

<https://sinhp.github.io/Poly> <https://github.com/sinhp/Poly> <https://sinhp.github.io/Poly/docs>

Polynomial Functors in Lean4

Sina Hazratpour

April 3, 2025

Chapter 1

Locally Cartesian Closed Categories

To be written

Chapter 2

Univariate Polynomial Functors

In this section we develop some of the definitions and lemmas related to polynomial endofunctors that we will use in the rest of the notes.

Definition 2.0.1 (Polynomial endofunctor). *UvPoly* Let \mathbb{C} be a locally Cartesian closed category (in our case, presheaves on the category of contexts). This means for each morphism $t : B \rightarrow A$ we have an adjoint triple

$$\begin{array}{ccc} & \mathbb{C}/B & \\ t_! \left(\begin{array}{c} \uparrow \\ \dashv \quad t^* \quad \dashv \\ \downarrow \end{array} \right) & & t_* \\ & \mathbb{C}/A & \end{array}$$

where t^* is pullback, and $t_!$ is composition with t .

Let $t : B \rightarrow A$ be a morphism in \mathbb{C} . Then define $P_t : \mathbb{C} \rightarrow \mathbb{C}$ be the composition

$$P_t := A_! \circ t_* \circ B^*$$

$$\mathbb{C} \xrightarrow{B^*} \mathbb{C}/B \xrightarrow{t_*} \mathbb{C}/A \xrightarrow{A_!} \mathbb{C}$$

Proposition 2.0.2 (Characterising property of Polynomial Endofunctors). *The data of a map into the polynomial applied to an object in \mathbb{C}*

$$\Gamma \longrightarrow P_t Y$$

corresponds to a pair of morphisms

$$\alpha : \Gamma \rightarrow A \quad \text{and} \quad \beta : \Gamma \cdot \alpha \rightarrow Y$$

and this correspondance is natural in both Γ and Y .

Given any such ϕ we can extract $\alpha : \Gamma \rightarrow A$ by composition

$$\begin{array}{ccc}
\Gamma & \xrightarrow{\phi} & P_t Y \\
& \searrow \alpha & \swarrow t_* B^* Y \\
& A &
\end{array}$$

Applying the adjunction $A_! \dashv A^*$, and viewing $\phi : \alpha \rightarrow t_* B^* Y$ as a map in the slice, this corresponds to

$$\begin{array}{ccc}
\alpha : \Gamma \rightarrow A & \text{and} & \\
B_! t^* \alpha \dashrightarrow B \times Y & \xrightarrow{\tilde{\phi}} & \\
\downarrow t^* \alpha & & \downarrow B^* Y \\
& B &
\end{array}$$

Applying the adjunction $t^* \dashv t_*$, this corresponds to

$$\begin{array}{ccc}
\alpha : \Gamma \rightarrow A & \text{and} & \\
\Gamma \cdot \alpha := B_! t^* \alpha \dashrightarrow Y & \xrightarrow{\beta} &
\end{array}$$

Henceforth we will write

$$(\alpha, \beta) : \Gamma \rightarrow P_t Y$$

for this map, since it is uniquely determined by this data.

This is natural in Γ . Precomposition by $\sigma : \Delta \rightarrow \Gamma$, acts on such a pair by

$$\begin{array}{ccc}
\Delta & & \\
\sigma \downarrow & \searrow (\alpha \circ \sigma, \beta \circ t^* \sigma) & \\
\Gamma & \xrightarrow{(\alpha, \beta)} & P_t Y
\end{array}$$

It is also natural in $f : X \rightarrow Y$, meaning the morphism $P_t f$ acts on such a pair by

$$\begin{array}{ccc}
\Gamma & \xrightarrow{(\alpha, \beta)} & P_t X \\
& \searrow (\alpha, f \circ \beta) & \downarrow P_t f \\
& & P_t Y
\end{array}$$

Lemma 2.0.3. *prop:UvPoly.equiv* Use R to denote the fiber product

$$\begin{array}{ccc}
R & \xrightarrow{\rho_{\text{Im}}} & B \\
\rho_P \downarrow & \lrcorner & \downarrow t \\
P_t Y & \xrightarrow{t_* B^* Y} & A
\end{array}$$

