

Writeup

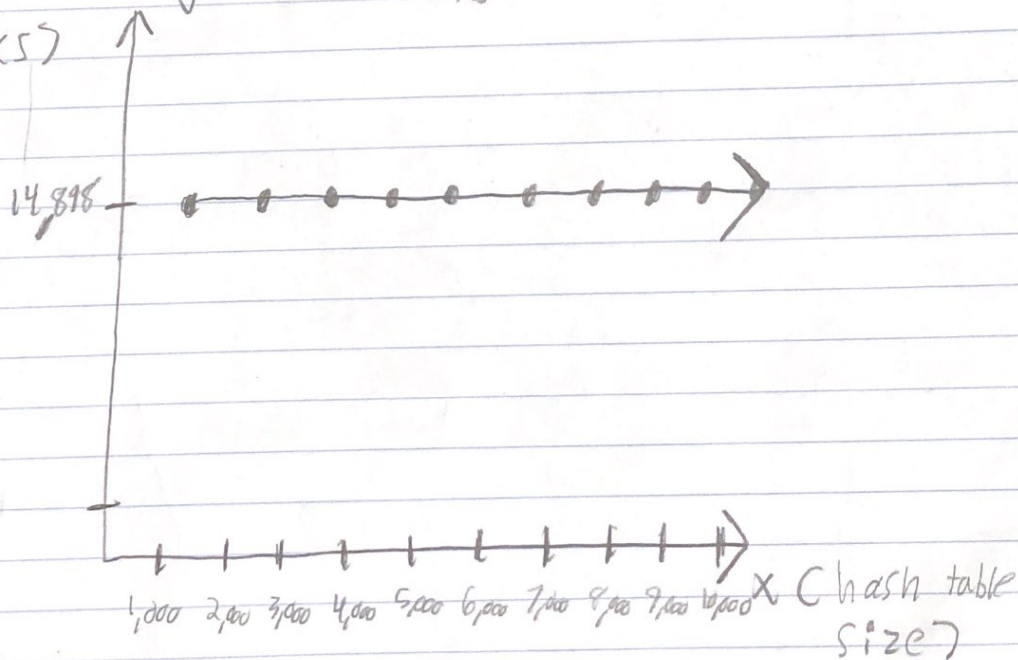
To analyze the results of assignment 7, the variables "seeks" & "links" were created to track the amount of times we called a linked list lookup & the amount of node links traversed,

The following graphs display the total number of seeks and the average seek length as we vary the hash table and bloom filter size. Numbers are based on the same file "banhammer.c"

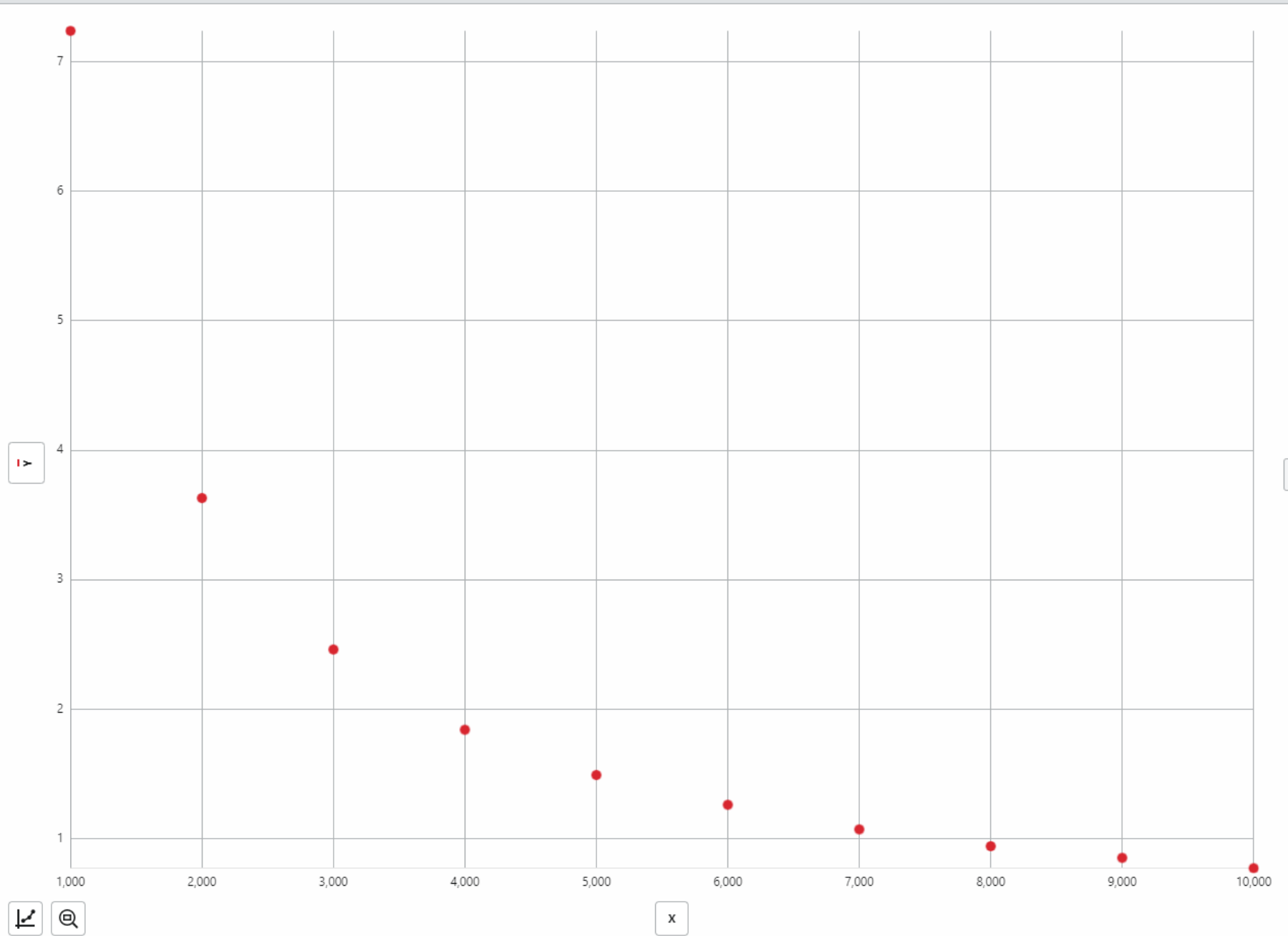
NOTE:

