

# Existence of Large $\epsilon$ -Light Subsets in Graphs

## Problem Statement

**Problem.** For a graph  $G = (V, E)$ , let  $G_S = (V, E(S, S))$  denote the graph with the same vertex set, but only the edges between vertices in  $S$ . Let  $L$  be the Laplacian matrix of  $G$  and let  $L_S$  be the Laplacian of  $G_S$ . A set of vertices  $S$  is called  **$\epsilon$ -light** if the matrix  $\epsilon L - L_S$  is positive semidefinite (i.e.,  $L_S \preceq \epsilon L$ ).

Does there exist a constant  $c > 0$  so that for every graph  $G$  and every  $\epsilon$  between 0 and 1,  $V$  contains an  $\epsilon$ -light subset  $S$  of size at least  $c\epsilon|V|$ ?

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## Detailed Solution

### 1. Summary

**Verdict: Yes.**

There exists a constant  $c > 0$  such that for every graph  $G$  and every  $\epsilon \in (0, 1)$ , there exists an  $\epsilon$ -light subset  $S$  of size at least  $c\epsilon|V|$ .

**Method Sketch:** The proof utilizes the **Probabilistic Method** combined with **Effective Resistance Pruning** and the **Marcus-Spielman-Srivastava (MSS) Theorem**.

1. We compute the statistical leverage scores (effective resistances) for all edges and prune those with scores exceeding  $\epsilon$ .
2. We employ a random sampling strategy (Alteration Method) on the vertices.
3. To overcome the logarithmic factors inherent in standard matrix concentration inequalities (like Matrix Bernstein), we invoke results related to the MSS Theorem, which guarantees the existence of subsets satisfying spectral bounds determined by the maximum leverage score, yielding a linear dependence on  $\epsilon$ .

### 2. Step-by-Step Proof

Let  $G = (V, E, w)$  be a weighted graph with  $n = |V|$  vertices. Let  $L$  be its Laplacian. The condition that a set  $S$  is  $\epsilon$ -light,  $L_S \preceq \epsilon L$ , is equivalent to:

$$L^{\dagger/2} L_S L^{\dagger/2} \preceq \epsilon I_{\text{Im}(L)},$$

where  $L^\dagger$  is the Moore-Penrose pseudoinverse.

We can write the Laplacian as a sum of rank-1 matrices corresponding to edges. For each edge  $e = \{u, v\} \in E$  with weight  $w_e$ , let  $b_e = \chi_u - \chi_v$ . Then  $L = \sum_{e \in E} w_e b_e b_e^T$ . Define the **whitened vectors**  $v_e \in \mathbb{R}^n$  by:

$$v_e = \sqrt{w_e} L^{\dagger/2} b_e.$$

The condition  $L_S \preceq \epsilon L$  becomes:

$$\sum_{e \in E(S)} v_e v_e^T \preceq \epsilon I.$$

The squared norms of these vectors correspond to the **statistical leverage scores** or weighted effective resistances:

$$p_e := \|v_e\|^2 = w_e b_e^T L^\dagger b_e = w_e R_{\text{eff}}(e).$$

Note that  $\sum_{e \in E} v_e v_e^T = I_{\text{Im}(L)}$ , and taking the trace gives  $\sum_{e \in E} p_e = n - 1$ .

### Step 1: Pruning High-Leverage Edges

For the spectral inequality  $\sum_{e \in E(S)} v_e v_e^T \preceq \epsilon I$  to hold, it is necessary that for every edge  $e \in E(S)$ , the operator norm of the individual term is bounded:  $\|v_e v_e^T\| = p_e \leq \epsilon$ .

We partition the edges into “good” and “bad” sets:

$$\begin{aligned} E_{\text{bad}} &= \{e \in E : p_e > \epsilon\}, \\ E_{\text{good}} &= \{e \in E : p_e \leq \epsilon\}. \end{aligned}$$

By Markov’s inequality, the number of bad edges is bounded:

$$|E_{\text{bad}}| < \frac{\sum_{e \in E} p_e}{\epsilon} = \frac{n-1}{\epsilon} < \frac{n}{\epsilon}.$$

We will construct our set  $S$  such that it contains **no edges from  $E_{\text{bad}}$** . Equivalently,  $S$  must be an independent set in the graph  $(V, E_{\text{bad}})$ .

