

# Graphs of simplex codes

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- $V = \mathbb{F}_q^n$ ,  $q = p^m$ ;
- $\alpha$ -primitive element of  $\mathbb{F}_q$ ;
- $[n, k]_q$ -code,  $k$ -dimensional subspace of  $V$ ;
- $[n]_q = \frac{q^n - 1}{q - 1}$ ;
- $C_i = \{(x_1, \dots, x_{i-1}, 0, x_{i+1}, \dots, x_n); x_j \in \mathbb{F}_q\}$ , (Coordinate hyperplane)  
Kernel of the coordinate functional  $c_i((x_1, x_2, \dots, x_n)) = x_i$ ;
- $X = \langle v_1, \dots, v_k \rangle - [n, k]_q$ -code,  $G_x$ -generator matrix;

$$G_x = \begin{bmatrix} v_1 \\ \vdots \\ v_k \end{bmatrix}$$

- $C_{i,j} = C_i \cap C_j$
- Code  $X$  such that  $\dim(X \cap C_{i,j}) = k - 2$  for all  $i, j$  is called a projective code.
- The columns of a generator matrix of a projective code are distinct points of a projective space i.e. every two columns are linearly independent.

$$\begin{bmatrix} 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 2 \end{bmatrix}$$

$$\begin{bmatrix} 2 & 0 & 2 & 2 \\ 0 & 1 & 1 & 2 \end{bmatrix}$$

Projective codes with maximal length  $n = [k]_q$  are called **simplex codes**.

$$\begin{bmatrix} 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 2 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 & 1 & 1 & 1 \\ 0 & 1 & 1 & \alpha & \alpha^2 \end{bmatrix}$$

Generator matrices of simplex codes of dimension 2 over  $\mathbb{F}_3$  and  $\mathbb{F}_4$ .

- We say that  $x \in V$  is a *simplex vector* if the Hamming weight of this vector is  $q^{k-1}$ .
- A non-zero vector is simplex if and only if it is a codeword of a certain  $q$ -ary simplex code of dimension  $k$ .
- A subspace of  $V$  is a  $q$ -ary simplex code of dimension  $k$  if and only if it is maximal with respect to the property that all non-zero vectors are simplex.

## Theorem (K, Pankov, Pasini)

*The simplex vectors form the algebraic variety defined by the equations*

$$\sum_{i_1 < \dots < i_{p^j}} x_{i_1}^{q-1} \dots x_{i_{p^j}}^{q-1} = 0$$

*for  $j \in \{0, \dots, mk - m - 1\}$ , where  $q = p^m$  and  $p$  is a prime number.*

*This variety is a quadric only if  $q = 2$ ,  $k = 3$  or  $q = 3$ ,  $k = 2$ .*

The group of monomial linear automorphisms of  $V$  acts transitively on the set of simplex codes and contains precisely  $n!(q-1)^n$  elements. There are precisely

$$(q^k - 1)(q^k - q) \dots (q^k - q^{k-1})$$

(the number of elements in  $\text{GL}(k, q)$ ) monomial linear automorphisms of  $V$  which preserve a fixed simplex code. Therefore the number of  $q$ -ary simplex codes of dimension  $k$  is equal to

$$\frac{n!(q-1)^n}{(q^k - 1)(q^k - q) \dots (q^k - q^{k-1})}$$

## Proposition (K, Pankov)

*If  $X$  is an  $m$ -dimensional subcode of a simplex code, then every generator matrix  $M$  of  $X$  satisfies the following condition:*

*$(*)_m$   $M$  contains precisely  $[k - m]_q$  zero columns and any non-zero column of  $M$  is proportional to precisely  $q^{k-m}$  columns including itself.*

*If a generator matrix  $M$  of an  $m$ -dimensional code  $X \subset V$  satisfies  $(*)_m$ , then  $X$  is a subcode of a simplex code.*

$$M = \begin{bmatrix} 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 2 & 2 & 2 \end{bmatrix}$$

$M$  is a generator matrix of a subcode of simplex code with generator matrix

$$\begin{bmatrix} 1 & 0 & 1 & 2 & 0 & 1 & 2 & 0 & 1 & 2 & 0 & 1 & 2 \\ 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 2 & 2 & 2 \end{bmatrix}$$

## Theorem (K, Pankov)

Let  $m \in \{0, 1, \dots, k-1\}$ . Every  $m$ -dimensional subcode of a simplex code is contained in precisely

$$\frac{[k-m]_q!(q-1)^{[k-m]_q}(q^{k-m}!)^{[m]_q}}{(q^{k-m}-1)(q^{k-m}-q)\dots(q^{k-m}-q^{k-m-1})q^{m(k-m)}}$$

distinct simplex codes. Furthermore, there are two simplex codes whose intersection is precisely this subcode except the case when  $q = k = 2$ .

