

Problem 3

Due on November 16th

Consider the RANDOMIZED HYPOTHESIS ELIMINATION algorithm defined for a finite hypothesis set H as follows:

Algorithm 1 RANDOMIZED HYPOTHESIS ELIMINATION

Initialization: Let $V_1 = H$.

For all $t = 1, 2, \dots, n$, **repeat**

- Choose an $h_t \in V_t$ uniformly at random.
 - Observe \mathbf{x}_t .
 - Predict $\hat{y}_t = h_t(\mathbf{x}_t)$.
 - Update $V_{t+1} = \{h \in V_t : h(\mathbf{x}_t) = y_t\}$.
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The following theorem establishes a performance guarantee for this algorithm:

Theorem 3. *Under the realizability assumption, the total expected loss of RANDOMIZED HYPOTHESIS ELIMINATION over n rounds satisfies*

$$\mathbb{E} \left[\sum_{t=1}^n \ell(\mathbf{x}_t, y_t, \hat{y}_t) \right] \leq 1 + \ln |H|.$$

Your homework is to answer the following questions about RANDOMIZED HYPOTHESIS ELIMINATION and the other algorithms we discussed in class:

Question 1: When is the bound of Theorem 3 better than the bound of Theorem 2 presented in the lecture? (What is the size of the smallest hypothesis class for which the former is better than the latter?)

Question 2: What is the computational complexity of RANDOMIZED HYPOTHESIS ELIMINATION? (Hints: How expensive is it to generate a random hypothesis and test if it is in the version space? How many random hypotheses do we need to generate until we find a consistent one?)

Question 3: Construct a hypothesis class and a sequence of examples for which HYPOTHESIS ELIMINATION (the non-randomized one) makes exactly $|H| - 1$ mistakes. (Hint: a possible good idea is choosing the \mathbf{x}_t 's from the natural numbers $1, 2, \dots, |H|$ and consider the hypothesis class H consisting of functions of the form $h_k(x) = \mathbb{I}\{x \leq k\}$ (i.e., $h_k(x) = 1$ if $x \leq k$ and $h_k(x) = 0$ otherwise) for all $k = 1, 2, \dots, |H|$.)

Question 4: Construct a hypothesis class and a sequence of examples for which HALVING makes exactly $\log_2 |H|$ mistakes. (Hint: The same hypothesis class as above and the same pool of \mathbf{x}_t 's may work here too.)

Question 5*: Prove Theorem 3! (See hints in the slides and below.)

Hints for Question 5

It is enough to consider a sequence $(\mathbf{x}_t, y_t)_{t=1}^n$ such that in every round t , there is exactly one hypothesis $g_t \in H$ making a mistake. (For the curious reader, this follows from the argument used for proving Corollary 4.5 in Cesa-Bianchi and Lugosi, 2006.) At any round t , what is the maximal probability of the algorithm choosing the hypothesis t ? Hint: choosing a hypothesis h_t uniformly at random from a set V means that each hypothesis in V is chosen with equal probability $\frac{1}{|V|}$.

References

Cesa-Bianchi, N. and Lugosi, G. (2006). *Prediction, Learning, and Games*. Cambridge University Press, New York, NY, USA.