

LECTURE 8:

DATA STRUCTURES

Agenda

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Previously in 2110:

- Encapsulation
- Inheritance
- Polymorphism

Today:

- Data structures
- Abstract data types (ADTs) and interfaces
- ADT List: array lists, Linked lists (intro to A3)
- Abstract classes

EQUALITY CONTINUED...

What Is Equality?

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- **Referential equality:** two references point to the same object in memory

- Java: `o1 == o2`

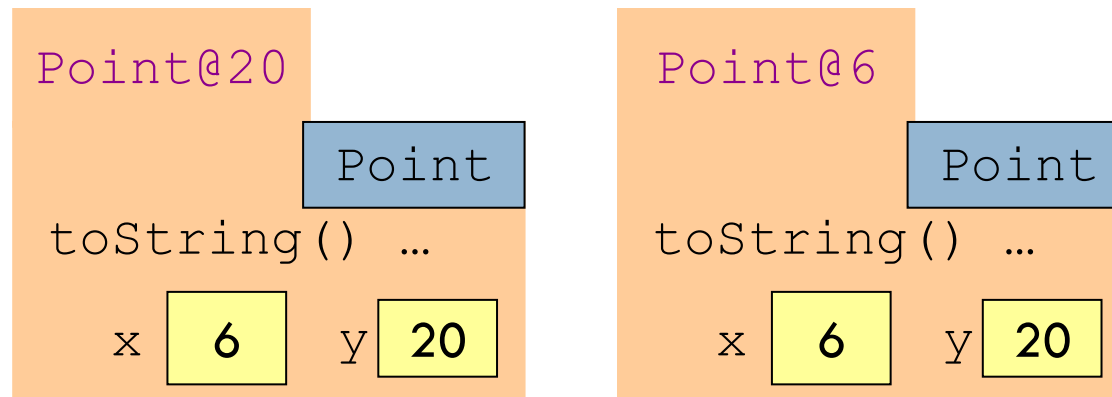
... and not this (except with primitives)

- **Equivalence:** two different objects are deemed “equivalent” according to the programmer

- Java: `o1.equals(o2)`

Almost always want this!

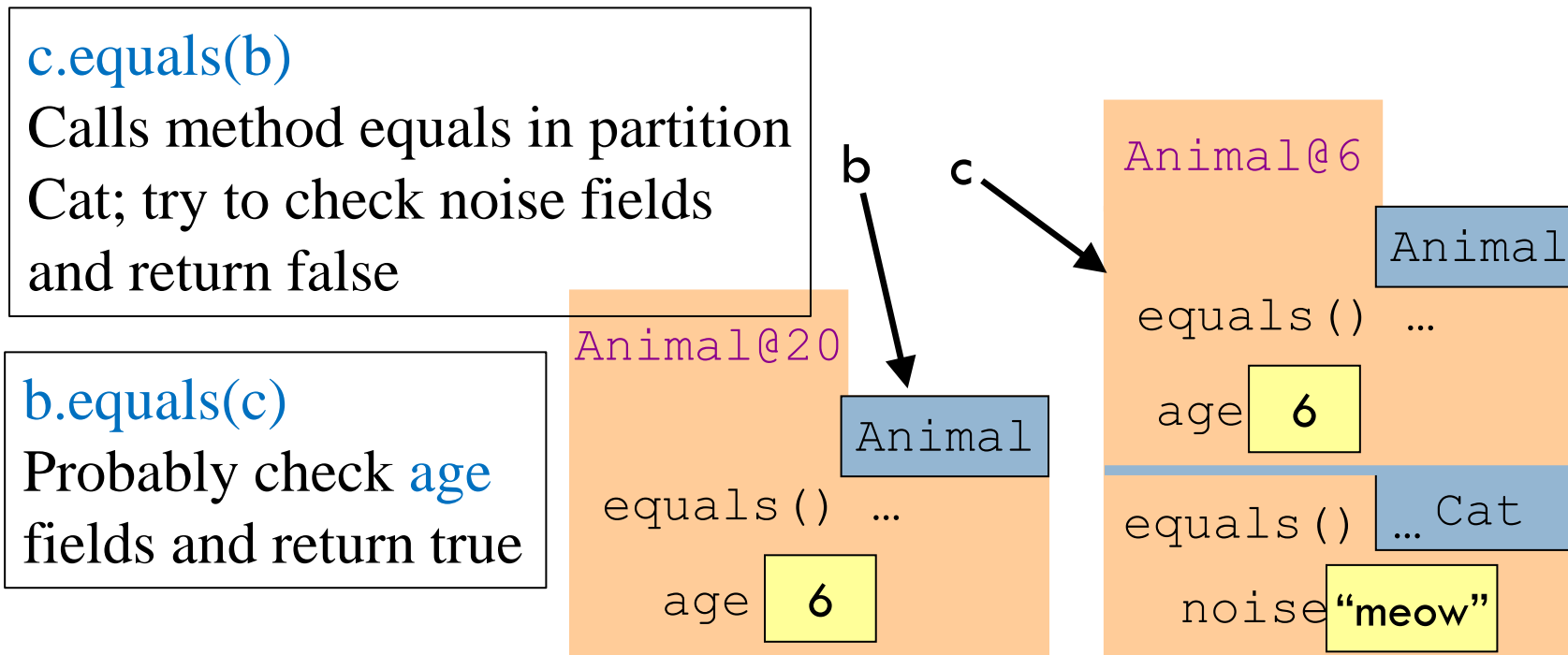
These objects could be considered equivalent, or equal



Overriding function equals

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We will override function equals in classes Animal and Cat.
(Note: it wouldn't make sense to think of objects **b** and **c** as being equal).