The *Grassmann graph*  $\Gamma_k(V)$  is the simple graph whose points are  $k$ -dimensional subspaces (codes) of  $V$  and two such subspaces are adjacent if their intersection is  $(k - 1)$ -dimensional. We assume that  $n = [k]_q$ ,  $k \geq 2$  and denote by  $\Gamma^s(k, q)$  the subgraph of  $\Gamma_k(V)$  induced by the set of  $q$ -ary simplex codes of dimension  $k$ .

Examples:

- $\Gamma^s(2, 2)$  is a single vertex;
- $\Gamma^s(2, 3)$  is the complete bipartite graph  $K_{4,4}$ ,
- $\Gamma^s(3, 2)$  is isomorphic to the graph  $\Gamma_{1,3}(\mathbb{F}_2^4)$  formed by 1-dimensional and 3-dimensional subspaces of  $\mathbb{F}_2^4$ , where distinct subspaces are connected by an edge if they are incident.



A *clique* is a complete subgraph. A clique  $\mathcal{X}$  is said to be *maximal* if every clique containing  $\mathcal{X}$  coincides with  $\mathcal{X}$ . Every maximal clique of  $\Gamma_k(V)$  is of one of the following type:

- the *star*  $\mathcal{S}(X)$  consisting of all  $k$ -dimensional subspaces containing a certain  $(k-1)$ -dimensional subspace  $X$ ;
- the *top*  $\mathcal{T}(Y)$  consisting of all  $k$ -dimensional subspaces contained in a certain  $(k+1)$ -dimensional subspace  $Y$ .

The intersections of  $\mathcal{S}(X)$  and  $\mathcal{T}(Y)$  with the set of simplex codes are denoted by  $\mathcal{S}^s(X)$  and  $\mathcal{T}^s(Y)$ , respectively. Every such intersection is a clique in  $\Gamma^s(k, q)$  (if it is non-empty), but we cannot assert that this clique is maximal. We say that  $\mathcal{S}^s(X)$  or

$\mathcal{T}^s(Y)$  is a *star* or a *top* of the simplex code graph  $\Gamma^s(k, q)$  only in the case when it is a maximal clique of  $\Gamma^s(k, q)$ .

Since  $\Gamma^s(2, 3)$  and  $\Gamma^s(3, 2)$  are bipartite, every maximal clique in these graphs consists of two vertices which implies that it is a star and a top simultaneously. If  $X, Y$  are adjacent vertices in one of these graphs, then

$$\{X, Y\} = \mathcal{S}^s(X \cap Y) = \mathcal{T}^s(X + Y).$$

### Proposition (K, Pankov)

*Suppose that one of the following possibilities is realized:*

- $q = 2$  and  $k \geq 4$ ;
- $q = 3$  and  $k \geq 3$ ;
- $q \geq 4$ .

*Then  $S^s(X)$  is a star of  $\Gamma^s(k, q)$  if and only if  $X$  is a  $(k - 1)$ -dimensional subcode of a simplex code. Furthermore, there are no maximal cliques of  $\Gamma^s(k, q)$  which are stars and tops simultaneously.*

It is clear that  $S^s(X)$  is non-empty if and only if  $X$  is a  $(k - 1)$ -dimensional subcode of a simplex code. We have that

$$|S^s(X)| = \frac{(q!)^{[k-1]_q}}{q^{k-1}}$$

This number is greater than  $q + 1$ .

If  $k = 2$ ,  $q \geq 4$ , then  $\mathcal{S}^s(X)$  consists of  $(q - 1)!$  elements and

$$(q - 1)! \geq (q - 2)(q - 1) = q^2 - 3q + 2 > q + 1$$

(since  $q^2 - 4q + 1 > 0$  for  $q \geq 4$ ).

