

HOMEWORK 2 – Q1

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1. Given positive integers M and n compute M^n using only $O(\log(n))$ many multiplications. (15 pts)

Solution:

There are two ways to solve this question.

1. We first divide M^n into two parts, then we have

$M^n = M^{\left(\frac{n}{2}\right)^2}$. We solve it recursively, then we have

$$M * M = M^2, M^4 = (M^2)^2 = M * M * M *$$

$$M, \text{ and } ((M^2)^2)^2 = M * M * \dots *$$

M (there are 8 multiplications), but we only need

$\log_2(8) = 3$ multiplications to get a answer that need

8 multiplications. Thus, we can compute M^2 using

only $O(\log(n))$ many multiplications. However, if n is

a odd number, we have $M^n = M * (M^{n-1})$, we need

to do the same thing with M^{n-1} , and multiple one

more M into the finally answer we compute (which is

M^{n-1}), it also using only $(\log(n) + 1) = O(\log(n))$

many multiplications.

2. We can write M^n in binary as $M^n = 2^{k_1} + 2^{k_2} +$

$2^{k_3} \dots 2^{k_m}$ where $k_1 > k_2 > k_3 > \dots k_m$ and $k_1 =$

$\log_2 n$; then $M^n = M^{2^{k_1}} M^{2^{k_2}} M^{2^{k_3}} \dots M^{2^{k_m}}$. This involves

at most $\log_2(n)$ multiplications. So it is enough to

compute all of M^{2^j} for all $1 \leq j \leq \log_2(n)$ with at

most $\log_2(n)$ multiplications. To do that, use repeated squaring, as show below:

```
result = 1;
base = 2;
while (b != 0){
    if (b&1 > 0){
        result *= base;
    }
    base *= base;
    b >>= 1;
}
return result
```