

Problem: Implement quick sort.

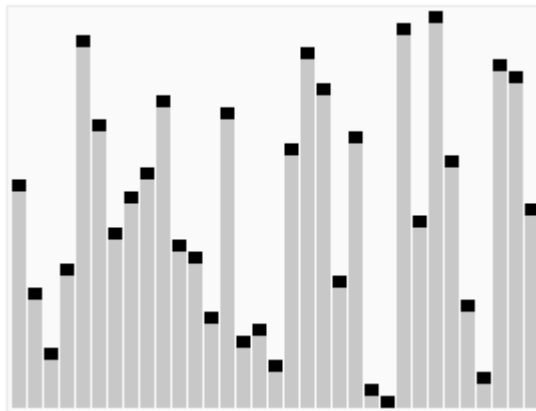
Constraints

- Is a naive solution sufficient (ie not in-place)?
 - Yes
- Are duplicates allowed?
 - Yes
- Can we assume the input is valid?
 - No
- Can we assume this fits memory?
 - Yes

Test Cases

- None -> Exception
- Empty input -> []
- One element -> [element]
- Two or more elements

Algorithm



Wikipedia's animation:

- Set pivot to the middle element in the data
- For each element:
 - If current element is the pivot, continue
 - If the element is less than the pivot, add to left array
 - Else, add to right array
- Recursively apply quicksort to the left array
- Recursively apply quicksort to the right array
- Merge the left array + pivot + right array

Complexity:

- Time: $O(n \log(n))$ average, best, $O(n^2)$ worst
- Space: $O(n)$

Misc:

- More sophisticated implementations are in-place, although they still take up recursion depth space
- Most implementations are not stable

See [Quicksort on wikipedia](#):

Typically, quicksort is significantly faster in practice than other $\Theta(n \log n)$ algorithms, because its inner loop can be efficiently implemented on most architectures [presumably because it has good cache locality], and in most real-world data, it is possible to make design choices which minimize the probability of requiring quadratic time.

See: [Quicksort vs merge sort](#)