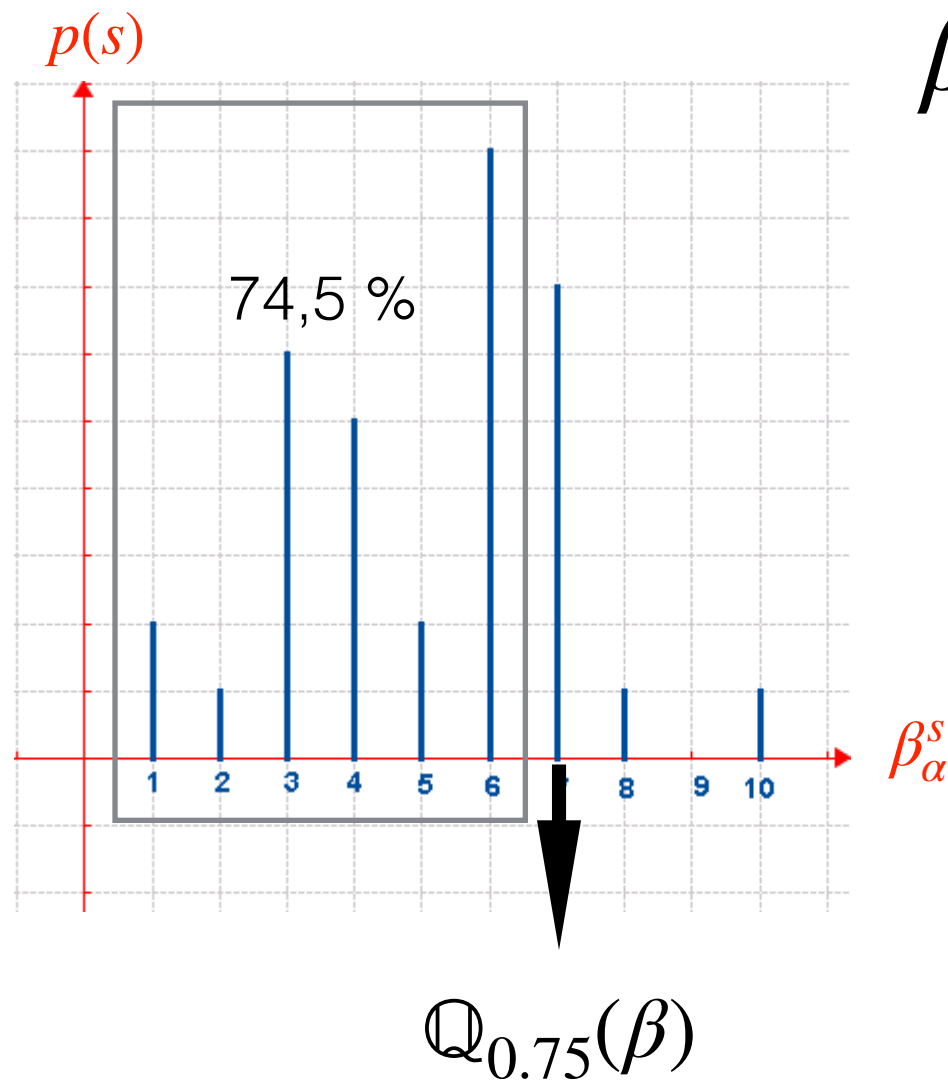
Xie, W., & Ahmed, S. (2018). On quantile cuts and their closure for chance constrained optimization problems. *Mathematical Programming*, 172, 621-646.

Song, Y., & Luedtke, J. R. (2013). Branch-and-cut approaches for chance-constrained formulations of reliable network design problems. *Mathematical Programming Computation*, 5(4), 397-432.

Song, Y., Luedtke, J. R., & Küçükyavuz, S. (2014). Chance-constrained binary packing problems. *INFORMS Journal on Computing*, 26(4), 735-747.