$$\text{avg. Seek length} = \frac{\text{links}}{\text{seeks}}$$

y(seeks)
seeks



"As we vary hash table size the amount of seeks does not change. This is likely due to the fact that we will always need to look up the same amount of words that we suspect are bad."



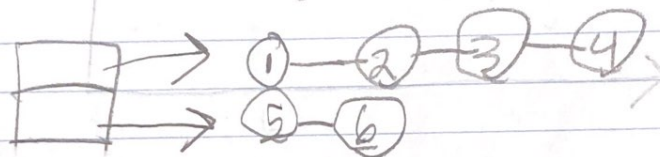
	Data Set 1		...	
	X	...	Y	...
1	1000		7.24	
2	2000		3.63	
3	3000		2.46	
4	4000		1.84	
5	5000		1.49	
6	6000		1.26	
7	7000		1.07	
8	8000		0.94	
9	9000		0.85	
10	10000		0.77	
11				

$$\text{avg. seek length} = \frac{\text{links}}{\text{seeks}}$$

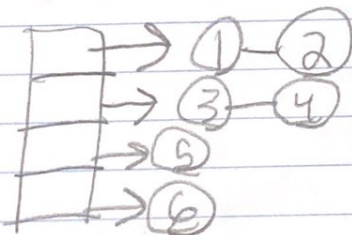
Avg

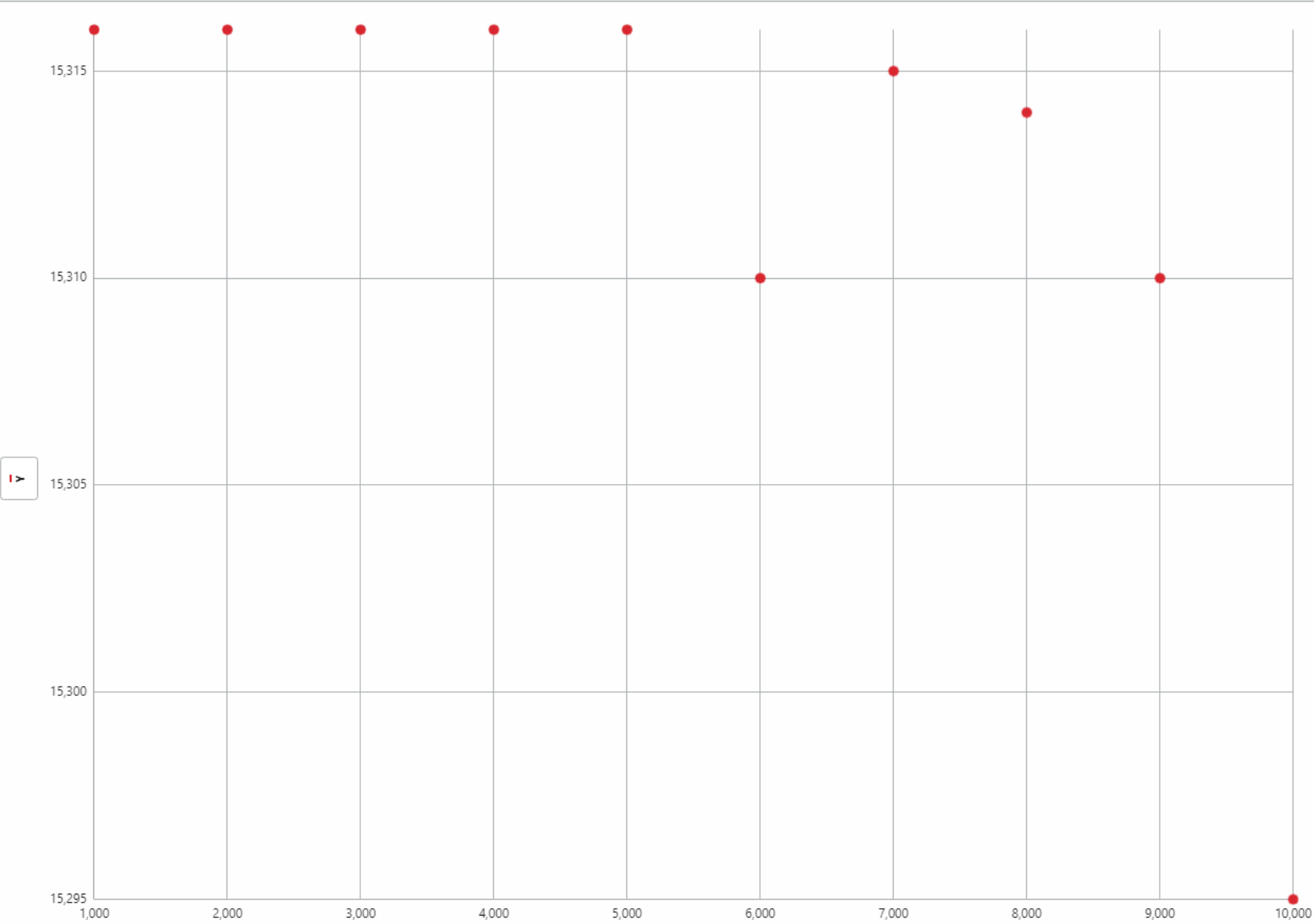
• As the Hash table size is increased by 1,000, we can clearly see a decrease in the average seek length. This is because the more space we have in our hash table the less "cramped" it becomes for linked lists. In other words, more hash table space means ~~more~~ less nodes to traverse per linked list.

Small Hash Table



Large Hash Table





	Data Set 1		...	
	X	...	Y	...
1	1000		15316	
2	2000		15316	
3	3000		15316	
4	4000		15316	
5	5000		15316	
6	6000		15310	
7	7000		15315	
8	8000		15314	
9	9000		15310	
10	10000		15295	
11				



- As we increase our bloom filter size by 1,000 we can clearly see that the amount of seeks remain at a constant (15216) until we reach a bloom filter size of 6,000. Then, our seeks noticeably decrease which is likely due to our system having less false-positives and thus decreasing the amount of times we must verify a word of being bad

