Session 23: Introduction to non-linear programming (NLP)

A generic non-linear program has the following form (*x* is a *n*-dimensional vector.)

Minimize:
$$f(x)$$

 $g_1(x) \le 0$
 \dots
 $g_m(x) \le 0$
 $x \in X$

In a linear program, all of the functions f, g_1 , \cdots , g_m are linear, and the constraint set X is n-dimensional real numbers. In a MIP, the constraint set X also requires certain variables to be integers.

Unlike with LP and MIPs, there are no generic algorithms that can solve arbitrary non-linear programs with large n or m. The strategy in practice is to find special structure in the functions f and g_i .

1. Sum of Squares

Let x, y and z be decision variables. Gurobi allows the following types of quadratic constraints:

- **Minimizing** a sum of squares plus a linear term, or an expression that can be expressed as a sum of squares plus a linear term.
- A sum of squares plus a linear term **less than** equal to a linear term.

Examples of what is allowed:

Minimize:
$$x^2 + (x+2y)^2 + 3(y+z)^2 - 5y$$

 $x^2 + 2y^2 + 3z^2 \le 5$
 $(x+2y)^2 \le 5z - 2x$
 $x^2 + 4y^2 + 4xy \le 5z - 2x$
 $6x + y^2 + 2z \le 10$

However, in the code, you cannot use the power ** operator or ^2, but must multiply out using *:

```
[1]: from gurobipy import Model, GRB
    mod=Model()
    x=mod.addVar(lb=-GRB.INFINITY)
    y=mod.addVar(lb=-GRB.INFINITY)
    z=mod.addVar(lb=-GRB.INFINITY)

mod.setObjective(x*x+(x+2*y)*(x+2*y)+3*(y+z)*(y+z)-5*y)
    mod.addConstr(x*x+2*y*y+3*z*z <= 5)
    mod.addConstr((x+2*y)*(x+2*y)<=5*z-2*x)
    mod.addConstr(x*x+4*y*y+4*x*y<=5*z-2*x)
    lhs4=6*x+y*y+2*z
    mod.addConstr(lhs4<=10)
    mod.optimize()

[2]: print(f'Optimal solution: obj={mod.objval}')
    print(f'\t x={x.x}')</pre>
```

```
print(f'\t y={y.x}')
print(f'\t z={z.x}')
print(f'\t lhs4={lhs4.getValue()}')

Optimal solution: obj=-2.0411051524261055
    x=-1.134296289724148
    y=0.7968655538079593
    z=-0.41150244505150585
    lhs4=-6.993787917602234
```

The following are **NOT allowed**:

• Subtracting a square instead of adding:

Minimize:
$$x^2 + (x + 2y)^2 - 3(y + z)^2 - 5y$$

• Maximizing a sum of squares:

Maximize:
$$x^2 + (x + 2y)^2 + 3(y + z)^2 - 5y$$

• A sum of squares larger than a linear expression:

$$x^2 + y^2 - 2xy \ge 5$$

Q1 (DMD Example 8.1)

Solve the following non-linear optimization formulation problem using Gurobi. The formulation maximizes expected returns of a portfolio subject to not exceeding a certain level of risk.

Decision variables: Let *A*, *G*, *D* denote the fraction of total investment to put in the assets Advent, GSS, and Digital.

Objective and constraints:

Maximize:
$$11A+14G+7D$$
 subset to:
$$(Fractions) \qquad \qquad A+G+D=1$$

$$(Target\;risk) \quad \sqrt{16A^2+22G^2+10D^2+6AG+2GD-10AD}\leq 3.1$$
 (Nonnegativity)
$$A,G,D\geq 0$$

[4]:

Optimal annual return: 12.250136110681433

A: 0.37801476764687053 G: 0.5340110057277011 D: 0.08797422662543453 risk: 3.099999192146312

2. Linearizing using Auxiliary Decision Variables

2.1 Max and Min

The non-linear objective

Minimize max(x, y)

is equivalent to

Minimize

z

subject to

$$\max(x, y) \le z$$

which is equivalent to the linear formulation:

Minimize

subject to

 $x \le z$

z

 $y \le z$

Similarly,

Minimize
$$max(x, y) - min(x, y)$$

is equivalent to

Minimize U - L

subject to:

 $L \le x \le U$

 $L \le y \le U$

2.2 Absolute Values

Similarly, the non-linear objective

Minimize
$$|x_1 - y_1| + |x_2 - y_2|$$

is equivalent

Minimize

 $z_1 + z_2$

subject to

(New constraint 1) $|x_1 - y_1| \le z_1$

(New constraint 2) $|x_2 - y_2| \le z_2$

Now, because $|x| = \max(x, -x)$, the above is equivalent to the linear formulation:

Minimize

 $z_1 + z_2$

subject to

$$x_1 - y_1 \le z_1$$

$$y_1 - x_1 \le z_1$$

$$x_2 - y_2 \le z_2$$

$$y_2 - x_2 \le z_2$$

2.3 Big-M Method

Suppose we want to either turn a continuous variable *X* on or off, we can do

where Z is a binary decision variable and M is a sufficiently large number that is guaranteed to be larger than the maximum possible value of X in any optimal soluiton.

2.4 Either/Or Constraint

Similar to the above, suppose that we want to force a continuous variable X to be either between A_1 and A_2 or between B_1 and B_2 , we can do

$$ZA_1 + (1-Z)B_1 \le X \le ZA_2 + (1-Z)B_2$$
.

Q2 (Portfolio Optimization with Complex Constraints)

Consider the following formulation of a portfolio optimization problem, use auxiliary decision variables to linearize the last four constraints.