If  $k \geq 3$ , then

$$[k - 1]_q = q^{k-2} + q^{k-3} \dots + 1 \geq q^{k-2} + k - 2 \geq k + 1$$

(we have  $q^{k-2} \geq 3$ , since  $k \geq 4$  if  $q = 2$ ). Therefore,

$$\frac{(q!)^{[k-1]_q}}{q^{k-1}} \geq \frac{q^{k+1}}{q^{k-1}} = q^2 > q + 1.$$

So,  $\mathcal{S}^s(X)$  contains more than  $q + 1$  elements and there is no  $\mathcal{T}^s(Y)$  containing  $\mathcal{S}^s(X)$  (since the intersection of a star and a top of the Grassmann graph is empty or contains precisely  $q + 1$  elements). This guarantees that  $\mathcal{S}^s(X)$  is a maximal clique of  $\Gamma^s(k, q)$ , i.e. it is a star of  $\Gamma^s(k, q)$  which is not a top.

There exist non-empty  $\mathcal{T}^s(Y)$  which is not a top of  $\Gamma^s(k, q)$ .

Let us take  $\mathbf{x}, \mathbf{y}, \mathbf{z} \in V$  such that

$$\begin{bmatrix} \mathbf{x} \\ \mathbf{y} \\ \mathbf{z} \end{bmatrix} = \begin{bmatrix} 0 & 1 & 1 & \alpha & \alpha^2 & \cdots & \alpha^{q-2} \\ 1 & 0 & 1 & 1 & 1 & \cdots & 1 \\ 1 & 0 & \alpha & \alpha & \alpha & \cdots & \alpha \end{bmatrix}.$$

Then  $\langle \mathbf{x}, \mathbf{y} \rangle$  and  $\langle \mathbf{x}, \mathbf{z} \rangle$  are adjacent  $q$ -ary simplex codes of dimension 2. Assume that there is a simplex code intersecting  $\langle \mathbf{x}, \mathbf{y} \rangle$  and  $\langle \mathbf{x}, \mathbf{z} \rangle$  in distinct 1-dimensional subcodes  $\langle \mathbf{x} + \alpha^j \mathbf{z} \rangle, \langle \mathbf{x} + \alpha^i \mathbf{y} \rangle$ , respectively. Its generator matrix is

$$\begin{bmatrix} \mathbf{x} + \alpha^j \mathbf{z} \\ \mathbf{x} + \alpha^i \mathbf{y} \end{bmatrix} = \begin{bmatrix} \alpha^j & 1 & 1 + \alpha^{j+1} & \alpha + \alpha^{j+1} & \alpha^2 + \alpha^{j+1} & \cdots & \alpha^{q-2} + \alpha^{j+1} \\ \alpha^i & 1 & 1 + \alpha^i & \alpha + \alpha^i & \alpha^2 + \alpha^i & \cdots & \alpha^{q-2} + \alpha^i \end{bmatrix}$$

and  $\alpha^i \neq \alpha^j$  (since the first and second columns are non-proportional).

We choose  $t \in \{0, 1, \dots, q-2\}$  such that

$$\alpha^t = \frac{\alpha^i(\alpha^{j+1} - \alpha^j)}{\alpha^j - \alpha^i}.$$

Then the determinant

$$\begin{vmatrix} \alpha^j & \alpha^t + \alpha^{j+1} \\ \alpha^i & \alpha^t + \alpha^i \end{vmatrix} = \alpha^t(\alpha^j - \alpha^i) - \alpha^i(\alpha^{j+1} - \alpha^j)$$

is zero, i.e. the first and  $(t+3)$ -th columns are proportional which is impossible. Therefore, every simplex code adjacent to both  $\langle \mathbf{x}, \mathbf{y} \rangle, \langle \mathbf{x}, \mathbf{z} \rangle$  belongs to the star  $\mathcal{S}^s(\langle \mathbf{x} \rangle)$ . Then  $\mathcal{T}^s(\langle \mathbf{x}, \mathbf{y}, \mathbf{z} \rangle)$  is a non-empty proper subset of  $\mathcal{S}^s(\langle \mathbf{x} \rangle)$  and, consequently, it is not a top.

Recall that every maximal clique of  $\Gamma^s(2, 3)$  and  $\Gamma^s(3, 2)$  is a star and a top simultaneously. Every maximal clique of  $\Gamma^s(2, 4)$  is a star. The graph  $\Gamma^s(k, q)$  contains tops in all remaining cases.

