

Problem

1. Minimize the following DFA using table filling algorithm where A is the start state. The states C, F and I are the final states.

δ	0	1
A	B	E
B	C	F
*C	D	H
D	E	H
E	F	I
F	G	B
G	H	B
H	I	C
*I	A	E

Table filling algorithm

	A	B	*C	D	E	*F	G	H
B								
*E	X	X						
D			X					
E			X					
*F	X	X		X	X			
G			X			X		
H			X			X		
*I	X	X		X	X	*F	G	H

Step 1: Cross the combinations of Final and non-final states.
 (A,C) (A,F) (A,I) (B,C) (B,F) (B,I) (C,H) (C,G) (C,D) (C,E) (D,I)
 (D,F) (E,I) (E,F) (F,G) (F,H) (G,I) (H,I) cannot be
 equivalent states.

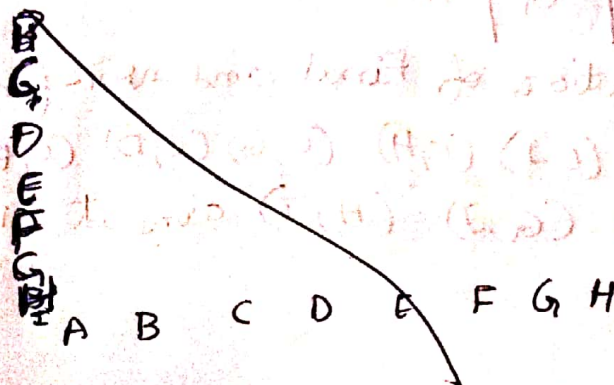
Boxes where combinations of final & final, non-final and non-final are left open.

Step 2:

Check the o/p and i/p combinations of A and B to start with.

δ	O	I
X(A,B)	(B,C)	(E,F)
(A,D)	(B,E)	(E,H)
X(A,E)	(B,F)	(E,I)
(A,G)	(B,H)	(E,B)
X(A,H)	(B,I)	(E,C)
(B,H)	(C,I)	(F,C)
X(B,G)	(C,H)	(F,B)
X(B,D)	(C,E)	(F,H)
(B,E)	(C,F)	(F,I)
(C,F)	(D,G)	(H,B)
(C,I)	(D,A)	(H,E)
X(D,H)	(E,I)	(H,C)
(D,G)	(E,H)	(H,B)
X(E,G)	(F,H)	(I,B)
(E,H)	(F,I)	(I,C)
(F,I)	(G,A)	(B,E)
X(G,H)	(H,I)	(B,C)

Table



B	X								
*C	X	X							
D			X	X					
E	X			X					
*F	X	X			X	X			
G			X	X		X	X		
H	X			X	X			X	
*I	X	X			X	X			X
A									X

Hence the remaining pairs are equivalent

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δ	0	1
(A,D)	(B,F)	(E,H)
(A,G)	(B,H)	(E,B)
(B,H)	(C,I)	(F,C)
(B,E)	(C,F)	(F,I)
* (C,F)	(D,G)	(H,B)
* (C,I)	(D,A)	(H,E)
(D,G)	(E,H)	(H,B)
(E,H)	(F,I)	(I,C)
(F,I)	(G,A)	(B,E)

From the above table

$$A = D \neq D = G \neq A = G$$

$$A = D = G$$

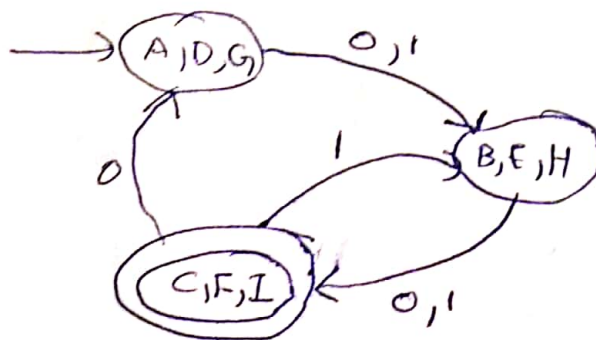
$$B = E \neq E \neq H \neq B = H$$

$$B = E = H$$

$$\Rightarrow C = I = F$$

δ	0	1
\Rightarrow (A,D,G)	(B,E,H)	(B,F,H)
(B,E,H)	(C,F,I)	(C,F,I)
(C,F,I)	(A,D,G)	(E,B,H)

Transition diagram of minimised DFA



Problem 2:

Consider the DFA given by the transition

δ	0	1
$\rightarrow q_1$	q_2	q_3
q_2	q_3	q_5
$*q_3$	q_4	q_3
q_4	q_3	q_5
$*q_5$	q_2	q_5

- (a) Draw the table of distinguishabilities for this automation.
 (b) Construct the minimum state equivalent DFA.

