

"No theorist —

" — Craig Hogan

" — Scott Aaronson

* Ordinary Prob Theory

States

$|H\rangle$

$|T\rangle$

$$\frac{1}{2}|H\rangle + \frac{1}{2}|T\rangle$$

$$\frac{1}{3} \cancel{|H\rangle} + \frac{2}{3}|T\rangle$$

$$\text{Any } P_0|H\rangle + P_1|T\rangle$$

$$\sim P_0, P_1 \geq 0 \quad P_0 + P_1 = 1$$

Transform

Transformations

Shake: b_3 : $P_0 |H\rangle + P_1 |T\rangle \xrightarrow{\left(\frac{1}{2}(P_0 + b_3 P_1)\right)} \left(\frac{1}{2}(P_0 + b_3 P_1)\right) |H\rangle + \frac{1}{2}b_3 |T\rangle$

$$\begin{aligned} & : P_0 |H\rangle + P_1 |T\rangle \rightarrow \\ & \qquad \qquad \qquad b_3 |H\rangle + \frac{2}{3}b_3 |T\rangle \end{aligned}$$

If $|H\rangle, b_m$ in

if $|T\rangle, \text{flip}$: $P_0 |H\rangle + P_1 |T\rangle \rightarrow$

$$(P_0 + \frac{1}{2}P_1) |H\rangle +$$

$$\frac{1}{2}P_1 |T\rangle$$

If $|H\rangle, \text{in } \alpha |H\rangle + \beta |T\rangle$

If $|T\rangle, \text{in } \alpha' |H\rangle + \beta' |T\rangle$

$$P_0 |H\rangle + P_1 |T\rangle \rightarrow$$

$$\alpha(P_0 + \alpha' P_1) |H\rangle +$$

$$(\beta_{P_0} \vdash \beta_{P_1}^{'}) \mid T$$

All linear

Measurement

Perfom a measurement

$$|P_0\rangle |H\rangle + |P_1\rangle |T\rangle$$

Measuring $|H\rangle$ w/ prob P_0

Ort Rules of ordinary problem w/ prob P_1 ,
Basis

Given H find set of basis states

$$\{ |x_1\rangle, \dots, |x_n\rangle \}$$

$$(\{ |H\rangle, |T\rangle \})$$

Set of all states is

$$\left\{ p_1 |x_1\rangle + \dots + p_n |x_n\rangle \mid p_1, \dots, p_n \geq 0, p_1 + \dots + p_n = 1 \right\}$$

Two upshots

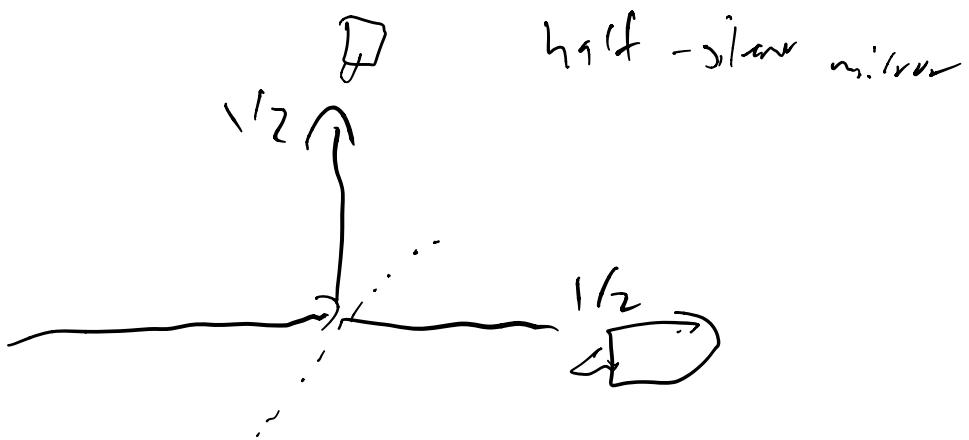
Transition : A linear maps from \mathbb{R}^n to \mathbb{R}^n