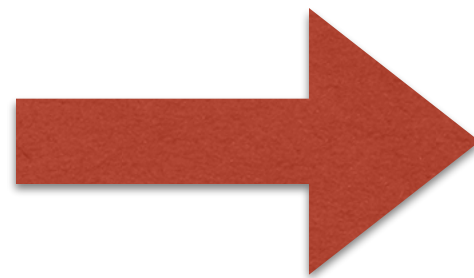
Cattaruzza, D., Labbé, M., Petris, M., Roland, M., & Schmidt, M. (2024). Exact and Heuristic Solution Techniques for Mixed-Integer Quantile Minimization Problems. *INFORMS Journal on Computing*.

# Quantile cut



$$\beta_{\alpha}^s = \min_x \{ \alpha^{\top} x : A^s x \geq b^s, x \geq 0 \}$$

$$\beta^{\alpha} = \mathbb{Q}_{1-\tau}(\beta_{\alpha}^s)$$



$$\alpha^{\top} x \geq \beta_{\alpha}$$

# Mixing set inequality (Günlük and Pochet 2001)

$$\alpha^\top x \geq \beta_\alpha + (\beta_k - \beta_\alpha)y_k$$

for  $\beta_k > \beta_\alpha$

# Relaxed Problem

$$\lambda^s \geq 0, \quad e^\top \lambda^s = 1, \quad s \in S$$

$$(\lambda^s)^\top A^s x \geq (\lambda^s)^\top \bar{b}^s + (\lambda^s)^\top (b^s - \bar{b}^s) y^s, \quad s \in S,$$

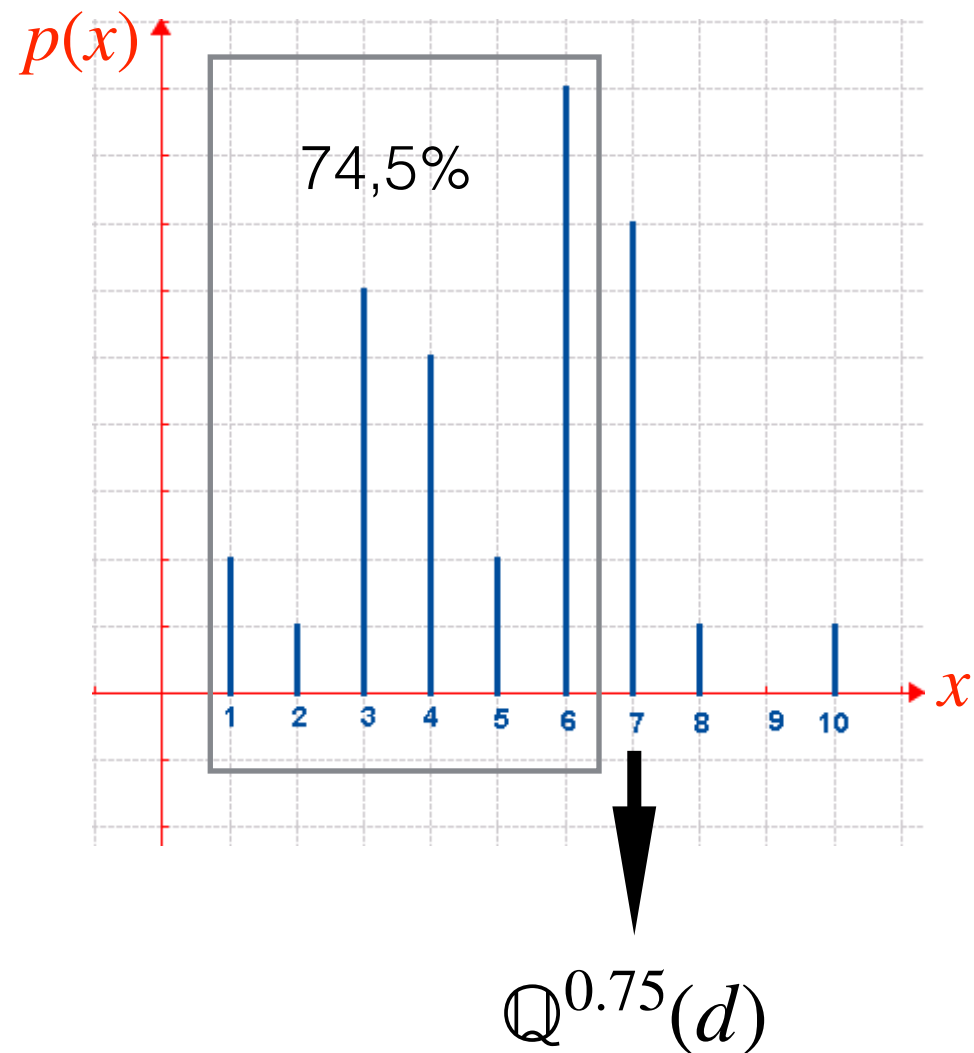
$$\sum_{s \in S} p^s y^s \geq 1 - \tau,$$

