

14/12/16
Masterclass, Angew

Hilbert Schemes

i.e. assume reducedness etc.

we gloss over schemey details

(1 was later) (wow)

?

Proposition: if p, p_2, p_3, p_4 are general points then $\text{codim}(I_2(X), S_2) = 4$

Proposition: if p_1, \dots, p_4 are on a line then $\text{codim}(I_2(X), S_2) = 3$

Problem: find a uniform degree which works for all X

proof: Suppose the line is $y=0$
the curve meets all four points and so contains the line $y=0$
possible equations: $y(ax+by+cz)$
so $\dim I_2 = 3$

question: $X(t) = (0,0) \vee (t,0)$
what is $\lim_{t \rightarrow 0} X(t)$?

if we define it to be $(0,0)$ then we lose a point (so no hope of continuity)

well $I(X(t)) = (y, (x-t)x)$
so $I(X_0) = (y, x^2)$

and $\dim k[x,y]/(y, x^2) = \dim k[x]/(x^2) = 2$

problem: need to consider ideals and not only geometric objects
i.e. schemes not varieties

Intuition: $\text{Hilb}^p(A^2) = \{ I \triangleleft k[x,y] \mid \dim k[x,y]/I = p \}$
 $\text{Hilb}^p(A^2) = \{ I \triangleleft k[x,y] \mid \dim k[x,y,z]/(I_d) = p \ \forall d \gg 0 \}$
lim and saturated

Proposition: $I \triangleleft S$ then $\exists P_I$ polynomial
and P_I is called the Hilbert polynomial of I
 $\text{codim}(I_d, S_d) = P_I(d)$
 $\forall d \gg 0$

for all polynomials $P \in k[x_0, \dots, x_n]$
 $(x_0 P, \dots, x_n P \in I \Rightarrow P \in I)$
(here $x_0 = x_0, y = x_1$)

example: $P^2 \subset P^4$, $P_I(d) = \binom{d+2}{2}$
(no of degree d monomials in three variables)

example: $P^2 \subset P^4$ then $P_I(P^2) = d+1$
Since $I(P^2) = (x_2, \dots, x_4)$ so
 $I_d(P^2) = \dots$? well
 $k[x_0, \dots, x_4]/I_d(P^2) = \frac{k[x_0^d, x_0^{d-1}x_1, \dots, x_1^d]}{d+1 \text{ elements}}$

Answer (to something): $I_1 \sim I_2 \iff P_{I_1} = P_{I_2}$

New question: what is / how do we construct $\text{Hilb}^p(P^n) = \{ I \subset S, P_I = p \}$

Proposition: $\{ C \subset P^n \mid C \text{ curve} \}$

Suppose $\exists P$ polynomial $\mathbb{A}^1 \rightarrow \mathbb{Z} \subset C \times \mathbb{A}^1 \times P^n$
and $\mathbb{Z} \subset \mathbb{A}^1 \times P^n$ closed

s.t. $\forall q \in C, P_I(C_q) = p$
If C smooth (near x_0) then

$\exists \tilde{C} \subset C \times P^n$ s.t. $P_I(\tilde{C}_x) = p$
and \tilde{C} extends \mathbb{Z}
(and \tilde{C} is the schematic closure of \mathbb{Z})

• unconstrained problem

$P = p_1(d) + \dots + p_r(d)$ is possible (\mathbb{P}^n)

what are the possible Hilbert polynomials?

• Proposition: if r, d, \dots, r_k then (by looking at \mathbb{P}^n)

example: $3d+1 = (d+1) + d + (d-1) + 1$
(can realize this $\sim \mathbb{P}^2$) } $P = p_1(d) + \dots + p_r(d)$ is possible (and \therefore true!)
so this is a Hilbert polynomial

• Definition: no. of terms in these sets of decompositions is the Gordan number of \mathbb{P} (imposed)

• example: $d^2 - d - 4 = \binom{d+2}{2} + \binom{d+1}{2} + 2d + \mu$ with $\mu \leq 0$

So NOT a Hilbert polynomial

• Theorem: $I \subseteq S$, $P = \sum p_i$
 $r =$ Gordan number of \mathbb{P}
then $\text{codim}(I_d, S_d) = P(d)$ for all $d \geq r$

• Definition: as a set,
 $\text{Hilb}^P(\mathbb{P}^n) = \{ I \subseteq S \mid P_I = P \}$

• Theorem:
 $\text{Hilb}^P(\mathbb{P}^n) \xrightarrow{(d, \mu) \mapsto P} G(\mathbb{P}^n, S_r)$
 $I \mapsto I_r$

is injective and $\{I_r\}$ can be defined by equations (in fact, determinants) (but also more recent, simpler ways) (over degree)

~~unconstrained problem~~
~~is equivalent to Hilb(\mathbb{P}^n)~~