Overriding equals in class Animal

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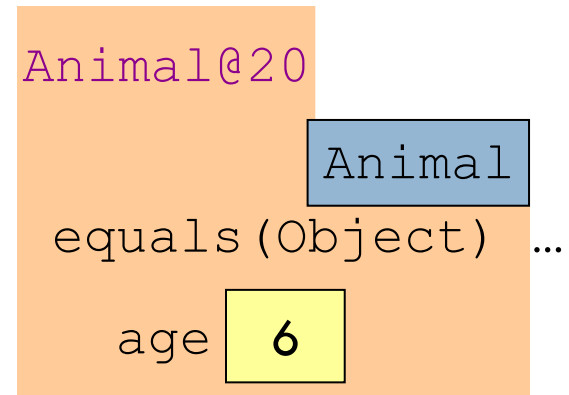
```
/** = "this object and ob are of the same class" and
    their ages are the same */

@Override
public boolean equals(Object ob) {
    if (ob == null || getClass() != ob.getClass()) return false;

    Animal oba= (Animal)ob;

    return age == oba.age;    // return age == ((Animal)ob).age;
}
```

By compile-time reference rule,
`ob.age` is illegal



Overriding equals in class Cat

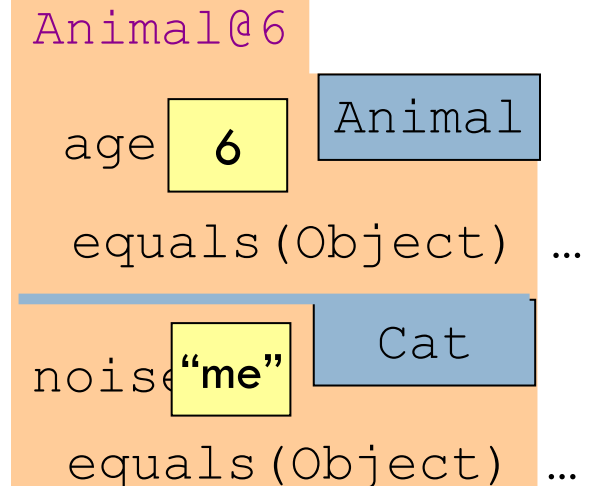
10

In class Cat

```
/** = "this object and ob are of the same class" and
    their ages and noises are the same */
public boolean equals(Object ob) {
    if (!super.equals(ob)) return false;
    Cat oba= (Cat) ob;
    return noise.equals(oba.noise);
}
```

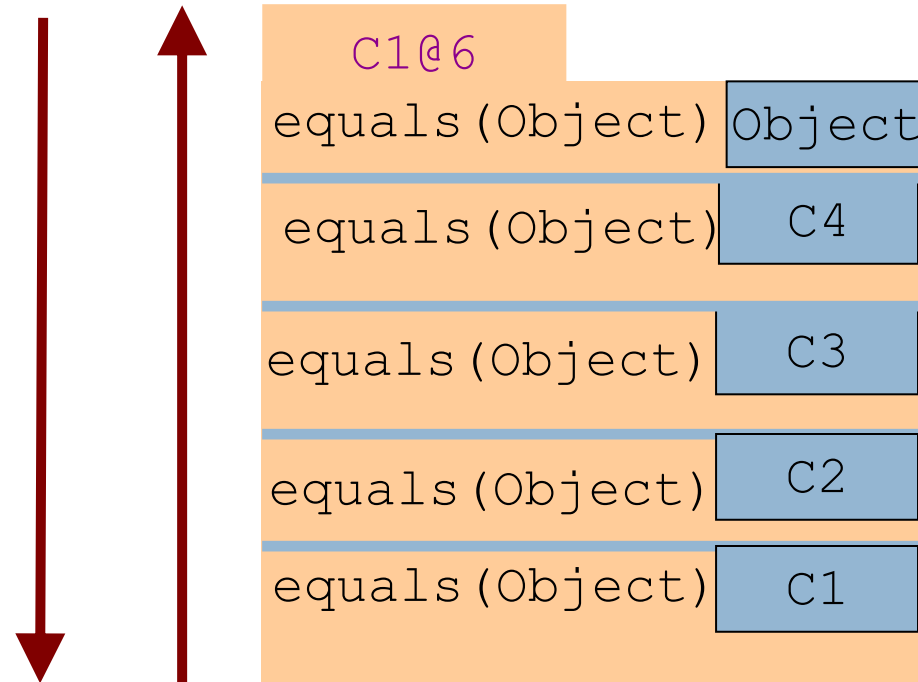
In class Animal

```
/** = "this object and ob are of the same class"
    and their ages are the same */
public boolean equals(Object ob)
```



OOP: process superclass partitions first

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Where have we seen this before?