$$y^s \in \{0, 1\}, \quad s \in S.$$



$$\mathbb{Q}_{1-\tau}[(\lambda^s)^\top (b^s - A^s x)] \leq 0$$

# Quantile



$$\begin{array}{ll} \max_{u, q} & (1 - \tau)q - \sum_{s \in \mathcal{S}} p^s u^s \\ \text{s.t.} & q - u^s \leq d^s, \quad u^s \geq 0, \quad s \in \mathcal{S}, \\ & q \in \mathbb{R}. \end{array}$$

$$\begin{array}{ll} \min_w & \sum_{s \in \mathcal{S}} d^s w^s \\ \text{s.t.} & \sum_{s \in \mathcal{S}} w^s = 1 - \tau, \\ & 0 \leq w^s \leq p^s, \quad s \in \mathcal{S}. \end{array}$$

$$(1 - \tau)q - \sum_{s \in \mathcal{S}} p^s u^s \geq \sum_{s \in \mathcal{S}} d^s w^s$$

# Primal-Dual (PD) inequality

$$\sum_{s \in \hat{S}} p^s (\lambda^s)^\top (A^s x - b^s) - \sum_{i=1}^n U_i(\lambda) x_i + L(\lambda) \leq 0$$

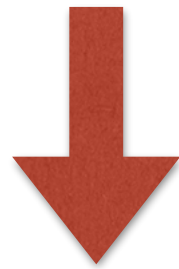
$$U_i(\lambda) = \max_w \left\{ \sum_{s \in S} (\lambda^s)^\top A_{\cdot i}^s w^s : \sum_{s \in S} w^s = 1 - \tau, \ 0 \leq w^s \leq p^s, \ s \in S \right\}$$

and

$$L(\lambda) = \min_w \left\{ \sum_{s \in S} (\lambda^s)^\top b^s w^s : \sum_{s \in S} w^s = 1 - \tau, \ 0 \leq w^s \leq p^s, \ s \in S \right\}$$

# multi-disjunctive (multi-D) inequality

Set  $\bar{S}$  of scenarios s.t.  $p(\bar{S}) \geq \tau$



$$(\lambda^s)^\top (A^s x - b^s) \geq 0, \quad \text{for } s \in \hat{S} \subseteq \bar{S} : p(\hat{S}) \geq p(\bar{S}) - \tau$$

$$\sum_{s \in \hat{S}} w^s (\lambda^s)^\top A^s x \geq \sum_{s \in \hat{S}} w^s (\lambda^s)^\top b^s$$

