

ESTRUCTURA DE DATOS 1 Código ST0245

Laboratory practice No. 1 Recursion

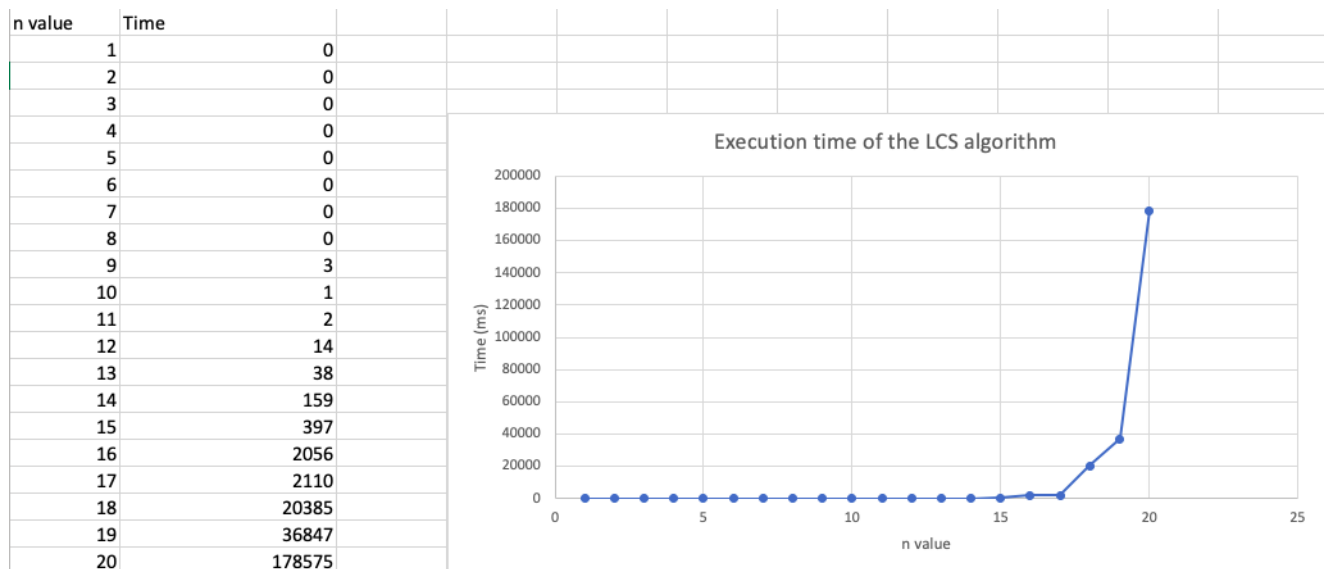
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3) Practice for final project defense presentation

3.1 $T(n+m) = c_3 T((m+n)-1) + T((m+n)-1)$
 $T(n+m) = c_3 T((2^m+n)-1) + c_1(2^{(m+n-1)})$
 $T(n+m) = 2^{(m+n)}$
 $T(n+m) = T(p)$
 $T(p) = O(2^p)$ complexity

3.2 With these results we can see how data grow exponentially, showing us the worst case of their performance.



3.3 It wouldn't be an optimal and recommended option because being an exponential function means that its complexity is going to double in every single value. For that reason, finding the DNA subsequence algorithm would take too much time because each DNA sequence has approximately 300,000 strings.

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3.4 How groupSum5 works?

It receives an array and execute the condition to verify there is a number multiple of 5. If the number is a multiple jump into the next condition that verify there is a number 1, then add the multiple of 5 and omit the 1. If the number is not a multiple of 5, rejects the number. At the end we use a two recursive call to continue with the next number if the conditions are correct

3.5 CODINGBAT

RECURSION-1

- **Bunny Ears**

```
public int bunnyEars(int bunnies) {

    if(bunnies==0){                //c1
        return bunnies;            //c2
    }else{
        return 2+bunnyEars(bunnies-1); // t(n)=c2 + t(n-1)
    }
}
```

Model:

$$T(n) = \begin{cases} 1 & \text{if } n == 0 \\ 2 + t(n-1) & \text{if } n > 0 \end{cases}$$

Recurrence equation solution :

$$t(n) = c_2 * n + c_1$$

Calculate complexity:

$$O(C_2 * n + c_1)$$

$$O(C_2 * n) \quad // \text{ By addition rule}$$

$$O(n) \quad // \text{ By multiplication rule}$$

- **Fibonacci**

```
public int fibonacci(int n) {                //
    if(n<=1)                                //c1
        return n;                            //c2
    return fibonacci(n-1) + fibonacci(n-2); //t(n)= t(n-1)+t(n-2)
}
```