Table filling algorithm

q_2				
$*q_3$	X	X		
q_4			X	
$*q_5$	X	X		X
	q_1	q_2	$*q_3$	q_4

Step 1: Cross the combinations of final and non-final states
 (q_1, q_5) (q_1, q_3) (q_2, q_5) (q_2, q_3) (q_3, q_4) (q_4, q_5) cannot be
 equivalent states

Boxes where combinations of final & final, non-final and non-final are left open.

Step 2:

Check the a/p and i/p combinations of A and B to start with.

	δ	0	1
$\times (q_1, q_2)$	(q_2, q_3)	(q_3, q_5)	
$\times (q_1, q_4)$	(q_2, q_3)	(q_3, q_5)	
	(q_2, q_4)	(q_3, q_3)	(q_5, q_5)
	(q_3, q_5)	(q_4, q_2)	(q_3, q_5)

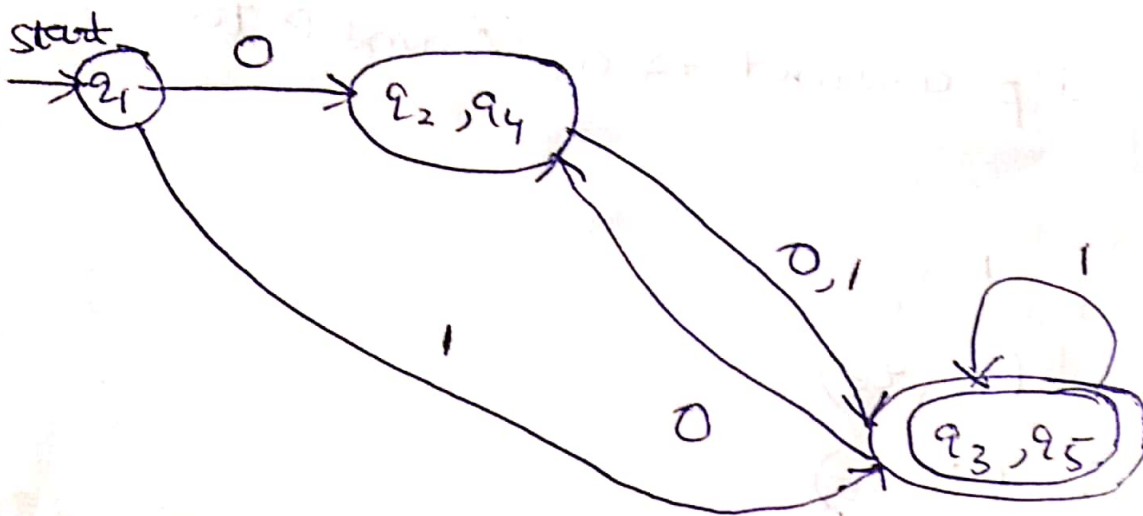
q_2	X			
* q_3	X	X		
q_4	X		X	
* q_5	X	X		X
	q_1	q_2	* q_3	q_4

Step 3:-

Therefore only (q_2, q_4) and (q_3, q_5) form the equivalents for the given transition.

	δ	0	1
\rightarrow	(q_2, q_4)	(q_3, q_5)	(q_5, q_5)
	* (q_3, q_5)	(q_4, q_2)	(q_3, q_5)

Reduced DFA



Problem 3

Repeat exercise 4.4.1 for the DFA

δ	0	1
$\rightarrow q_1$	q_2	q_6
q_2	q_1	q_3
$*q_3$	q_2	q_4
q_4	q_4	q_2
q_5	q_4	q_5
$*q_6$	q_5	q_4

Table filling algorithm

q_2	* x					
* q_3	x	x				
q_4	* x	x	x			
q_5	* x	x	x	x		
* q_6	x	x	x	x	x	
	q_1	q_2	* q_3	q_4	q_5	

Step 1: Cross the combinations of final and non-final states.

(q_1, q_2) (q_1, q_4) ~~(q_1, q_5)~~ (q_2, q_3) (q_2, q_5) (q_3, q_4) (q_3, q_5)

(q_4, q_6) (q_5, q_6) cannot be equivalents.

Step 2:

δ	\emptyset	1
(q_1, q_2)	(q_2, q_1)	(q_6, q_3) x
(q_1, q_4)	(q_2, q_4)	(q_6, q_2) x
(q_1, q_5)	(q_2, q_4)	(q_6, q_5) x
(q_2, q_4)	(q_1, q_4)	(q_3, q_2) x
(q_2, q_5)	(q_1, q_4)	(q_3, q_5) x
(q_3, q_6)	(q_2, q_5)	(q_4, q_4) x
(q_4, q_5)	(q_4, q_4)	(q_2, q_5) x

After step 2 and step 3 we can say there are no equivalent states for the given table