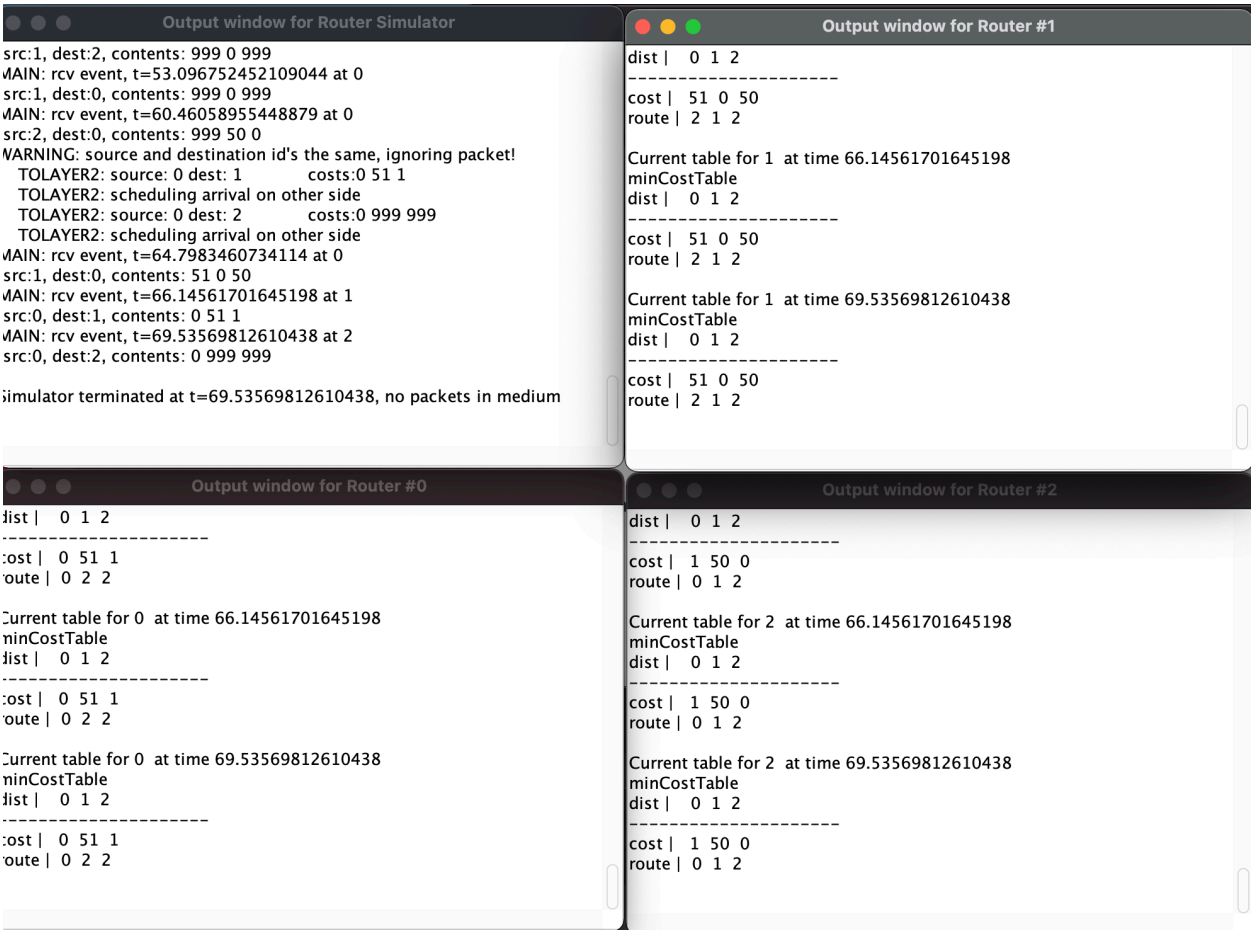
CPSC 441 A4 Report

1. Distance vector routing determines the best route to take based on distance. In order to determine the best route to take, it must take into account the cost of each path that it can take to get to a particular node and compute the distance. To calculate the best route, the Bellman-Ford algorithm is used. In the solution, it calculates the shortest distances that have at most one edge in the path. It then calculates the shortest paths with at most 2 edges and so on. The distances are summed for each path that the source vector needs to reach. The minimum cost is then taken into account for all vector routes, that the source vector can take to reach the neighbouring nodes. The algorithm must initialize all distances from the source to all vertices as infinite and distances to the source and themselves as 0. Then a distance vector must be created and record the values. To calculate the shortest distances, at each edge, we must determine if the cost at the edge is greater than the cost at the other edge plus the weight. If that is the case, then you must update the cost. This almost works like dijkstra's algorithm, but is less efficient in terms of runtime. With this solution, we also have a poisoned reverse to ensure that we don't reach the count to infinity problem.

