

Ve-LAB Part 2) I In the worst case, the shell sout is executively an insertion sort. It approaches O(n2) complexity, The awage complexity depends on the gop 5:205, because it changes the amount of companisons you have to do, Pre-Lab Part 31 I Quick sout is not doomed due to its worst case being of (na), this is because Quicksout USES midpoints or pivots which eliminates the need to compare we numbers left of pivat and Sert other Annay ardered way, It is vare for it to be ocurs blood for you would have to pick a minduak pilot out of all the numbers, Tre-Lab Party 1.7 I will intralize a counter for each Sorting algor, than that will be incremented for each comparison Likewise, a for counter will also be initialized and increment each time the algorithm swaps two numbers.

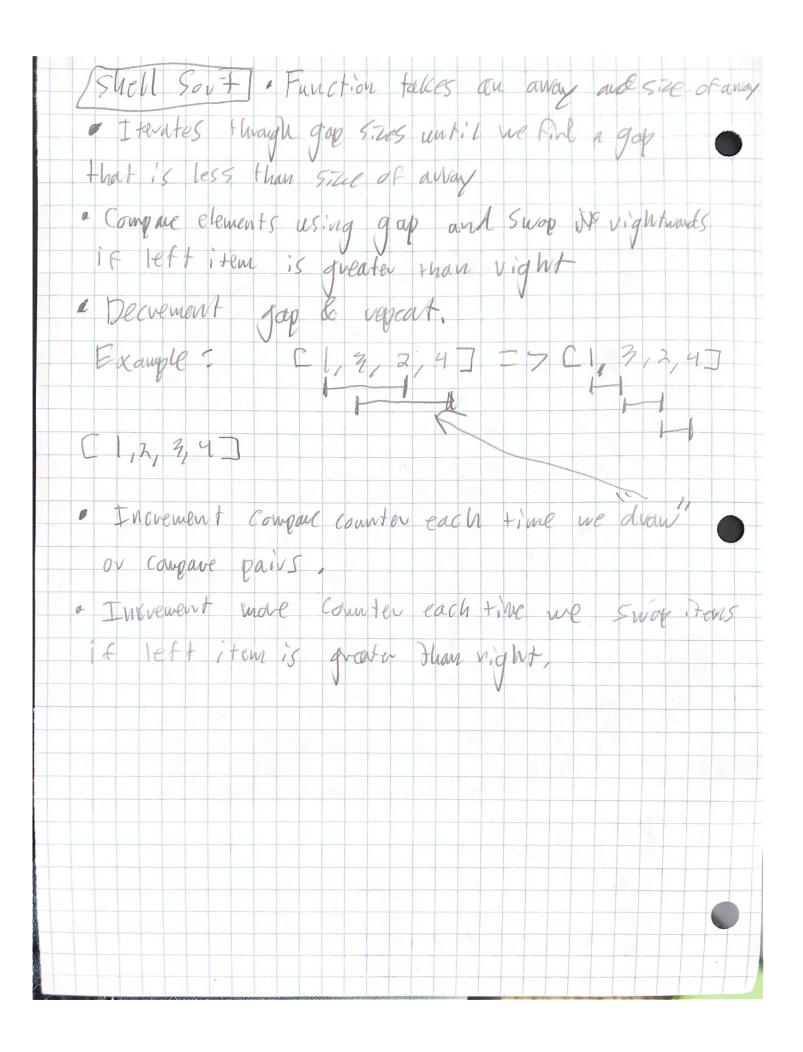
Struct Stack last in =>first out System · Use Stack. In as a interface for actual Structure · Actual definition will be in Stack. , Will have following components that we track top Cto track where to pop / push ? condity Cto track if we can push? i tems Cto stage hold actual values) Coente Function to instantiate a stack w/ given & Capacir, "Each Hem requires space of 4 bits, cuse called) to is zero " Capacity is capacity Stack Delete Function to dake the delete Stacke and make it null. · Use free () to we a pointer to a pointer (fred(\$5)) tack Empty/ Returns true it top is at index Fall it otherwise Stack-S.Zel - Returns index of top

Stack Push . Place given value at index of top Increment top Stack-Pop) Decrement top & Place Value of number of House array of items in given & and decrement top Stack-print Print I terate through items and grint Hall Stack is NULL print Stack's Structure? Has components - Head C to track where to place next item? C to track which item is to dequeve? - 5720 C tracks amount of items) - Capacity & tracks when eve me Full - items (Away to hold Items) Queue Coeate (Instantiate each item ne ubits si Callor head, tail, sile start at zero Threwe delete 1. To delete queue and merer won his · Use FreeC) w/ 28 pueul

+ Size = of them return queue empty frul otherwise, veturn False quere Full I + size = Capacity, veturn frue 1 Otherwise vetorn Palo anerue 5,20 / vetura 5,20 Enquewe I F full vetur false · Elce, increment size and place given value at The fail index in item away · Inquewent tail dequewel : If queue is empty veturn false ELSP decrement size place value at wead index in item away in given x and increment head queue print i terate though all items and print each are

ASSIGNMENT 3 VEST clopase Assignment 3 temants a gertable Shell and quick fort (we queuelstack), we will generate aways of varying 5:205 and have the USEV input commands to control Size print size, Sorting algorithm Output will print statistics Layout/ Structure Bubble Sout Bubble Sout will be a function in an away and a size of that away, 1 It will iterate a take pairs within away.

I To one item is queater more it vightwards * This eventually quarantees the greatect element to be the wast eightwards coxclusting the previous numbers we souted) 1 3, 2, FD souted, do not touch =>C1,2,3 conster to track comparisons, Each time we find item greater, we snow and increment move counter



Quick Sort. Function takes away and aways s. 20 · We use a pivot point; it left item of pilot is greater we more it vight of givet. It item vight of pivot is loss than givet we more it left of pivot. . Now we don't need to compare items left of givot we items vight of pivot, Halving comparison Repeat w/ a pivot within the pivot. Example [64, 34, 25, 12, 22, 11, 90] Cas [64, 34, 25, 22, 90] Left Pivot Right =>MERGE => C11,12,64,34,25,22,90] => (32) [25] [64,34,90] Left Pivot Right MERGE => (11,12,22,29,64-94,90) MERGE -> (34) (647) (907) -> (1/4) ax, 2534, 64, 907) a Increment comparisons each comparison w/ prot . Increment mark each surp to the text or vight of prot.