### Theorem (K, Pankov)

*Suppose that one of the following possibilities is realized:*

- $k = 2$  and  $q \geq 5$ ,
- $k \geq 4$  and  $q = 2$ ,
- $k \geq 3$  and  $q \geq 3$ .

*Then  $\Gamma^s(k, q)$  contains tops. If  $k \geq 4$  and  $q \geq 3$ , then there are tops of  $\Gamma^s(k, q)$  containing different numbers of elements. If  $k \geq 5$  and  $q \geq 3$ , then there is a top of  $\Gamma^s(k, q)$  consisting of precisely three elements.*

Consider  $\mathbf{x}, \mathbf{y}, \mathbf{z} \in V$  such that

$$\begin{bmatrix} \mathbf{x} \\ \mathbf{y} \\ \mathbf{z} \end{bmatrix} = \begin{bmatrix} 0 & 1 & 1 & 1 & 1 & \cdots & 1 & 1 \\ 1 & 0 & \alpha^0 & \alpha^1 & \alpha^2 & \cdots & \alpha^{q-3} & \alpha^{q-2} \\ 1 & \frac{1}{1+0} & \frac{1}{1+\alpha^0} & \frac{1}{1+\alpha^{-1}} & \frac{1}{1+\alpha^{-2}} & \cdots & \frac{1}{1+\alpha^{-(q-3)}} & \frac{1}{1+\alpha^{-(q-2)}} \end{bmatrix};$$

if  $1 + \alpha^{-i} = 0$ , then we put 0 instead of  $\frac{1}{1+\alpha^{-i}}$ . It is clear that the columns of  $\begin{bmatrix} \mathbf{x} \\ \mathbf{y} \end{bmatrix}$  and  $\begin{bmatrix} \mathbf{x} \\ \mathbf{z} \end{bmatrix}$  are mutually non-proportional, and  $\langle \mathbf{x}, \mathbf{y} \rangle$  and  $\langle \mathbf{x}, \mathbf{z} \rangle$  are simplex codes. With some calculations  $\langle \mathbf{y}, \mathbf{z} \rangle$  is also a simplex code. And  $\mathcal{T}^s(\langle \mathbf{x}, \mathbf{y}, \mathbf{z} \rangle)$  is a top.

For  $\mathbb{F}_3$  and  $\mathbb{F}_4$  this vectors are linearly dependent.

$$A = \begin{bmatrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{bmatrix} \quad \text{and} \quad B = \begin{bmatrix} 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \end{bmatrix}$$

For every  $i \in \{1, \dots, [k-2]_2\}$  we denote by  $D_i$  the  $(k-2, 4)$ -matrix whose columns are  $i$ -th non-zero vectors of  $\mathbb{F}_2^{k-2}$  (in lexicographical order). and consider the  $(k+1, n)$ -matrix

$$M = \begin{bmatrix} A & B & \dots & B \\ \mathbf{0} & D_1 & \dots & D_{[k-2]_2} \end{bmatrix}.$$

Let  $\mathbf{v}_1, \mathbf{v}_2$  and  $\mathbf{v}_3$  be the first, second and third rows of  $M$  (respectively). Denote by  $C$  the  $(k-2)$ -dimensional subspace whose generator matrix is  $[\mathbf{0}, D_1, \dots, D_{[k-2]_2}]$ . then  $X_1 = \langle C, \mathbf{v}_2, \mathbf{v}_3 \rangle, X_2 = \langle C, \mathbf{v}_1, \mathbf{v}_3 \rangle, X_3 = \langle C, \mathbf{v}_1, \mathbf{v}_2 \rangle$  are simplex codes  $C$  is a subcode of every  $X_i$  and  $X_i \cap X_j = C$   
 $\mathcal{T}^s(\langle X_1, X_2, X_3 \rangle)$  is a top.



$$\left[ \begin{array}{c|c|c|c} 011 & 1001 & 1001 & 1001 \\ 101 & 0101 & 0101 & 0101 \\ 110 & 0011 & 0011 & 0011 \\ \hline 000 & 0000 & 1111 & 1111 \\ 000 & 1111 & 0000 & 1111 \end{array} \right]$$

$$\left[ \begin{array}{c|c} 011 & 1001 \\ 101 & 0101 \\ 110 & 0011 \\ \hline 000 & 1111 \end{array} \right]$$

