www.locuz.com **Shared Memory Programming with OpenMP** Presented By: Mandeep Kumar Converge to the Cloud

Parallelism







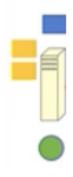
Why Parallelism Is Good

- > The Trees: We like parallelism because, as the number of processing units working on a problem grows, we can solve the same problem in less time
- > The Forest: We like parallelism because, as the number of processing units working on a problem grows, we can solve bigger problems

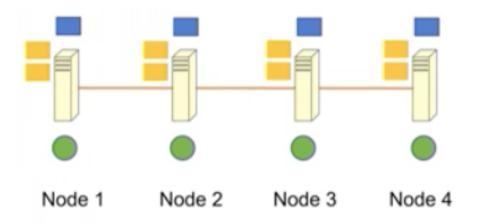


Shared Memory Vs Distributed Memory Model

Shared Memory Model



Distributed Memory Model





Process

Memory



OpenMP vs MPI

OpenMP	MPI
Shared memory model	Distributed memory model
On Multi-core processors	On Distributed network
Directive based	Message based
Easier to program and debug	Flexible and expressive



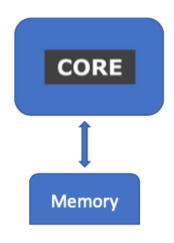
Agenda

- Core concepts
- What is OpenMP
- OpenMP Programming and Execution Model
- OpenMP Constructs
- Granularity of Parallelization
- Advantages and Disadvantages of OpenMP

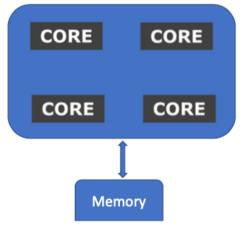


Basis System Architecture

Single Core Processor



Multi Core Processor





Sequential Program Execution

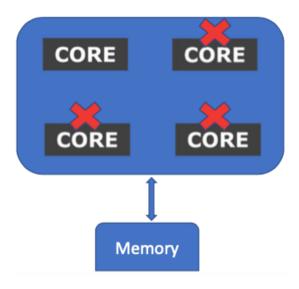
When you run sequential program

- Instructions executed in serial
- Other cores are idle

Waste of available resource... We want all core to be used to execution program.

How?

Multi Core Processor





Process and Thread

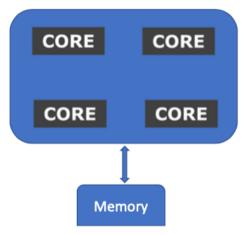
- An executing instance of a program is called a process
- Process has its independent memory space
- A thread is a subset of the process also called lightweight process allowing faster context switching
- Threads share memory space within process's memory
- Threads may have some (usually small) private data
- A thread is an independent instruction stream, thus allowing concurrent operation
- In OpenMP one usually wants no more than one thread per core



Shared Memory Model

- Multiple threads operate independently but share same memory resources
- Data is not explicitly allocated
- Changes in a memory location effected by one process is visible to all other processes
- Communication is implicit
- Synchronization is explicit

Multi Core Processor





What is OpenMP?

- Open Specification for Multi Processing
- Provides multi-threaded parallelism
- OpenMP is an Application Program Interface (API) for writing multi-threaded, shared memory parallelism.
- It is an specification for
 - Compiler Directives
 - Runtime Library Routines
 - Environment Variables
- Easy to create multi-threaded programs in C, C++ and Fortran.



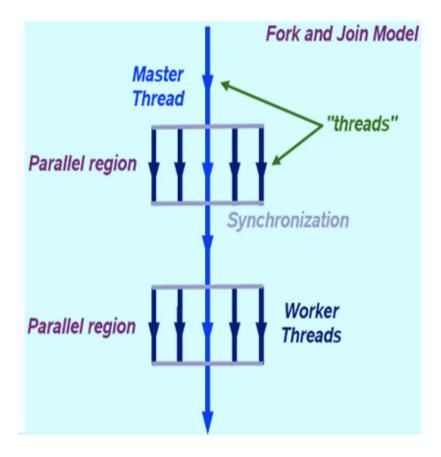
OpenMP History

- 1997: Version 1.0 for Fortran
- 1998: Version 1.0 for C/C++
- 2002-2005: Versions 2.0-2.5, Merger of Fortran and C/C++ specifications
- 2008: Version 3.0, Incorporates Task Parallelism
- 2013: Version 4.0, Support for Accelerators, SIMD support
- 2018: Version 5.0, C11/C++17/Fortran 2008 support



