Resource Allocation in an Open RAN System using Network Slicing

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***Abstract*—Taking advantage of both virtual RAN (v-RAN) and Cloud RAN (C-RAN), Open RAN (O-RAN) is introduced as the next generation of RAN systems which leads to increase flexibility, Openness, and reduce operational costs and allow them to add new capabilities to the network more quickly. O-RAN separate RAN into three different units, namely Radio Unit (O-RU), Distributed Unit (O-DU), and Central Unit (O-CU). In this paper, we study the problem of baseband resource allocation and virtual network function (VNF) activation in O-RAN architecture based on their service priority for different types of 5G services includes enhanced mobile broadband (eMBB), ultra-reliable low latency communications (URLLC) and massive Machine Type Communications (mMTC) services. According to the concept of network slicing, the isolation of different types of services in O-DU, O-CU, and user plane function (UPF) is performed. The limited fronthaul capacity and the restriction of end-to-end delay are considered in this problem. The optimization of baseband resources includes O-RU assignment, physical resource block (PRB), and power allocation. The main problem is the mixed-integer non-linear programming that is tremendously difficult. So we broke it down into two different steps that the iterative algorithm solves it. In the first step, we reformulated and simplified the problem to find the power allocation, PRB assignment, and the number of activated VNFs. In the second step, the O-RU association is carried out. The proposed method (IAPPVO) is confirmed by the simulation results which illustrate a higher achievable data rate than a baseline scheme that only optimizes one of the baseband resources and the FBDR algorithm described in other papers. Also, the simulation results achieved less end-to-end delay for the proposed method than the FBDR and baseline scheme.**

***Index Terms*—Open Radio Access Network (O-RAN), Virtual Network Function (VNF)**

I. Introduction

One of the fifth-generation goals of the wireless system is to achieve the desired QoS (such as rate, delay, power, ...) for different types of services. Network slicing is the best solution for this aim. A network slice is an end-to-end logical network that offers services with special needs. Multiple isolated network slices run, manage, and work independently on the same infrastructure. There are several implementations of network slicing, including network core slicing, RAN slicing, and slicing of both sections. Different type of services includes enhanced mobile broadband (eMBB), and ultra-reliable low latency communications (URLLC), and massive Machine Type Communications (mMTC) services are introduced in 5th generation of mobile network. Each type of service requires a particular slice of network based on its QoS [1]–[3].

Recently, RAN virtualization attracts significant attention from industry and academia since it has remarkable benefits that increase flexibility and reduce operating costs such as CAPEX and OPEX and allow them to add new capabilities to the network more quickly. In addition to RAN virtualization, openness and RAN intelligence are two other fundamental points that encourage the Open Radio Access Network (O-RAN) Alliance to establish O-RAN as the next generation of RAN systems.

The idea of O-RAN comes from the integration of virtual RAN (vRAN) and cloud RAN (CRAN), and it takes advantage of both. CRAN divides RAN into two parts radio remote head (RRH) and baseband unit (BBU). More than one distributed RRHs can be connected to a centralized BBU, which is named BBU-pool [4]. Unlike the previous generation of RAN that divides RAN into two parts, O-RAN separate RAN into three different units, namely Radio Unit (O-RU), Distributed Unit (O-DU), and Central Unit (O-CU). O-RU is a logical node that contains RF and lowers PHY. Moreover, the O-DU expresses another logical node that includes higher PHY, MAC, and RLC. In addition, the O-CU depicts the logical node contains two parts, which are the O-CU user plane (O-CU-UP) and O-CU control plane (O-CU-CP). O-CU-UP hosts PDCP-UP and SDAP, and O-CU-CP hosts PDCP-CP and RRC. O-DU and O-CU are connected via an open and well-defined interface F_1. Moreover, O-DU is connected to a radio unit (O-RU) with an open fronthaul interface. The architecture of O-RAN contains other principal logical nodes called Orchestration and Automation, RAN Intelligent Controller (RIC)- Near Real-Time and O-Cloud. One of the necessities of the new generation of wireless networks is its intelligence. Based on the requirement of an intelligent wireless network, O-RAN offers machine learning techniques. The two logical nodes RIC-Non Real-Time (which is placed in Orchestration and Automation node) and RIC- Near Real Time, implement the algorithms for network intelligence [5]–[11].

The separation of network software and hardware elements has been done and introduced as network function virtualization (NFV), and virtual network functions (VNF) are system function blocks. This technology improves the

system performance in the fifth generation of telecommunications. The key idea of the implementation of NFV is to decouple software from physical hardware, dynamic scaling, and the deployment of flexible network functions. A usual NFV offer is to execute VNFs on virtual machines or containers in a cloud system [12], [13]. As a result, some O-RAN components that include user plane function (UPF), O-CU, O-DU, and RAN Intelligent Controller (RIC)-near real-time, are virtualized and implemented as a VNF that can be run on virtual machines (VMs) or containers.

A. Related Works

Network slicing is increasingly receiving research attention. Many researchers studied the problem of resource allocation in network slicing for multitenant cellular networks [14]–[16]. In [15], dynamic network slicing in multitenant heterogeneous CRAN (H-CRAN) is considered. The process of allocating network resources to users is discussed. The network slicing scheme includes a higher level, which manages user acceptance control, user communication includes radio unit association (RRH association to maximize user rates and allocate baseband resource capacity), the allocation of BBU capacity, and a lower level, which is the allocation of power and physical resource blocks (PRB) among us In the article, [17], network slicing in the radio section is considered for fog or F-RAN structure, in which two network slices are set for hotspots and vehicle scenarios with related infrastructure. In [18], [19] the implementation of RAN level slicing is discussed in mobile network operator (MNO). Also, the problem of resource allocation is considered. Moreover, the challenges facing RAN slicing have also been explored, one of which involves designing and managing multiple slices in the shared infrastructure in an efficient manner while guaranteeing the agreed service level agreement (SLA) for each.

Multiplexing eMBB and URLLC services on the same RAN and sharing the resources for these services is challenging because many researchers pay attention to this topic. In [2], [20] the problem of resource allocation in the coexistence of these two services (URLLC and eMBB) is considered based on their QoS. In [21], the problem of resource allocation for joint eMBB and URLLC is formulated and solved by deep reinforcement learning. In [22] the problem of power minimization for these two types of services (URLLC and eMBB) is presented for non-orthogonal multiple access (NOMA) and orthogonal multiple access (OMA). In [23], the authors proposed to allocate RAN resources for the network slicing system in the coexistence of eMBB and URLLC services. The system guarantees the latency, the service rate, and the maintenance of reliability.

Virtualization technique for RAN and core is one of the exciting topics. In [24], [25], the authors solve the problem of obtaining beamforming and VMs activation in a C-RAN system with limited fronthaul capacity. This paper aims to minimize the energy cost with the system delay, fronthaul capacity, and rate constraint. Also, Transmission and processing delay are modeled based on M/M/1 queuing theory to guarantee delay for the services. In [26], [27], the problem of joint virtual computing resource allocation with beamforming is formulated; Also, the association of RRH to UE is considered and solved using novel methods. In [28]–[30], the problem of joint power allocation and RRH association in the H-CRAN system is considered to maximize energy efficiency.

B. Main Contribution

This study aims to optimize VNF activation, power allocation, PRB allocation, and O-RUs association to develop an isolated network slicing outline for different types of services in an O-RAN platform. In this paper, as depicted in Figure 1, the downlink of the ORAN system is studied. The main contributions of this paper are summarized as follow:

- In this paper, a network slicing model is depicted for three different types of services introduced in 5G (URLLC, eMBB, and mMTC). The problem of radio resource allocation and VNF activation is studied in this paper in the O-RAN architecture. We formulate the problem of baseband resource allocation to maximize the weighted throughput of the O-RAN system for a different types of services with specific QoS.
- The desired QoS conditions such as delay and throughput are considered for different types of services. We formulate the end-to-end mean delay of the system based on the activated VNFs and each user's throughput. We take into account the limited fronthaul capacity and accurately obtain the power and the capacity of each O-RU based on the quantization noise. We model the interference of neighboring O-RU and formulate the actual throughput for different types of services. We model the throughput of URLLC and mMTC based on their short packet transmission.
- The main problem is mixed-integer non-linear programming that is extremely difficult to solve. We perform a two-step iterative algorithm to solve it. The two-step iterative algorithm is presented for the resource management framework with the first-step VNF activation, power allocation, PRB association, and the second-step O-RU association.
- We reformulated and simplified the main problem for the first step to find an upper and lower bound for the number of activated VNFs and use lagrangian function and KKT conditions to find optimal power and PRB allocation. For the second step, the problem of O-RU association can be converted to a multiple knapsack problem and solved by the Greedy algorithm.