We take  $\mathbf{x}, \mathbf{y} \in \mathbb{F}_q^{q+1}$  spanning a  $q$ -ary simplex code of dimension 2. For any non-zero scalars  $a, b \in \mathbb{F}$  the  $(3, q+1)$ -matrix

$$\begin{bmatrix} \mathbf{x} \\ \mathbf{y} \\ a\mathbf{x} + b\mathbf{y} \end{bmatrix}$$

will be denote by  $A_{a,b}$ . and denote by  $B_{a,b}$  the  $(3, q^2)$ -matrix

$$[A_{a,b} \ \alpha A_{a,b} \ \dots \ \alpha^{q-2} A_{a,b} \ \mathbf{0}]$$

For every  $i \in \{1, \dots, [k-2]_2\}$  we denote by  $D_i$  the  $(k-2, 4)$ -matrix whose columns are  $i$ -th non-zero vectors of  $\mathbb{F}_q^{k-2}$  with first non-zero coordinate equal to 1 (in lexicografical order).

Let us take any collection of non-zero scalars

$$a_0, b_0, a_1, b_1, \dots, a_{[k-2]_q}, b_{[k-2]_q} \in \mathbb{F}$$

such that for some  $i, j$  we have  $(a_i, b_i) \neq (a_j, b_j)$

consider the  $(k + 1, n)$ -matrix

$$M = \begin{bmatrix} A_{a_0, b_0} & B_{a_1, b_1} & \cdots & B_{a_{[k-2]_q}, b_{[k-2]_q}} \\ \mathbf{0} & D_1 & \cdots & D_{[k-2]_q} \end{bmatrix}.$$

Let  $\mathbf{v}_1, \mathbf{v}_2$  and  $\mathbf{v}_3$  be the first, second and third rows of  $M$  (respectively). Denote by  $C$  the  $(k - 2)$ -dimensional subspace whose generator matrix is  $[\mathbf{0}, D_1, \dots, D_{[k-2]_q}]$ . then  $X_1 = \langle C, \mathbf{v}_2, \mathbf{v}_3 \rangle, X_2 = \langle C, \mathbf{v}_1, \mathbf{v}_3 \rangle, X_3 = \langle C, \mathbf{v}_1, \mathbf{v}_2 \rangle$  are simplex codes  $C$  is a subcode of every  $X_i$  and  $X_i \cap X_j = C$

### Lemma

*A non-zero vector of  $\langle \mathbf{v}_1, \mathbf{v}_2, \mathbf{v}_3 \rangle$  is not a code word of a simplex code if and only if it is a scalar multiple of*

$$a_i \mathbf{v}_1 + b_i \mathbf{v}_2 - \mathbf{v}_3$$

*for a certain  $i \in \{0, 1, \dots, [k - 2]_q\}$ .*

$\Gamma^s(2, q)$ -the restriction of the Grassmann graph to the set of simplex lines.  
We will use a projective terminology (1-dimensional subspaces are points and 2-dimensional are lines).

### Proposition

- (1) *Every simplex point is contained in precisely  $(q - 1)!$  simplex lines.*
- (2) *The degree of every vertex of the graph  $\Gamma^s(2, q)$  is equal to  $(q + 1)[(q - 1)! - 1]$ .*

A simplex point  $\langle 1, a_1, \dots, a_q \rangle$  is adjacent to simplex point  $P = \langle 0, 1, \dots, 1 \rangle$  if and only if all columns of the matrix

$$\begin{bmatrix} 0 & 1 & \dots & 1 \\ 1 & a_1 & \dots & a_q \end{bmatrix}$$

are mutually non-proportional; the latter is equivalent to the fact that  $a_1, \dots, a_q$  are mutually distinct. Therefore, there are precisely  $q!$  simplex points adjacent to  $P$ .

Every simplex line passing through  $P$  contains  $q$  points distinct from  $P$  which means that  $P$  is contained in precisely  $(q - 1)!$  simplex lines.

Each simplex line  $L$  consists of  $q + 1$  simplex points. Every such point is contained in precisely  $(q - 1)! - 1$  simplex lines distinct from  $L$ .

For  $q = 4, k = 2$  the degree is 25

For  $q = 3, 4$  the sum of  $q - 1$  non-zero elements is zero if and only if these elements are mutually distinct.