Execution Model

- OpenMP program starts single threaded
- To create additional threads, user starts a parallel region
 - additional threads are launched to create a team
 - original (master) thread is part of the team
 - threads "go away" at the end of the parallel region
- Repeat parallel regions as necessary
 Fork-join model



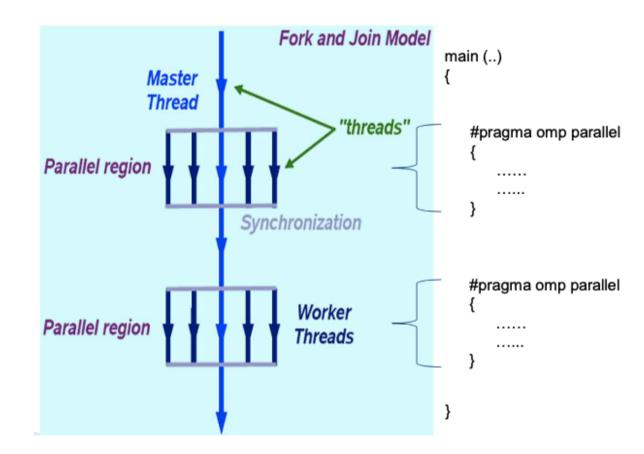


OpenMP Basic Syntax

Header file: #include <omp.h>

```
Parallel Region:

C:
#pragma omp construct [clauses...]
{
    // ... Do some work here
} // end of parallel region/block
```





Parallel Region

Fork a team of N threads {0...N-1}

Without it, all code are sequential

```
// sequential code here (master thread)

#pragma omp parallel [clauses]
{

A team of N threads start here

// parallel computing here

// ...
}

implicit barrier; no thread can progress until all others have reached here
```



OpenMP Directives

- OpenMP directives are comments in source code that specify parallelism for shared memory parallel (SMP)
 machines
- C/C++ compiler directives begin with the sentinel **#pragma omp**
- FORTRAN compiler directives begin with one of the sentinels !\$OMP, C\$OMP, or *\$OMP
 - use **!\$OMP** for free-format F90

C/C++

```
#pragma omp parallel
{
    ...
}

#pragma omp parallel for
  for(...) { ...
}
```

FORTRAN

```
!$OMP parallel
...
!$OMP end parallel
!$OMP parallel do
   Do ...
!$OMP end parallel do
```

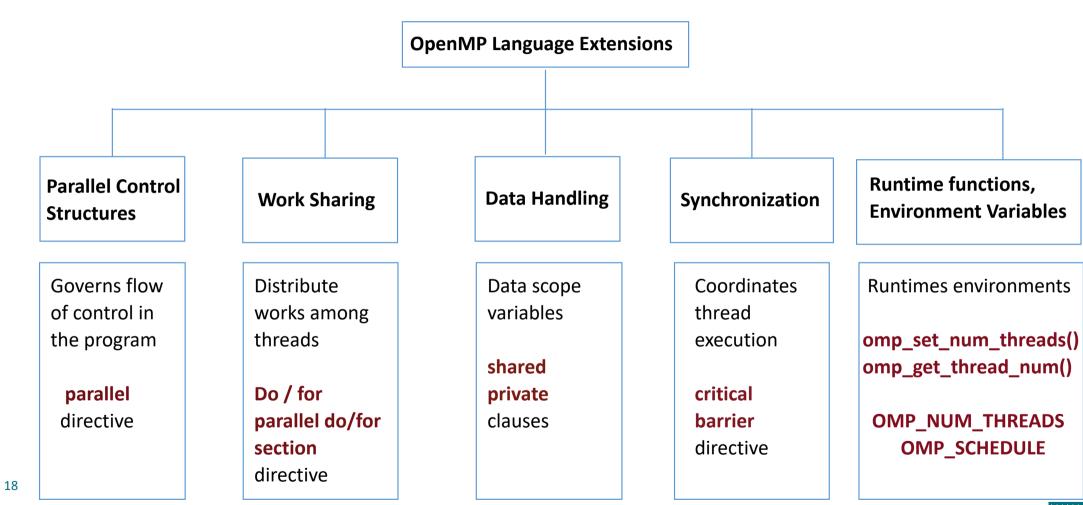


How do Threads Interact?