The rest of this paper is organized as follows. The system model and the problem formulation are described in Section II. The details of our proposed resource management algorithm are introduced in Section III. In Section IV, numerical results are provided to evaluate the performance of the proposed algorithm. Section V, concludes the whole paper.

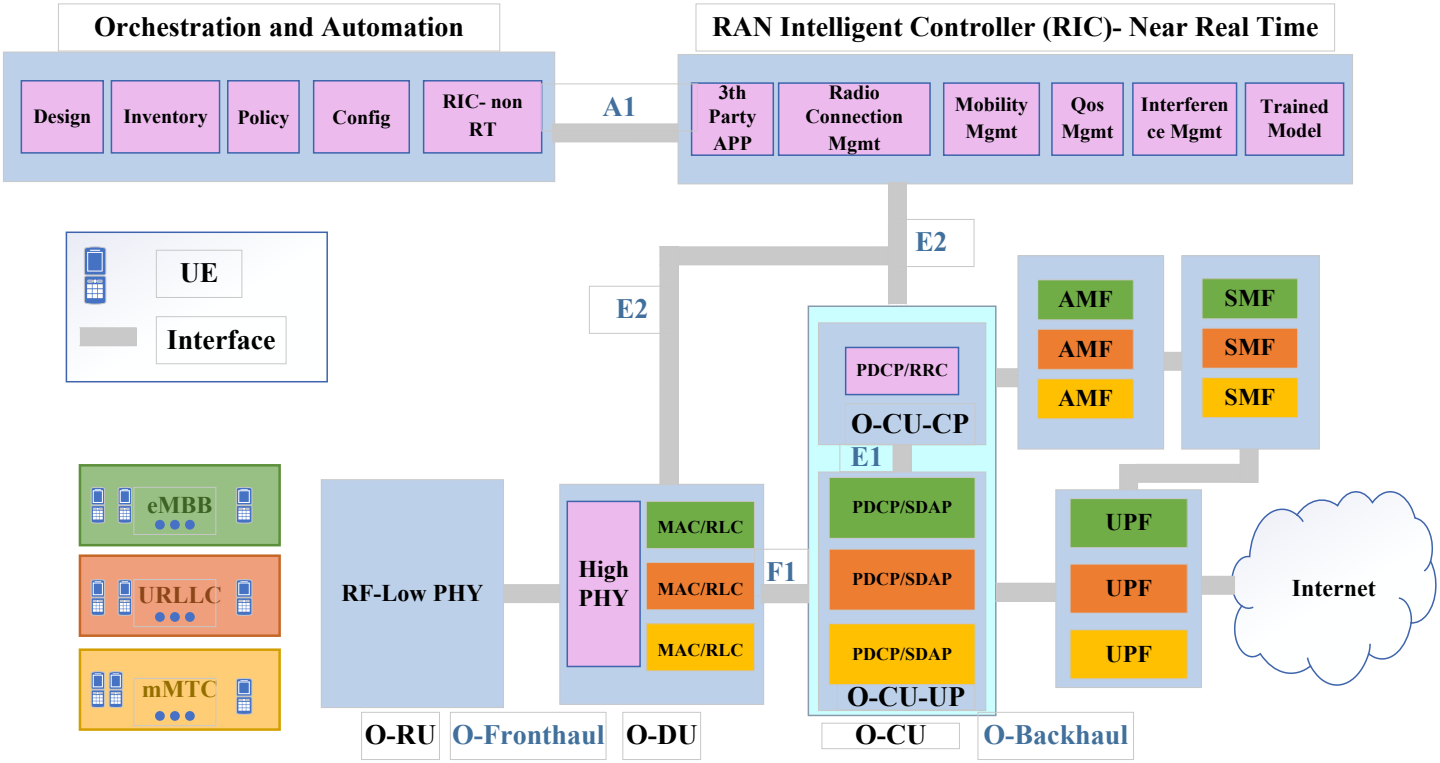

Fig. 1. Network sliced ORAN system

II. System Model and Problem Formulation

In this section, we describe the downlink system in O-RAN slicing as depicted in Figure 1. Here, firstly, we present the system model. Then, we obtain achievable data rates, power of O-RU, and the fronthaul capacity for the downlink (DL) of the ORAN system. Afterward, we discuss the mean delay and the power of VNFs. Finally, the main problem is expressed.

A. System Model

Suppose we have three service types includes mMTC, eMBB, and URLLC, which support different applications.

Assume we have S_1, S_2 and S_3 different applications for the first, second and third service type, respectively ($S = S_1 + S_2 + S_3$). So, we have S preallocated slices serving these S services; There are S_1 slices for the first service type (eMBB), S_2 slices for the second service type (URLLC), and S_3 slices for the third service type (mMTC). So each service request s is served by its corresponding slice.

Each Service $s_j \in \{1, 2, ..., S_j\}$ consists of U_s request from the single-antenna UEs which require certain QoS to be able to use the requested program($j \in \{1, 2, 3\}$ indicate service type). There are different application requests which fall into one of these service categories. Each application request requires a specific QoS. Based on the application and QoS request, UE may be admitted and allocated to the resources. Each slice $s_j \in \{1, 2, ..., S_j\}$, $j \in \{1, 2\}$ consists of preallocated virtual resource blocks that are mapped to the Physical Resource Blocks (PRBs), M_s^d VNFs for the processing of O-DU, M_s^c VNFs for the processing of O-CU-UP and M_s^u VNFs for the processing of UPF.

All K PRBs can be assigned to all UE in each service. Also, each VNF instance runs on a virtual machine (VM) that uses resources from the data centers. Each VM requires enough resources of CPU, memory, storage, and network bandwidth.

In addition, there are R multi-antenna O-RU that are shared between slices. Each O-RU $r \in \{1, 2, ..., R\}$ has J antenna for transmitting and receiving data. Also $\mathcal{R} = \{r | r \in 1, 2, ..., R\}$ depicts the set of O-RUs. Moreover, all O-RUs, have access to all PRBs.

B. The Achievable Rate

The SNR of i^{th} UE served at slice s on PRB k is obtained from

$$\rho^k_{r,u(s,i)} = \frac{|p^k_{r,u(s,i)} \mathbf{h}^{H\,k}_{r,u(s,i)} \mathbf{w}^k_{r,u(s,i)}|^2}{BN_0 + I^k_{r,u(s,i)}}, \quad (1)$$

where $p^k_{r,u(s,i)}$ represents the transmission power from O-RU r to i^{th} UE served at slice s on PRB k. $\mathbf{h}^k_{r,u(s,i)} \in \mathbb{C}^J$ is the vector of channel gain of a wireless link from r^{th} O-RU to the i^{th} UE in s^{th} slice. In addition, $\mathbf{w}^k_{r,u(s,i)} \in \mathbb{C}^J$ depicts the transmit beamforming vector from r^{th} O-RU to the i^{th} UE in s^{th} slice that is the zero forcing beamforming vector to minimize the interference which is indicated as below

$$\mathbf{w}^k_{r,u(s,i)} = \mathbf{h}^k_{r,u(s,i)} (\mathbf{h}^{H\,k}_{r,u(s,i)} \mathbf{h}^k_{r,u(s,i)})^{-1} \quad (2)$$

Moreover, $g^r_{u(s,i)} \in \{0,1\}$ is the binary variable that illustrates whether O-RU r served the i^{th} UE that is allocated to s^{th} slice or not. Also, BN_0 denotes the power of Gaussian additive noise, and $I^k_{r,u(s,i)}$ is the power of interfering signals represented as follow.

$$I^k_{r,u(s,i)} = \underbrace{\sum_{\substack{l=1\\l\neq i}}^{U_s} \gamma_1 p^k_{u(s,l)} \sum_{\substack{r'=1\\r'\neq r}}^{R} |\mathbf{h}^{H\,k}_{r',u(s,i)} \mathbf{w}^k_{r',u(s,l)} g^{r'}_{u(s,l)}|^2}_{\text{(intra-slice interference)}} + \underbrace{\sum_{\substack{n=1\\n\neq s}}^{S} \sum_{l=1}^{U_s} \gamma_2 p^k_{u(n,l)} \sum_{\substack{r'=1\\r'\neq r}}^{R} |\mathbf{h}^{H\,k}_{r',u(s,i)} \mathbf{w}^k_{r',u(n,l)} g^{r'}_{u(n,l)}|^2}_{\text{(inter-slice interference)}} + \underbrace{\sum_{j=1}^{R} \sigma^2_{q_{r_j}} |\boldsymbol{h}_{r,u(s,i)}|^2}_{\text{(Quantization Noise Interference)}} \tag{3}$$

where $\gamma_1 = e^k_{u(s,i)} e^k_{u(s,l)}$ and $\gamma_2 = e^k_{u(s,i)} e^k_{u(n,l)}$.
$e^k_{u(s,i)}$ is the binary variable to show whether the k^{th} PRB is allocated to the UE i in slice s, assigned to r^{th} o-RU.