LECTURE 8: DATA STRUCTURES

PART 1: WHY DATA STRUCTURES



Search Results

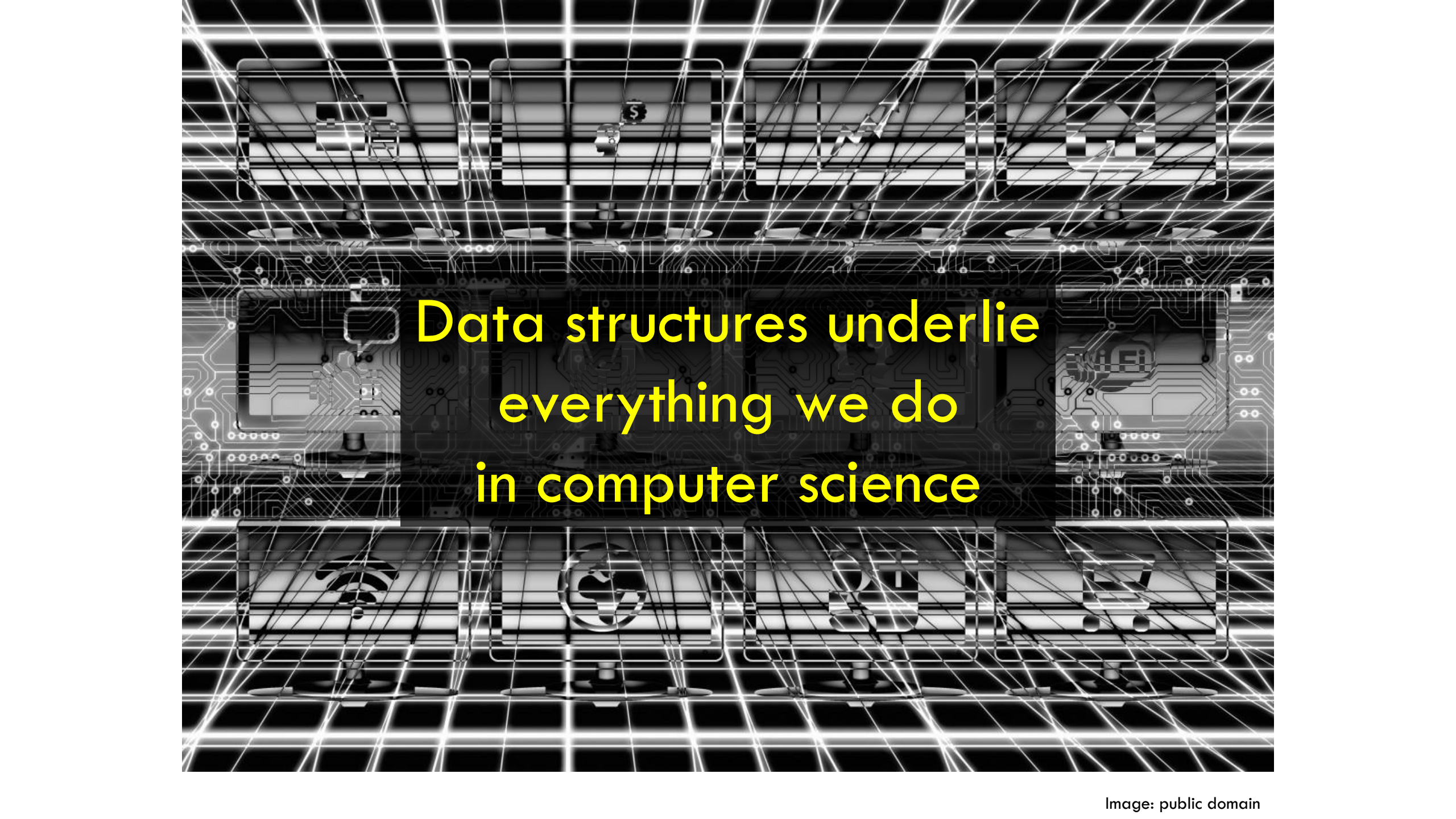
Showing 2 results.

Course descriptions provided by the [Courses of Study 2019-2020](#).

CS 2110

Object-Oriented Programming and Data Structures

Intermediate programming in a high-level language and introduction to computer science.



**Data structures underlie
everything we do
in computer science**

What is a Data Structure?

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- A format for storing data, and
- Algorithms for operations that manage the data

Format: what data,
and where to put it



Image: <https://novoresume.com/career-blog/resume-formats>

Operations: how to
transform data

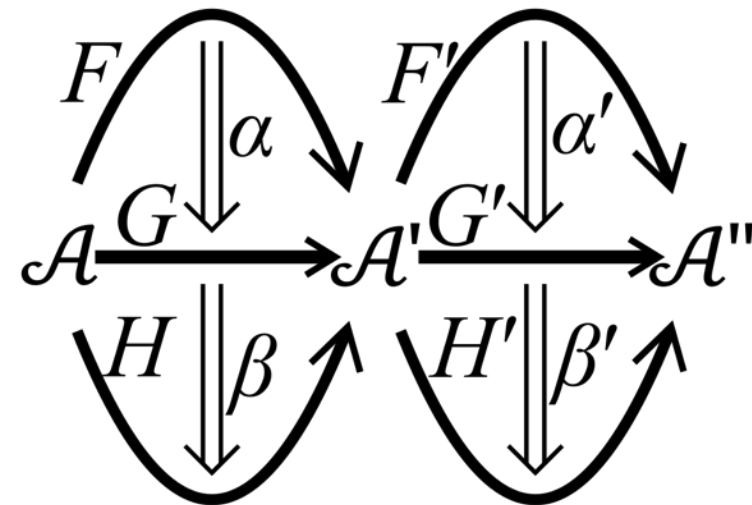


Image: public domain

Why Study Data Structures?

