

**Theorem 2.1.1 Logical Equivalences**

Given any statement variables  $p$ ,  $q$ , and  $r$ , a tautology  $\mathbf{t}$  and a contradiction  $\mathbf{c}$ , the following logical equivalences hold.

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|--|---|---|
| 1. <i>Commutative laws:</i>  | $p \wedge q \equiv q \wedge p$                              | $p \vee q \equiv q \vee p$                                |
| 2. <i>Associative laws:</i>  | $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$        | $(p \vee q) \vee r \equiv p \vee (q \vee r)$              |
| 3. <i>Distributive laws:</i>   | $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$ | $p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$ |
| 4. <i>Identity laws:</i>   | $p \wedge \mathbf{t} \equiv p$                              | $p \vee \mathbf{c} \equiv p$                              |
| 5. <i>Negation laws:</i>   | $p \vee \sim p \equiv \mathbf{t}$                           | $p \wedge \sim p \equiv \mathbf{c}$                       |
| 6. <i>Double negative law:</i>   | $\sim(\sim p) \equiv p$                                     |   |
| 7. <i>Idempotent laws:</i>   | $p \wedge p \equiv p$                                       | $p \vee p \equiv p$                                       |
| 8. <i>Universal bound laws:</i>  | $p \vee \mathbf{t} \equiv \mathbf{t}$                       | $p \wedge \mathbf{c} \equiv \mathbf{c}$                   |
| 9. <i>De Morgan's laws:</i>  | $\sim(p \wedge q) \equiv \sim p \vee \sim q$                | $\sim(p \vee q) \equiv \sim p \wedge \sim q$              |
| 10. <i>Absorption laws:</i>  | $p \vee (p \wedge q) \equiv p$                              | $p \wedge (p \vee q) \equiv p$                            |
| 11. <i>Negations of <math>\mathbf{t}</math> and <math>\mathbf{c}</math>:</i> | $\sim \mathbf{t} \equiv \mathbf{c}$                         | $\sim \mathbf{c} \equiv \mathbf{t}$                       |

<b>Modus Ponens</b>	$p \rightarrow q$ $p$ $\therefore q$	<b>Elimination</b>	<b>a.</b> $p \vee q$ $\sim q$ $\therefore p$	<b>b.</b> $p \vee q$ $\sim p$ $\therefore q$
<b>Modus Tollens</b>	$p \rightarrow q$ $\sim q$ $\therefore \sim p$	<b>Transitivity</b>	$p \rightarrow q$ $q \rightarrow r$ $\therefore p \rightarrow r$	
<b>Generalization</b>	<b>a.</b> $p$ $\therefore p \vee q$	<b>Proof by Division into Cases</b>	$p \vee q$ $p \rightarrow r$ $q \rightarrow r$ $\therefore r$	
<b>Specialization</b>	<b>b.</b> $q$ $\therefore p \vee q$			
<b>Conjunction</b>	<b>a.</b> $p \wedge q$ $\therefore p$	<b>Contradiction Rule</b>	$\sim p \rightarrow \mathbf{c}$ $\therefore p$	
	<b>b.</b> $p \wedge q$ $\therefore q$			

**Valid Argument Forms**

Figure 3: Theorem 2.1.1 and Theorem 2.3.1