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Model :

$$T(n) = \begin{cases} c_1 + c_2 & \text{if } n \leq 1 \\ T(n-1) + t(n-2) & \text{if } n \geq 2 \end{cases}$$

Recurrence equation solution :

$$t(n) = c_1 * 2^n + c_2$$

Calculate complexity:

$$O(C_1 * 2^n + c_2)$$

$$O(C_1 * 2^n) \quad // \text{ By "Laboratory practice No. 1: Recursion" point 4.4.1}$$

$$O(2^n)$$

- **BunnyEars2**

```
public int bunnyEars2(int bunnies) {
    if(bunnies==0)
        return bunnies;                // c1
    if (bunnies%2==1)
        return 2+bunnyEars2(bunnies-1); // c2=2+t(n-1)
    return 3+bunnyEars2(bunnies-1);    // t(n) = c2 + t(n - 1)
}
```

Model :

$$T(n) = \begin{cases} 0 & \text{if } n == 0 \\ 2 + t(n-1) & \text{if } n == 1 \\ 3 + t(n-1) & \text{if } n \geq 2 \end{cases}$$

Recurrence equation solution :

$$t(n) = c_1 + 3n$$

Calculate complexity:

$$O(C_1 + 3 * n)$$

$$O(3*n) \quad // \text{ By addition rule}$$

$$O(n) \quad // \text{ By multiplication rule}$$

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- **AllStar:**

```
public String allStar(String str) {
    if(str.length() <=1)                //c1
    return str;                          //c2
    return str.charAt(0) + "*" + allStar(str.substring(1)); //t(n)=c2+t(n-1)
}
```

Model :

$$T(n) = \begin{cases} c_1 + c_2 & \text{if } n.length \leq 1 \\ c_2 + t(n-1) & \text{if } n.length \geq 2 \end{cases}$$

Recurrence equation solution :

$$t(n) = c_2 * n + c_1$$

Calculate complexity:

$$O(c_2 * n + c_1)$$

$$O(c_2 * n) \quad // \text{ By addition rule}$$

$$O(n) \quad // \text{ By multiplication rule}$$

- **Triangle**

```
public int triangle(int rows){
    if(rows <=1)                //c1
    return rows;                //c2
    return rows+triangle(rows-1); //t(n)=c2+t(n-1)
}
```

Model :

$$T(n) = \begin{cases} 1 & \text{if } n \leq 1 \\ c_2 + t(n-1) & \text{if } n \geq 2 \end{cases}$$

Recurrence equation solution :

$$t(n) = c_2 n + c_1$$

Calculate complexity:

$$O(c_2 * n + c_1)$$

$$O(c_2 * n) \quad // \text{ By addition rule}$$

$$O(n) \quad // \text{ By multiplication rule}$$

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RECURSION 2

- **groupSum5**

```
public boolean groupSum5(int start, int[] nums, int target) {
    if(start==nums.length)
        return target==0; //c1
    if(nums[start] % 5 == 0){
        if(start<nums.length-1 && nums[start+1]==1)
            return groupSum5(start+2,nums,target-nums[start]); //c2
        else
            return groupSum5(start+1,nums,target-nums[start]); //c3+t(n-2)
    }
    ReturngroupSum5(start+1,nums,target-nums[start])||groupSum5(start+1,nums,target);
    //c4+2(tn-1)
}
```

Recurrence equation solution:

$$T(n) = c_4 (2^n - 1) + c_1 2^{(n-1)}$$

Calculate complexity:

$$O(c_4 (2^n - 1) + c_1 2^{(n-1)})$$

$$O(c_4(2^n) + c_1 2^n) \quad // \text{By addition rule}$$

$$O((2^n) + 2^n) \quad // \text{By multiplication rule}$$

$$O(2 * 2^n) \quad // \text{By addition rule}$$

$$O(2^n) \quad // \text{By multiplication rule}$$

- **groupSum6**

```
public boolean groupSum6(int start, int[] nums, int target) {
    if (start >= nums.length)
        return target == 0; //C1
    if (groupSum6(start + 1, nums, target - nums[start]))
        return true; //C2 + T(n-2)
    if (nums[start] == 6)
        target -= 6;
    return groupSum6(start + 1, nums, target); //C3+2t(n-1)
}
```

