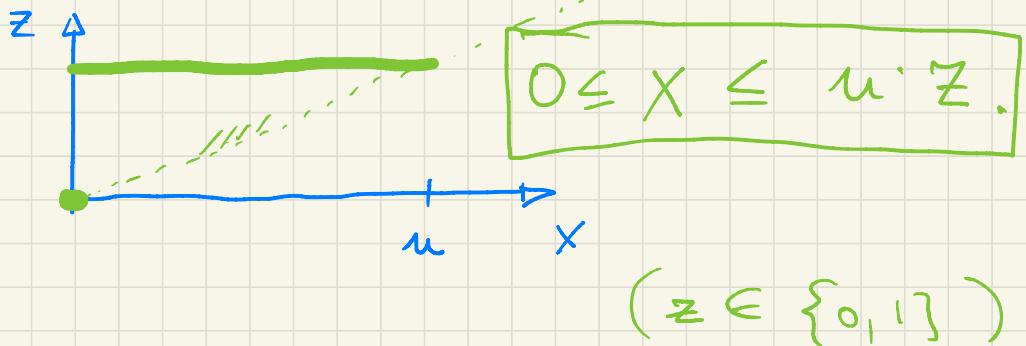


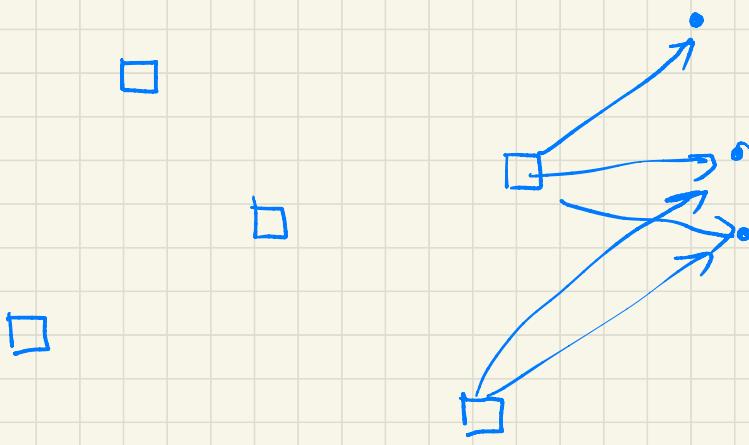
Forcing constraint :

| If binary variable  $z = 0$ ,  
| then the continuous variable  $x$  ( $x \leq u$ )  
must be zero.



Facility location:

Given sites for potential facilities  
to serve customers



Graph  $(V, A)$   $V = V_1 \cup V_2$

$\uparrow$  facilities       $\uparrow$  customers

Demands  $b_j$  for customer  $j \in V_2$   
(for a single good)

Each facility has a capacity  
(potential supply)  $u_i$  for  $i \in V_1$ .

Costs per unit of good  $c_{ij}$

Amount of good shipped from  
facility  $i$  to customer  $j$ :  $x_{ij}$ .

Cost of opening facility  $i$ :  $d_i$ .

Logic constraint:

A facility can only serve a  
customer if it is open.

Decision variable:

$$z_i = \begin{cases} 1 & \text{if we open facility } i \\ 0 & \text{o/w.} \end{cases}$$

$$\min \sum_{i \in V_0} d_i z_i + \sum_{(i,j) \in A} c_{(i,j)} x_{(i,j)}$$

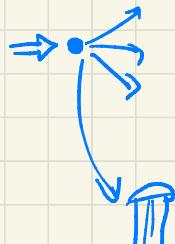
s.t.

$$\sum_{\substack{i \in V_0 \\ (i,j) \in A}} x_{(i,j)} = b_j \quad \forall j \in V_2$$

(disaggregated)  
forcing  
constraint

$$x_{(i,j)} \leq [u_i] \cdot z_i \quad \forall i \quad \forall (i,j) \in A$$

$$\boxed{\sum_{\substack{j \in V_2 \\ (i,j) \in A}} x_{(i,j)} \leq u_i}$$



total amount of goods supplied by facility  $i$ .