Data:

- *S*: the set of stocks.
- w_i : the old weight of stock $i \in S$ before optimization. (The "weight" of a stock is % of total funds invested in the stock; weights of all stocks should add to one.)
- R_i : the expected annual return of stock $i \in S$.
- C_{ij} : the estimated covariance between stocks $i, j \in S$.
- σ_{target} : the maximum volatility of the final portfolio.
- Δ : the total movement allowed between the old weights and the new weights.
- *k*: the maximum # of stocks allowed in the portfolio.
- ϵ : the minimum non-zero weight allowed.
- λ : penalty for different weights for stocks 1, 2 and 3.

Decision variables:

- x_i : the new weight of stock i. (Continuous)
- *y*: variation of weights among stocks 1, 2 and 3. (Continuous)

Formulation: All summations are over the set *S* of stocks.

Maximize:
$$\sum_{i} R_{i}x_{i} - \lambda y \qquad \text{(Average Return)}$$
 subject to:
$$\text{(Valid weights)} \qquad \sum_{i} x_{i} = 1$$

$$\text{(Risk tolerance)} \qquad \sum_{i,j} C_{ij}x_{i}x_{j} \leq \sigma_{target}^{2}$$

$$\text{(Change in weights)} \qquad \sum_{i} |x_{i} - w_{i}| \leq 2\Delta$$

$$\text{(Non-negligible weights)} \qquad \text{If } x_{i} > 0 \text{ then } x_{i} \geq \epsilon \qquad \text{for each stock } i.$$

$$\text{(Simplicity)} \qquad \text{(# of stock } i \text{ with } x_{i} > 0) \leq k$$

$$\text{(Similar weights for stocks 1-3)} \qquad \max(x_{1}, x_{2}, x_{3}) - \min(x_{1}, x_{2}, x_{3}) \leq y$$

$$x_{i} \geq 0$$

3. (Optional) Special Non-Linear Constraints Supported in Gurobi

The following information will not be tested in any exam or quiz, but may be helpful for the final project.

In the following table, x, y, z, and w are decision variables. Moreover, the code assume that you have imported all of the functions.

```
from gurobipy import Model, GRB, max_, min_, abs_, and_, or_
mod=Model()
```

Non-linear relationship	Sample Constraint as Math Expression	Gurobi Command
Maximum of arbitrary variables	x = max(y, z, 5)	mod.addConstr(x==max_(y,z,5))
Minimum of arbitrary variables	x = min(y, z, 5)	<pre>mod.addConstr(x==min_(y,z,5))</pre>
Absolute value of arbitrary variables	x = y = max(y, -y)	<pre>mod.addConstr(x==abs_(y))</pre>
AND of binary variables	x = min(y, z, w)	<pre>mod.addConstr(x==and_(y,z,w))</pre>
OR of binary variables	x = max(y, z, w)	<pre>mod.addConstr(x==or_(y,z,w))</pre>
At most one (arbitrary	$1(x \neq 0) + 1(y \neq$	<pre>mod.addSOS(GRB.SOS_TYPE1,[x,y,z</pre>
variable) non-zero	$0) + 1(z \neq 0)) \le 1$	

Example:

```
[5]: from gurobipy import Model, GRB, max_, min_, and_, abs_
     mod=Model()
     x=mod.addVar(lb=-GRB.INFINITY)
     y=mod.addVar(lb=-GRB.INFINITY)
     z=mod.addVar()
     l=mod.addVar(lb=-GRB.INFINITY)
     a=mod.addVar()
     mod.addConstr(z==max_(x,5,y))
     mod.addConstr(l==min_(x,y))
     mod.addConstr(z<=1+5)</pre>
     mod.addConstr(a==abs_(y))
     mod.addSOS(GRB.SOS_TYPE1, [x,y])
     mod.setObjective(z-1-a, sense=GRB.MAXIMIZE)
     mod.optimize()
     print('\nObjective', mod.objval)
     print(f'x=\{x.x\} y=\{y.x\} z=\{z.x\} 1=\{1.x\} a=\{a.x\}')
Optimize a model with 1 rows, 5 columns and 2 nonzeros
Model has 1 SOS constraint
Model has 3 general constraints
Variable types: 5 continuous, 0 integer (0 binary)
Coefficient statistics:
 Matrix range
                    [1e+00, 1e+00]
  Objective range [1e+00, 1e+00]
 Bounds range
                   [0e+00, 0e+00]
 RHS range
                   [5e+00, 5e+00]
```

Presolve added 6 rows and 6 columns

Presolve time: 0.00s

Presolved: 7 rows, 11 columns, 19 nonzeros Presolved model has 4 SOS constraint(s)

Variable types: 7 continuous, 4 integer (4 binary)

Root relaxation: objective 5.000000e+00, 5 iterations, 0.00 seconds

Nodes		Current Node			Object	Objective Bounds		Work			
E	xpl Une	xpl	Obj	Depth	IntI	nf	Incumbent	BestBd	Gap	It/Node	Time
	0	0	5.00	0000	0	1	_	5.00000	-	-	0s
Н	0	0					5.0000000	5.00000	0.00%	_	0s

Explored 1 nodes (5 simplex iterations) in 0.01 seconds Thread count was 4 (of 4 available processors)

Solution count 1: 5

Optimal solution found (tolerance 1.00e-04)
Best objective 5.000000000000e+00, best bound 5.00000000000e+00, gap 0.0000%

Objective 5.0 x=5.0 y=0.0 z=5.0 1=0.0 a=0.0