By the universal property of pullbacks and ??, The data of a map $\Gamma \rightarrow R$ corresponds to the data of $\beta : \Gamma \rightarrow B$ and $(t \circ \beta, y) : \Gamma \rightarrow P_t Y$, or just $\beta : \Gamma \rightarrow B$ and $y : \Gamma \cdot t \circ \beta \rightarrow Y$

$$\begin{array}{ccccc}
\Gamma & & \xrightarrow{\beta} & & B \\
& \searrow (\beta, y) & & \searrow \rho_B & \\
& & R & \xrightarrow{\rho_B} & B \\
& & \downarrow \rho_P \lrcorner & & \downarrow t \\
& & P_t Y & \xrightarrow{t_* B^* Y} & A \\
& \searrow (t \circ \beta, y) & & &
\end{array}$$

By uniqueness in the universal property of pullbacks and ??, Precomposition by a map $\sigma : \Delta \rightarrow \Gamma$ acts on such a pair by

$$\begin{array}{ccc}
\Delta & & \\
\sigma \downarrow & \searrow (\beta \circ \sigma, y \circ t^* \sigma) & \\
\Gamma & \xrightarrow{(\beta, y)} & R
\end{array}$$

Definition 2.0.4 (Evaluation). *defn:UvPoly* Let $\text{counit} : \rho_B \rightarrow B \rightarrow B^*Y$ denote the counit of the adjunction $f^* \dashv f_*$ at the object B^*Y , recalling that $\rho_B = t^* t_* B^*Y$. Then viewing the object B^*Y in the slice as the object $Y \times B$ in the ambient category, we define $\text{ev} : R \rightarrow Y$ as the composition

$$\begin{array}{ccccc}
& & \text{ev} & & \\
& \nearrow & & \searrow & \\
R & \xrightarrow{\text{counit}} & Y \times B & \xrightarrow{\pi_Y} & Y \\
\rho_B \downarrow & & \nwarrow & & \\
& & B & &
\end{array}$$

Lemma 2.0.5 (Evaluation Computation). *defn:UvPolyEv* Suppose $(\beta, y) : \Gamma \rightarrow R$, as in ??

$$\beta : \Gamma \rightarrow B \quad \text{and} \quad y : \Gamma \cdot t \circ \beta \rightarrow Y$$

Then the evaluation of y at β can be computed as

$$\text{ev} \circ (\beta, y) = y \circ b$$

where

$$\begin{array}{ccccc}
\Gamma & & \xrightarrow{\beta} & & B \\
& \searrow b \text{ (dashed)} & & \searrow v & \\
& & \Gamma \cdot t \circ \beta & \xrightarrow{v} & B \\
& & \downarrow d \lrcorner & & \downarrow t \\
& & \Gamma & \xrightarrow{t \circ \beta} & A
\end{array}$$

and

$$\begin{array}{c}
\Gamma \xrightarrow{(t \circ \beta, y)} P_t Y \\
\downarrow (y \circ b, \beta) \quad \downarrow \text{counit} \quad \downarrow \lrcorner \quad \downarrow t_* B^* Y \\
Y \xleftarrow{\pi_Y} Y \times B \xrightarrow{\pi_B} B \xrightarrow{t} A
\end{array}$$

Proof. It suffices to show $(\text{counit} \circ (\beta, y)) = (y \circ b, \beta)$ instead.

$$\begin{aligned}
& \text{counit} \circ (\beta, y) \\
&= \text{counit} \circ (v \circ b, y \circ t^* d \circ t^* b) && ?? \\
&= \text{counit} \circ (v, y \circ t^* d) \circ b && ??, ?? \\
&= \text{counit} \circ t^* (t \circ \beta, y) \circ b && ?? \\
&= \overline{(t \circ \beta, y)} \circ b && ?? \\
&= (y, v) \circ b && ?? \\
&= (y \circ b, v \circ b) \\
&= (y \circ b, \beta)
\end{aligned}$$

$$\begin{array}{ccccc}
\Gamma \cdot t \circ \beta & \xrightarrow{t^* b} & \Gamma \cdot t \circ \beta \cdot t \circ \beta & \xrightarrow{t^* d} & \Gamma \cdot t \circ \beta \xrightarrow{v} B \\
\downarrow d & \lrcorner & \downarrow & \lrcorner & \downarrow d \lrcorner \downarrow t \\
\Gamma & \xrightarrow{b} & \Gamma \cdot t \circ \beta & \xrightarrow{d} & \Gamma \xrightarrow{t \circ \beta} A
\end{array}$$