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- Data management is a primary use of **computers**
- **Libraries** provide them for our use
- Choices affect performance, usability of software
- Implementing them ourselves gives us:
 - ▣ excellent coding **practice**
 - ▣ deep **understanding** of the data structures

LECTURE 8: DATA STRUCTURES

PART 2: ADTS AND THE JAVA INTERFACE

Type: a set of values together with operations on them.

Java: type **int**, **double**, ...

class types, like **String**, **Animal**, **Account**

Abstract Data Type, or **ADT**

A type:

(1) defines a data structure

(2) It's abstract —no implementation is given

List:

The Cornell Daily Sun

161 Things Every Cornellian Should Do

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0. [Redacted]
1. Finally meet the dazzling Denice Cassaro
2. Go to the Cornell-Harvard men's hockey game and throw fish on the ice
3. Take off to NYC for Fall Break, being sure to post on Instagram about it at least twice
4. Sled down Libe Slope during a snowstorm
5. Take Hotel Administration 4300: Introduction to Wines
- ...
158. Tell a professor what you really think of their class
159. Attend a Sun meeting
160. Climb all 161 steps to the top of McGraw Tower

<https://cornellsun.com/161-things-every-cornellian-should-do/>

An ADT: List

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A **list**, or **sequence**, is an ordered collection of items

s1: (5, 8, -2, 8, 7)

s2: (“who”, “what”, “when”, “where”, “why”)

s3: ('h', 'o', 'n', 'e', 's', 't', 'y')

s4: (“word”, 4, “the”, 3, “and”, 3)

Notation for lists (*not* Java)

s[0], s[1], s[2], ..., s[s.length-1]

s[i..j]: stands for sublist consisting of s[i], s[i+1], ..., s[j]

s[i..]: stands for sublist consisting of s[i], s[i+1], ..., s[s.length-1]

s1 + s2: list catenation

An ADT: List —its operations

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In introducing ADTs and a Java feature that can be used to define an ADT, we limit our operations on a List to 3:

`prepend(e)`: insert item e at the beginning
`get(i)`: return item s[i]
`size()`: number of elements

Examples:

Prepend 5 to (4, 5, 7, 2): Result is (5, 4, 5, 7, 2)

`get(3)` from the list (5, 4, 5, 7, 2) is 7

`size()` of (5, 4, 5, 7, 2) is 5

Interfaces

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- An **interface** describes the operations of an ADT
- The notion of “interface” can be found in many languages
- Java actually provides a syntax for it ...

Java Interface

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```
public interface List<T> {  
    void prepend(T val);  
    T get(int i);  
    int size();  
}
```

No method bodies.

Declares what operations must be provided, not how they are to be implemented.

No constructors. Cannot instantiate —not a class.

Description of values and method specifications given in Javadoc comments

Methods are automatically public. No need to put in access modifier

Demo

Java Interface

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```
public interface List<T> {  
    void prepend(T val);  
    T get(int i);  
    int size();  
}
```

How do use this interface? What good is it?

In part 3 of this lecture, we will see that we can write several different implementations of the ADT defined by interface `List`. Then, we can choose one of them to use in any program, and switch to a different one if it performs better in our program.

LECTURE 8: DATA STRUCTURES

PART 3: TWO IMPLEMENTATIONS OF ADT LIST

An implementation of List: ArrayList

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```
public class ArrayList<T>                                implements List<T>
    /** The items of the list are in values[0..size-1] */
    T[] values;
    /** number of items in the list */
    int size;