s.t. States get mapped to sum

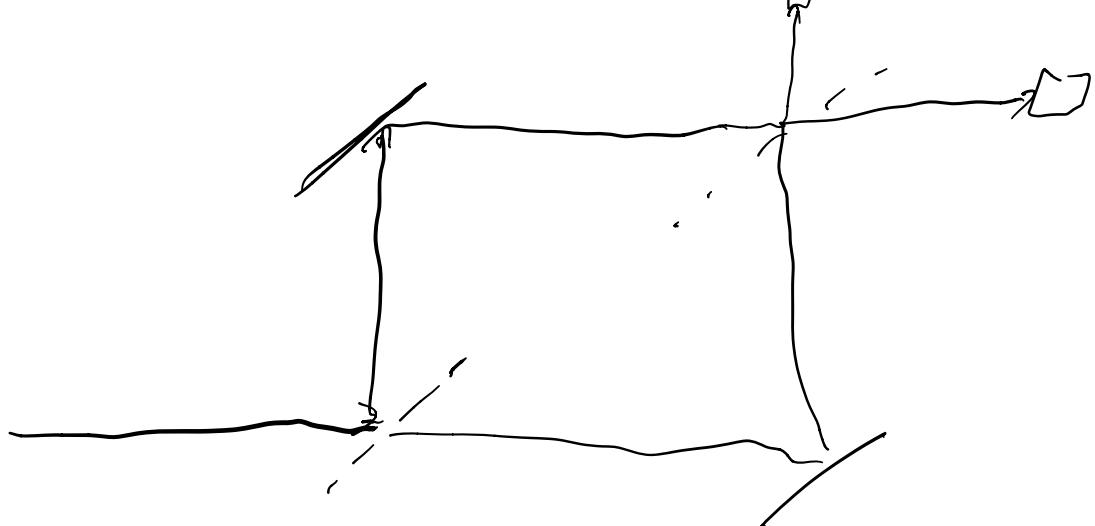
Measurement : $p_1 |x_1\rangle + \dots + p_n |x_n\rangle$

