

Semantic Theory 2025: Exercise 1 Key

Question 1

Translate the following sentences into first-order predicate logic. You can freely introduce predicates, but try to retain as much of the structure as possible. Also provide the key to the translation.

- a. Every student uses their computer.

$$\forall x(student(x) \rightarrow \exists y(computer(y) \wedge own(x, y) \wedge use(x, y)))$$

- b. John doesn't see anything.

$$\neg \exists x(see(j, x)) \text{ (equivalently: } \forall x(\neg see(j, x)) \text{)}$$

- c. The pen is the only instrument that is mightier than the sword.

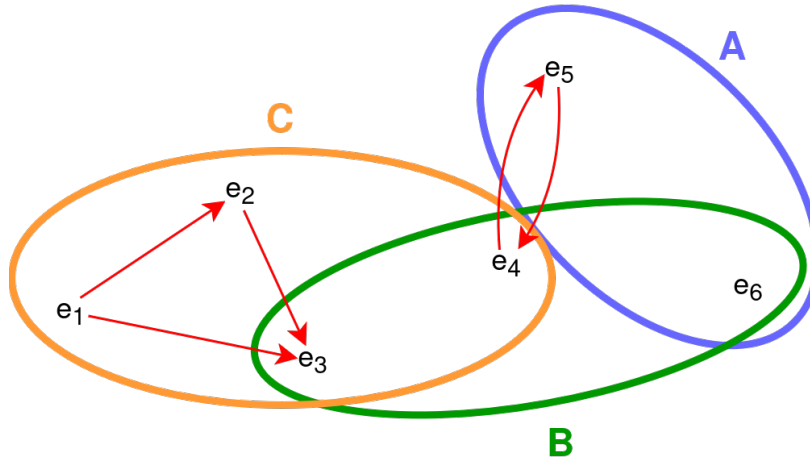
$$\forall x((instrument(x) \wedge \forall y(sword(y) \rightarrow mightier-than(x, y))) \rightarrow pen(x))$$

- d. If the Prime Minister of the country is defeated, the city will celebrate.

$$defeated(pm') \rightarrow celebrate(c')$$

Question 2

Consider the following model $M_1 = (U_1, V_1)$, with $U_1 = \{e_1, e_2, e_3, e_4, e_5, e_6\}$.



The interpretation function V_1 is defined as follows:

- $V_1(j) = e_1$
- $V_1(m) = e_4$
- $V_1(b) = e_6$
- $V_1(A) = \{e_5, e_6\}$
- $V_1(B) = \{e_3, e_4, e_6\}$
- $V_1(C) = \{e_1, e_2, e_3, e_4\}$
- $V_1(R) = \{(e_1, e_2), (e_1, e_3), (e_2, e_3), (e_4, e_5), (e_5, e_4)\}$

Let the assignment function g_1 be defined as follows:

$g_1(x) = e_4$, $g_1(x') = e_2$, $g_1(x'') = e_3$ and for all other variables x'^* : $g_1(x'^*) = e_5$.

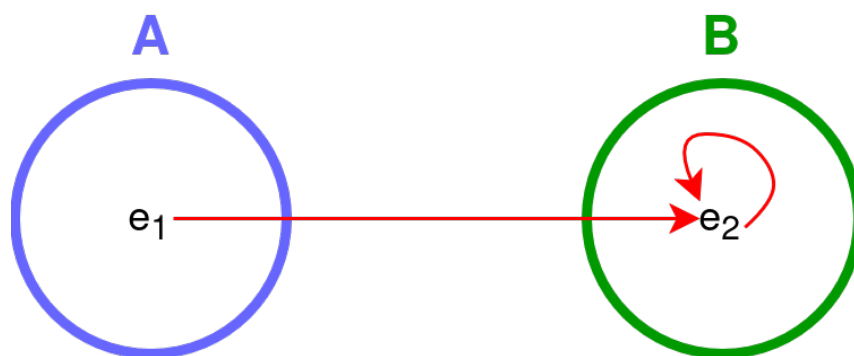
2.1 Evaluate the following formulas in model M_1 , with respect to assignment function g_1 . First, derive the truth conditions (showing all relevant steps of the derivation) and then evaluate these truth conditions with respect to M_1 and g_1 .