    /** Constructor: An empty list with capacity c. */
    public ArrayList(int c) {
        values= (T[]) new Object[c];
    }
```

... more to come ...

Demo in Eclipse

Implements clause: Requires that the class implements all methods declared in interface List

An implementation of List: LinkedList

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We introduce the notion of a linked list, as an implementation of abstract data type [List](#).

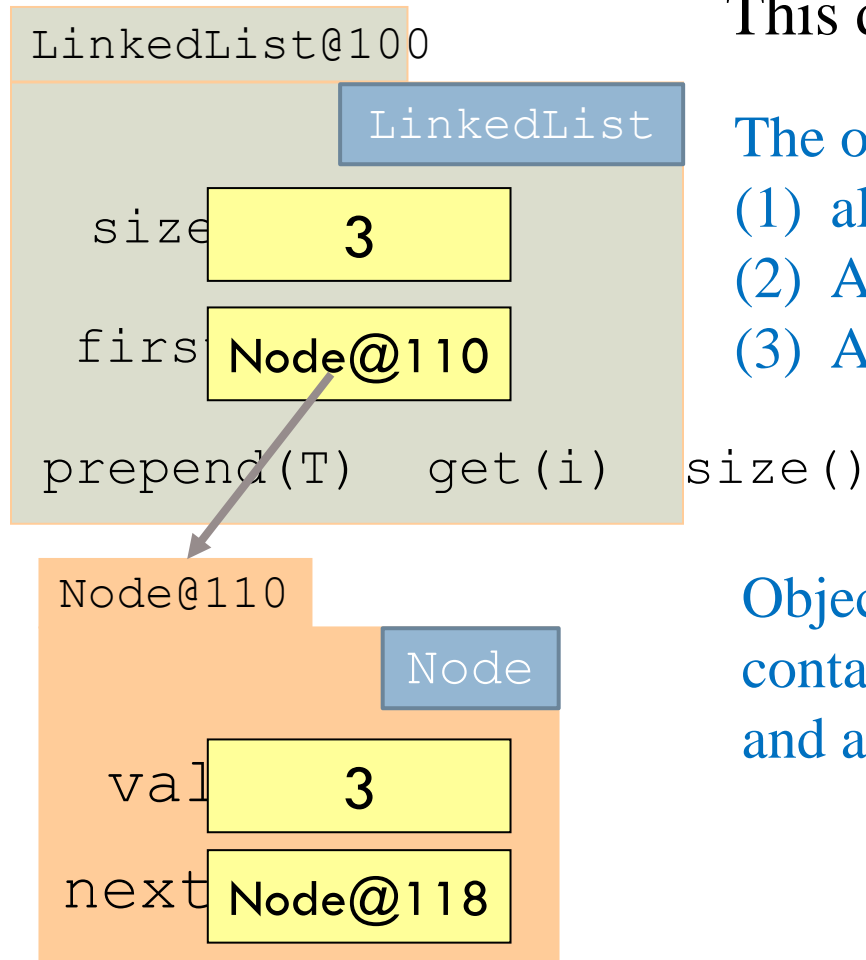
[Assignment A4](#) is all about linked lists. For A4, you are expected to read about linked lists in [JavaHyperText](#), entry [linked list](#)

The emphasis is on learning about a data structure by reading about it. This short intro will help.

You will learn about a new Java feature: the [inner class](#).

Singly Linked List: How To Store

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This data structure contains the list (3, 5, 2)

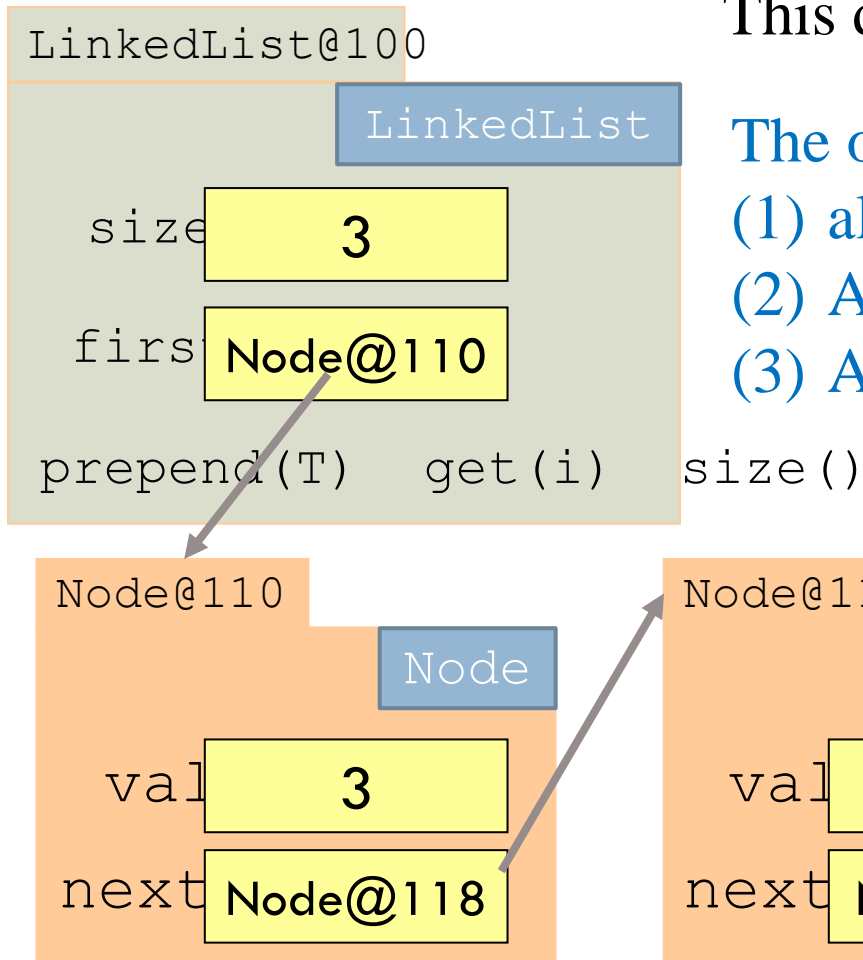
The object of class `LinkedList` contains

- (1) all three methods of interface/ADT `List`
- (2) A field that contains the size of the list
- (3) A pointer to an object of class `Node`:

Object `Node@110`, of class `Node`, contains the first value of the list, 3, and a pointer to another `Node` object.

Singly Linked List: How To Store

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This data structure contains the list (3, 5, 2)

The object of type `LinkedList` contains

- (1) all three methods of linked lists,
- (2) A field that contains the size of the list
- (3) A pointer to an object of class `Node`:

LinkedList

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```
public class LinkedList<T> {  
    private class Node {  
        T val;  
        Node next;  
        Node(T val, Node next) {  
            this.val= val;    this.next= next;  
        }  
    }  
  
    private Node first;  
    private int size;  
    ...  
}
```

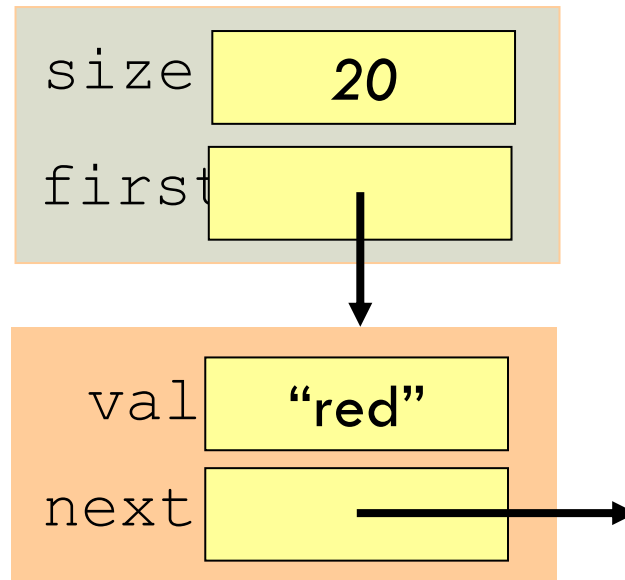
Inner class: Node nested inside
LinkedList. No reason to
expose to outside world.