- Threads read and write shared variable
 - hence communication is implicit
- Unintended sharing of data causes race conditions
 - race condition can lead to different outputs across different runs
- Use synchronisation protect against race conditions
- Synchronization is expensive
 - change data storage attributes for minimizing synchronization and improving cache reuse



OpenMP Constructs





OpenMP Constructs

- Parallel region #pragma omp parallel
- Worksharing #pragma omp for #pragma omp sections

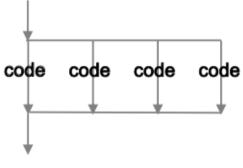
- Data Environment #pragma omp parallel shared/private (...)
- Synchronization
 #pragma omp barrier
 #pragma omp critical



Parallel Directives

- Replicated executed by all threads
- Worksharing divided among threads

```
#pragma omp parallel
    code();
```



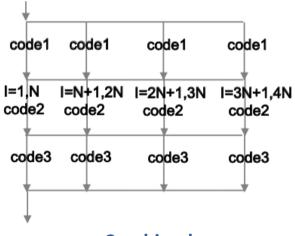
```
Replicated
```

```
#pragma omp for
 for(I=1; I<=4*N; I++)
      code();
```

```
I=1,N I=N+1,2N I=2N+1,3N I=3N+1,4N
code
       code
                 code
                           code
```

Worksharing

```
#pragma omp parallel
    code1();
    #pragma omp for
      for(I=1; I<=4*N; I++)
         code2();
     code3();
 }
```

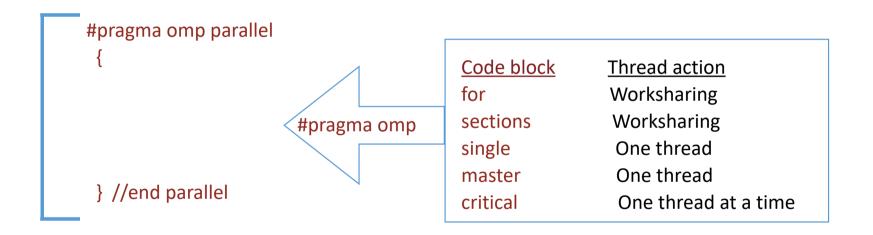


Combined



Worksharing

Use OpenMP directives to specify worksharing in a parallel region, as well as synchronization



parallel do/for parallel sections

Directives can be combined, if a parallel region has just one worksharing construct.



First OpenMP Program

```
#include <stdio.h>
#include <omp.h>

int main()
{
    #pragma omp parallel
    {
        printf("Parallel hello to you!\n");
    }
    return(0);
}
```



First OpenMP Program

- Check your system support: locate omp.h /usr/lib/gcc/x86_64-redhat-linux/4.8.2/include/omp.h
- Compilation: gcc -fopenmp first OpenMP.c

```
• Flags:
```

```
GNU: -fopenmp for Linux, Solaris, AIX, MacOSX, Windows. IBM: -qsmp=omp for Windows, AIX and Linux. Sun: -xopenmp for Solaris and Linux. Intel: -qopenmp on Linux or Mac, /Qopenmp on Windows. PGI: -mp
```

• Conditional Compilation:

```
$ gcc -fopenmp first_OpenMP.c
$ ./a.out
Parallel hello to you!
```



Control the Number of Threads

- Parallel region clause #pragma omp parallel num_threads(integer)
- Run-time function omp_set_num_threads(integer)
- Environment Variable OMP_NUM_THREADS





Setting Environment Variable and Executing

```
$ echo $SHELL //Know your shell
$ /bin/bash
$ export OMP_NUM_THREADS=4 // Setting environment variables
$ gcc -fopenmp first_OpenMP.c //Compilation
$ ./a.out // Execution
Parallel hello to you!
Parallel hello to you!
Parallel hello to you!
Parallel hello to you!
```



Second Example

```
#include <stdio.h>
#include <omp.h>
int main(int argc, char *argv[])
   int id:
   double wtime;
   printf("Number of processors available: %d\n",omp_get_num_procs());
   printf("Number of threads: %d\n".omp get max threads());
   wtime=omp get wtime();
   printf("Outside the parallel region\n");
   id=omp get thread num();
   printf("Hello from process %d\nGoing inside the parallel region\n",id);
   #pragma omp parallel private(id)
      id=omp get thread num();
      printf("Hello from process %d\n".id);
   wtime=omp get wtime() - wtime;
   printf("Back from the parallel region. Normal execution\n");
   printf("Elapsed wall clock time %f\n",wtime);
   return 0;
```