$\rightarrow |x_j\rangle$ w/ prob., p_j

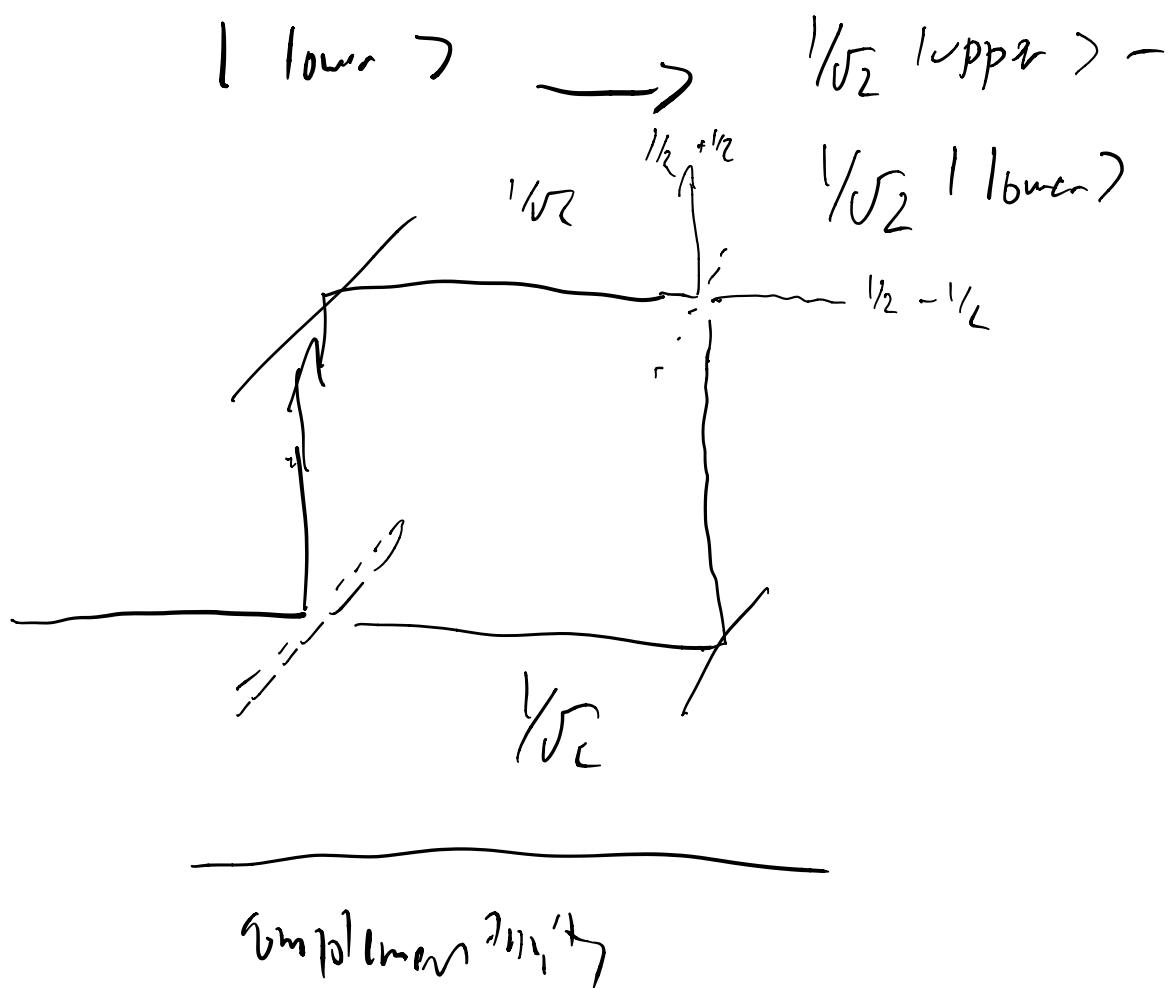
Interference



Mach-Zehnder Interferometer



$$| \text{Upper} \rangle \rightarrow \frac{1}{\sqrt{2}} | \text{Upper} \rangle + \frac{1}{\sqrt{2}} | \text{Lower} \rangle$$



partially
 partially
 electric h_{α} α spin up
 1. can be measured with my detector
 2. the output of measure will be
 even $+ \frac{1}{2}$ or $- \frac{1}{2}$

Ψ Suppose now \rightarrow

\rightarrow (Right)
 (Left)

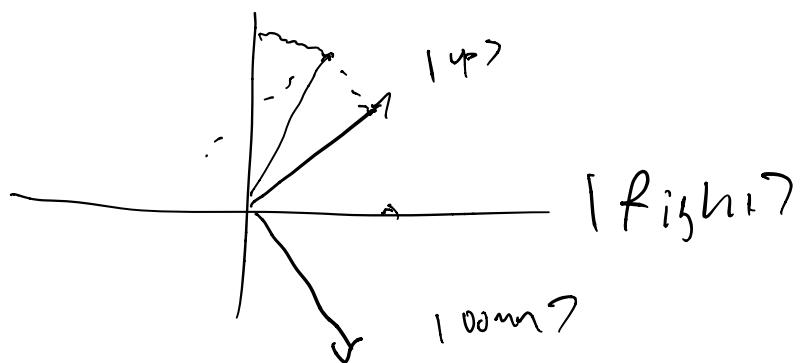
$$\alpha |Left\rangle + \beta |Right\rangle$$

$$|\alpha|^2 + |\beta|^2 = 1$$

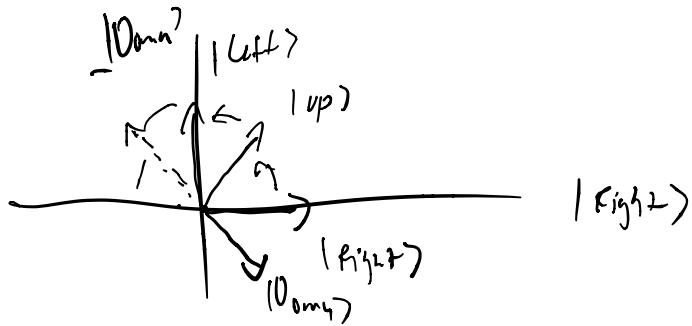
$$|Up\rangle = \frac{1}{\sqrt{2}} |Left\rangle + \frac{1}{\sqrt{2}} |Right\rangle$$