To obtain SNR as formulated in (1), let $y_{u(s,i)}$ be the received signal user i in s^{th} service

$$y_{u(s,i)} = \textstyle\sum_{r=1}^{R} \sum_{k=1}^{K_s} \boldsymbol{h}^{H\,k}_{r,u(s,i)} g^r_{u(s,i)} e^k_{r,u(s,i)} \mathfrak{y}^k_{r,u(s,i)} + z_{u(s,i)}, \tag{4}$$

where $\mathfrak{y}^k_{r,u(s,i)} = \boldsymbol{w}^k_{r,u(s,i)} p^{k\,\frac{1}{2}}_{r,u(s,i)} x_{u(s,i)} + \boldsymbol{q}_r$ and $x_{u(s,i)}$ depicts the transmitted symbol vector of UE i in s^{th} set of service, $z_{u(s,i)}$ is the additive Gaussian noise $z_{u(s,i)} \backsim \mathcal{N}(0, N_0)$ and N_0 is the noise power. In addition, $\boldsymbol{q}_r \in \mathbb{C}^J$ indicates the quantization noise, which is made from signal compression in O-DU.

The achievable data rate for the i^{th} UE request in the s^{th}_1 application of service type 1 (eMBB) can be written as $\mathcal{R}_{u(s_1,i)}$ that is formulated as below.

$$\begin{aligned} \mathcal{R}^k_{r,u(s_1,i)} &= B \log_2(1 + \rho^k_{r,u(s_1,i)}), \\ \mathcal{R}^r_{u(s_1,i)} &= \sum_{k=1}^{K} B \log_2(1 + \rho^k_{r,u(s_1,i)} e^k_{r,u(s_1,i)}), \\ \mathcal{R}_{u(s_1,i)} &= \sum_{r=1}^{R} \mathcal{R}^r_{u(s_1,i)} g^r_{u(s_1,i)} \end{aligned} \tag{5}$$

where B is the bandwidth of system. $\mathcal{R}^r_{u(s_1,i)}$ is the achievable rate of each RU r to UE i in slice s_1. Since the blocklength in URLLC and mMTC is finite, the achievable data rate for the i^{th} UE request in the s^{th}_j ($j \in \{2,3\}$) application of service type 2 (URLLC) and 3 (mMTC) is not achieved from Shannon Capacity formula. So, for the short packet transmission, the achievable data rate is approximated from following

$$\begin{aligned} \mathcal{R}^k_{r,u(s_j,i)} &= B \log_2(1 + \rho^k_{r,u(s_j,i)} - \zeta^k_{u(s_j,i)}) e^k_{u(s_j,i)}, \\ \mathcal{R}^r_{u(s_j,i)} &= \sum_{k=1}^{K} B(\log_2(1 + \rho^k_{u(s_2,i)}) - \zeta^k_{u(s_j,i)}) e^k_{u(s_j,i)} \\ \mathcal{R}_{u(s_j,i)} &= \sum_{r=1}^{R} \mathcal{R}^r_{u(s_j,i)} g^r_{u(s_j,i)} \end{aligned} \tag{6}$$

Where $j \in \{1,2\}$. Also we have

$$\zeta^k_{u(s_j,i)} = log_2(e) Q^{-1}(\epsilon) \sqrt{\frac{\mathfrak{C}^k_{u(s_j,i)}}{N^k_{u(s_j,i)}}}) \tag{7}$$

Where, ϵ is the transmission probability, Q^{-1} is the inverse of Q- function (Gaussian), $\mathfrak{C}^k_{u(s_j,i)} = 1 - \frac{1}{(1+\rho^k_{u(s_j,i)})^2}$ depicts the channel dispersion of UE i at slice s_j, experiencing PRB k and $N^k_{u(s_j,i)}$ represents the blocklength of it. $\mathcal{R}^{e,r}_{u(s_j,i)}$ is the achievable rate of each O-RU r to UE i in slice s_j.

If we replace $p^k_{u(s,l)}$ and $p^k_{u(n,l)}$ in (3) by P_{max}, an upper bound $\bar{I}^k_{r,u(s,i)}$ is obtained for $I^k_{r,u(s,i)}$. Therefore, $\bar{\mathcal{R}}_{u_{(s,i)}} \forall s, \forall i$ is derived by using $\bar{I}^k_{r,u(s,i)}$ instead of $I^k_{r,u(s,i)}$ in (6) and (5).

C. Power of O-RU and Fronthaul Capacity

Let P_r denote the power of the transmitted signal from the r^{th} O-RU to UEs served by it. From (4), we have,

$$P_r = \sum_{s=1}^{S} \sum_{k=1}^{K_s} \sum_{i=1}^{U_s} |\mathbf{w}^k_{r,u(s,i)}|^2 p^k_{r,u(s,i)} g^r_{u(s,i)} e^k_{r,u(s,i)} + \sigma^2_{q_r}. \tag{8}$$

Since we have a fiber link between O-RU and O-DU, the rate of users on the fronthual link between O-DU and the r^{th} O-RU is formulated as

$$C_r = \log\,(1 + \frac{\sum_{s=1}^{S} \sum_{k=1}^{K_s} \sum_{i=1}^{U_s} |\mathbf{w}^k_{r,u(s,i)}|^2 \alpha^k_{r,u(s,i)}}{\sigma^2_{q_r}}), \tag{9}$$

Where, $\alpha^k_{r,u(s,i)} = p^k_{r,u(s,i)} g^r_{u(s,i)} e^k_{r,u(s,i)}$ and $\sigma^2_{q_r}$ is the power of quantization noise.

D. Mean Delay

In this part, the end-to-end mean delay for a service is obtained. Suppose the mean total delay is depicted as T_{tot}.

$$\begin{aligned} T_{tot} &= T_{process} + T_{transmission} + T_{propagation} \\ T_{process} &= T_{RU} + T_{DU} + T_{CU} + T_{UPF} \\ T_{transmission} &= T_{front} + T_{mid} + T_{back} + T_{trans2net} \\ T_{propagation} &= T_{front} + T_{mid} + T_{back} + T_{trans2net} \end{aligned} \tag{10}$$

Total delay is the sum of processing delay, transmission delay, and propagation delay. The propagation delay is the time takes for a signal to reach its destination. So it has a constant value based on the length of the fiber link ($T = L/c$, where L is the length of the link and c is the speed of signal). Also, the transmission delay is the amount of time required to push all the packets into the fiber link. ($T = \frac{\alpha}{R}$

Where R is the rate of transmission in each link and α is the mean arrival data rate of each link which is constant in this model.) Here we assume the value of propagation delay and transmission is negligible compared to the rest.

$$T_{tot} \approx T_{process} \tag{11}$$

1) Processing Delay: Assume the packet arrival of UEs follows a Poisson process with arrival rate $\lambda_{u(s,i)}$ for the i^{th} UE of the s^{th} slice. Therefore, the mean arrival data rate of the s^{th} slice in the UPF layer is $\alpha_s^U = \sum_{u=1}^{U_s} \lambda_{u(s,i)}$. Assume the mean arrival data rate of the UPF layer for slice s (α_s^U) is approximately equal to the mean arrival data rate of the O-CU-UP layer (α_s^C) and O-DU (α_s^D). so $\alpha_s = \alpha_s^U \approx \alpha_s^C \approx \alpha_s^D$. since, by using Burkes Theorem, the mean arrival data rate of the second and third layers, which are processed in the first layer, is still Poisson with rate α_s It is assumed that there are load balancers in each layer for each service to divide the incoming traffic to VNFs equally. Suppose the baseband processing of each VNF is depicted as M/M/1 processing queue. Each packet is processed by one of the VNFs of a slice. So, the mean delay for the s^{th} slice in the first and the second layer, modeled as M/M/1 queue, is formulated as follows, respectively.