Alternative:

Combine them

$$\sum_{\substack{j \in V_2 \\ (i,j) \in A}} x_{(i,j)} \leq u_i \cdot z_i$$

## (aggregated forcing constraint)

- if  $z_i = 1$ , then this is the facility capacity constraint.
- if  $z_i = 0$ , then it forces the aggregated flow to 0, together with  $x_{(i,j)} \geq 0$  it implies that  $x_{(i,j)} = 0 \quad \forall j$ .

Aggregated formulation has fewer constraints, but it is a "weaker" formulation.

Relationship between strong and weak formulations of the same problem.

Revisit modeling of logical AND.

Correct formulation :

$$\begin{array}{ccc} 0 & 0 & 1 \end{array} \quad x_i + x_j + (1 - y_{ij}) \geq 1$$

$$\Leftrightarrow x_i + x_j \geq y_{ij}$$

$$\begin{array}{ccc} 0 & 1 & 1 \end{array} \quad x_i + (1 - x_j) + (1 - y_{ij}) \geq 1$$

$$\Leftrightarrow x_i - x_j - y_{ij} \geq -1$$

$$\begin{array}{ccc} 1 & 0 & 1 \end{array} \quad (1 - x_i) + x_j + (1 - y_{ij}) \geq 1$$

$$\begin{array}{ccc} 1 & 1 & 0 \end{array} \quad (1 - x_i) + (1 - x_j) + y_{ij} \geq 1$$

$$\Leftrightarrow y_{ij} \geq -x_i + x_j - 1$$



1st constraint and 2nd constraint together:

Sum these constraints:

$$x_i + x_j \geq y_{ij}$$

$$x_i - x_j - y_{ij} \geq -1$$

$$2x_i - y_{ij} \geq y_{ij} - 1$$

$$2x_i - 2y_{ij} \geq -1$$

$$x_i - y_{ij} \geq -\frac{1}{2}$$

- valid inequality for all points that satisfy the given ineqs.

$$x_i \in \mathbb{Z}, y_{ij} \in \mathbb{Z}$$



$$x_i - y_{ij} \in \mathbb{Z}.$$

uses integrality

$$x_i - y_{ij} \geq 0$$



$$y_{ij} \leq x_i.$$

(by combining 1st & 3rd ineq.) get  $y_{ij} \leq x_j.$ )

More general : Chvátal–Gomory cut.

Take integer variables  $x \in \mathbb{Z}_{\geq 0}^n$

and an inequality system

$$Ax \leq b. \quad \left( \sum_{j=1}^m a_{ij} x_j \leq b_i \right) \text{ in } \boxed{A}$$

Any nonnegative vector of multipliers  $\gamma \in \mathbb{R}_{\geq 0}^m$  gives a valid inequality

$$\gamma^T A x \leq \gamma^T b$$

$$\sum_{j=1}^m \left( \sum_{i=1}^m \gamma_i a_{ij} \right) x_j = \sum_{i=1}^m \gamma_i \sum_{j=1}^m a_{ij} x_j \leq \sum \gamma_i b_i$$

If  $\gamma^T A \in \mathbb{Z}^m$  ( $\forall j, \sum_{i=1}^m \gamma_i a_{ij} \in \mathbb{Z}$ ),

then for  $x \in \mathbb{Z}_{\geq 0}^n$

$$\gamma^T A x \in \mathbb{Z}, \quad \text{so} \quad \gamma^T A x \leq \gamma^T b$$

implies  $\gamma^T A x \leq [\gamma^T b]$ .

↑  
a new valid inequality.

If  $\gamma^T A \notin \mathbb{Z}^n$ , then b/c  $x \geq 0$

$$\underbrace{[\gamma^T A]}_{\in \mathbb{Z}^n} x \leq \gamma^T A x \leq \gamma^T b$$

implies

$$[\gamma^T A] x \leq [\gamma^T b].$$

Chvátal-Gomory cut.

There are infinitely many of these cuts but the set

$$\{ x \in \mathbb{R}_{\geq 0}^m : Ax \leq b,$$

$$\forall \gamma \in \mathbb{R}_{\geq 0}^m : [\gamma^T A] x \leq [\gamma^T b] \}$$

is a polyhedron.

"1st Chvátal closure".