$$|Down\rangle = \frac{1}{\sqrt{2}} |Left\rangle - \frac{1}{\sqrt{2}} |Right\rangle$$

|Left>



Consider $\text{rotate } -b_y - 90^\circ$



$$|\text{Right}\rangle \rightarrow |\text{Up}\rangle \rightarrow |\text{Left}\rangle \rightarrow \\ -|\text{0mm}\rangle \rightarrow |\text{Right2}\rangle$$

$$\text{rotate } -b_y - 360^\circ (\varphi) = -\varphi$$

Performing rotation on φ is same as

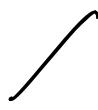
$-\varphi$

at θ $\ell - \varphi$ (roll) and sliding

In this form of superposition

$\boxed{\text{down 100\%}}$

$\square \text{ to } 25^\circ$



$\wedge 360^\circ \wedge 720^\circ$

↳



Feynman's trick

from Prof. No-Cloning Thm

"Ordinary" Probability Theory

Suppose two systems

$|A\rangle, |B\rangle$ A; w/ basis states $|a_1\rangle, \dots, |a_n\rangle$

$|C\rangle, |D\rangle$ B; w/ basis states $|b_1\rangle, \dots, |b_m\rangle$

then, the composite system ~~A ⊗ B~~^{A ⊕ B} is w/ basis states

$(a_1 b_1), (a_1 b_2), \dots, (a_1 b_n)$

$(a_2 b_1), (a_2 b_2), \dots, (a_2 b_n)$

\vdots — {

$(a_n b_1), \dots, (a_n b_m)$

⊗ expands to n_m basis states

$$\begin{aligned}
 & \left(\frac{1}{2} |H\rangle + \frac{1}{2} |T\rangle \right) \otimes \left(\frac{1}{3} |H\rangle + \frac{2}{3} |T\rangle \right) \\
 = & \frac{1}{6} |H\rangle\langle H| + \frac{1}{6} |T\rangle\langle H| \\
 & + \frac{1}{3} |H\rangle\langle T| + \frac{1}{3} |T\rangle\langle T|
 \end{aligned}$$

If a state ψ of A \otimes B can

be written as $\psi_A \otimes \psi_B$ it's called

Simple or Product or A & B are independent. Only the two parts

depend.

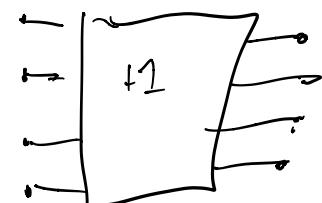
~~Possible to duplicate information~~

$(C^{1\text{qubit}})_B$ is "0 or 1"

$C^{1\text{qubit}}$ operates

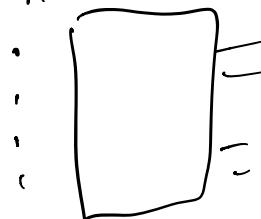
independently on

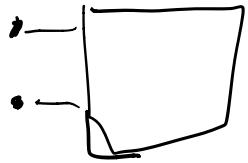
symmetries of bits



Qubit is any quantum system A w/ states $\alpha|0\rangle + \beta|1\rangle$

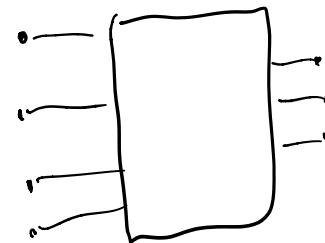
$$A = C \otimes A \otimes A \otimes A \rightarrow A \otimes A$$





Circuits
 (Classical) bit is 0 or 1 $\{0, 1\} \xrightarrow{\text{any fn fn}} \{0, 1\}$

and a classical circ



$$0 \rightarrow 1$$

$$1 \rightarrow 0$$

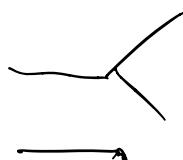


$$0, 0$$

$$0, 1$$

$$1, 0$$

$$1, 1 \rightarrow 1$$



Quantum bit (qubit)