$$\begin{aligned} T_{DU}^s &= \frac{1}{\mu_s^d - \alpha_s/M_s^d}, \\ T_{CU}^s &= \frac{1}{\mu_s^c - \alpha_s/M_s^c} \\ T_{UPF}^s &= \frac{1}{\mu_s^u - \alpha_s/M_s^u} \end{aligned} \tag{12}$$

Where M_s^d, M_s^c and M_s^u are the variables that depict the sum of VNFs in O-DU, O-CU-UP and UPF, respectively. Moreover, $1/\mu_s^d$, $1/\mu_s^c$, and $1/\mu_s^u$ are the mean service time of the O-DU, O-CU, and the UPF layers, respectively. Besides, α_s is the arrival rate which is divided by load balancer before arriving to the VNFs. The arrival rate of each VNF in each layer for each slice s is α_s/M_s^i $i \in \{d, c, u\}$.

In addition, $T_{RU}^{u(s,i)}$ is the mean transmission delay of i^{th} UE in s^{th} service on the wireless link. The arrival data rate of wireless link for each UE i in service s is $\lambda_{u(s,i)}$ As a result we have, $\sum_{i=1}^{U_s} \lambda_{u(s,i)} = \alpha_s$. Moreover, The service time of transmission queue for UE i requesting service s has an exponential distribution with mean $1/R_{u(s,i)}$ and can be modeled as a M/M/1 queue.

Therefore, the mean delay of the transmission layer for UE i in slice s is

$$T_{RU}^{u(s,i)} = \frac{1}{R_{u(s,i)} - \lambda_{u(s,i)}}; \tag{13}$$

So the mean processing delay for each UE i in slice s is

$$T_{process}^{u(s,i)} = T_{RU}^{u(s,i)} + T_{DU}^s + T_{CU}^s + T_{UPF}^s \tag{14}$$

Hence, $T_{tot}^{u(s,i)} \approx T_{process}^{u(s,i)}$

E. VNF Power

Assume the power consumption of baseband processing at each DC d that is connected to VNFs of a slice s is depicted as ϕ_s. So the system's total power for all active DCs connected to slices can be represented as follows.

$$\phi_{tot} = \textstyle\sum_{s=1}^{S} \phi_s.$$

Where, ϕ_s is obtained from below

$$\phi_s = M_s^u \phi_s^u + M_s^c \phi_s^c + M_s^d \phi_s^d \tag{15}$$

Moreover, ϕ_s^u, ϕ_s^c, and ϕ_s^d are the fixed cost of energy in UPF, O-CU, and O-DU, respectively.

F. Problem Statement

Suppose each slice s has priority factor δ_s where $\sum_{s=1}^{S} \delta_s = 1$. The optimization problem is formulated as follow. This paper aims to maximize the sum rate of all UEs with the presence of constraints as follows.

$$\max_{\boldsymbol{P},\boldsymbol{E},\boldsymbol{M},\boldsymbol{G}} \quad \sum_{s=1}^{S}\sum_{i=1}^{U_s} \delta_s \bar{\mathcal{R}}_{u_{(s,i)}} \tag{16a}$$

$$\text{subject to} \quad P_r \le P_{max} \quad \forall r \tag{16b}$$

$$p_{r,u(s,i)}^k \ge 0 \quad \forall i, \forall r, \forall s, \forall k, \tag{16c}$$

$$\bar{\mathcal{R}}_{u_{(s_j,i)}} \ge \mathcal{R}_{min}^{s_j} \quad \forall s, j \in \{1,2,3\}, \tag{16d}$$

$$C^r \le C_{max}^r \quad \forall r, \tag{16e}$$

$$T_{tot}^{u(s,i)} \le T_{max}^s \quad \forall i, \forall s, \tag{16f}$$

$$\mu_s \ge \alpha_s / M_s \quad \forall s, \tag{16g}$$

$$\bar{\mathcal{R}}_{u_{(s,i)}} \ge \lambda_{u_{(s,i)}} \quad \forall i, \forall s, \tag{16h}$$

$$0 \le M_s \le M^{max} \quad \forall s, \tag{16i}$$

$$\sum_{r} g_{u(s,i)}^r = 1 \quad \forall s, \forall i, \tag{16j}$$

$$\sum_{k=1}^{K_s} g_{u(s,i)}^r e_{r,u(s,i)}^k \ge 1 \quad \forall s, \forall i, \forall r \tag{16k}$$

$$\sum_{s=1}^{S}\sum_{i=1}^{U_s} g_{u(s,i)}^r e_{r,u(s,i)}^k \le 1 \quad \forall s, \forall i, \forall r \tag{16l}$$

$$\phi_{tot} \le \phi_{max}, \tag{16m}$$

$$g_{u(s,i)}^r \in \{0,1\} \quad \forall s, \forall i, \tag{16n}$$

$$e_{r,u(s,i)}^k \in \{0,1\} \quad \forall s, \forall i, \tag{16o}$$

where $\boldsymbol{P} = [p_{r,u(s,i)}^k] \ \forall s, \forall i, \forall r, \forall k$, is the matrix of power for UEs, $\boldsymbol{E} = [e_{r,u(s,i)}^k] \ \forall s, \forall i, \forall r, \forall k$ indicate the binary variable for PRB association. Moreover, $\boldsymbol{G} = [g_{u(s,i)}^r] \forall s, \forall i, \forall r$ is a binary variable for O-RU association. Furthermore, $M = [M_s^d, M_s^c, M_s^u] \ \forall s$ is the matrix that shown the number of VNFs in each layer of slice. (16b), and (16c), indicate that the power of each RU does not exceed the maximum power, and the power of each UE is a positive integer value, respectively. Also, (16d) shows that the rate of each UE requesting eMBB, URLLC, and mMTC is more than a threshold, respectively. (16e) and (16f) expressed the limited capacity of the fronthaul link and the limited delay of receiving the signal, respectively.

(16g) and (16h) denoted the stability of the M/M/1 queue model. (16i) restricted the number of VNF in each slice due to the limited resources. (16j) and (16k) guarantee that O-RU and PRB are associated with the UE, respectively. Also, (16l) ensures that each PRB can not be assigned to more than one UE associated with the same O-RU. In addition, (16m) indicates that the fixed cost of energy of VNFs in each slice does not exceed the threshold. Moreover, (16n) and (16o) depict that $\boldsymbol{E}$ and $\boldsymbol{G}$ are matrix of binary variables.

III. Proposed Algorithm Scheme

In this section, we first apply some simplifications to the system; Solving problem (16) is complicated since this problem is non-convex and it is a mixed-integer non-linear problem (MINLP) with a binary variable and an integer variable. We simplified and reformulated MINL parts and used an iterative heuristic algorithm to solve the problem. We solve this problem in two-level iteratively until it converges. In the first level, parameters $(\boldsymbol{P}, \boldsymbol{E}, \boldsymbol{M})$ are obtained by relaxing and reformulating parameters and turn them into a convex problem; Afterward, we solve it by dual optimization problem. In the second level, finding optimal O-RU association ($\boldsymbol{G}$) is concerned with the fixed parameter of power, PRB allocation, and number of VNFs. We repeat this procedure until the algorithm converges.

A. Sub-Problem 1

Suppose that $\boldsymbol{G}$ is fixed, we want to obtain $\boldsymbol{P}$, $\boldsymbol{E}$ and $\boldsymbol{M}$. Here, we first simplify and relax the parameters to convexify the problem.