Figure 2.1: $t^* d \circ t^* b = \text{id}_{\Gamma \cdot t \circ \beta}$

$$\begin{array}{ccc}
\Gamma & & \\
\downarrow b & \searrow (v \circ b, y \circ t^* d \circ t^* b) & \\
\Gamma \cdot t \circ \beta & \xrightarrow{(v, y \circ t^* d)} & R
\end{array}$$

Figure 2.2: $(v, y \circ t^* d) \circ b = (v \circ b, y \circ t^* d \circ t^* b)$

□

$$\begin{array}{ccccc}
& & v & & \\
& \searrow & & \nearrow & \\
\Gamma \cdot t \circ \beta & \xrightarrow{(v, y \circ t^* d)} & R & \longrightarrow & B \\
\downarrow d & \lrcorner & \downarrow & \lrcorner & \downarrow t \\
\Gamma & \xrightarrow{(t \circ \beta, y)} & P_t Y & \xrightarrow{t_* B^* Y} & A \\
& \nearrow & & \searrow & \\
& & t \circ \beta & &
\end{array}$$

$$\begin{array}{ccc}
\Gamma \cdot t \circ \beta & & \\
\downarrow d & \searrow (t \circ \beta \circ d, y \circ t^* d) & \\
\Gamma & \xrightarrow{(t \circ \beta, y)} & P_t Y
\end{array}$$

using ??, ??

Figure 2.3: $t^*(t \circ \beta, y) = (v, y \circ t^* d)$

$$\begin{array}{ccc}
t^*(t \circ \beta) & & t \circ \beta \\
\downarrow t^*(t \circ \beta, y) & \searrow \overline{(t \circ \beta, y)} & \downarrow (t \circ \beta, y) \\
t^* t_* B^* Y & \xrightarrow{\text{counit}} & B^* Y \\
& & \text{counit}
\end{array}
\quad
\begin{array}{ccc}
\parallel & & \\
t^* \dashv t_* & & t_* B^* Y \xrightarrow{\text{counit}} t_* B^* Y
\end{array}$$

Figure 2.4: $\text{counit} \circ t^*(t \circ \beta, y) = \overline{(t \circ \beta, y)}$

Definition 2.0.6 (Polynomial composition). Let $f : B \rightarrow A$ and $g : D \rightarrow C$. Define the *polynomial composition* $f \triangleleft g : Q \rightarrow P_f C$ as the composition of the two vertical maps in the following

$$\begin{array}{ccccc}
D & \longleftarrow & Q & & \\
\downarrow g & & \downarrow & \searrow f \triangleleft g & \\
C & \xleftarrow{\text{ev}} & R & \longrightarrow & B \\
& & \downarrow & \lrcorner & \downarrow f \\
& & P_f C & \xrightarrow{f_* B^* C} & A
\end{array}$$

Then the two functors

$$P_{f \triangleleft g} \cong P_f \circ P_g$$

are naturally isomorphic.

Proof.

□

Definition 2.0.7 (Mate). Suppose

$$\begin{array}{ccc}
\Gamma \cdot t \circ \beta \overset{(y,v)}{\dashrightarrow} Y \times B & & \Gamma \overset{(t \circ \beta, y)}{\rightarrow} P_t Y \\
\downarrow v = t^*(t \circ \beta) & \swarrow B^* Y & \downarrow t \circ \beta \\
B & & A \\
& \searrow t^* \dashv t_* & \swarrow t_* B^* Y
\end{array}$$

Figure 2.5: $\overline{(t \circ \beta, y)} = (y, v)$

$$\begin{array}{ccc}
C & \xrightarrow{\rho} & B \\
& \searrow s & \swarrow t \\
& & A
\end{array}$$