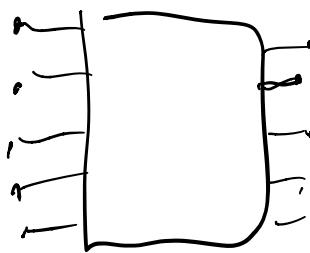
is a quantum system with

basis states $|0\rangle \dots |1\rangle$ - $\underbrace{A \otimes A \otimes A \dots A}_{n \text{ times}}$

$\alpha|0\rangle + \beta|1\rangle$

1 quantum circ

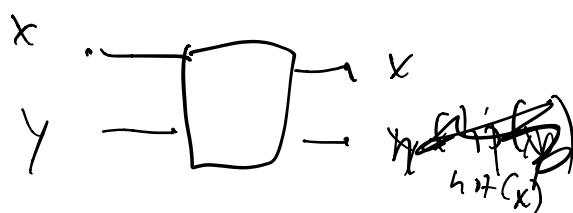
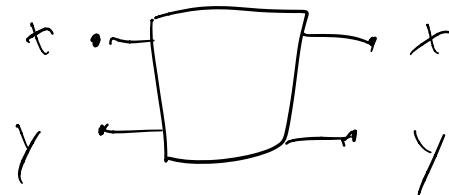
$A^{\otimes n} \rightarrow A^{\otimes n}$ unitary



Deutsch's Algorithm ("Fast or slow")

way: Artificial

4 datum items

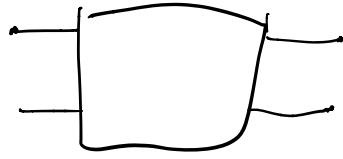


$|00\rangle \rightarrow |01\rangle$

$|01\rangle \rightarrow |10\rangle$

$\beta|10\rangle \rightarrow |11\rangle$

$|11\rangle \rightarrow |10\rangle$



Flips y if
 x is $|0\rangle$

$$|00\rangle \rightarrow |01\rangle$$

$$|01\rangle \rightarrow |10\rangle$$

$$|10\rangle \rightarrow |10\rangle$$

!

Determine which of the circuit is in the first set or second set (with looking out about

$$0\ 0 \rightarrow 0\ 0$$

$$\text{Suppose } |0\rangle = \text{Left} \quad (R_{jL})$$

$$|1\rangle = \text{Right} \quad (L_{iR})$$

$$x = |U_p\rangle = \frac{|Left\rangle + |Right\rangle}{\sqrt{2}}$$

$$y = |Down\rangle = \frac{-|Left\rangle + |Right\rangle}{\sqrt{2}}$$

$$-|11\rangle - |01\rangle + |01\rangle + |10\rangle$$

$$-|10\rangle - |00\rangle + |01\rangle + |11\rangle$$

$$= (|11\rangle + |0\rangle)(-|0\rangle +$$

$$-|11\rangle -|01\rangle + |10\rangle + |00\rangle$$

$$-|10\rangle -|01\rangle + |00\rangle + |11\rangle$$

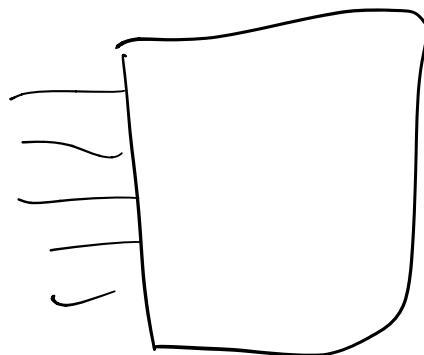
$$\underline{-|0\rangle} \\ -|10\rangle + |11\rangle - |01\rangle + |00\rangle$$

$$|11\rangle(-|0\rangle + |1\rangle) \quad * -|0\rangle(|11\rangle - |00\rangle)$$

$$-|0\rangle + |1\rangle$$

Graur's algorithm

in support size size arbitrary
black box com



Think of it as a function or lookup table

e.g., may it map from to
phon #s.

Problem: Given α find #

Classical: $O(n)$, where n is #
of perm

Quantum: $O(\sqrt{n})$

Empirical:

Throughput

How not good it is: suggests
Quantum computers can solve NP
problems efficiently

Complexity Class,
PSPACE