As we mentioned before, by replacing $p^k_{u(s,l)}$ and $p^k_{u(n,l)}$ in (3) by P_{max}, an upper bound $\bar{I}^k_{r,u(s,i)}$ for $I^k_{r,u(s,i)}$, the lower bound $\bar{\rho}^k_{u(s,i)}$ for $\rho^k_{u(s,i)}$ and the lower bound $\bar{\mathcal{R}}_{u_{(s,i)}} \forall s, \forall i$ for $\mathcal{R}_{u_{(s,i)}}$ is obtained by replacing with $I^k_{r,u(s,i)}$ $\bar{I}^k_{r,u(s,i)}$ in (6) and (5) and make them concave.

Suppose $\hat{\rho}^k_{r,u(s,i)} = \frac{|P_{max}\mathbf{h}^{H\,k}_{r,u(s,i)}\mathbf{w}^k_{r,u(s,i)}g^r_{u(s,i)}|^2}{BN_0}$. To convexify (6) (for the short packet transmission), we replace $\rho^k_{r,u(s,i)}$ with $\hat{\rho}^k_{r,u(s,i)}$ in (7). So, a lower bound for (6) is given that is a concave function.

$$
\begin{aligned}
\bar{\mathcal{R}}^r_{u(s_j,i)} &= \sum_{k=1}^{K_{s_j}} B(\log_2(1+\bar{\rho}^k_{u(s_2,i)}) - \hat{\zeta}^k_{u(s_j,i)}) e^k_{u(s_j,i)} \\
\bar{\mathcal{R}}_{u(s_j,i)} &= \sum_{r=1}^{R} \bar{\mathcal{R}}^r_{u(s_j,i)} \\
\hat{\zeta}^k_{u(s_j,i)} &= log_2(e) Q^{-1}(\epsilon) \sqrt{\frac{\hat{\mathfrak{C}}^k_{u(s_j,i)}}{N^k_{u(s_j,i)}}}) \\
\hat{\mathfrak{C}}^k_{u(s_j,i)} &= 1 - \frac{1}{(1+\hat{\rho}^k_{u(s_j,i)})^2}
\end{aligned}
\tag{17}
$$

Consider UPF, O-CU and O-DU have the same processor (for simplification), so we have $\mu_s = \mu^u_s \approx \mu^c_s \approx \mu^d_s$. Moreover, as mentioned before, the mean arrival data rate of the UPF layer for a service s (α^U_s) is approximately equal to the mean arrival data rate of the O-CU-UP layer (α^C_s) and O-DU (α^D_s). so $\alpha_s = \alpha^U_s \approx \alpha^C_s \approx \alpha^D_s$. So the given assumption leads to have same energy for each layer $\phi^u_s = \phi^c_s = \phi^d_s$. As a result of these assumption, for simplicity, we can assume that $M_s = M^u_s = M^c_s = M^d_s$ Using the above assumption, we have $T^s_{DU} = T^s_{CU} = T^s_{UPF}$

$$
\begin{aligned}
T^s_{process} &= T^s_{RU} + T^s_{DU} + T^s_{CU} + T^s_{UPF} \\
T^s_{process} &= T^s_{RU} + 3 \times T^s_{DU}.
\end{aligned}
\tag{18}
$$

Lemma 1. *In problem* (16)*, the constraint* (16f) *can be reformulated as below* $\forall i, \forall s$

$$
\begin{aligned}
T^s_{max} &\geq \frac{1}{R_{u(s,i)} - \lambda_{u(s,i)}} + \frac{3}{\mu_s - \alpha_s/M_s} \\
M_s &\geq \frac{\alpha_s(T^s_{max}R_{u(s,i)} - T^s_{max}\lambda_{u(s,i)} - 1)}{(T^s_{max}\mu_s - 3)(R_{u(s,i)} - \lambda_{u(s,i)}) - \mu_s}
\end{aligned}
\tag{19}
$$

Also from equation (16m)*,* (16g) *and* (16i) *we have*

$$
0 \leq M_s \leq \min\{M^{max}, \alpha_s/\mu_s, \phi_{max}/3\phi_s\} \tag{20}
$$

We denote $\mathfrak{M}_s = \min\{M^{max}, \alpha_s/\mu_s, \phi_{max}/3\phi_s\}$*. Thus, if we restrict* (16f) *to equality we have*

$$
0 \leq \frac{\alpha_s(T^s_{max}R_{u(s,i)} - T^s_{max}\lambda_{u(s,i)} - 1)}{(T^s_{max}\mu_s - 3)(R_{u(s,i)} - \lambda_{u(s,i)}) - \mu_s} \leq \mathfrak{M}_s \tag{21}
$$

In (21)*,* $0 \leq \frac{\alpha_s(T^s_{max}R_{u(s,i)} - T^s_{max}\lambda_{u(s,i)} - 1)}{(T^s_{max}\mu_s - 3)(R_{u(s,i)} - \lambda_{u(s,i)}) - \mu_s}$ *is established due to the fact that the numerator and the denominator will both have same sign. Using* (16h)*, in numerator,* $\alpha_s \geq 0$*,* $R_{u(s,i)} - \lambda_{u(s,i)} \geq 0$ *and to simplify the problem, assume* $(R_{u(s,i)} - \lambda_{u(s,i)})T^s_{max} \geq 1$ *since the order of* T^s_{max} *is about milli second and the difference between achievable rate and packet rate can be more than* $1/T^s_{max}$*. Therefore, we restrict constraint* (16h) *to* $R_{u(s,i)} \geq \lambda_{u(s,i)} + 1/T^s_{max}$*. So the numerator is positive. In denominator, it can be said approximately that* $(T^s_{max}\mu_s)(R_{u(s,i)} - \lambda_{u(s,i)}) - \mu_s \geq 0$*, since,* $(R_{u(s,i)} - \lambda_{u(s,i)}) \geq 1/T^s_{max}$ *as mentioned above. Therefore, we need to have constraints below*

$$
\frac{\alpha_s(T^s_{max}R_{u(s,i)} - T^s_{max}\lambda_{u(s,i)} - 1)}{(T^s_{max}\mu_s - 3)(R_{u(s,i)} - \lambda_{u(s,i)}) - \mu_s} \leq \mathfrak{M}_s \tag{22}
$$

So by reformulating the equation (22)*, we have a new constraint* $\forall i, \forall s$ *as below,*

$$
\begin{aligned}
\mathcal{R}_{u(s,i)} &\geq \frac{\mathfrak{M}_s((T^s_{max}\mu_s - 3)\lambda_{u(s,i)} + \mu_s) - \alpha_s(T^s_{max}\lambda_{u(s,i)} + 1)}{\mathfrak{M}_s(T^s_{max}\mu_s - 3) - \alpha_s T^s_{max}}, \\
\varpi_{u(s,i)} &= \frac{\mathfrak{M}_s((T^s_{max}\mu_s - 3)\lambda_{u(s,i)} + \mu_s) - \alpha_s(T^s_{max}\lambda_{u(s,i)} + 1)}{\mathfrak{M}_s(T^s_{max}\mu_s - 3) - \alpha_s T^s_{max}}, \\
\mathcal{R}_{u(s,i)} &\geq \varpi_{u(s,i)}.
\end{aligned}
\tag{23}
$$

In addition, we denote $\mathtt{M}_{u(s,i)} = \frac{\alpha_s(T^s_{max}R_{u(s,i)} - T^s_{max}\lambda_{u(s,i)} - 1)}{(T^s_{max}\mu_s - 3)(R_{u(s,i)} - \lambda_{u(s,i)}) - \mu_s}$*. So we have,*

$$
M_s = \max\{\mathtt{M}_{u(s,i)} | i \in 1, 2, ..., U_s\} \quad \forall s. \tag{24}
$$

Despite simplifying the problem (16), it is still non-convex and hard to be solved. So the simplest approach is to relax $\mathbf{E}$ into continuous value $e^k_{r,u(s,i)} \in [0,1]\, \forall s, \forall i, \forall r, \forall k$ Furthermore, the problem can be solved using the Lagrangian function and iterative algorithm.