Second Example Output

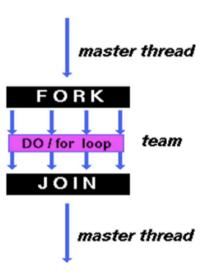
```
Number of processors available: 10
Number of threads: 10
Outside the parallel region
Hello from process 0
Going inside the parallel region
Hello from process 8
Hello from process 5
Hello from process 4
Hello from process 1
Hello from process 2
Hello from process 7
Hello from process 6
Hello from process 9
Hello from process 0
Hello from process 3
Back from the parallel region. Normal execution
Elapsed wall clock time 0.006727
```

```
$ gcc -fopenmp second_OpenMP.c
$ ./a.out
```



Loop Constructs: Parallel for

```
In C/C++:
#pragma omp parallel for
    for(i=0; i<n; i++)
        {
            c[i] = a[i] + b[i];
        }</pre>
```





OpenMP Clauses

- Directives dictate what the OpenMP thread team will do
- Examples:
 - Parallel regions are marked by the parallel directive
 - Work sharing loops are marked by do, for directives (Fortran, C/C++)
- Clauses control the behaviour of any particular OpenMP directive
- Examples:
- 1. Scoping of variables: private, shared, default
- 2. Initialization of variables: copyin, firstprivate
- 3. Scheduling: static, dynamic, guided
- 4. Conditional application: if
- 5. Number of threads in team: num threads



Scheduling of loop iterations

- Schedule clause
 - specifies how loop iteration are divided among team of threads
- Supporting Scheduling types
 - Static
 - Dynamic
 - Guided
 - Runtime



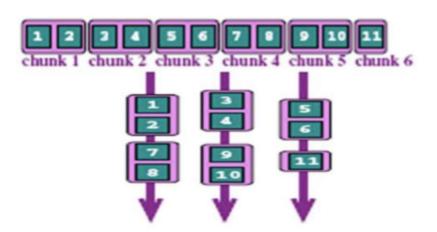
schedule Clause

schedule(static, [n])

- Each thread is assigned chunks in "round robin" fashion, known as block cyclic scheduling
- If n has not been specified, it will contain CEILING(number_of_iterations/number_of_threads) iterations
- Deterministic

Example:

loop of length 11, with 3 threads, and chunk size of 2:





Example

```
#include <stdio.h>
#include <omp.h>
#include <unistd.h>
#define THREADS 8
int main()
   int i,N=10;
   printf("First part....\n");
   #pragma omp parallel for num threads(THREADS)
   for (i = 0; i < N; i++) {
       printf("Thread %d is doing iteration %d.\n", omp get thread num(), i);
   printf("\n\nSecond part...\n");
   #pragma omp parallel for schedule(static) num_threads(THREADS)
   for (i = 0; i < N; i++) {
      sleep(i); /* wait for i seconds */
       printf("Thread %d has completed iteration %d.\n", omp_get_thread_num( ), i);
   return 0;
```

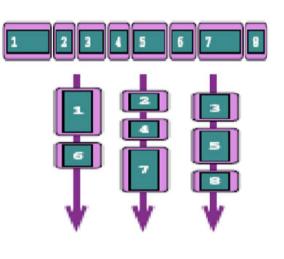


schedule Clause

schedule(dynamic, [n])

- Iteration of loop are divided into chunks containing n iterations each
- Default chunk size is 1
- Iterations picked by threads depends upon the relative speeds of thread execution

```
#pragma omp parallel for schedule (dynamic)
  for(i=0; i<8; i++)
  {
     ... (loop body)
  }</pre>
```



```
#pragma omp parallel for schedule(dynamic) num_threads(THREADS)
    for (i = 0; i < N; i++) {
        sleep(i); /* wait for i seconds */
        printf("Thread %d has completed iteration %d.\n", omp_get_thread_num(), i);
}</pre>
```



schedule Clause

schedule (guided, [n])

- If you specify n, that is the minimum chunk size that each thread should get
- Size of each successive chunks is decreasing chunk size = max((num_of_iterations_remaining/2*num_of_threads),n)
 - the formula may differ across compiler implementations

```
#pragma omp parallel for schedule(guided) num_threads(THREADS)
   for (i = 0; i < N; i++) {
      sleep(i); /* wait for i seconds */
      printf("Thread %d has completed iteration %d.\n", omp_get_thread_num(), i);
}</pre>
```

schedule (runtime)

Determine the scheduling type at run time by the OMP_SCHEDULE environment variable export OMP_SCHEDULE="static, 4"



Data Scoping in OpenMP

#pragma omp parallel [data scope clauses ...]