Then we have a mate $\mu_! : \rho_! \circ s^* \Rightarrow t^*$. This is given by the universal property of pullbacks: given $f : x \rightarrow y$ in the slice \mathbb{C}/A we have

$$\begin{array}{ccccc}
\bullet & \xrightarrow{\mu_! x} & \bullet & \longrightarrow & X \\
s^* f \downarrow & \lrcorner & \downarrow t^* f & \lrcorner & \downarrow f \\
\bullet & \xrightarrow{\mu_! y} & \bullet & \longrightarrow & Y \\
s^* y \downarrow & \lrcorner & \downarrow t^* y & \lrcorner & \downarrow y \\
C & \xrightarrow{\rho} & B & \xrightarrow{t} & A
\end{array}
\quad \begin{array}{c} \curvearrowright x \end{array}$$

By the calculus of mates we also have a reversed mate between the right adjoints $\mu^* : t_* \rightarrow s_* \circ \rho^*$. Explicitly μ^* is the composition

$$t_* \xrightarrow{\text{unit } t_*} s_* \rho^* \rho_! s^* t_* \xrightarrow{s_* \rho^* \mu_! t_*} s_* \rho^* t^* t_* \xrightarrow{s_* \rho^* \text{ counit}} s_* \rho^*$$

Definition 2.0.8 (Contravariant action of P_- on a slice). Let $P_- : (\mathbb{C}/A)^{\text{op}} \rightarrow [\mathbb{C}, \mathbb{C}]$ be defined by taking $s \mapsto P_s$ on objects and act on a morphism by

$$\begin{array}{ccc}
\begin{array}{ccc} & B & \\ t \swarrow & \uparrow \rho & \\ A & & C \\ s \swarrow & & \end{array} & \longmapsto & \begin{array}{c} P_t \\ \downarrow \rho^* \\ P_s \end{array}
\end{array}$$

where

$$\rho^* := A_!(s_* \eta \circ \mu B^*) : P_t \rightarrow P_s$$

$$\begin{array}{ccc}
& \mathbb{C} & \\
& \downarrow C^* & \searrow B^* \\
& \mathbb{C}/C \xleftarrow{\rho^*} \mathbb{C}/B & \\
& \downarrow s_* & \swarrow t_* \\
& \mathbb{C}/A & \\
& \downarrow A_! & \\
& \mathbb{C} &
\end{array}
\begin{array}{c}
\curvearrowright P_s \quad \quad \quad \curvearrowleft P_t
\end{array}$$

where $\mu = \mu^*$ is the mate from ??, and η is the natural isomorphism given by pullback pasting.

Pointwise, this natural transformation acts on a pair $(\alpha, \beta) : \Gamma \rightarrow P_t X$ by

$$\begin{array}{ccc}
\Gamma & \xrightarrow{(\alpha, \beta)} & P_t X \\
& \searrow (\alpha, \beta \circ \alpha^* \rho) & \downarrow \rho_X^* \\
& & P_s X
\end{array}$$

where $\alpha^* \rho$ is defined as

$$\begin{array}{ccc}
\Gamma \cdot_s \alpha & \xrightarrow{s^* \alpha} & C \\
\alpha^* \rho \downarrow & \lrcorner & \downarrow \rho \\
\Gamma \cdot_t \alpha & \xrightarrow{t^* \alpha} & B \\
\downarrow & \lrcorner & \downarrow t \\
\Gamma & \xrightarrow{\alpha} & A
\end{array}$$

We prove this now.

Proof. Firstly $\rho_X^* = A_!(s_* \eta_X \circ \mu_{B^* X})$, so the first component $\alpha : \Gamma \rightarrow A$ is preserved by ρ_X^* and it suffices to show, in \mathbb{C}/A

$$\begin{array}{ccc}
\alpha & \xrightarrow{(\alpha, \beta)} & t_* B^* X \\
& \searrow (\alpha, \beta \circ \alpha^* \rho) & \downarrow s_* \eta_X \circ \mu_{B^* X} \\
& & s_* C^* X
\end{array}$$

By the adjunction $s^* \dashv s_*$, it suffices to show, in \mathbb{C}/C

$$\begin{array}{ccc}
s^* \alpha & \xrightarrow{s^* (\alpha, \beta)} & s^* t_* B^* X \\
& \searrow (\alpha, \beta \circ \alpha^* \rho) & \downarrow \overline{s_* \eta_X \circ \mu_{B^* X}} \\
& & C^* X
\end{array}$$