For every  $q \geq 5$  there are non-zero  $a_1, \dots, a_{q-1} \in \mathbb{F}_q$  whose sum is zero and  $a_i = a_j$  for some distinct  $i, j$ . In the case when  $q$  is odd, we take any  $a \neq 0$  and  $b \neq 0, a, -a$  and replace the pair  $a, -a$  in the sum of all non-zero elements by the pair  $b, -b$ . If  $q = 2^m$  with  $m > 2$ , then  $1 + \alpha + \alpha^2 \neq 0$  and we replace the pair  $1, \alpha + \alpha^2$  in the sum of all non-zero elements by the pair  $1 + \alpha, \alpha^2$ .

Two distinct simplex points are said to be *adjacent* if they are connected by a simplex line.

### Proposition (K, Pankov)

Let  $P = \langle x_1, \dots, x_{q+1} \rangle$  be a simplex point and  $x_i = 0$ . If a simplex point  $Q = \langle y_1, \dots, y_{q+1} \rangle$  is adjacent to  $P$ , then

$$a_1 y_1 + \dots + a_{q+1} y_{q+1} = 0, \text{ where } a_j = x_j^{-1} \text{ for } j \neq i \text{ and } a_i = 0. \quad (1)$$

In the case when  $q = 4$ , a simplex point  $Q = \langle y_1, \dots, y_{q+1} \rangle$  is adjacent to  $P$  if and only if it satisfies (1) and  $y_i \neq 0$ .

The simplex points  $P, Q$  are adjacent if and only if  $y_i \neq 0$  and the columns of the matrix

$$\begin{bmatrix} x_1 & \dots & x_{q+1} \\ y_1 & \dots & y_{q+1} \end{bmatrix}$$

are mutually non-proportional. i.e. the determinants

$$\begin{vmatrix} x_i & x_j \\ y_i & y_j \end{vmatrix}$$

are non-zero. The latter is equivalent to the fact that  $y_i \neq 0$  and

$$x_1^{-1}y_1, \dots, x_{i-1}^{-1}y_{i-1}, x_{i+1}^{-1}y_{i+1}, \dots, x_{q+1}^{-1}y_{q+1}$$

are mutually distinct elements of  $\mathbb{F}_q$  (one of them is zero). The sum of all  $q - 1$  non-zero elements of  $\mathbb{F}_q$  is zero and we obtain (1).

In the case when  $q = 4$ , the sum of three non-zero elements of  $\mathbb{F}_q$  is zero if and only if these elements are mutually distinct. This implies the second statement.

## Theorem (K, Pankov)

Let  $q = 4$ . Then  $\Gamma^s(2, 4)$  is a connected graph of diameter 3 consisting of 162 simplex lines and the degree of every vertex of  $\Gamma$  is equal to 25. Furthermore, for each simplex line  $L$ :

- (1) There are precisely 6 simplex lines  $L_1, \dots, L_6$  which are at distance 3 from  $L$  in the graph  $\Gamma^s(2, 4)$ .
- (2) There are precisely 130 simplex lines which are at distance 2 from  $L$  in the graph  $\Gamma^s(2, 4)$ . The set of all such lines is the union of three mutually disjoint subsets denoted by  $\mathcal{X}_{20}^3, \mathcal{X}_{90}^1, \mathcal{X}_{20}^0$ , where
  - $\mathcal{X}_{20}^3$  is formed by 20 simplex lines and each of these lines is adjacent to precisely three distinct  $L_i$ ,
  - $\mathcal{X}_{90}^1$  consists of 90 lines and every such line is adjacent to a unique  $L_i$ ,
  - $\mathcal{X}_{20}^0$  consists of 20 simplex lines disjoint with all  $L_i$ .
- (3)  $\{L, L_1, \dots, L_6\} \cup \mathcal{X}_{20}^0$  is a spread of the set of all 135 simplex points, i.e. this set consists of 27 mutually disjoint lines which cover the set of simplex points.



For every simplex line  $L$  we denote by  $G(L)$  the group of all projective transformations induced by monomial semilinear automorphisms of  $V$  preserving  $L$ , i.e. the extensions of automorphisms of the corresponding simplex code (recall that every such extension is unique).

The group  $G(L)$  is isomorphic to  $P\Gamma L(2, 4)$  (since the automorphism group of the corresponding code is isomorphic to  $\Gamma L(2, 4)$ )

$P\Gamma L(2, 4)$  is a subgroup of the permutation group  $S_5$  acting on the points of  $L$ . These groups both are of order 120 which means that they are coincident.