- shared
- private
- firstprivate
- lastprivate
- default



shared Clause (Data Scope)

- Shared data among team of threads
- Each thread can modify shared variables
- Data corruption is possible when multiple threads attempt to update the same memory location
- Data correctness is user's responsibility



private Clause (Data Scope)

- The values of private data are undefined upon entry to and exit from the specific construct.
- Loop iteration variable is private by default

```
#pragma omp parallel for private(tid)
for(i=0; i<n; i++)
{
  tid = omp_get_thread_num();
  printf(" My rank is %d ", tid)
}</pre>
```



firstprivate Clause (Data Scope)

• The clause combines behaviour of private clause with automatic initialization of the variables in its list with values prior to parallel region

```
int a, i, b=51, n=100 ;
printf("Before parallel loop: b=%d ,n=%d\n",b,n)

#pragma omp parallel for private(i), firstprivate(b)

for(i=0; i<n; i++)

{
    a = i + b;
}</pre>
```



lastprivate Clause (Data Scope)

• Performs finalization of private variables Each thread has its own copy

```
Example:
int a, i, b=51,n=100;
printf("Before parallel loop: b=%d ,n=%d\n",b,n)

#pragma omp parallel for private(i), firstprivate(b), lastprivate(a)
    for(i=0; i<n; i++)
    {
        a=i+b;
    }</pre>
```

//After parallel region: a = 150



default Clause (Data Scope)

- Defines the default data scope within parallel region
- default (private | shared | none)



More clauses for parallel directive

```
#pragma omp parallel [clause , clause, ...]
```

- nowait
- if
- reduction



nowait Clause

#pragma omp parallel nowait

- By default there is implicit barrier at the end of parallel region
- Allows threads that finish earlier to proceed without waiting
- If specified, then threads do not synchronize at the end of parallel loop



if Clause

if (integer expression)

- Determines if the region should be parallelized
- Useful option when data is too small



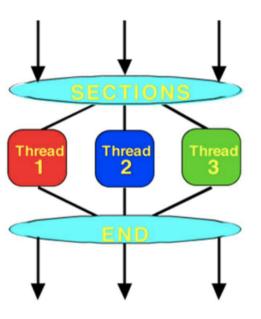
Reduction Clause

- Performs a collective operation on variables according to the given operators
 - built-in reduction operations such as +, *, -, max, min, logical operators
 - user can define his/her own operations
- Makes reduction variable as private
 - The variable is initialized according to reduction operator e.g. 0 for addition
- Each thread will perform the operation in its local variable
- Finally local results are combined into global result in shared variable



Work sharing: Section Directive

- One thread executes one section
- Each section is executed exactly once and





Work sharing: Single Directive

• Designated section is executed by single thread only.

```
#pragma omp single
{
    // read value of "a" from file
}
#pragma omp for
for (i=0;i<N;i++)
b[i] = a;</pre>
```



Work sharing: Master

• Similar to single, but code block will be executed by the master thread only

```
#pragma omp master
{
    // reading and writing data etc.
}
```

```
#pragma omp master
- - block of code - -
```



Synchronization Constructs

• CRITICAL: Mutual Exclusion

• ATOMIC: Atomic Update

• BARRIER: Barrier Synchronization



Example

```
#include <stdio.h>
int main()
{
    int x=0;
    x+=1;
    printf("%d\n",x);
    return 0;
}
```

```
$ gcc
example.c
$ ./a.out
1
$
```



Example

Race Condition

```
#include <stdio.h>
#include <omp.h>

int main()
{
    int x=0, size=8;
    omp_set_num_threads(size);
    #pragma omp parallel shared(x)
    {
        x+=1;
    }
        printf("%d\n",x);
        return 0;
}
```

```
$ gcc -fopenmp example_OpenMP.c
$ ./a.out
5
$ ./a.out
8
$ ./a.out
4
$ ./a.out
4
$ ./a.out
3
```



Race Condition

Intended

Thread 0		Thread 1		Value
read	←			0
increment				0
write	\rightarrow			1
		read	←	1
		increment		1
		write	\rightarrow	2

Possible...