Now we calculate $\overline{s_*\eta_X \circ \mu_{B^*X}} = \eta_X \circ \overline{\mu_{B^*X}}$. So that our goal is to show

$$\begin{array}{ccc} s^*\alpha & \xrightarrow{s^*(\alpha,\beta)} & s^*t_*B^*X \xrightarrow{\overline{\mu_{B^*X}}} \rho^*B^*X \\ & \searrow (\alpha,\beta \circ \alpha^*\rho) & \swarrow \sim_{\eta_X} \\ & & C^*X \end{array}$$

Since η_X is an isomorphism between two limits of the same diagram, namely $X \times C \cong C_!C^*X \cong C_!\rho^*B^*X$, it suffices to show that both $\overline{\mu_{B^*X}} \circ s^*(\alpha,\beta)$ and $(\alpha,\beta \circ \alpha^*\rho)$ are uniquely determined by the same two maps into X and C .

By the characterising property of polynomial endofunctors (??) we calculate

$$\overline{(\alpha,\beta \circ \alpha^*\rho)} = (\beta \circ \alpha^*\rho, s^*\alpha)$$

$$\alpha \xrightarrow{(\alpha,\beta \circ \alpha^*\rho)} s_*C^*X \quad \parallel \quad s^*\alpha \xrightarrow{(\alpha,\beta \circ \alpha^*\rho)}_{(\beta \circ \alpha^*\rho, s^*\alpha)} C^*X \quad \parallel \quad C_!s^*\alpha \xrightarrow{\beta \circ \alpha^*\rho} X$$

More formally, this means $\beta \circ \alpha^*\rho : C_!s^*\alpha \rightarrow X$ and $s^*\alpha : C_!s^*\alpha \rightarrow C$ are the two maps that uniquely determine the map $C_!\overline{\alpha,\beta \circ \alpha^*\rho} : C_!s^*\alpha \rightarrow X \times C$.

On the other hand,

$$\begin{array}{c} \alpha \xrightarrow{(\alpha,\beta)} t_*B^*X \xrightarrow{\text{unit}_{t_*B^*X}} s_*\rho^*\rho_!s^*t_*B^*X \xrightarrow{s_*\rho^*\mu_!t_*B^*X} s_*\rho^*t^*t_*B^*X \xrightarrow{s_*\rho^*\text{counit}_{B^*X}} s_*\rho^*B^*X \\ \searrow \mu_{B^*X} \nearrow \\ \hline s^*\alpha \xrightarrow{s^*(\alpha,\beta)} s^*t_*B^*X \xrightarrow{\overline{\text{unit}_{t_*B^*X}}} \rho^*\rho_!s^*t_*B^*X \xrightarrow{\rho^*\mu_!t_*B^*X} \rho^*t^*t_*B^*X \xrightarrow{\rho^*\text{counit}_{B^*X}} \rho^*B^*X \\ \searrow \overline{\mu_{B^*X}} \nearrow \\ \hline \rho_!s^*\alpha \xrightarrow{\rho_!s^*(\alpha,\beta)} \rho_!s^*t_*B^*X \xrightarrow{\overline{\text{unit}_{t_*B^*X}}} \rho_!s^*t_*B^*X \xrightarrow{\mu_!t_*B^*X} t^*t_*B^*X \xrightarrow{\text{counit}_{B^*X}} B^*X \\ \parallel \quad \parallel \quad \parallel \quad \parallel \\ \Gamma \cdot_s \alpha \quad S \quad R \quad X \times B \end{array}$$

The mate $\mu_!$ is calculated via the universal map into the pullback R (dotted below).

$$\begin{array}{ccccc}
\Gamma \cdot_s \alpha & \longrightarrow & \Gamma \cdot_t \alpha & \longrightarrow & \Gamma \\
\downarrow s^*(\alpha, \beta) & & \downarrow & & \downarrow (\alpha, \beta) \\
S & \overset{\text{dotted}}{\longrightarrow} & R & \longrightarrow & P_t X \\
\downarrow s^* t_* B^* X & & \downarrow \mu_! t_* B^* X & & \downarrow t_* B^* X \\
C & \xrightarrow{\rho} & B & \xrightarrow{t} & A
\end{array}
\begin{array}{l}
\swarrow s^* \alpha \\
\searrow \alpha
\end{array}$$