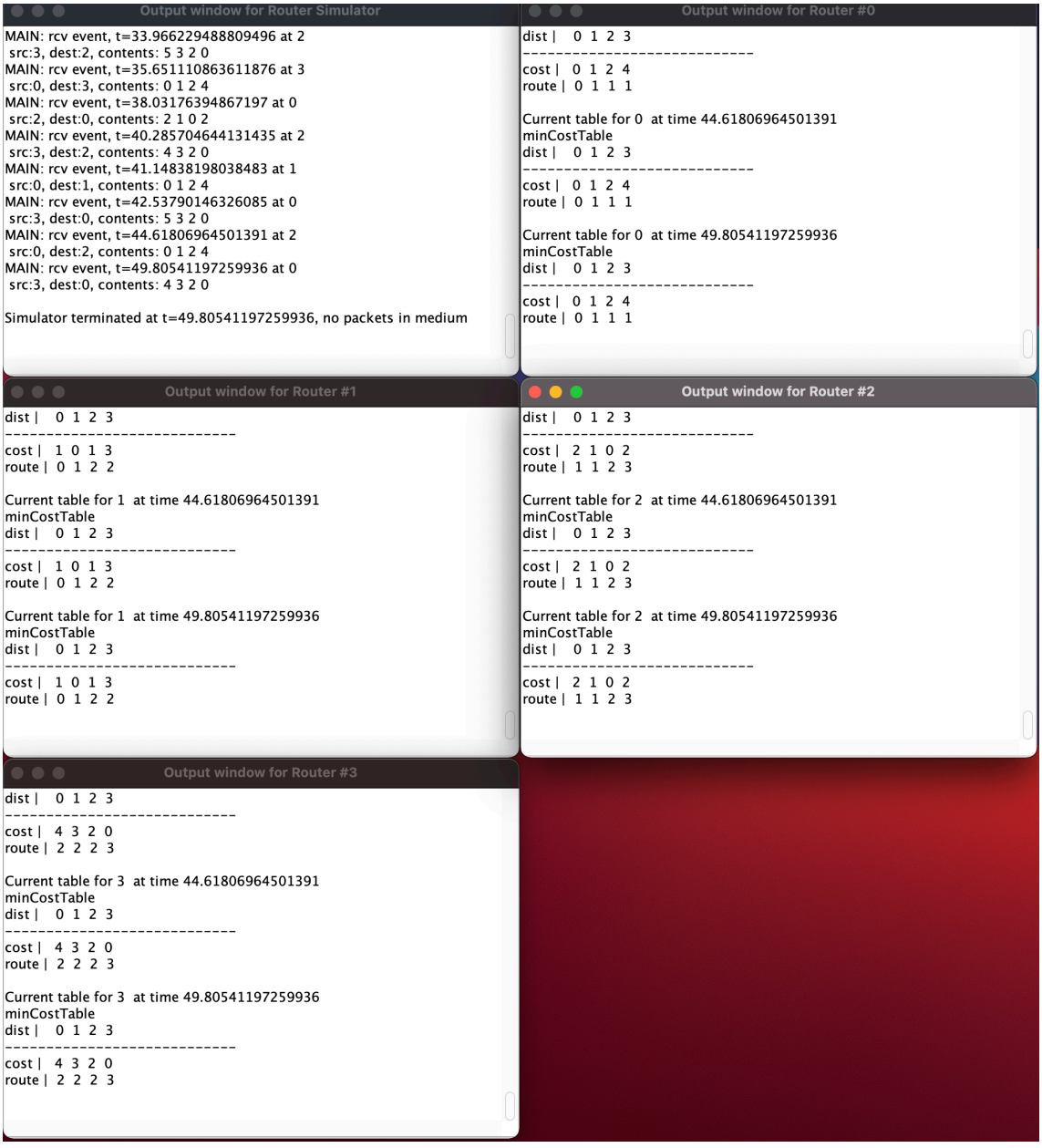
2. The solution was tested by running the three tests as per the assignment description, as well as comparing the values computed with the values that were computed by the instructor and the TA. All of which resulted in the same values. Additionally, I had also calculated some of the values by hand, tracing the costs and distances to ensure accuracy.