In order to make (16) as a standard form of a convex optimization problem, it is required to change the variable of equations (9) to $P_r = \sigma_{q_r}^2 \times 2^{C^r}$ so the constraint (16e) is changed to $P_r \leq \sigma_{q_r}^2 \times 2^{C^r_{max}}$. The combination of equations (16d) and (16e) leads to the following equation

$$\begin{aligned} &\zeta_r = min\{P_{max}, \sigma_{q_r}^2 \times 2^{C^r_{max}}\}, \\ &P_r \leq \zeta_r. \end{aligned} \tag{25}$$

Moreover, the combination of equations (16d), (16h) and (23) leads to the following equation

$$\begin{aligned} &\eta_{u_{(s,i)}} = max\{\mathcal{R}^{max}_{u_{(s,i)}}, \lambda_{u_{(s,i)}} + 1/T^s_{max}, \varpi_{u(s,i)}\}, \\ &\bar{\mathcal{R}}_{u_{(s,i)}} \geq \eta_{u_{(s,i)}}. \end{aligned} \tag{26}$$

Assume $\mathfrak{v}$, $\mathfrak{m}$, $\mathfrak{h}$, $\boldsymbol{\xi}$, $\boldsymbol{\chi}$ and $\boldsymbol{\kappa}$ are the matrix of Lagrangian multipliers that have non-zero positive elements.

The Lagrangian function is written as follow

$$\mathcal{L}(\boldsymbol{P}, \boldsymbol{E}; \mathfrak{v}, \boldsymbol{\chi}, \mathfrak{h}, \boldsymbol{\xi}, \boldsymbol{\kappa}, \mathfrak{m}) = \sum_{s=1}^{S}\sum_{i=1}^{U_s} \delta_s \bar{\mathcal{R}}_{u_{(s,i)}} \tag{27a}$$

$$+ \sum_{s=1}^{S}\sum_{i=1}^{U_s} \mathfrak{h}_{u_{(s,i)}} (\bar{\mathcal{R}}_{u_{(s,i)}} - \eta_{u_{(s,i)}}) \tag{27b}$$

$$- \sum_{r=1}^{R} \mathfrak{m}_r (P_r - \zeta_r) \tag{27c}$$

$$+ \sum_{s=1}^{S}\sum_{i=1}^{U_s}\sum_{k=1}^{K}\sum_{r=1}^{R} \kappa^k_{r,u(s,i)} p^k_{r,u(s,i)} \tag{27d}$$

$$+ \sum_{r=1}^{R}\sum_{s=1}^{S}\sum_{i=1}^{U_s} \chi_{r,u(s,i)} (\sum_{k=1}^{K_s} e^k_{r,u(s,i)} - 1) \tag{27e}$$

$$- \sum_{s=1}^{S}\sum_{i=1}^{U_s}\sum_{k=1}^{K}\sum_{r=1}^{R} \mathfrak{v}^k_{r,u(s,i)} (e^k_{r,u(s,i)} - 1) \tag{27f}$$

$$+ \sum_{s=1}^{S}\sum_{i=1}^{U_s}\sum_{k=1}^{K}\sum_{r=1}^{R} \xi^k_{r,u(s,i)} e^k_{r,u(s,i)}. \tag{27g}$$

Lemma 2. *By taking derivatives of* (27) *(the lagrangian function), with respect to the* $\boldsymbol{P}$ *and the* $\boldsymbol{E}$*, these two variables are obtained. Assume,* $e^k_{r,u(s,i)} = 1$

$$\frac{\partial \mathcal{L}}{\partial p^k_{r,u(s,i)}} = (\delta_s + \mathfrak{h}_{u_{(s,i)}}) \mathfrak{B}^k_{r,u(s,i)} + (\kappa^k_{r,u(s,i)} - \mathfrak{m}_r \mathfrak{D}^k_{r,u(s,i)}) = 0 \tag{28}$$

Where

$$\begin{aligned} &\mathfrak{D}^k_{r,u(s,i)} = |\mathbf{w}^k_{r,u(s,i)}|^2 g^r_{u(s,i)} e^k_{r,u(s,i)}, \\ &\mathfrak{B}^k_{r,u(s,i)} = \frac{B |\mathbf{h}^{H\,k}_{r,u(s,i)} \mathbf{w}^k_{r,u(s,i)}|^2 g^r_{u(s,i)} e^k_{r,u(s,i)}}{\ln(2)} \mathfrak{S}^k_{r,u(s,i)}, \\ &\mathfrak{S}^k_{r,u(s,i)} = \frac{1}{|\mathbf{h}^{H\,k}_{r,u(s,i)} \mathbf{w}^k_{r,u(s,i)}|^2 \mathfrak{k}^k_{r,u(s,i)} + BN_0 + I^k_{r,u(s,i)}}. \end{aligned} \tag{29}$$

where $\mathfrak{k}^k_{r,u(s,i)} = g^r_{u(s,i)} e^k_{r,u(s,i)} p^k_{r,u(s,i)}$*. Thus, from equation* (28)*, optimal power is obtained and power is allocated. We denote* $\mathfrak{j}^k_{r,u(s,i)} = g^r_{u(s,i)} e^k_{r,u(s,i)}$*.*

$$p^k_{r,u(s,i)} = [\frac{(\delta_s + \mathfrak{h}_{u_{(s,i)}}) B \mathfrak{j}^k_{r,u(s,i)}}{\kappa^k_{r,u(s,i)} - \mathfrak{m}_r \mathfrak{D}^k_{r,u(s,i)}} - \frac{BN_0 + I^k_{r,u(s,i)}}{|\mathbf{h}^{H\,k}_{r,u(s,i)} \mathbf{w}^k_{r,u(s,i)}|^2 \mathfrak{j}^k_{r,u(s,i)}}]^+. \tag{30}$$

Also $[a]^+ = \max(0, a)$*. In addition, PRB assignment is obtained as follow*

$$\begin{aligned} \frac{\partial \mathcal{L}}{\partial e^k_{r,u(s,i)}} &= \bar{\mathcal{R}}^k_{r,u(s,i)} (\delta_s + \mathfrak{h}_{u_{(s,i)}}) \\ &- \mathfrak{m}_r |\mathbf{w}^k_{r,u(s,i)}|^2 p^k_{r,u(s,i)} g^r_{u(s,i)} \\ &+ (\xi^k_{r,u(s,i)} - \mathfrak{v}^k_{r,u(s,i)} + \chi_{r,u(s,i)}) = 0. \end{aligned} \tag{31}$$

Using KKT conditions, we have

$$e^k_{r,u(s,i)} \times (\mathfrak{F}^k_{r,u(s,i)} - \mathfrak{v}^k_{r,u(s,i)} - \mathfrak{m}_r |\mathbf{w}^k_{r,u(s,i)}|^2 p^k_{r,u(s,i)} g^r_{u(s,i)}) = 0. \tag{32}$$

Where $\mathfrak{F}^k_{r,u(s,i)} = \bar{\mathcal{R}}^k_{r,u(s,i)} (\delta_s + \mathfrak{h}_{u_{(s,i)}}) + (\xi^k_{r,u(s,i)} + \chi_{r,u(s,i)})$*. Hence, from equation* (31) *and* (32)*, PRB assignment is performed as follow.*

$$e^k_{r,u(s,i)} = \begin{cases} 1 & u(s,i) = argmax \mathfrak{F}^k_{r,u(s,i)} \forall s, \forall r, \forall k \\ 0 & otherwise \end{cases} \tag{33}$$

Thus the user in each slice s that have the the most considerable value of $\mathfrak{F}^k_{r,u(s,i)}$, should be allocated to the PRB k. Since just one PRB can be allocated to a UE between those UEs (regardless to the services) that are associated to the same O-RU.

B. Sub-Problem 2

After power allocation and PRB assignment, the remaining problem is to assign O-RU to the UE in each service.

Assume $\boldsymbol{P}$ and $\boldsymbol{E}$ are fixed, we want to find $\boldsymbol{G}$. Next, we introduce a greedy algorithm that assigns one O-RU to each UE.