Small Bloom Filter

1	—
1	—
1	—

"Python"
"Java"

- "Python" and "Java" are forced to share the same bits bc filter is too small... we aren't sure which one is the real positive transgression.

Large Bloom Filter

1	—
1	—
1	—
0	—
0	—
0	—

"Python"
"Java"

- Now we have more space & have more certainty as to which word might be a transgression

y

x



	Data Set 1		...	
	X	...	Y	...
1	1000		0.8	
2	2000		0.8	
3	3000		0.8	
4	4000		0.8	
5	5000		0.8	
6	6000		0.8	
7	7000		0.8	
8	8000		0.8	
9	9000		0.8	
10	10000		0.8	
11				

- As we increase the blank filter size, there is a minimal decrease in our avg. seek length. This makes sense because our seeks or the amount of times we need to call `lookup()` is decreased.

$$\text{avg. seek length} = \frac{\text{links}}{\text{seeks}}$$

Now, when we compare our results w/ the move-to-front option enabled compared to when it is not we ~~can~~ will notice a difference.

Move To Front Enabled

Seeks: 14898
Avg. Seek Length: 0.740099

Move-To Front Disabled

Seeks: 14898
Avg. Seek Length: 0.771781

- It is clear that the avg. seek length is the only difference w/ the disabled version being greater than when move to front was enabled.
- This means that the amount of links we are traversing is greater which makes sense because humans tend to be repetitive in their speech & some words are used more often i.e. "the" or "something".