LinkedList

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```
public class LinkedList<T> {  
    ...  
    public void prepend(T val) {  
        Node n= new Node(val, first);  
        first= n;  
        size++ ;  
    }  
}
```

prepend("big")

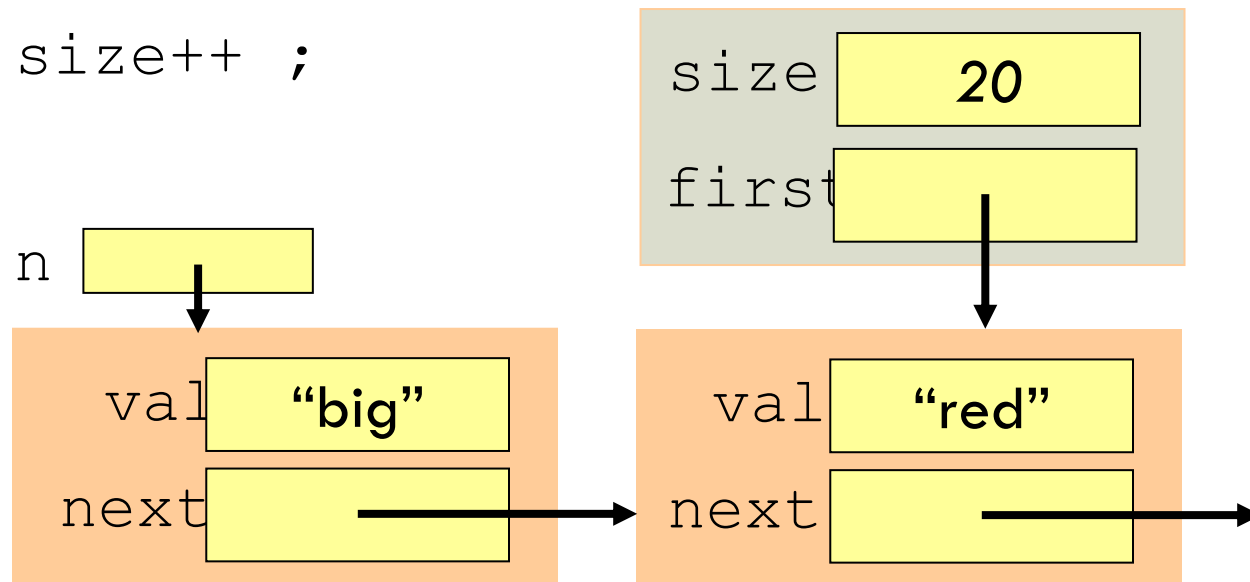


LinkedList

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```
public class LinkedList<T> {  
    ...  
    public void prepend(T val) {  
        Node n= new Node(val, first);  
        first= n;  
        size++ ;  
    }
```

```
prepend("big")
```

