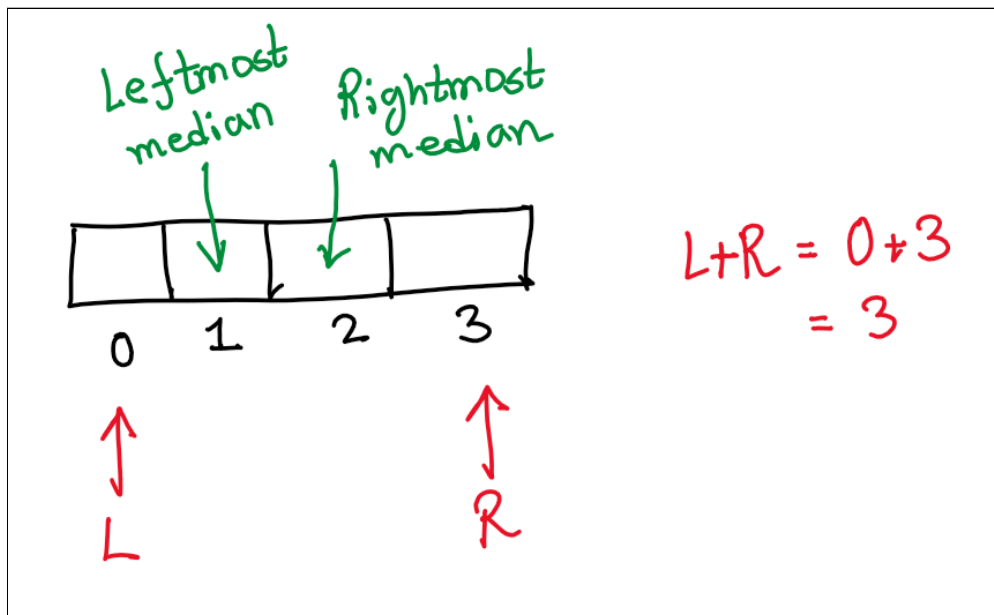


Before jumping into variation of binary search, let's clear some basic concepts at first:

1.  $M = (L + R) / 2$  is risky. It can create overflow issues.
2.  $M = L + (R - L) / 2$  is safe. It is computationally the same as no 1, but it has no overflow issue.
3. If  $(L+R)$  is ODD, then there will be an even number of elements in total. So, there will be two median values. Let's look at the following figure for better understanding:



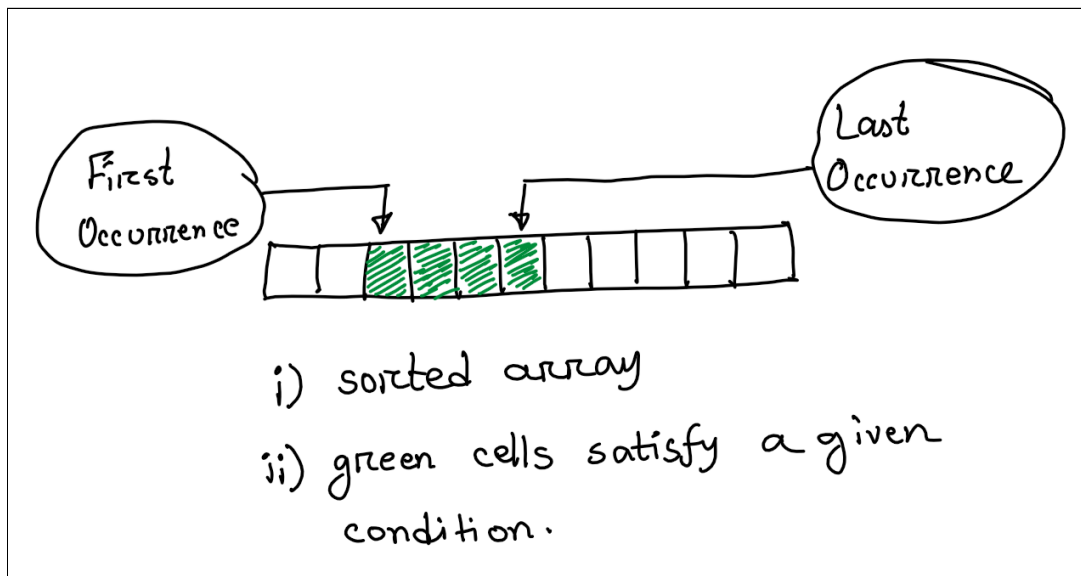
If you divide  $(L+R)$  by 2, you will always get the leftmost median in such scenarios. But what if we want to get the rightmost median? :D

The answer is :  $M = (L+R+1)/2$  :D

## Types of Binary Search:

1. Find the target element from an array.
2. Find the first occurrence of a number that satisfies a given condition.
3. Find the last occurrence of a number that satisfies a given condition.

To understand type 2 and 3, refer to the following figure:



## Sample Code Template

1. Find the target element from an array:

```
bool doesExist(vector<int>&vec, int target)
{
    int L = 0;
    int R = vec.size() - 1;

    while(L <= R)
    {
        int M = L + (R - L) / 2;
        if(vec[M] == target) return 1;
        if(vec[M] < target)
        {
            L = M + 1;
        }
        else
        {
            R = M - 1;
        }
    }
    return 0;
}
```

2. Find the first occurrence of a number that satisfies a given condition (considering at least one number exists in the array that satisfies the given condition) :

```
int first_occurence(vector<int>&vec, int target)
{
    int L = 0;
    int R = vec.size() - 1;

    while(L < R)
    {
        int M = L + (R - L) / 2;
        if(vec[M] < target)
        {
            L = M + 1;
        }
        else
        {
            R = M;
        }
    }
    return L;
}
```

3. Find the last occurrence of a number that satisfies a given condition (considering at least one number exists in the array that satisfies the given condition) :

```
int last_occurence(vector<int>&vec, int target)
{
    int L = 0;
    int R = vec.size() - 1;

    while(L < R)
    {
        int M = L + (R - L + 1) / 2;
        /**
         * this is equivalent to M = (L + R + 1) / 2
         * we are using this because we want to go as
         * right as we can. :D
         */
        if(vec[M] > target)
        {
            R = M - 1;
        }
        else
        {
            L = M;
        }
    }
    return R;
}
```

**Note:**

1. lower\_bound = first occurrence of a number smaller or equal to a given number.
2. upper\_bound = first occurrence of a number smaller or equal to a given (number + 1)
3. Hence, lower\_bound(value + 1) == upper\_bound(value)
4. type 2 and 3 do not entirely reflect the idea to find the lower and upper bound. But one can easily change the code of type 2 , to find out lower\_bound from an array. If you have the code for lower\_bound, you don't need to write another function for upper\_bound. :) . Just change R = vec.size() in type 2 code and you're done.