## Theorem (K, Pankow)

Suppose that  $q = 4$  and  $L$  is a simplex line. Let  $L_1, \dots, L_6$  and  $\mathcal{X}_{20}^3, \mathcal{X}_{90}^1, \mathcal{X}_{20}^0$  be as in previous theorem and let  $\mathcal{A}$  be the set of all simplex lines adjacent to  $L$ . The action of the group  $G(L)$  on the set of all simplex lines has the following properties:

- (1) The sets  $\{L_1, \dots, L_6\}$ ,  $\mathcal{X}_{20}^3$  and  $\mathcal{X}_{20}^0$  are orbits of this action; moreover, the action of  $G(L)$  on the set  $\{L_1, \dots, L_6\}$  is sharply 3-transitive.
- (2) The set  $\mathcal{A}$  is the union of two orbits consisting of 10 and 15 simplex lines.
- (3) The set  $\mathcal{X}_{90}^1$  is the union of two orbits consisting of 30 and 60 simplex lines.

$$f(x, y) = x_1 y_1^{q-2} + x_2 y_2^{q-2} + \dots + x_{q+1} y_{q+1}^{q-2}$$

- $f(ax + bx', y) = af(x, y) + bf(x', y)$
- $f(x, ay) = a^{q-2} f(x, y)$
- $f(y, x) = f(x^{q-2}, y^{q-2})$
- $f(x, x) = 0$  gives first equation defining simplex vectors.
- $f(x, y) = 0$  for simplex vectors  $x, y$  that give a simplex line.

for  $q = 4$  we get a Hermitian form

- Geometry  $S(k,q)$ : point-line geometry whose maximal singular subspaces correspond to  $q$ -ary simplex codes of dimension  $k$ .
- Points: 1-dimensional subcodes of simplex codes.
- Lines: 2-dimensional subcodes of simplex codes.
- **Collinearity graph**: vertices = simplex points, edges = collinear points.

## Lemma (Fisher's inequality)

*if  $X_1, \dots, X_m$  are subsets of  $\{1, \dots, n\}$  such that  $|X_i \cap X_j|$  is constant for all pairs of distinct  $i, j$ , then  $m \leq n$ .*

## Proposition (K,Pankov,Tyc)

*Every clique of the collinearity graph of  $S(k, q)$  contains no more than  $n$  elements.*

## Proof

- For a clique of simplex points  $\langle v_1 \rangle, \dots, \langle v_m \rangle$ , define  $X_i$  as the set of indices where the  $i$ -th coordinate of  $v_i$  is zero.
- Then  $|X_i \cap X_j| = [k - 2]_q$  for distinct  $i, j$ .
- By Fisher's inequality,  $m \leq n$ .

## Definition (Power Transformation)

Let  $\mathbb{F}_q$  be a finite field and  $s$  a positive integer.

- The  $s$ -th power map on  $\mathbb{F}_q$ :

$$\phi_s : a \mapsto a^s, \quad a \in \mathbb{F}_q.$$

- $\phi_s$  is bijective iff  $\gcd(s, q-1) = 1$ .
- Induces a map on  $\mathbb{F}_q^n$ :

$$F_s(x_1, \dots, x_n) = (x_1^s, \dots, x_n^s).$$

## Theorem (K,Pankov,Tyc)

*If  $F_s$  is bijective and  $s \neq p^m$  (not a Frobenius automorphism), then  $F_s$  sends every maximal singular subspace of  $S(k, q)$  to an  $n$ -clique of the collinearity graph which is **not** a singular subspace.*

No three points of these cliques are collinear. Under additional conditions those cliques form normal rational curves.

## Theorem (K,Pankov,Tyc)

Let  $k = 2$ ,  $q = p^r$ , and let  $F_s$  be the  $s$ -th power map on  $\mathbb{F}_q^{q+1}$ :

$$F_s(x_1, \dots, x_{q+1}) = (x_1^s, \dots, x_{q+1}^s).$$

If the following hold:

- ①  $\gcd(s, q - 1) = 1$  (so  $F_s$  is bijective),
- ②  $s \neq p^m$  (not a Frobenius automorphism),
- ③ **Condition (A):** the  $p$ -cyclotomic coset of  $s$  modulo  $q - 1$  contains  $up^m - 1$  for some  $0 < u < p$ ,

then  $F_s$  sends every line of  $S(2, q)$  to a **normal rational curve** in a projective space  $\text{PG}(s, q)$ .

