The Alternating Method of Multipliers

From the past, with Vengeance

Operator Splitting Quadratic Programming

We will tackle our QP using the OSQP solver by Oxford University

OSQP is a modern solver for Quadratic Programs in the form:

$$\arg\min_{x} \left\{ \frac{1}{2} x^{T} P x + q^{T} x \mid l \le A x \le u \right\}$$

The solver:

- is <u>very fast</u>, especially for problems with sparse matrices
- is available under a (very permissive) Apache 2.0 license
- has API for many programming languages

The solver relies on the Alternating Direction Method of Multipliers (ADMM)

- ...Plus <u>a bunch of clever "tricks"</u> to improve speed
- Here we will discuss only the basic ADMM, to provide an intuition

The Alternating Direction Method of Multipliers

The <u>ADMM</u> solves numerical constrained optimization problems in the form:

argmin
$$f(x) + g(z)$$

subject to: $Ax + Bz = c$

lacktriangle Where f and $oldsymbol{g}$ are assumed to be convex

The methods relies on a so-called augmented Lagrangian

This is a reformulation where the constraints are turned into penalty terms:

$$\mathcal{L}_{\rho}(x, z, \lambda) = f(x) + g(z) + \lambda^{T}(Ax + Bz - c) + \frac{1}{2}\rho ||Ax + Bz - c||_{2}^{2}$$

- The algorithm idea is to optimize the augmented Lagrangian
- ...And to encourage constraint satisfaction via the penalty terms
- In practice, this is done by adjusting the multiplier vector λ

The Alternating Direction Method of Multipliers

The ADMM operates as follows

We start from an initial assignment x^0, z^0, λ^0 , then we iterate:

$$x^{k+1} = \operatorname{argmin}_{x} \mathcal{L}_{\rho}(x, z^{k}, \lambda^{k})$$

$$z^{k+1} = \operatorname{argmin}_{z} \mathcal{L}_{\rho}(x^{k+1}, z, \lambda^{k})$$

$$\lambda^{k+1} = \lambda^{k} + \rho(Ax^{k+1} + Bz^{k+1} - c)$$

In other words:

- We keep everything fixed and we optimize over x to obtain x^{k+1}
- lacktriangle We replace x^k with x^{k+1} , keep everything fixed and optimize over z
- Finally, we update the multiplier vector

The switch between x and z optimization is the "alternating" part

...While the use of the multipliers λ explains the rest of the name

Let's try to understand better the multiplier update

...Which consists in the rule:

$$\lambda^{k+1} = \lambda^k + \rho(Ax^{k+1} + Bz^{k+1} - c)$$

- The term $Ax^{k+1} + Bz^{k+1} c$ is just the current constraint violation
- ...In particular both its amount and direction

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If $(Ax^{k+1} + Bz^{k+1})_i > c_i$ for some constraint i:

- Then we increase the corresponding multiplier λ_i
- So that the penalty term $\lambda_i(Ax^{k+1}+Bz^{k+1}-c)_i$ grows
- This will push the next iteration to reduce the degree of violation

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- Then we decrease the corresponding multiplier λ_i
- So that the penalty term $\lambda_i(Ax^{k+1}+Bz^{k+1}-c)_i$ grows (again)
- This will push the next iteration to reduce the degree of violation (again)

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- ...In particular both its amount and direction

If
$$(Ax^{k+1} + Bx^{k+1})_i = c_i$$
 for some constraint i :

- Then we keep the corresponding multiplier λ_i as it is
- The constraint is not violated, so there is nothing to do

Main Advantages of the Method

The method has two major advantages:

1) The x and z variables can be handled in isolation

- This results into simpler problems
- ...And in some cases enables massive parallelization

2) The ADMM converges under relatively mild conditions

- In the classical formulation, f and g need to be closed, proper, convex functions
 - They do not need to be differentiable
 - They can take the value $+\infty$
 - We will see why that matters in the next slides
- The second condition is that $\mathcal{L}_0(x,z,\lambda)$ should have a saddle point
 - This one is way trickier to check...

The full convergence proof can be found e.g. <u>here</u>

The ADMM and QP

That was the whole point, right?

QP Reformulation

Let's see these advantages at work on Quadratic Programs

We need to solve:

$$\operatorname{argmin}_{x} \left\{ \frac{1}{2} x^{T} P x + q^{T} x \mid l \leq A x \leq u \right\}$$

...Which we reformulate to:

argmin
$$x^T P x + q^T x$$

subject to: $z = Ax$
 $l \le z \le u$

- We have introduced a new variables z
- ...And posted the inequality constraints over that

QP Reformulation

Then, we turn the inequality constraints into a function

argmin
$$x^T P x + q^T x + \chi_{l \le z \le u}(z)$$

subject to: $z = Ax$

Where:

- $\chi_{l < z < u}$ is the characteristic function of $l \le z \le u$
 - It's value is $+\infty$ when the constraint is violated and 0 elsewhere
 - In this case, it is non-differentiable, but closed, proper, and convex!
- $x^T P x + q^T x$ is our usual cost term
 - It is differentiable
 - lacktriangleright ...And closed, proper, and convex if $m{P}$ is semi-definite positive

We can now proceed to apply the ADMM!

The ADMM Steps for QP

We need to start from a feasible x^0 , z^0 , λ^0 :

• That's easy, we get it by setting $\lambda^0=0$, $z^0=l$, then solving $Ax^0=l$

The x minimization step for $\hat{z} = z^k$ becomes:

$$\operatorname{argmin}_{x} x^{T} P x + q^{T} x + \chi_{l \le z \le u}(\hat{z}) + \lambda^{T} (\hat{z} - A x) + \frac{1}{2} \rho ||\hat{z} - A x||_{2}^{2}$$

And then, since \hat{z} is fixed and feasible:

$$\operatorname{argmin} x^{T} P x + q^{T} x + \lambda^{T} (\hat{z} - A x) + \frac{1}{2} \rho ||\hat{z} - A x||_{2}^{2}$$

This is a convex, differentiable, quadratic minimization problem

- It can be tackled via gradient descent
- ...Or by solving a linear system of equations

The ADMM Steps for QP

The z minimization step for $\hat{x} = x^{k+1}$ becomes

$$\operatorname{argmin}_{z} \hat{x}^{T} P \hat{x} + q^{T} \hat{x} + \chi_{l \leq z \leq u}(z) + \lambda^{T} (z - A \hat{x}) + \frac{1}{2} \rho \|z - A \hat{x}\|_{2}^{2}$$

Since \hat{x} is fixed, this can be reformulated as:

$$\operatorname{argmin} \lambda^T z + \frac{1}{2} \rho \|z - A\hat{x}\|_2^2$$

subject to: $l \le z \le u$

...And finally separated in to n problems (one per variable) in the form:

$$\operatorname{argmin}_{z_j} \left\{ \lambda_j z_j + \frac{1}{2} \rho (z_j - A_j \hat{x})^2 \mid l \leq z_j \leq u \right\}$$

These are all very easy to solve

Some Considerations

We used the ADMM to break QP into a sequence of simpler problems

The method can be used in other clever ways:

- Optimization with non-differentiable reguralizers
- Parallel training, by splitting examples into multiple problems
- ...And using constraints to reach a consensus

The ADMM is best used for convex problems

- Classical results are for convex problems only
- There are some (local) results non non-convex problems (e.g. this one)
- ...But in practice it's less reliable

About the convergence pace

- It's very fast in the first iterations, but much slower later
- You can high-quality solutions early, but reaching the optimum takes long
- All in all, it's best to use the ADMM as an approximate method