Thread 0		Thread 1		Value
				0
read	←			0
increment		read	←	0
write	\rightarrow	increment		1
		write	\rightarrow	1
				1

- In a critical section, need mutual exclusion to get intended result
- The following OpenMP directives prevent this race condition:

```
#pragma omp critical – for a code block (C/C++)#pragma omp atomic – for single statement
```



Race Condition: Real Life Example



• So, how do we ensure that there is no contention and both person get to access the bathroom?

Use Lock and Key



Critical Construct

```
#pragma omp critical [name]
   structured-block
```



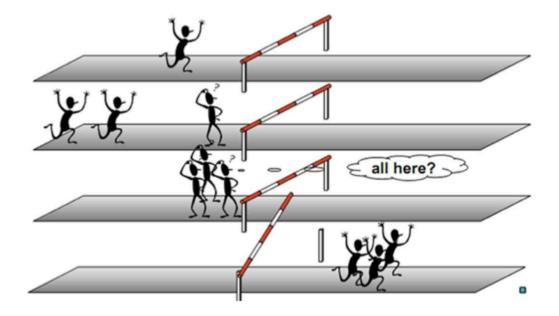
Example

```
$ gcc -fopenmp example_OpenMP.c
$ ./a.out
8
```



Barrier Directive

Synchronizes all the threads in a team





Barrier Directive

#pragma omp barrier //Threads wait until all threads reach this point

```
#include <stdio.h>
#include <omp.h>
int main()
{
   int x=0;
    #pragma omp parallel
   {
        #pragma omp master
        x=1;
        #pragma omp barrier
        printf("%d\n",x);
   }
   return 0;
}
```

Be careful not to cause deadlock:

No barrier inside of critical, master, sections, single!



Atomic Directive

- Mini Critical section
- Specific memory location must be updated atomically

#pragma omp atomic

- Single line code -



Example

```
#include <stdio.h>
int main()
{
    int i, n=5, sum;
    for (i = 1; i <= n ; i++)
    {
        sum += i ;
    }
    printf("%d\n",sum);
    return 0;
}</pre>
```

Race Condition

```
$ gcc example.c
$ ./a.out
15
$
```



Reduction Clause

Recall previous example,

- Simple parallel-for doesn't work due to race condition on shared sum
- Best solution is to apply OpenMP's reduction clause
- Doing private partial sums is fine too; add a critical section for sum of i



Example

```
#include <stdio.h>
#include <omp.h>

int main()
{
    int i, n=5, sum;
    omp_set_num_threads(8);
    #pragma omp parallel for reduction(+:sum)
        for(i=1; i<=n;i++)
        {
            sum+=i;
        }
        printf("%d\n",sum);
        return 0;
}</pre>
```

```
$ gcc -fopenmp example_OpenMP.c
$ ./a.out
15
$ ./a.out
15
$ ./a.out
15
$ ./a.out
```



Table: OpenMP Directives

Name	Parallel	Worksharing	Synchronization	Data environment	Description
					Defines a parallel region, which is
parallel	Yes	No	No	No	code that will be executed by mul-
paraner	165	110	110	110	tiple threads in parallel
					Causes the work done in a for loop
for	No	Yes	No	No	inside a parallel region to be di-
101	110	100	110	110	vided among threads
					Identifies code sections to be di-
sections	No	Yes	No	No	vided among all threads
					Lets you specify that a section of
					code should be executed on a single
single	No	Yes	No	No	thread, not necessarily the master
					thread
					Specifies that only the master
master	No	No	Yes	No	thread should execute a section of
			200		the program
					Specifies that code is only executed
critical	No	No	Yes	No	on one thread at a time
					Synchronizes all threads in a team;
					all threads pause at the barrier,
barrier	No	No	Yes	No	until all threads execute the bar-
					rier
					Specifies that a memory location
atomic	No	No	Yes	No	that will be updated atomically
					Specifies that all threads have the
flush	No	No	Yes	No	same view of memory for all shared
					objects
					Specifies that code under a paral-
ordered	No	No	Yes	No	lelized for loop should be executed
					like a sequential loop
		.,	3.7	37	Specifies that a variable is private
threadprivate	No	No	No	Yes	to a thread