- $q = p$ ,  $s = p - 2$ ;  $F_{p-2}$  sends lines to normal rational curves in  $\text{PG}(p - 2, p)$ .
- For  $k = 2$ ,  $q = 5$ : only  $s = 3$  works and we have 2 types of maximal cliques: simplex lines and their images under  $F_3$ .
- For  $k = 2$ ,  $q = 7$ : only  $s = 5$  works but we have additional cliques that are not obtained by the power map.



## Definition

A **near orthomorphism** of a group  $G$  is a bijection

$$\theta : G \setminus \{a\} \longrightarrow G \setminus \{b\}$$

such that the map

$$\delta(x) = x^{-1}\theta(x)$$

is also a bijection

$$\delta : G \setminus \{a\} \longrightarrow G \setminus \{c\}$$

for some  $a, b, c \in G$ .

The elements  $(a, b, c)$  are called the *ex elements* of  $\theta$ .

## Theorem (K)

*Every top can be obtained from near-orthomorphism of the multiplicative group of the field.*

## Setup:

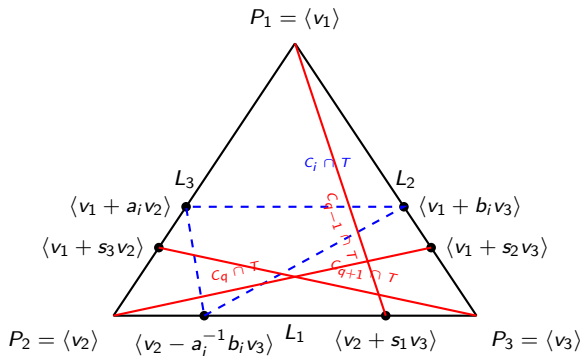
- Top  $T$  contains 3 lines  $L_1 = \langle v_2, v_3 \rangle$ ,  $L_2 = \langle v_1, v_3 \rangle$ ,  $L_3 = \langle v_1, v_2 \rangle$ .
- Lines intersect at points  $P_1 = L_2 \cap L_3 = \langle v_1 \rangle$ ,  $P_2 = L_1 \cap L_3 = \langle v_2 \rangle$ ,  $P_3 = L_1 \cap L_2 = \langle v_3 \rangle$ .

## Hyperplane intersections:

- For  $i \in \{1, \dots, q-2\}$ :
  - $C_i \cap L_3 = \langle v_1 + a_i v_2 \rangle$ ,  $a_i \neq a_j$  for  $i \neq j$
  - $C_i \cap L_2 = \langle v_1 + b_i v_3 \rangle$ ,  $b_i \neq b_j$  for  $i \neq j$
  - $C_i \cap L_1 = \langle v_2 - b_i a_i^{-1} v_3 \rangle$ ,  $-b_i a_i^{-1} \neq -b_j a_j^{-1}$  for  $i \neq j$
- Remaining intersections:
  - $C_{q-1} \cap T = \langle v_1, v_2 + s_1 v_3 \rangle$ ,  $s_1 \neq -b_i a_i^{-1}$
  - $C_q \cap T = \langle v_2, v_1 + s_2 v_3 \rangle$ ,  $s_2 \neq b_i$
  - $C_{q+1} \cap T = \langle v_3, v_1 + s_3 v_2 \rangle$ ,  $s_3 \neq a_i$

## Generator Matrix for $T$ :

$$\begin{bmatrix} 1 & 1 & \cdots & 1 & 0 & 1 & 1 \\ -a_1^{-1} & -a_2^{-1} & \cdots & -a_{q-2}^{-1} & 1 & 0 & -s_3^{-1} \\ -b_1^{-1} & -b_2^{-1} & \cdots & -b_{q-2}^{-1} & -s_1^{-1} & -s_2^{-1} & 0 \end{bmatrix}$$



If we put  $\theta(a_i) = -b_i$  then  $\delta(a_i) = a_i^{-1}\theta(a_i) = -a_i^{-1}b_i$  and  $\theta$  is a near orthomorphism of the multiplicative group of the field with ex elements  $(s_3, s_2, s_1)$ .