

What if we have inequalities.

$$\begin{array}{ll}
 \max & 5x_1 + 4x_2 + 3x_3 \\
 \text{s.t.} & 2x_1 + 3x_2 + x_3 \leq 5 \\
 & 4x_1 + x_2 + 2x_3 \leq 11 \\
 & 3x_1 + 4x_2 + 2x_3 \leq 8 \\
 & x_1, x_2, x_3 \geq 0
 \end{array}$$

"subject to (constraints)"

nonnegative RHS

Standard form LP.

( $\geq$  variables,  
 $\leq$  constraints,

( $\in \mathbb{R}$ )

Idea: Introduce a variable for the "slack" in each  $\leq$  constraint. "a slack variable".

Names:  $x_4 \geq 0$  - slack in 1st constraint

$x_5 \geq 0$  2nd

$x_6 \geq 0$  3rd

Equivalent formulation with slack variables added:

Alternative notation:

$$\begin{array}{ll}
 \max & 5x_1 + 4x_2 + 3x_3 \\
 \text{s.t.} & 2x_1 + 3x_2 + x_3 + \boxed{x_4} = 5 \\
 & 4x_1 + x_2 + 2x_3 + \boxed{x_5} = 11 \\
 & 3x_1 + 4x_2 + 2x_3 + \boxed{x_6} = 8
 \end{array}$$

$$\begin{array}{l}
 || \quad || \quad x_1 + x_2 + 2x_3 \\
 || \quad || \quad 3x_1 + 4x_2 + 2x_3 \\
 || \quad || \quad x_1, x_2, x_3, x_4, x_5, x_6 \geq 0.
 \end{array}
 \quad
 \begin{array}{l}
 [x_5] \quad w_3 = 11 \\
 + [x_6] = 8
 \end{array}$$

This is an LP in standard <sup>equation</sup> form.

$$A = \begin{array}{c} x_1 \quad x_2 \quad x_3 \quad x_4 \quad x_5 \quad x_6 \\ \left( \begin{array}{cccccc} 2 & 3 & 1 & 1 & . & . \\ 4 & 1 & 2 & . & 1 & . \\ 3 & 4 & 2 & . & . & 1 \end{array} \right) \end{array} \Rightarrow \text{full row rank (3).}$$

identity matrix

Can set all "original variables"  $x_1, x_2, x_3$  to 0.

$\Rightarrow$  feasible solution. b/c this vector  $b \geq 0$ .

► Initial assumption (to get started w/ algorithm)!

After introducing slack variables:

$x_1, x_2, x_3$  — set to 0.

$x_4, x_5, x_6$  — set them to "right values"

... by writing  $x_4, x_5, x_6$  as functions

positive

of  $x_1, x_2, x_3$ .

max

$$0 + 5x_1 + 4x_2 + 3x_3$$

$$x_4 = 5 - 2x_1 - 3x_2 - x_3$$

$$x_5 = 11 - 4x_1 - x_2 - 2x_3$$

$$x_6 = 8 - 3x_1 - 4x_2 - 2x_3$$

The "dictionary" corresponding to the "basis" with variables  $x_4, x_5, x_6$

"basic" variables  $x_i$  for  $i \in B$

Claim / Invariant:  
The system of

"nonbasic variables"  $x_j$  for  $j \in N$

A dictionary defines a particular solution  
"the basic solution"  
— set nonbasic variables

linear equations  
is equivalent to  
the original system.

to 0

set basic variables  
to right-hand side  
constant

In notation with variables

$\underbrace{x_1, \dots, x_n}_{\text{decision variables}}, \underbrace{x_{n+1}, \dots, x_{n+m}}_{\text{slack variables}}$

Use index set  
for all variables:

$\mathcal{B} \subseteq \{1, \dots, n+m\}$

Index sets for basic variables:  
nonbasic variables:

Partition of index set:  $\{1, \dots, n+m\} = \mathcal{B} \sqcup \mathcal{N}$

$\mathcal{N}$  disjoint.

By def. of "basic solution", it satisfies the  
equations.

is it always feasible? (The one in example is.)

Feasible if right-hand side constants are  $\geq 0$

• "Feasible dictionary" — "feasible basic solution":

Is the solution in the example optimal?

idea: Find a nonbasic variable with positive obj coeff in the dictionary.

Try to increase it from 0.