Using the characterization of maps into R from ?? we can calculate

$$\mu_! t_* B^* X \circ s^*(\alpha, \beta) = (\rho \circ s^* \alpha, \beta \circ t^* \alpha^* s)$$

since the first component is simply the map $\Gamma \cdot_s \alpha \rightarrow B$ and the second component is the second component of the map

$$(\alpha \circ \alpha^* s, \beta \circ t^* \alpha^* s) = (\alpha, \beta) \circ \alpha^* s : \Gamma \cdot_s \alpha \rightarrow P_t X$$

Then using ??

$$\overline{\mu_{B^* X} \circ s^*(\alpha, \beta)} \quad (2.1)$$

$$= \text{counit}_{B^* X} \circ \mu_! t_* B^* X \circ s^*(\alpha, \beta) \quad (2.2)$$

$$= \text{counit}_{B^* X} \circ (\rho \circ s^* \alpha, \beta \circ t^* \alpha^* s) \quad (2.3)$$

$$= (\beta \circ t^* \alpha^* s \circ r, \rho \circ s^* \alpha) \quad (2.4)$$

$$= (\beta \circ \alpha^* \rho, \rho \circ s^* \alpha) \quad (2.5)$$

$$: \Gamma \cdot_s \alpha \rightarrow X \times B \quad (2.6)$$

where

$$\begin{array}{ccc}
\Gamma \cdot_s \alpha & \xrightarrow{\rho \circ s^* \alpha} & B \\
\downarrow r & \searrow & \downarrow t \\
\Gamma \cdot_s \alpha \cdot_t \alpha \circ \alpha^* s & \longrightarrow & B \\
\downarrow & & \downarrow \\
\Gamma \cdot_s \alpha & \xrightarrow{\alpha \circ \alpha^* s} & A
\end{array}$$

and

$$\left(\begin{array}{ccccc}
\Gamma \cdot_s \alpha & \xlongequal{\quad} & \Gamma \cdot_s \alpha & \xrightarrow{s^* \alpha} & C \\
\downarrow r & & \downarrow \alpha^* \rho & & \downarrow \rho \\
\Gamma \cdot_s \alpha \cdot_t \alpha \circ \alpha^* s & \overset{\text{dotted}}{\longrightarrow} & \Gamma \cdot_t \alpha & \longrightarrow & B \\
\downarrow & & \downarrow & & \downarrow t \\
\Gamma \cdot_s \alpha & \xrightarrow{\alpha^* s} & \Gamma & \xrightarrow{\alpha} & A
\end{array} \right) s$$

Moving back along the adjunction $\rho_! \dashv \rho^*$?? tells us that

$$\begin{array}{ccccc}
 & & s^* \alpha & & \\
 & \Gamma \cdot_s \alpha & \xrightarrow{\quad} & X \times C & \longrightarrow C \\
 & \searrow \overline{\mu_{B^* X} \circ s^*}(\alpha, \beta) & \searrow & \downarrow \lrcorner & \downarrow \rho \\
 & \searrow \overline{\mu_{B^* X} \circ s^*}(\alpha, \beta) & \searrow & X \times B & \longrightarrow B \\
 & \searrow \beta \circ \alpha^* \rho & \searrow & \downarrow \lrcorner & \downarrow \\
 & & & X & \longrightarrow 1
 \end{array}$$

So that, as required, $\overline{\mu_{B^* X} \circ s^*}(\alpha, \beta)$ and $\overline{(\alpha, \beta \circ \alpha^* \rho)}$ are uniquely determined by the same two maps into X and C . \square

Definition 2.0.9 (Covariant action of P_- on a cartesian square). We can also view taking polynomial endofunctors as a covariant functor on the category of arrows with cartesian squares as morphisms

$$P_- : \text{CartArr}(\mathbb{C}) \rightarrow [\mathbb{C}, \mathbb{C}]$$

where the action on a cartesian square is

$$\begin{array}{c}
 \begin{array}{ccc}
 \mathbb{C} & \xlongequal{\quad} & \mathbb{C} \\
 D^* \downarrow & \xRightarrow{\eta^{-1}} & \downarrow B^* \\
 \mathbb{C}/D & \xleftarrow{\rho^*} & \mathbb{C}/B \\
 s_* \downarrow & \xRightarrow{\mu^{*-1}} & \downarrow t_* \\
 \mathbb{C}/C & \xleftarrow{\theta^*} & \mathbb{C}/A \\
 C_! \downarrow & \xRightarrow{\mu_!} & \downarrow A_! \\
 \mathbb{C} & \xlongequal{\quad} & \mathbb{C}
 \end{array} \\
 \left(\begin{array}{c} \downarrow \\ \downarrow \\ \downarrow \\ \downarrow \end{array} \right) P_s \quad \left(\begin{array}{c} \downarrow \\ \downarrow \\ \downarrow \\ \downarrow \end{array} \right) P_t
 \end{array}$$