3.

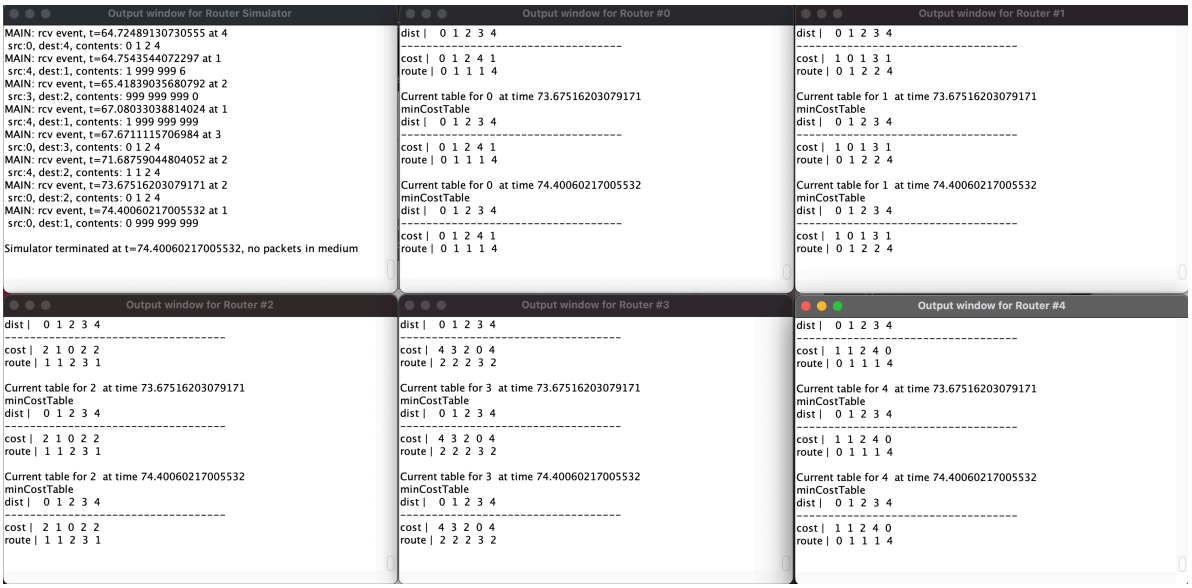
Test case: make install3



Test case: make install4

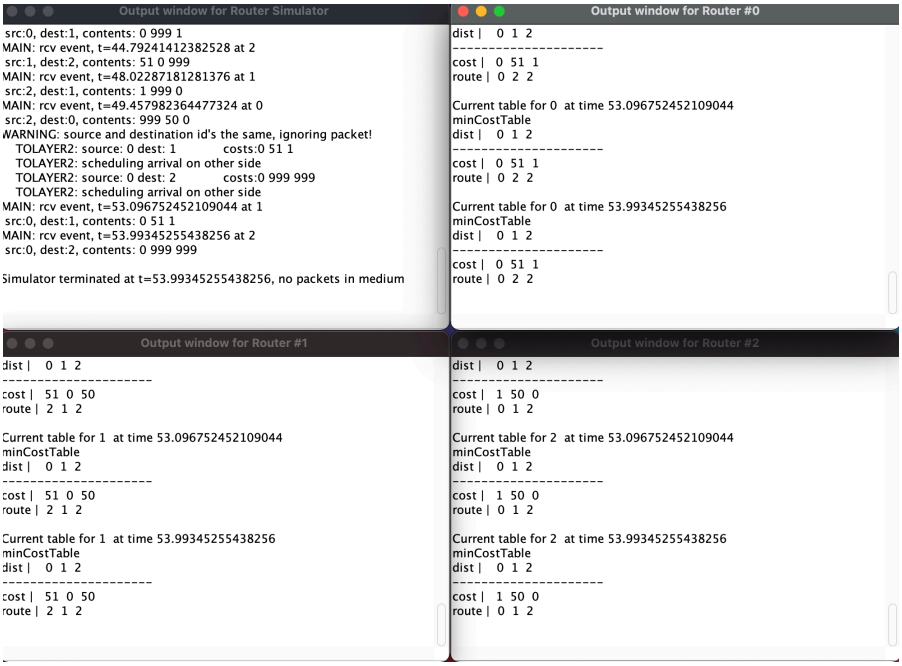


Test case: make install5

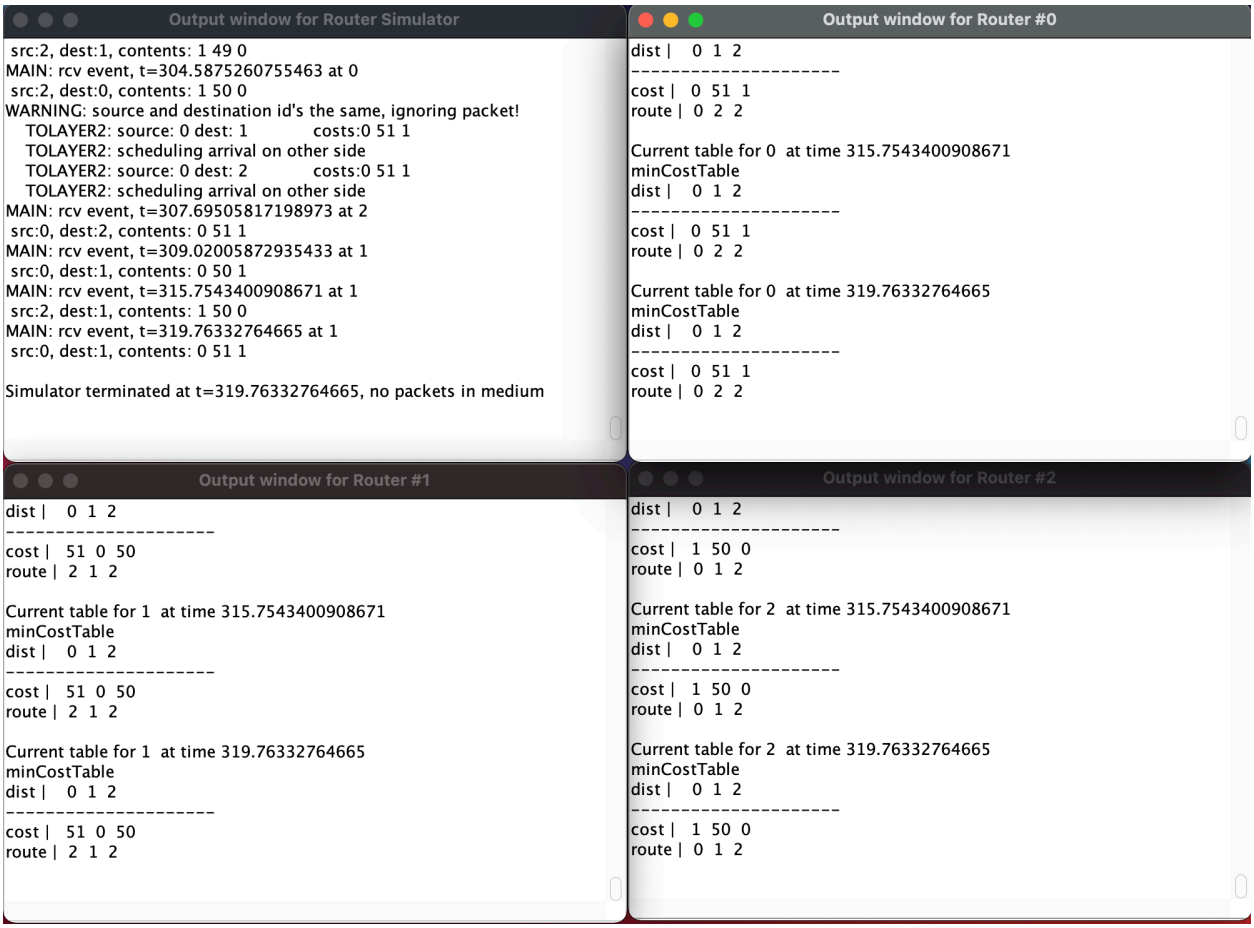


4.

Poisoned Reverse On



Poisoned Reverse Off



From the test case above, we can see that poisoned reverse is necessary because the Bellman-Ford algorithm does not have a solution to get rid of routing loops. In the figure with the poisoned reverse toggled off, we can see a much higher time. This is due to the fact the program increments the cost by two for each iteration, resulting in a much longer time.

5. Poisoned reverse can communicate with the other nodes, indicating that it has reached a case in which it continuously loops back upon itself creating a count to infinity problem. When it reaches this, the node that that is affected will no longer route to the node in which the loop was created. This gets rid of the routing loop when an update is successfully transmitted. However, poison reverse is not capable of getting rid of loops greater than two.