Greedy Algorithm for Non-Comp O-RU Assignment: The problem can be reformulated as follow

$$\max_{\boldsymbol{G}} \quad \sum_{s=1}^{S}\sum_{i=1}^{U_s}\sum_{r=1}^{R} \delta_s g^r_{u(s,i)} \bar{\mathcal{R}}^r_{u_{(s,k)}} \tag{34a}$$

$$\text{subject to} \quad \sum_{s=1}^{S}\sum_{i=1}^{U_s} g^r_{u(s,i)} \psi_{r,u(s,i)} \leq \mathfrak{t}_r \quad \forall r \tag{34b}$$

$$\sum_{r} g^r_{u(s,i)} = 1 \quad \forall s, \forall i, \tag{34c}$$

$$g^r_{u(s,i)} \in \{0, 1\} \quad \forall s, \forall i, \tag{34d}$$

Where $\psi_{r,u(s,i)} = \sum_{k=1}^{K_s} |\mathbf{w}^k_{r,u(s,i)}|^2 p^k_{r,u(s,i)} e^k_{r,u(s,i)}$ and $\mathfrak{t}_r = \zeta_r - \sigma_r$ Since we obtained (26) in (III-A), we can ignore this constraint in (34). The problem (34) is an NP-complete 0-1 multiple knapsack problem. We solve this problem using GAAOU, which is a greedy algorithm (1) as follows [15], [31]. Firstly, we set all variables $g^r_{u(s,i)} = 0, \quad \forall s, \forall i, \forall r$. Then we define $\mathfrak{B}^{rem}_{u_{(s,i)}} = \mathcal{R}$ $\forall s, \forall i$ and $\mathfrak{C}_r = \mathfrak{t}_r, \forall r$ as a set of all O-RUs and value of each O-RU, respectively. Next, we sort all slices based on their priority. Afterward, we assign the O-RU that provides the highest achievable rate for each UE (we start from the UEs on the slices with the highest priority) on the condition that it does not exceed the value of each O-RU (that is, a function of maximum power and capacity of O-RU). If it exceeds the value of O-RU, then O-RU with the next highest achievable rate is selected. The

complexity of sorting S slices based on their priority is $O(Slog(S))$. Depict $\mathfrak{N} = \sum_{s=1}^{S} \sum_{i=1}^{U_s} 1$. The complexity of this algorithm is about $O(Slog(S)) + O(R \times \mathfrak{N})$.

Algorithm 1 Greedy Algorithm for Assignment of O-RU to UEs (GAAOU)

Set $g^r_{u(s,i)} = 0, \quad \forall s, \forall i, \forall r.$
Set $\mathfrak{C}_r = \mathfrak{t}_r, \forall r$
Set $\mathfrak{B}^{rem}_{u(s,i)} = \mathcal{R}\ \forall s, \forall i$
Sort slices according to their priority factor (δ_s) in descending order
for $s \leftarrow 1$ to S **do**
  for $i \leftarrow 1$ to U_s **do**
    $RU = 0$
    for $r \leftarrow 1$ to R **do**
      Acquire $\mathfrak{G}^r_{u_{(s,i)}} = \bar{\mathcal{R}}^r_{u_{(s,i)}}$
    end for
    Obtain $r^* = \text{argmax}_{r \in \mathfrak{B}^{rem}_{u(s,i)}} \mathfrak{G}^r_{u_{(s,i)}}$
    while $RU == 0$ **do**
      if $\mathfrak{C}_{r^*} \geq \psi_{r^*,u(s,i)}$ **then**
        Set $g^{r^*}_{u(s,i)} = 1$
        Set $\mathfrak{C}_{r^*} = \mathfrak{C}_{r^*} - \psi_{r^*,u(s,i)}$
        Set $RU = 1$
      else
        $\mathfrak{B}^{rem}_{u(s,i)} = \mathcal{R} \setminus \{r^*\}$
      end if
    end while
  end for
end for

C. Iterative Proposed Algorithm

In (III-A) and (III-B), the details of solving each sub-problem are depicted. Here, the iterative algorithm for the whole problem is demonstrated. Firstly, we fixed $\boldsymbol{G}$, then $\boldsymbol{P}$ and $\boldsymbol{E}$ is achieved using the Lagrangian method. Afterward, $\boldsymbol{G}$ is updated using the GAAOU algorithm. This process is repeated until it converges. The whole algorithm is depicted as follows (Algorithm (2)).

Algorithm 2 Iterative Algorithm for Power Allocation, PRB, VNF and O-RU Association (IAPPVO)

Set the maximum number of iterations $Iter_{max}$, convergence condition $\epsilon > 0$
Assign Users to O-RU randomly (Initialize $\boldsymbol{G}$)
for $i \leftarrow 1$ to $Iter_{max}$ **do**
  Acquire $\boldsymbol{P}^{(i)}$, $\boldsymbol{E}^{(i)}$ and $\boldsymbol{M}^{(i)}$ using Lagrangian function and sub-gradient method based on (III-A)
  Update $\boldsymbol{G}^{(i)}$ based on algorithm GAAOU (1) in (III-B)
  if the algorithm converged with the tolerence of ϵ **then**
    Break
  else
    Continue the algorithm
  end if
end for

IV. Numerical Results

In this section, numerical results for the problem are depicted to evaluate the performance of the algorithms. We consider three network slices, one for eMBB and two for URLLC and mMtc. Assume we have six 4-antenna O-RU (MISO) located in a cell with a diameter of 500 meters. We consider 25 PRB in the network. The maximum number of VNF for each slice is 20 and the mean arrival rate for eMBB and URLLC is $\lambda = 300KHz$ and for mMTC $\lambda = 600KHz$. The other parameters of these simulations are depicted in Table I. Two different methods compare the performance of the method. The first one is a baseline scheme, which used random PRB allocation. The association of O-RU is based on the distance, channel quality, and the fronthaul capacity; also, the number of activated VNF for each slice is fixed. The second one is similar to the FBDR algorithm proposed in [15]. In this method, PRB and power are dynamically allocated, the number of VNFs is fixed, and the UEs are associated with O-RU based on the quality of their channels, the fronthaul capacity, and channel distance and quality. In Fig. 2 the aggregate throughput is demonstrated versus the different number of UEs in each service for these three methods. Suppose we have one service instance for each type of service, so we have three various services in this figure. Here, we did not consider the priority. The figure presented that the proposed algorithm is 18.6% higher than the baseline scheme. As the number of UEs increases in each service, the aggregated throughput initially increases. Still, due to the interference and the power constraint, it will be saturated from 12 UEs in each service.

TABLE I
Simulation Parameter

Parameter	Value
Noise power	-174dBm
Bandwidth	180 KHz
Maximum transmit Power of each O-RU	38dBm
Maximum delay for eMBB	4msec
Maximum delay for URLLC	1msec
Maximum delay for mMTC	20msec
Maximum fronthaul capacity	200 bits/sec/Hz
Minimum data rate for eMBB	20 bits/sec/Hz
Minimum data rate for URLLC and mMTC	5 bits/sec/Hz
Maximum received power for mMTC	20 dBm
Maximum received power for eMBB and URLLC	33 dBm

Figure 3, depicts the number of activated VNF for the five different mean service times of one URLLC service vs. the mean arrival time for 12 UEs. This figure presented that as the mean arrival rate increases, the number of activated VNF increases. Moreover, the number of activated VNFs decreased when the mean service rate increases.

In figure 4, the aggregate throughput is depicted vs. the maximum power of UE for three different instances of eMBB service using proposed method. Here, we suppose

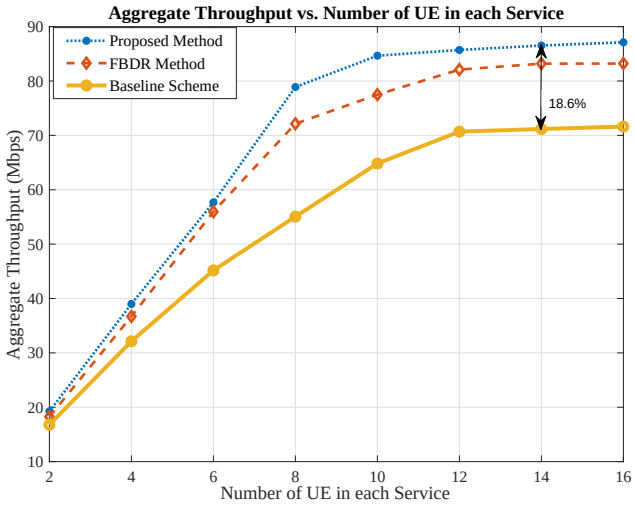

Fig. 2. Aggregate Throughput vs. Number of UE in each Service

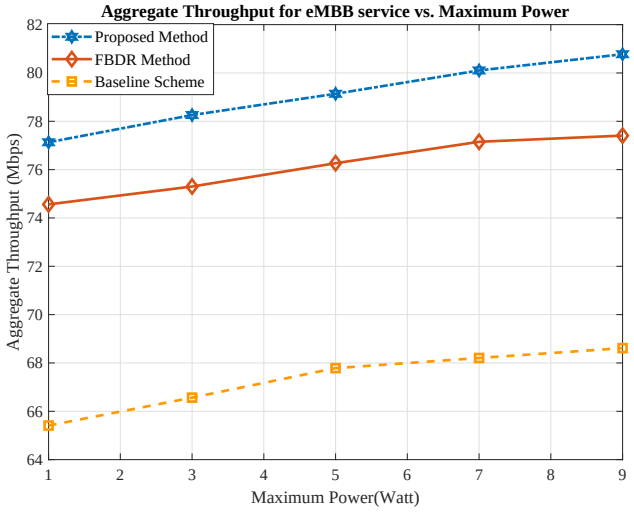

Fig. 4. Aggregate Throughput for eMBB vs. Maximum Transmit power for various number of UEs