# multi-disjunctive (multi-D) inequality

$$\sum_{i=1}^n U_i(\lambda, \bar{S}) x_i \geq L(\lambda, \bar{S}), \quad \bar{S} \subseteq \mathcal{S}, \quad p(\bar{S}) > \tau,$$

with

$$U_i(\lambda, \bar{S}) = \max_w \left\{ \sum_{s \in \bar{S}} (\lambda^s)^\top A_{\cdot i}^s w^s : \sum_{s \in \bar{S}} w^s = p(\bar{S}) - \tau, \quad 0 \leq w^s \leq p^s, \quad s \in \bar{S} \right\}$$

and

$$L(\lambda, \bar{S}) = \min_w \left\{ \sum_{s \in \bar{S}} (\lambda^s)^\top b^s w^s : \sum_{s \in \bar{S}} w^s = p(\bar{S}) - \tau, \quad 0 \leq w^s \leq p^s, \quad s \in \bar{S} \right\}$$

dominates P-D inequality



# Separation

- All VIs are hard to separate: Determine  $\bar{S}$  and  $\lambda^s, s \in S$
- For given  $\lambda^s, s \in S$ , PD inequalities are easy to separate
- For given  $\lambda^s, s \in S$ , multi-D inequalities are hard to separate
- For given  $\bar{S}$ , PD and multi-D inequalities are easy to separate (LPs)

# Chance Constraint- Covering LP

$$\begin{array}{ll}\min_{x,y} & c^\top x \\ \text{s.t.} & A^s x \geq y^s \mathbf{1}, \quad s \in S, \\ & x \geq 0, \\ & \sum_{s \in S} p^s y^s \geq 1 - \tau, \\ & y^s \in \{0, 1\}, \quad s \in S\end{array}$$

# Closures for CC-Covering-LP

- $\mathcal{Q} = \cap$  all quantile cuts
- $\text{multi-}\mathcal{D} = \cap$  all multi-disjunctive inequalities
- $\text{single-}\mathcal{D} = \cap$  all single-disjunctive inequalities  
(  $p(\bar{S}) > \tau$  and  $p(\bar{S} \setminus \{s\}) < \tau$ , for all  $s \in \bar{S}$  )

Theorem:  $\text{multi-}\mathcal{D} \subseteq \text{single-}\mathcal{D} = \mathcal{Q}$

- There exists instances where the inclusion is strict

# Cut impact: single-covering

TABLE 1. Results of the CP-tests with MD-VIs, mixing-set inequalities and quantile cuts on the instances of the CKVLP.

Instances			MD-VIs			Mixing-set			Quantile		
$n-m$	Gen.	$ \mathcal{S} $	T	it.	$v$ (%)	T	it.	$v$ (%)	T	it.	$v$ (%)
20-1	lit	100	0.00	3.80	<b>58.90</b>	0.00	2.40	17.32	0.00	2.40	15.82
		500	0.01	3.40	<b>50.97</b>	0.01	2.00	23.10	0.01	2.00	22.30
		1000	0.02	3.40	<b>34.43</b>	0.01	3.60	19.87	0.01	2.40	19.03
		3000	0.08	3.20	<b>36.53</b>	0.05	6.40	20.92	0.04	3.40	17.88
	new	100	0.00	3.00	<b>32.61</b>	0.00	2.40	30.57	0.00	2.20	30.11
		500	0.01	3.00	37.24	0.01	2.20	<b>41.82</b>	0.01	2.00	41.72
		1000	0.02	3.00	34.01	0.01	3.20	<b>48.76</b>	0.01	2.00	48.76
		3000	0.09	3.00	31.72	0.05	6.20	<b>37.67</b>	0.04	2.20	37.46
80-1	lit	100	0.00	4.40	<b>35.99</b>	0.00	3.00	6.14	0.00	2.40	3.51
		500	0.01	4.00	<b>33.03</b>	0.01	3.60	7.81	0.01	2.60	5.88
		1000	0.03	4.40	<b>22.91</b>	0.02	3.80	6.38	0.02	3.20	3.77
		3000	0.10	3.20	<b>42.22</b>	0.07	4.00	8.36	0.07	2.20	7.21
	new	100	0.00	3.20	<b>27.04</b>	0.00	3.40	17.37	0.00	3.20	15.39
		500	0.01	3.00	<b>31.13</b>	0.01	3.00	26.55	0.01	2.60	24.79
		1000	0.03	3.00	<b>30.64</b>	0.02	4.40	28.29	0.02	3.40	27.46
		3000	0.12	3.00	<b>33.62</b>	0.08	6.00	30.99	0.07	2.60	30.80

# Cut impact: multi- knapsack

TABLE 2. Results of the CP with MD-VIs, Mixing-set inequalities and quantile cuts on the instances of the CCMKP.