Recurrence equation solution:

$$T(n) = c_3(2^n - 1) + c_1 2^{n-1}$$

Calculate complexity:

$$O(c_3(2^n - 1) + c_1 2^{n-1})$$

$$O(c_3(2^n) + c_1 2^n) \quad // \text{By addition rule}$$

$$O(2^n) + 2^n \quad // \text{By multiplication rule}$$

$$O(2 * 2^n) \quad // \text{By addition rule}$$

$$O(2^n) \quad // \text{By multiplication rule}$$

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- **groupNoAdj**

```
public boolean groupNoAdj(int start, int[] nums, int target) {
    if (start >= nums.length)
        return (target == 0); //c
    if (groupNoAdj(start+1, nums, target))
        return true; //c+t(n-1)
    if (groupNoAdj(start+2, nums, target-nums[start]))
        return true; //c2+t(n-1)+t(n-2)
    return false;
}
```

Recurrence equation solution:

$$T(n) = c_1 \cdot 2^n + c_2$$

Calculate complexity:

$$O(C_1 \cdot 2^n + c_2)$$

$$O(C_1 \cdot 2^n) \quad // \text{ By "Laboratory practice No. 1: Recursion" point 4.4.1}$$

$$O(2^n)$$

- **SplitOdd10**

```
public boolean splitOdd10Aux(int start, int[] nums, int mult, int odd) {
    if (start >= nums.length)
        return mult % 10 == 0 && odd % 2 == 1; //c
    if (splitOdd10Aux(start+1, nums, mult + nums[start], odd))
        return true; //c+(tn-1)
    if (splitOdd10Aux(start+1, nums, mult, odd + nums[start]))
        return true; //c2+2t(n-1)
    return false;
}
```

Recurrence equation solution:

$$t(n) = c_2 (2^n - 1) + c_1 2^{n-1}$$

Calculate complexity:

$$O(c_2(2^n - 1) + c_1 2^{n-1})$$

$$O(c_2(2^n) + c_1 2^n) \quad // \text{ By addition rule}$$

$$O(2^n) + 2^n \quad // \text{ By multiplication rule}$$

$$O(2 \cdot 2^n) \quad // \text{ By addition rule}$$

$$O(2^n) \quad // \text{ By multiplication rule}$$

- **Split53**

```
public boolean split53(int[] nums) {
    return splitAux(nums, 0, 0, 0);
}

public boolean splitAux(int[] nums, int prefix, int sum1, int sum2){
```

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```

if(prefix==nums.length)return sum1==sum2;
if(nums[prefix]%5==0){                                     //c1
    return splitAux(nums, prefix+1, sum1, sum2+nums[prefix]);
}else if(nums[prefix]%3==0){                                //c2+t(n-1)
    return splitAux(nums, prefix+1, sum1+nums[prefix], sum2);
}else{
    return splitAux(nums, prefix+1, sum1, sum2+nums[prefix])||
    splitAux(nums, prefix+1, sum1+nums[prefix], sum2);        //c3+2t(n-1)
}
}

```

Recurrence equation solution:

$$T(n) = c_3(2^n - 1) + c_1 2^{n-1}$$

Calculate complexity:

$$O(c_3(2^n - 1) + c_1 2^{n-1})$$

$$O(c_3(2^n) + c_1 2^n) \quad // \text{By addition rule}$$

$$O(2^n) + 2^n \quad // \text{By multiplication rule}$$

$$O(2 * 2^n) \quad // \text{By addition rule}$$

$$O(2^n) \quad // \text{By multiplication rule}$$

3.6

- **BunnyEars:** The variable **Bunnies** sets the length and complex of the problema, **n** represents the number of bunnies
- **Fibonacci:** The variable **n** is the Fibonacci number to calculate
- **BunnyEars2:** The variable **Bunnies** represents the number of bunnies
- **AllStar:** The **str** is a string the algorithm needs to separate by * all the letters
- **Triangle:** The variable **rows** sets the amount of blocks we need to complete the triangle
- **Groupsum5:** The variable **n** is the size of the array
- **Groupsum6:** The variable **n** is the size of the array
- **Groupoadj:** The variable **n** is the size of the array
- **SplitOdd10:** The variable **n** is the size of the array
- **Split53:** the variable **n** is the size of the array

4) Practice for midterms

4.1 1)a

2)c

3)a

4.2 1)b

2)a-c

4.3 b

4.4 c

4.5 1) a

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2) *b*

- 4.6** 1) Complete la línea 10=
 `return sumaAux(n, i+1);`
2) Complete la línea 12=
 `return (n.charAt(i)- '0' + sumaAux(n, i+1));`

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