Adjust the basic variables:

$$x_1 = \lambda \geq 0$$

$$x_1 = 0$$

$$x_2 = 0$$

$$x_3 = 0$$

param.  
eqn.  
of a  
ray!

$$x_4 = 5 - 2\lambda \Rightarrow \lambda \leq \frac{5}{2}$$

$$x_5 = 11 - 4\lambda \Rightarrow \lambda \leq \frac{11}{4}$$

$$x_6 = 8 - 3\lambda \Rightarrow \lambda \leq \frac{8}{3}$$

$$x(\lambda) = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 5 \\ 11 \\ 8 \end{pmatrix} + \lambda \begin{pmatrix} 1 \\ 0 \\ 0 \\ -2 \\ -4 \\ -3 \end{pmatrix}$$

Walk along the ray.

How far? As far as possible — until leaving the

feasible region. (Nonnegativities)  
of the basic variables.

Take minimum of ratios  $\Rightarrow$  maximum  $\lambda$ .

$$\min \left\{ \frac{5}{2}, \frac{11}{4}, \frac{8}{3} \right\} = \frac{5}{2}.$$

basic variable  $x_4$  attains the min. ratio.

By choosing  $\lambda$  as max. step length that keeps a point in the ray feasible, we guarantee: 1 basic variable will change to 0.

Idea: Make this basic variable non-basic  
& the chosen non-basic variable basic  
("pivot")

← "leaving"  
← "entering".

Use Gaussian elimination to rewrite the dictionary.

nonbasic

max  $0 + 5x_1 + 4x_2 + 3x_3$  ← entering

leaving  $x_4 = 5 - 2x_1 - 3x_2 - x_3$  ←

$x_5 = 11 - 4x_1 - x_2 - 2x_3$

$x_6 = 8 - 3x_1 - 4x_2 - 2x_3$

$2x_1 = 5 - 3x_2 - x_3 - x_4$

$x_1 = \frac{5}{2} - \frac{3}{2}x_2 - \frac{1}{2}x_3 - \frac{1}{2}x_4$

Rewrite so that basic variables are on LHS, nonbasic RHS.  
isolated with coefficient 1

$x_5 =$	$11$	$+ 5x_2$	$+ 0x_3$	$+ 2x_4$
$x_6 =$	$\frac{1}{2}$	$+ \frac{1}{2}x_2$	$- \frac{1}{2}x_3$	$+ \frac{3}{2}x_4$
$x_1 =$	$\frac{5}{2}$	$- \frac{3}{2}x_2$	$- \frac{1}{2}x_3$	$- \frac{1}{2}x_4$

Use equation for  $x_1$  to remove  $x_1$  from the equation of  $x_5$ :

$$x_5 = 11 - 4 \left( \frac{5}{2} - \frac{3}{2}x_2 - \frac{1}{2}x_3 - \frac{1}{2}x_4 \right) - x_2 - 2x_3$$

$$x_6 = 8 - 3 \cdot \left( \frac{5}{2} - \frac{3}{2}x_2 - \frac{1}{2}x_3 - \frac{1}{2}x_4 \right) - 4x_2 - 2x_3$$

Also rewrite objective function:

$$\begin{aligned} \max \quad z &= 5 \left( \frac{5}{2} - \frac{3}{2}x_2 - \frac{1}{2}x_3 - \frac{1}{2}x_4 \right) + 4x_2 + 3x_3 \\ &= \boxed{\frac{25}{2}} - \frac{7}{2}x_2 + \frac{1}{2}x_3 - \frac{5}{2}x_4 \quad \checkmark \end{aligned}$$

function of nonbasic variables.

New dictionary — corresponding to a new basis

$$B = \{5, 6, 1\} \Rightarrow N = \{2, 3, 4\}$$

$\Rightarrow$  Read off a basic solution:



$$\begin{pmatrix} x_5 \\ x_6 \\ x_1 \end{pmatrix} = \begin{pmatrix} 1 \\ \frac{1}{2} \\ \frac{5}{2} \end{pmatrix}$$

$$x_2 = 0$$

$$x_3 = 0$$

$$x_4 = 0.$$

$$x = \begin{pmatrix} 5/2 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1/2 \end{pmatrix}.$$

Read off obj. value:

$$z = \frac{25}{2}.$$

This new feasible solution — optimal?  
basic

Repeat the same procedure.