- a. $\llbracket R(x', x'') \wedge R(x''', b) \rrbracket^{M_1, g_1} = 1$
 iff $\llbracket R(x', x'') \rrbracket^{M_1, g_1} = 1$ and $\llbracket R(x''', b) \rrbracket^{M_1, g_1} = 1$
 iff $\llbracket (x', x'') \rrbracket^{M_1, g_1} \in \llbracket R \rrbracket^{M_1, g_1}$ and $\llbracket (x''', b) \rrbracket^{M_1, g_1} \in \llbracket R \rrbracket^{M_1, g_1}$
 iff $(g_1(x'), g_1(x'')) \in V_1(R)$ and $(g_1(x'''), V_1(b)) \in V_1(R)$
 $\Rightarrow (e_2, e_3) \in \{(e_1, e_2), (e_1, e_3), (e_2, e_3), (e_4, e_5), (e_5, e_4)\}$ (true)
 and $(e_5, e_6) \in \{(e_1, e_2), (e_1, e_3), (e_2, e_3), (e_4, e_5), (e_5, e_4)\}$ (false)
 $\Rightarrow \llbracket R(x', x'') \wedge R(x''', b) \rrbracket^{M_1, g_1} = 0$
- b. $\llbracket \exists x(A(x) \rightarrow R(x'', j)) \rrbracket^{M_1, g_1} = 1$
 iff $\llbracket \exists y(A(y) \rightarrow R(x'', j)) \rrbracket^{M_1, g_1} = 1$
 iff there is a $d \in U_1$ such that $\llbracket (A(y) \rightarrow R(x'', j)) \rrbracket^{M_1, g_1[y/d]} = 1$
 iff $\llbracket A(y) \rrbracket^{M_1, g_1[y/d]} = 0$ or $\llbracket R(x'', j) \rrbracket^{M_1, g_1[y/d]} = 1$
 iff $\llbracket y \rrbracket^{M_1, g_1[y/d]} \notin \llbracket A \rrbracket$ or $\llbracket (x'', j) \rrbracket^{M_1, g_1[y/d]} \in \llbracket R \rrbracket^{M_1, g_1[y/d]}$
 iff $g_1[y/d](y) \notin V_1(A)$ or $(g_1[y/d](x''), V_1(j)) \in V_1(R)$
 Pick any $d \in U_1 - V_1(A)$: let's set $d = e_1$
 $e_1 \notin V_1(A) = \{e_5, e_6\} \Rightarrow \llbracket \exists x(A(x) \rightarrow R(x'', j)) \rrbracket^{M_1, g_1} = 1$
- c. $\llbracket \forall x(B(x) \rightarrow (A(x) \vee \exists x'(R(x', x)))) \rrbracket^{M_1, g_1} = 1$
 iff for all $d \in U_1$: $\llbracket B(x) \rightarrow (A(x) \vee \exists x'(R(x', x))) \rrbracket^{M_1, g_1[x/d]} = 1$
 iff $\llbracket B(x) \rrbracket^{M_1, g_1[x/d]} = 0$ or $\llbracket A(x) \vee \exists x'(R(x', x)) \rrbracket^{M_1, g_1[x/d]} = 1$
 ... we know how to do the rest from (a) and (b)
 Evaluation w.r.t. M_1 :
 for all $d \in U_1 - V_1(B) = \{e_1, e_2, e_5\}$: $\llbracket B(x) \rrbracket^{M_1, g_1[x/d]} = 0$, so
 $\llbracket B(x) \rightarrow (A(x) \vee \exists x'(R(x', x))) \rrbracket^{M_1, g_1[x/d]} = 1$
 for all $d \in \{e_3, e_4\}$: $\llbracket \exists x'(R(x', x)) \rrbracket^{M_1, g_1[x/d]} = 1$, so
 $\llbracket B(x) \rightarrow (A(x) \vee \exists x'(R(x', x))) \rrbracket^{M_1, g_1[x/d]} = 1$
 for $d = e_6$: $\llbracket A(x) \rrbracket^{M_1, g_1[x/d]} = 1$, so
 $\llbracket B(x) \rightarrow (A(x) \vee \exists x'(R(x', x))) \rrbracket^{M_1, g_1[x/d]} = 1$
 $\Rightarrow \llbracket \forall x(B(x) \rightarrow (A(x) \vee \exists x'(R(x', x)))) \rrbracket^{M_1, g_1} = 1$

2.2 Provide a graphical representation of a model that satisfies the following formulas (NB: c_1 and c_2 are constants):

- $R(x, x')$
- $\forall x(A(x) \vee \exists x'(R(x, x')))$
- $\neg \exists x(R(x, c_1))$
- $\exists x''(B(x'') \wedge \neg \exists x'(A(x') \wedge R(x'', x')))$
- $\forall x'(A(x') \rightarrow (B(x') \vee R(x', c_2)))$

The simplest model I could think of:



Where:

$$V_M(c_1) = e_1, V_M(c_2) = e_2$$

$$g(x) = e_1, g(x') = e_2$$

2.3 (Bonus) Can you think of a sensible (or: funny) interpretation for the predicates A, B and R , and the constants c_1 and c_2 in your model of the previous exercise? Given this interpretation, what is the natural language translation of the formulas given in exercise 2.2?

Up to you!