Instances			MD-VIs												
			ADM				C-wise			Mixing-set			Quantile		
$n-m$	Gen.	$ \mathcal{S} $	$s(\lambda)$	T	it.	$v$ (%)	T	it.	$v$ (%)	T	it.	$v(\%)$	T	it.	$v$ (%)
20-10	lit	100	0.87	0.18	18.80	<b>30.21</b>	0.01	4.00	28.05	0.11	4.80	25.07	0.11	4.20	21.18
		500	0.73	0.42	13.80	<b>25.50</b>	0.04	4.00	23.47	0.52	6.20	22.10	0.51	3.80	19.84
		1000	0.89	0.81	15.00	<b>27.00</b>	0.07	3.80	25.14	1.09	9.00	24.47	1.04	5.00	21.89
		3000	0.88	3.27	11.80	<b>25.80</b>	0.26	3.40	24.77	3.88	17.80	22.99	3.23	4.80	20.82
	new	100	0.33	0.23	21.60	<b>24.03</b>	0.01	3.60	4.30	0.12	5.20	16.30	0.12	2.40	0.92
		500	0.33	0.43	30.80	<b>27.27</b>	0.04	4.20	10.97	0.60	8.00	18.60	0.58	4.00	1.55
		1000	0.29	0.79	15.80	<b>25.10</b>	0.09	3.00	11.31	1.17	8.40	15.92	1.09	4.20	2.65
		3000	0.31	3.04	22.20	<b>25.20</b>	0.37	3.80	5.21	3.72	18.20	16.67	3.40	3.80	2.11
40-30	lit	100	0.77	0.62	39.80	<b>27.65</b>	0.05	4.20	25.05	1.36	6.40	22.78	1.37	5.60	15.81
		500	0.79	1.27	21.00	<b>26.13</b>	0.19	3.80	23.66	6.72	7.40	19.72	6.67	4.60	16.04
		1000	0.73	3.17	47.40	<b>28.19</b>	0.50	4.20	25.46	13.55	12.40	22.74	11.83	6.20	18.94
		3000	0.79	8.96	20.80	<b>25.59</b>	1.66	3.60	23.25	41.06	17.20	19.05	38.16	5.20	15.85
	new	100	0.23	0.76	63.80	<b>26.08</b>	0.07	2.60	4.18	1.25	6.80	17.78	1.37	3.00	0.99
		500	0.26	1.15	31.80	<b>23.30</b>	0.28	3.20	6.42	6.89	6.80	14.27	6.74	4.00	1.67
		1000	0.22	2.64	39.80	<b>24.47</b>	0.67	3.00	4.40	13.18	10.80	15.73	12.66	4.20	2.14
		3000	0.25	8.75	32.20	<b>25.08</b>	2.80	3.80	7.13	40.25	19.80	14.93	36.87	4.00	1.77

# MIP tests: single-covering

TABLE 3. Results of the MIP with MD-VIs, the mixing-set inequalities, and the quantile cuts for CKVLP instances.