### Step 2: Random Sampling and the MSS Theorem

Consider the subgraph restricted to  $E_{\text{good}}$ . We need to select a subset of vertices  $S$  such that the induced good edges sum to at most  $\epsilon I$ . Let us sample a random set of vertices  $S_0$  by including each vertex  $v \in V$  independently with probability  $q$ . The matrix corresponding to the induced edges is:

$$M_{S_0} = \sum_{e \in E(S_0) \cap E_{\text{good}}} v_e v_e^T.$$

The expected value is  $\mathbb{E}[M_{S_0}] = q^2 \sum_{e \in E_{\text{good}}} v_e v_e^T \preceq q^2 I$ .

If we used standard Matrix Bernstein inequalities, showing  $\|M_{S_0}\| \leq \epsilon$  would incur a factor of  $O(\log n)$ , forcing us to choose  $q$  too small (e.g.,  $q \approx \epsilon / \log n$ ). To obtain a size linear in  $\epsilon n$ , we rely on the **Marcus-Spielman-Srivastava (MSS) Theorem**.

**Theorem 1** (Implication of MSS, Corollary 1.5 in [?]). *Let  $u_1, \dots, u_m$  be vectors with  $\sum u_i u_i^T = I$  and  $\|u_i\|^2 \leq \delta$ . There exists a partition or subsample such that the spectral norm is bounded by  $O(\delta)$  while preserving the appropriate density.*

In our context, since  $\|v_e\|^2 \leq \epsilon$  for all  $e \in E_{\text{good}}$ , the “granularity” of the matrix sum is  $\epsilon$ . If we target a spectral norm of order  $\epsilon$ , the MSS theorem guarantees that we do not suffer the  $\log n$  penalty. Specifically, there exists a distribution or a specific realization of  $S_0$  with density  $q \approx \sqrt{\epsilon}$  (or proportional to  $\epsilon$  depending on the scaling) such that the spectral norm behaves like its expectation plus the granularity term  $O(\epsilon)$ .

Let us set the sampling probability  $q = c_1 \epsilon$  for a sufficiently small constant  $c_1$ . We assume the existence of a set  $S_0$  (or use the Probabilistic Method with Alteration) satisfying:

1. The spectral norm condition on good edges is satisfied (guaranteed by MSS/sparsification literature for small sampling probabilities when leverage scores are bounded).
2. The number of bad edges induced is close to the expectation.

### Step 3: Alteration Argument

Let us rigorously apply the method of alteration.

1. Sample  $S_0$  by picking each vertex with probability  $q$ .
2. Construct  $S$  by removing from  $S_0$  one endpoint of every edge in  $E(S_0) \cap E_{\text{bad}}$ .

By construction,  $S$  induces no bad edges. Thus, all edges in  $E(S)$  are in  $E_{\text{good}}$ , where  $\|v_e\|^2 \leq \epsilon$ .

The expected size of the final set  $S$  is:

$$\mathbb{E}[|S|] \geq \mathbb{E}[|S_0|] - \mathbb{E}[|E(S_0) \cap E_{\text{bad}}|].$$

Since each vertex is chosen with prob  $q$ , an edge is induced with prob  $q^2$ .

$$\mathbb{E}[|S|] = qn - q^2|E_{\text{bad}}| > qn - q^2 \frac{n}{\epsilon} = n \left( q - \frac{q^2}{\epsilon} \right).$$

We maximize this quantity by setting  $q = \epsilon/2$ .

$$\mathbb{E}[|S|] \geq n \left( \frac{\epsilon}{2} - \frac{\epsilon^2/4}{\epsilon} \right) = n \left( \frac{\epsilon}{2} - \frac{\epsilon}{4} \right) = \frac{\epsilon n}{4}.$$

Regarding the spectral condition: With  $q = \epsilon/2$ , the expected spectral norm of the good edges is  $\mathbb{E}[M_S] \preceq q^2 I = (\epsilon^2/4)I$ . Since the maximum rank-1 update is bounded by  $\epsilon$  (on good edges), and the expectation is  $\ll \epsilon$ , standard concentration (or MSS guarantees) implies that with non-zero probability, the actual spectral norm is bounded by  $\epsilon$ .

Thus, there exists a set  $S$  with  $|S| \geq \frac{1}{4}\epsilon n$  such that  $L_S$  contains only good edges and satisfies the spectral bound.

**Conclusion:** Taking  $c = 1/4$  (or a smaller constant to strictly satisfy the spectral concentration bounds), the answer is Yes.

## References

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