Table: OpenMP Functions

Name	Environment Execution	Lock	Description	
$omp_set_num_threads$	Yes	No	Sets the number of threads in upcoming parallel regions, unless overridden by a num_threads clause.	
$omp_get_num_threads$	Yes	No	Returns the number of threads in the parallel region.	
$omp_get_max_threads$	Yes	No	Returns an integer that is equal to or greater than the number of threads that would be available if a parallel region without num_threads were defined at that point in the code.	
$omp_get_thread_num$	Yes	No	Returns the thread number of the thread executing within its thread team.	
omp_get_num_procs	Yes	No	Returns the number of processors that are available when the function is called.	
omp_in_parallel	Yes	No	Returns nonzero if called from within a parallel region.	
$omp_set_dynamic$	Yes	No	Indicates that the number of threads available in upcoming parallel regions can be adjusted by the run time.	
omp_get_dynamic	Yes	No	Returns a value that indicates if the number of threads available in upcoming parallel regions can be adjusted by the run time.	
omp_set_nested	Yes	No	Enables nested parallelism.	
omp_get_nested	Yes	No	Returns a value that indicates if nested parallelism is enabled.	
omp_init_lock	No	Yes	Initializes a simple lock.	
omp_init_nest_lock	No	Yes	Initializes a lock.	
omp_destroy_lock	No	Yes	Uninitializes a lock.	
$omp_destroy_nest_lock$	No	Yes	Uninitializes a nestable lock.	
omp_set_lock	No	Yes	Blocks thread execution until a lock is available.	
omp_set_nest_lock	No	Yes	Blocks thread execution until a lock is available.	
omp_unset_lock	No	Yes	Releases a lock.	
$omp_unset_nest_lock$	No	Yes	Releases a nestable lock.	
omp_test_lock	No	Yes	Attempts to set a lock but doesn't block thread execution.	
$omp_test_nest_lock$	No	Yes	Attempts to set a nestable lock but doesn't block thread execution.	



Table: OpenMP Datatypes

Name	Environment Execution	Lock	Description
omp_lock_t	No	Yes	A type that holds the status of a lock, whether the lock is available or if a thread owns a lock.
$omp_nest_lock_t$	No	Yes	A type that holds one of the follow- ing pieces of information about a lock: whether the lock is available, and the identity of the thread that owns the lock and a nesting count.



Table: OpenMP Environmental Variables

Environment Variable	Description		
OMP_SCHEDULE	Modifies the behavior of the schedule clause when $schedule(runtime)$ is		
	specified in a for or parallel for directive.		
OMP_NUM_THREADS	Sets the maximum number of threads in the parallel region, unless over-		
	ridden by omp_set_num_threads or num_threads.		
OMP_DYNAMIC	Specifies whether the OpenMP run time can adjust the number of		
	threads in a parallel region.		
OMP_NESTED	Specifies whether nested parallelism is enabled, unless nested parallelism		
OMF_NESTED	is enabled or disabled with omp_set_nested.		



Granularity of Parallelization

Coarse-grain parallelism vs. Fine grain parallelism

```
#pragma omp parallel for
for(i=0, i<n; i++)
{
    // work1;
}

#pragma omp parallel for
for(i=0, i<n; i++)
    {
        // work;
}

#pragma omp for
for(i=0, i<n; i++)
        {
        // work;
}

#pragma omp for
for(i=0, i<n; i++)
        {
        // work;
}
</pre>
```



Some Tips

- Identify Loop-level parallelism: Run the loop backwards and see if same results are produced
- Load imbalance due to branching statements, sparse matrices: schedule(dynamic)
- Parallelization of less compute intensive loops: Use small number of threads e.g.
 #pragma omp parallel num_threads(4)
- Parallelize initialization of input data speedup and data locality



Advantages and Disadvantages

Advantages

- Shared address space provides user friendly programming
- Ease of programming
- Data sharing between threads is fast and uniform (low latency)
- Incremental parallization of sequential code
- Leaves thread management to compiler
- Directly support by compiler

Disadvantages

- Internal details are hidden
- Programmer is responsible for specifying synchronisation, e.g. locks
- Cannot run across distributed memory
- Performance limited by memory architecture
- Lack of scalability between memory and CPUs
- Requires compiler which supports OpenMP
- Bigger machines are heavy on budget



References

The contents of the presentation have been adapted from several sources. Some of the sources are as following:

www.openmp.org/ https://computing.llnl.gov/tutorials/openMP/





Thanks!