Instances			MIP	MIP-MD		MIP-MS		MIP-Q	
$n-m$	Gen.	$ \mathcal{S} $	gap/time	cuts	gap/time	cuts	gap/time	cuts	gap/time
20-1	lit	100	<b>0.24s</b>	2.80	0.73s	24.80	0.48s	1.40	0.37s
		500	3.35s	2.40	<b>1.01s</b>	100.00	4.74s	1.00	2.49s
		1000	0.03%(4)	2.40	0.10%(4)	199.40	0.06%(4)	1.40	<b>272.61s</b>
		3000	3.18%(2)	2.20	<b>1.61%(3)</b>	501.00	4.43%(1)	2.40	4.75%(0)
	new	100	<b>0.13s</b>	2.00	0.28s	24.40	0.34s	1.20	0.28s
		500	10.01s	2.00	1.18s	104.00	2.07s	1.00	<b>0.91s</b>
		1000	20.57s	2.00	1.67s	151.60	1.65s	1.00	<b>1.22s</b>
		3000	<b>0.72%(4)</b>	2.00	0.97%(4)	466.00	0.97%(3)	1.20	1.25%(4)
80-1	lit	100	0.98s	3.40	<b>0.91s</b>	30.60	1.31s	1.40	0.97s
		500	0.58%(4)	3.00	<b>0.03%(4)</b>	213.60	0.61%(4)	1.60	0.58%(4)
		1000	3.01%(1)	3.40	<b>1.94%(1)</b>	252.80	2.62%(2)	2.20	3.29%(0)
		3000	0.90%(1)	2.20	<b>0.13%(4)</b>	271.20	2.14%(2)	1.20	0.52%(3)
	new	100	<b>0.59s</b>	2.20	<b>0.59s</b>	44.00	1.19s	2.20	0.68s
		500	12.42s	2.00	<b>4.95s</b>	161.60	22.85s	1.60	6.06s
		1000	0.01%(4)	2.00	<b>472.05s</b>	281.80	0.21%(4)	2.40	0.34%(4)
		3000	1.02%(3)	2.00	0.76%(4)	459.60	1.55%(2)	1.60	<b>0.53%(4)</b>



# MIP tests: multi- knapsack

TABLE 5. Results of the MIP with MD-VIs, the mixing-set inequalities, and the quantile cuts for the CCMKP instances, where the  $x$ -variables are continuous.

Instances			MIP-MD								
			MIP	ADM		C-wise		MIP-MS		MIP-Q	
$n-m$	Gen.	$ \mathcal{S} $	gap/time	cuts	gap/time	cuts	gap/time	cuts	gap/time	cuts	gap/time
20-10	lit	100	204.87s	17.80	172.64s	18.60	<b>114.96s</b>	202.80	202.04s	15.60	117.55s
		500	3.35%(0)	12.80	2.76%(0)	20.80	<b>2.20%(0)</b>	380.80	2.60%(0)	14.60	2.31%(0)
		1000	11.92%(0)	14.00	8.61%(0)	17.20	<b>6.99%(0)</b>	674.00	7.12%(0)	18.20	8.85%(0)
		3000	32.96%(0)	10.80	<b>9.40%(0)</b>	17.80	10.69%(0)	1538.00	11.27%(0)	17.60	14.21%(0)
	new	100	313.56s	20.60	<b>227.12s</b>	10.20	305.02s	166.20	368.67s	3.60	292.92s
		500	10.10%(0)	29.80	10.83%(0)	12.40	<b>9.47%(0)</b>	503.60	9.73%(0)	6.60	9.87%(0)
		1000	16.77%(0)	14.80	14.60%(0)	7.20	13.25%(0)	593.40	<b>13.17%(0)</b>	6.80	13.21%(0)
		3000	44.58%(0)	21.20	<b>19.84%(0)</b>	9.40	26.31%(0)	1552.40	29.96%(0)	7.40	24.86%(0)
40-30	lit	100	0.15%(2)	38.80	0.21%(3)	46.60	0.12%(3)	214.80	0.14%(4)	44.00	<b>0.08%(4)</b>
		500	10.35%(0)	20.00	8.49%(0)	41.20	8.67%(0)	430.00	6.95%(0)	35.60	<b>6.56%(0)</b>
		1000	24.77%(0)	46.40	<b>13.21%(0)</b>	46.80	14.44%(0)	801.00	15.64%(0)	52.60	14.38%(0)
		3000	35.04%(0)	19.80	<b>11.91%(0)</b>	38.80	15.44%(0)	1412.20	18.89%(0)	34.80	20.17%(0)
	new	100	1.47%(1)	62.80	<b>0.80%(1)</b>	13.60	1.26%(1)	185.00	1.31%(1)	6.60	0.94%(1)
		500	16.61%(0)	30.80	15.14%(0)	11.00	14.36%(0)	387.00	<b>13.36%(0)</b>	10.00	13.60%(0)
		1000	30.26%(0)	38.80	<b>19.85%(0)</b>	16.20	21.40%(0)	700.40	24.04%(0)	16.80	21.83%(0)
		3000	41.97%(0)	31.20	<b>18.78%(0)</b>	17.20	33.37%(0)	1694.80	36.25%(0)	16.00	30.81%(0)