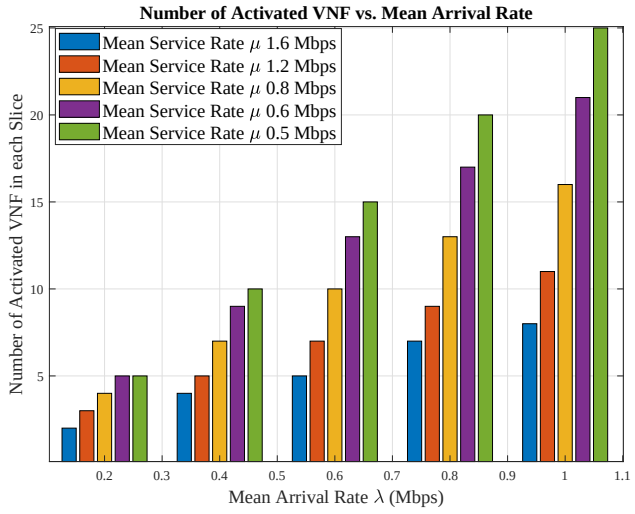

Fig. 3. Number of activated VNF in each Service vs. Mean Arrival Rate(Mbps)

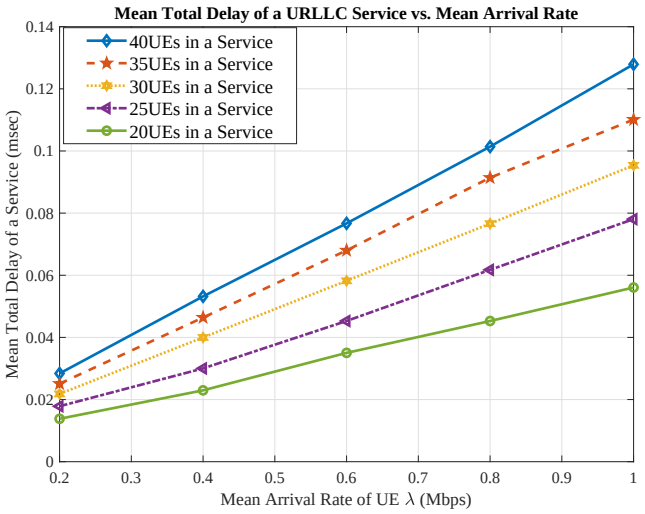

Fig. 5. Mean Total Delay of a URLCC Service vs. the Mean Arrival Rate of a UE in the Service for various number of UEs

that we have 12 UEs in each service. We assume that these three services require 5bits/sec/Hz, 10bits/sec/Hz, and 15bits/sec/Hz. In addition, We suppose that each O-RU can transmit three times of the maximum power of UE. As you can see in the figure, increasing the maximum power increases the aggregate throughput.

Figure 5, illustrates the mean total delay of a UE in a URLLC service regarding the mean arrival rate of the UE and the number of UEs in the service for proposed method. It is shown that the delay is an ascending function of the mean arrival rate and the number of UEs in the service. In this figure, we assume that the maximum number of VNF for each slice is 50 and the maximum delay of each UE in a URLLC service is $0.5ms$. Also, the maximum number of PRB is considered to be 50. Moreover, we can see that the mean delay of a URLLC service does not reach the maximum threshold of the delay. Figure 6 is the same as figure 6 that presented the mean total delay of a UE in a URLLC service regarding the mean arrival rate of the UE for 20 UEs using three different methods. As you can see, the proposed method (IAPPVO) outperforms the other scenarios.

Figure 7, represents the aggregate throughput concerning the number of UEs and the maximum power for three different mMTC services. mMTC services include a large number of UEs with low data rates and low power. Assume each service instance requires $1bits/sec/Hz$ data rate and is not sensitive to end-to-end delay. The figure depicts that by increasing the number of UEs in each instance of the service, the aggregate throughput increases. Also, by increasing the maximum power of each UE in each instance of mMTC service, the aggregate throughput rises too.

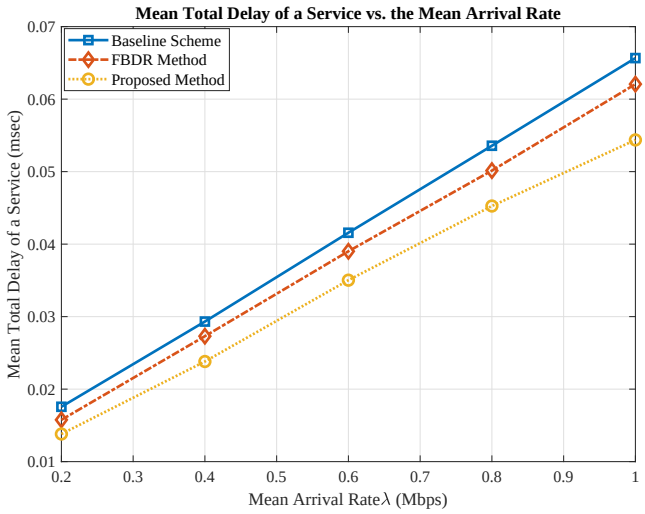

Fig. 6. Mean Total Delay of a URLCC Service vs. the Mean Arrival Rate of a UE in the Service for different methods

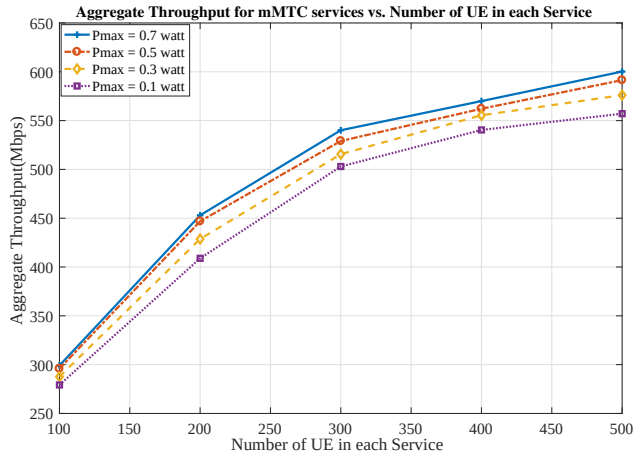

Fig. 7. Aggregate Throughput for mMTC services vs. Number of UE in each Service for three different mMTC service instances

V. Conclusion

This paper proposed the downlink of the O-RAN system using network slicing for different types of 5G services (eMBB, mMTC, and URLC). The isolation of various types of services (eMBB, mMTC, and URLC) in O-DU, O-CU, and user plane function (UPF) is accomplished. Also, the paper aims to obtain the number of activated VNFs in each service, RU association, power, and PRB allocation to maximize the sum rate. The limited fronthaul capacity and the mean end-to-end delay for each service are considered. The problem is mixed-integer non-linear programming that is solved by the two-step iterative algorithm. In the first step, we reformulated the problem to achieve the number of activated VNFs as a function of data rate. Then we obtain PRB association and power allocation using the Lagrangian method. Then in the second step, O-RU association is acquired. The performance of our proposed method (IAPPVO) is compared with the baseline scheme and FBDR in [15]. Also, we assume distinct scenarios for each service (eMBB, URLLC, and mMTC) based on their requirement QoS. Simulation results depicts that the proposed method (IAPPVO) achieved 18.6% higher data rate than the baseline scheme. Moreover, simulation results illustrate less delay for the proposed method (IAPPVO) than FBDR and baseline scheme.

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