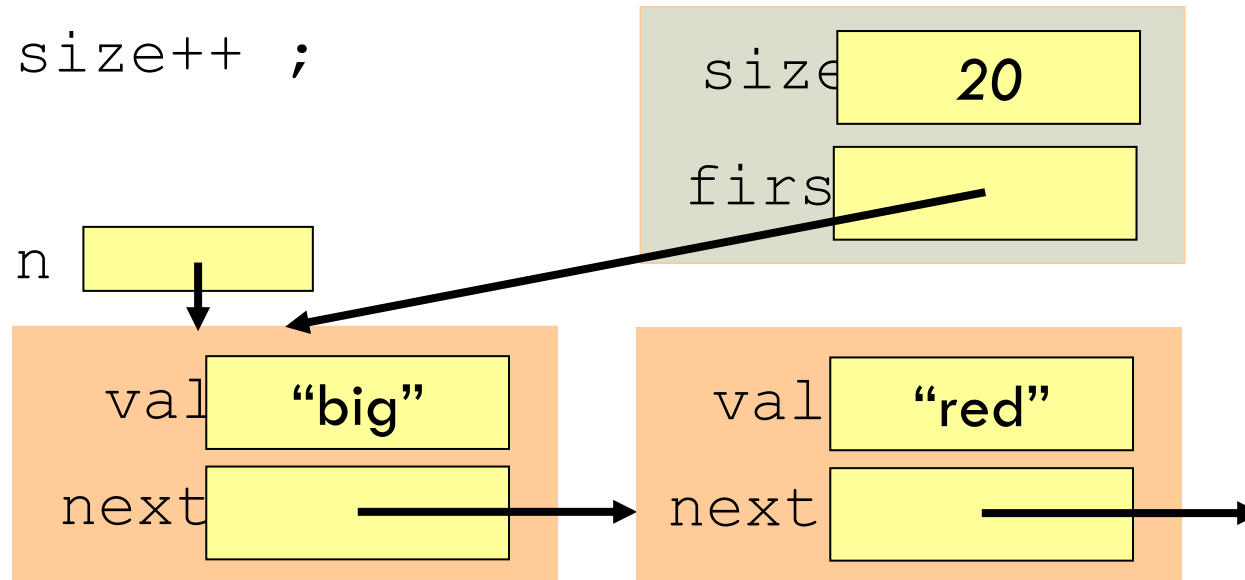


LinkedList

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```
public class LinkedList<T> {  
    ...  
    public void prepend(T val) {  
        Node n= new Node(val, first);  
        first= n;  
        size++ ;  
    }
```

prepend("big")

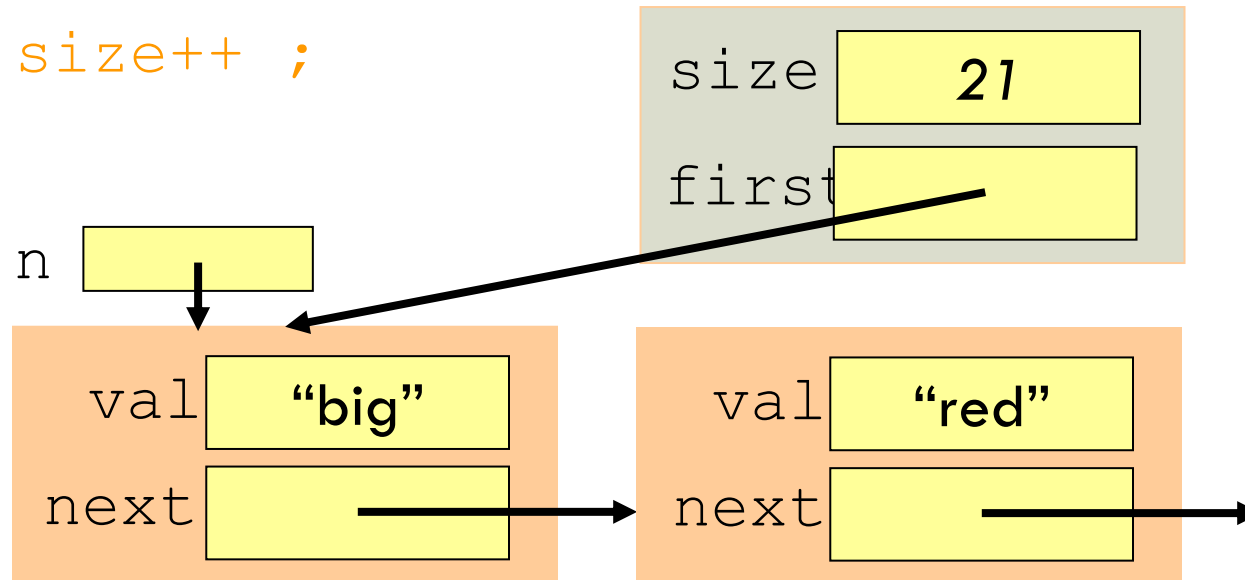


LinkedList

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```
public class LinkedList<T> {  
    ...  
    public void prepend(T val) {  
        Node n= new Node(val, first);  
        first= n;  
        size++ ;  
    }  
}
```

prepend("big")