# MIP tests: multi- knapsack

TABLE 4. Results of the MIP-tests with MD-VIs, the mixing-set inequalities, and the quantile cuts for the CCMKP instances, where the  $x$ -variables are binary.

Instances			MIP-MD								
			MIP	ADM		C-wise		MIP-MS		MIP-Q	
$n-m$	Gen.	$ S $	gap/time	cuts	gap/time	cuts	gap/time	cuts	gap/time	cuts	gap/time
20-10	lit	100	9.28s	17.80	7.47s	18.60	<b>4.31s</b>	202.80	6.34s	15.60	5.49s
		500	110.42s	12.80	46.33s	20.80	<b>21.53s</b>	380.80	60.40s	14.60	39.13s
		1000	642.60s	14.00	<b>111.55s</b>	17.20	120.40s	674.00	260.68s	18.20	164.71s
		3000	12.03%(1)	10.80	2.19%(4)	17.80	<b>696.69s</b>	1538.00	4.60%(4)	17.60	1500.19s
	new	100	14.43s	20.60	13.01s	10.20	8.40s	166.20	<b>8.01s</b>	3.60	13.11s
		500	91.75s	29.80	<b>62.04s</b>	12.40	70.99s	503.60	<b>72.04s</b>	6.60	91.51s
		1000	388.66s	14.80	<b>254.66s</b>	7.20	296.24s	593.40	356.02s	6.80	442.39s
		3000	41.46%(0)	21.20	<b>32.92%(0)</b>	9.40	33.83%(1)	1552.40	37.82%(0)	7.40	52.73%(0)
40-30	lit	100	10.17%(0)	38.80	9.16%(1)	46.60	<b>8.32%(1)</b>	214.80	11.73%(0)	44.00	12.05%(0)
		500	29.22%(0)	20.00	5.73%(2)	41.20	<b>1.81%(4)</b>	430.00	14.44%(0)	35.60	15.76%(0)
		1000	52.17%(0)	46.40	<b>20.85%(0)</b>	46.80	24.03%(0)	801.00	29.35%(0)	52.60	34.95%(0)
		3000	41.63%(0)	19.80	<b>15.07%(0)</b>	38.80	18.14%(0)	1412.20	24.66%(0)	34.80	29.13%(0)
	new	100	19.70%(0)	62.80	22.23%(0)	13.60	<b>18.79%(0)</b>	185.00	21.98%(0)	6.60	19.05%(0)
		500	39.24%(0)	30.80	<b>20.45%(0)</b>	11.00	32.42%(0)	387.00	27.21%(0)	10.00	34.72%(0)
		1000	48.89%(0)	38.80	<b>29.81%(0)</b>	16.20	55.00%(0)	700.40	39.56%(0)	16.80	52.94%(0)
		3000	49.05%(0)	31.20	<b>27.68%(0)</b>	17.20	48.18%(0)	1694.80	42.26%(0)	16.00	50.00%(0)