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Promela model for peterson leader election protocol.

The model initialized N processes using the example init process from leader0. Each node process has a input and output channel. The input and output channels are assigned to create a ring of nodes with unidirectional communication. Each process is connected to its clockwise neighbor through out and counter clockwise neighbor via in. There are three possible states for a process: Active, Relay, and Stop. In order to avoid invalid end states we use the end keyword to label relay as an end state. All processes start as active and has a unique identity or process number. When a process is active it initially sends it's own id which it assumes is the max. The process then waits to receive the value *e* from its neighbor. If *e* does not equal *d* then it also waits for the second message *f*. Next the process compares *d*, *e*, and *f*. If *e* is the max of the three values *d* is reassigned to the value of *e* otherwise the process goes to relay. Processes in relay wait to receive input *d*, check if *d* is their identity. If it is not they pass on *d* to the output.

2. Verify the following properties using the SPIN model checker:

- (a) There is always at most one leader. **PASS**
- (b) Eventually always a leader will be elected. **PASS**
- (c) The elected leader will be the process with the highest number. **PASS**
- (d) The maximum total amount of messages sent in order to elect a leader is at most  $2N \log_2 N + N$ . **PASS**

**/\*Peterson leader election protocol  
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```
*/  
#define N      6      /* nr of processes */  
#define I      3      /* node given the smallest number */  
#define L     12      /* size of buffer (>= 2*N) */  
#define X     38      /* 2*N*(Log2(N))+6 */  
ltl p1 {<>[] (nr_leaders==1)};  
ltl p2 {[] (messages<X)};  
ltl p3 {[] (nr_leaders <= 1)};  
ltl p4 {[] (leader_id==N)};
```

```
chan q[N] = [L] of {byte};
```

```
//counter for total leaders elected
```

```
byte nr_leaders = 0;
```

```
//counter for total messages sent
```

```
byte messages=0;
```

```
//used to check if leader has the largest ID which should be N
```

```
byte leader_id=N;
```

```
proctype node (chan in, out; byte ident)
```

```
{   byte d,e,f;
```

```
//Set exclusive read and write access for chan for this process
```

```
    xr in;
```

```
    xs out;
```

```
//    printf("MSC: %d\n", ident);
```

```
//Initially output my identity which I assume is the max until told otherwise
```

```
    activ:
```

```
//    printf("active[%d]\n",ident);
```

```
    d=ident;
```

```
    do :: true -> out!d;
```

```
        messages++;
```

```
        in?e;
```

```
        if :: (e==ident) ->
```

```
            //process is the leader
```

```
            goto stop
```

```
        :: else -> skip
```

```
    fi;
```

```
    if::d>e->out!d;
```

```
        ::else-> out!e;
```

```
    fi
```

```
    messages++;
```

```
    in?f;
```

```
    if :: (f==ident) ->
```

```
        //process is the leader
```

```
        goto stop
```

```
    ::else ->skip
```

```
    fi
```

```

        if::(e>=d)&&(e>=f) -> d =e
        ::else -> goto relay //put process into passive mode
    fi
od;

relay:
//relay mode
//printf("RelayNode[%d]\n",ident);
end:
    do :: in?d->
        if::(d != ident)->
            //process is not the leader
            out!d
            messages++;
        ::else -> goto stop //process is the leader
    fi
od;

stop:
printf("Leader: %d",ident);
nr_leaders++;
leader_id=ident;
    skip
}

/* initialize N processes with IDs*/
init {
    byte proc;
    atomic {
        proc = 1;
        do
            :: proc <= N ->
                run node (q[proc-1], q[proc%N], (N+1-proc)%N+1);
                proc++
            :: proc > N ->
                break
        od
    }
}

```