LinkedList

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```
public class LinkedList<T> {  
    ...  
    public void prepend(T val) {  
        Node n= new Node(val, first);  
        first= n;  
        size++ ;  
    }  
}
```

Cost: inexpensive! Create one Node object, do three assignments. Doesn't matter how many items are already in list. We call it a "constant time" operation.

LinkedList: Method get(i)

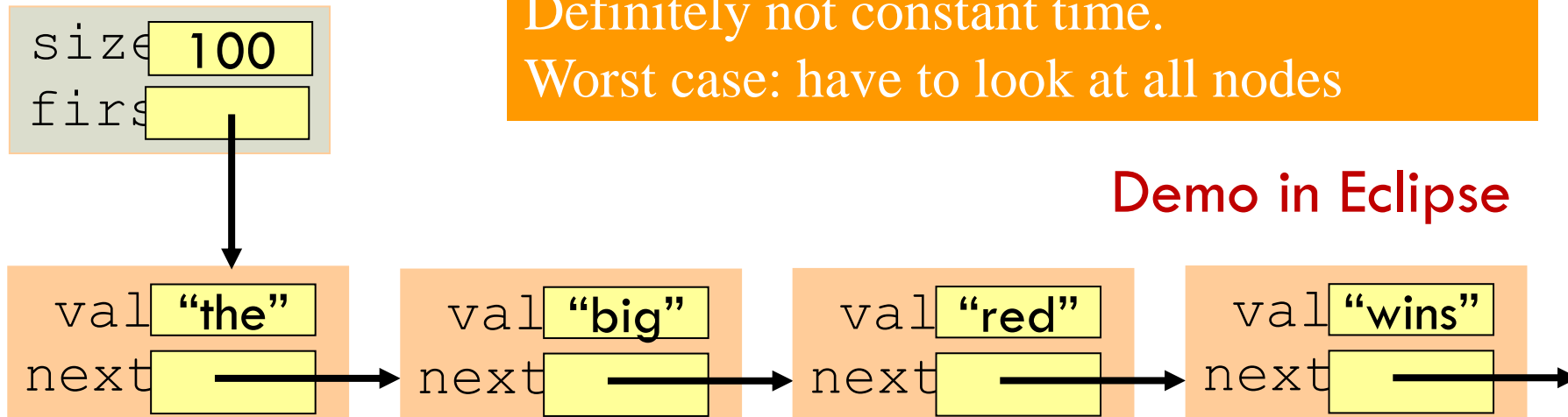
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The call `get(3)` has to search the list from the beginning to find item number 3 and return the value “wins”.

`get(3)`

The call `get(99)` has to look at all nodes to get to the last one!

The call `get(i)` takes time proportional to i .
It's a linear operation in i .
Definitely not constant time.
Worst case: have to look at all nodes



LinkedList vs. ArrayList

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Efficiency of operations

	LinkedList	ArrayList
prepend	constant time	linear time
get	linear time	constant time

constant time: inexpensive; independent of number of list elements

linear time: relatively expensive; proportional to number of
list elements

Next part of lecture: Illustrate how the use of interface **List** helps us easily choose one implementation and perhaps switch to the other later on.

LECTURE 8: DATA STRUCTURES

PART 4: INTERFACES AND SUBTYPING

Interface vs. Implementation

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Interface vs. Implementation

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Images:

<http://www.globalvendinggroup.com/products/National-637-Coffee-Machine.html>,

[https://en.wikipedia.org/wiki/Coffee_vending_machine#/media/File:Mechanisms inside a coffee vending machine.jpg](https://en.wikipedia.org/wiki/Coffee_vending_machine#/media/File:Mechanisms_inside_a_coffee_vending_machine.jpg)

Interface vs. Implementation

Interface: the operations of an ADT

- What you see in documentation web pages
- Method names and specifications
- Abstract from details: **what** to do, not **how** to do it
- Java syntax: **interface**

Implementation: the code for a data structure

- What you see in source code
- Fields and method bodies
- Provide the details: **how** to do operation
- Java syntax: **class**

Could be many implementations of an interface

e.g. List: `ArrayList`, `LinkedList`

LinkedList vs. ArrayList

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Both support the same operations: `prepend`, `get`

Always an engineering tradeoff: choose efficient data structure for operations of concern

But, some **clients** won't care about the different efficiency as long as they get the operations they want...

	LinkedList	ArrayList
<code>prepend</code>	constant time	linear time
<code>get</code>	linear time	constant time

Interfaces and Subtyping

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Recall: if S is a subtype of T , then anywhere a T is expected, an S can be used

Recall: if SC extends C , then SC is a subtype of C

New: if C implements I , then C is a subtype of I

Examples:

```
□ List<T> lst= new LinkedList<>();
```

```
□ void m(List<T> lst) { ... }  
  m(new ArrayList<>());
```

Compile-time reference rule:
Only operations allowed on `lst`:
prepend, get, size

Example

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```
int size= 50000;
```

```
List<Integer> al= new ArrayList<>(size);
```

```
long start= System.currentTimeMillis();
```

```
for (int k= 0; k < size; k= k + 1)  al.prepend(k);
```

```
long time= System.currentTimeMillis() - start;
```

```
System.out.println("Time for " + size + " prepends: " +  
                    time + " milliseconds");
```

Change implementation?
Change new-expression

Create ArrayList
object, store it in
a List variable

Current time, in
milliseconds*

Elapsed time to
prepend size values

* Time since midnight, January 1, 1970 UTC

Demo in Eclipse

Java List, ArrayList, LinkedList

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Java has in package `java.util`

`interface List<T>`

`class ArrayList<T> implements List<T>`

`class LinkedList<T> implements List<T>`

`class Stack<T> implements List<T>`

`class Vector<T> implements List<T>`

and more

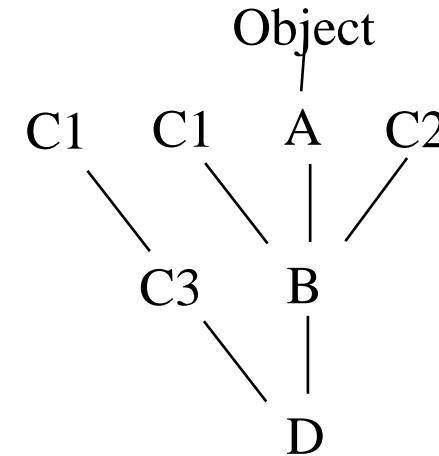
Your task: Learn about casting

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Homework, to complete your introduction to interfaces.

JavaHyperText, upper navigation bar,
item [Abstract classes and interfaces](#)

1. Read about [Three other components of an interface](#)
2. Study [3.75-minute video on Casting to learn about casting and drawing objects of classes that implement interfaces](#)



www.cs.cornell.edu/courses/JavaAndDS/abstractInterface/01ai.html

LECTURE 8: DATA STRUCTURES

PART 5: ABSTRACT CLASSES AND METHODS

Abstract Classes and Methods

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Make a class abstract so that it cannot be instantiated —objects of the class cannot be created.

Make a method in an abstract class (or interface abstract so it must eventually be overridden —defined in a subclass.

Watch the first, short, JHT tutorial on Abstract classes and interfaces:

<https://www.cs.cornell.edu/courses/JavaAndDS/abstractInterface/01ai.html>

Your Turn: Read in JavaHyperText

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- ❑ data structure
- ❑ list, linked list, doubly-linked list
- ❑ abstract data type
- ❑ interface, implements
- ❑ abstract method, abstract class