given by the whiskered natural transformations

$$\begin{array}{c}
 \begin{array}{ccc}
 \mathbb{C} & \xlongequal{\quad} & \mathbb{C} \\
 C^* \downarrow & \xRightarrow{\eta^{-1}} & \downarrow B^* \\
 \mathbb{C}/D & \xleftarrow{\rho^*} & \mathbb{C}/B \\
 s_* \downarrow & \xRightarrow{\mu^{*-1}} & \downarrow t_* \\
 \mathbb{C}/C & \xleftarrow{\theta^*} & \mathbb{C}/A \\
 C_! \downarrow & \xRightarrow{\mu_!} & \downarrow A_! \\
 \mathbb{C} & \xlongequal{\quad} & \mathbb{C}
 \end{array} \\
 \left(\begin{array}{c} \downarrow \\ \downarrow \\ \downarrow \\ \downarrow \end{array} \right) P_s \quad \left(\begin{array}{c} \downarrow \\ \downarrow \\ \downarrow \\ \downarrow \end{array} \right) P_t
 \end{array}$$

Furthermore, the natural transformation P_κ is cartesian. meaning each naturality square is a pullback square.

$$\begin{array}{ccc} P_s X & \xrightarrow{P_{\kappa Y}} & P_t X \\ P_s f \downarrow & \lrcorner & \downarrow P_s f \\ P_s Y & \xrightarrow{P_{\kappa Y}} & P_t Y \end{array}$$

The natural transformation P_κ computes in the following way

$$\begin{array}{ccc} \Gamma \cdot_t \theta \circ \alpha & & \\ \swarrow \text{dashed } i & \searrow & \\ \Gamma \cdot_s \alpha & \xrightarrow{\quad} & D \xrightarrow{\rho} B \\ \downarrow & \lrcorner & \downarrow s \lrcorner \downarrow t \\ \Gamma & \xrightarrow{\alpha} & C \xrightarrow{\theta} A \end{array} \quad \begin{array}{ccc} \Gamma & & \\ (\alpha, \beta) \downarrow & \searrow (\theta \circ \alpha, \beta \circ i) & \\ P_s X & \xrightarrow{P_{\kappa X}} & P_t X \end{array}$$

using the fact that $\Gamma \cdot_s \alpha$ and $\Gamma \cdot_t \theta \circ \alpha$ are limits of the same diagram.

Proof. We can use the computation of $P_{\kappa X}$ and $P_s f$ to show that the natural transformation P_κ is cartesian. Essentially, the first component of a map $\Gamma \rightarrow P_s X$ is determined by its composition with $P_s f$ and its second component is determined by its composition with $P_{\kappa X}$. \square

Corollary 2.0.10. *If we have*

$$\begin{array}{ccc} D' & \longrightarrow & B' \\ \left(\begin{array}{ccc} \downarrow \rho_1 \lrcorner & & \downarrow \rho_2 \\ D & \longrightarrow & B \\ \downarrow q_1 \lrcorner & & \downarrow q_2 \end{array} \right) & & \\ C & \xrightarrow{\theta} & A \end{array}$$

then the two possible ways of obtaining composing the covariant and contravariant actions of P_- form a (strictly commuting) pullback square in $[\mathbb{C}, \mathbb{C}]$.

$$\begin{array}{ccc} P_{q_1} & \xrightarrow{P_\kappa} & P_{q_2} \\ \rho_1^* \downarrow & \lrcorner & \downarrow \rho_2^* \\ P_{q'_1} & \xrightarrow{P_{\kappa'}} & P_{q'_2} \end{array}$$

Proof. To check that it commutes and is a pullback, it suffices to do this pointwise, for some $X \in \mathbb{C}$. Then we simply unfold the computation for each of P_κ and ρ^* . \square

Chapter 3

Multivariate Polynomial Functors

To be written