

# Goppa codes from a Singer cycle

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# Linear codes

## Definition

An  $[n, k]_q$ -linear code  $\mathcal{C}$  is a subspace of  $\mathbb{F}_q^n$  of dimension  $k$ .

## Definition

- The Hamming distance between two codewords  $x = (x_1, x_2, \dots, x_n)$  and  $y = (y_1, y_2, \dots, y_n)$  is the number of entries in which  $x$  and  $y$  differ:  
 $d(x, y) = |\{i | x_i \neq y_i\}|$ .
- The minimum distance of a code  $\mathcal{C}$  is  
 $d = d(\mathcal{C}) = \min\{d(x, y) | x, y \in \mathcal{C}, x \neq y\}$ .

In this case we say  $\mathcal{C}$  is a  $[n, k, d]_q$ -linear code.

## Theorem

Let  $\mathcal{C}$  be a  $[n, k, d]_q$ -linear code. Then,  $\mathcal{C}$  can correct  $\lfloor \frac{d-1}{2} \rfloor$  errors.  
If used for detection,  $\mathcal{C}$  can detect  $d - 1$  errors.

# Linear codes

## Dual codes

### Definition

Let  $\mathcal{C}$  be an  $[n, k]_q$ -linear code. Consider the standard inner product in  $\mathbb{F}_{q^n}$ :  $x \cdot y = \sum_{i=1}^n x_i y_i$ . The dual code  $\mathcal{C}^\perp$  is

$$\mathcal{C}^\perp = \{x \in \mathbb{F}_{q^n} \mid x \cdot c = 0, \forall c \in \mathcal{C}\}$$

### Theorem

$\mathcal{C}^\perp$  is a  $[n, n - k]_q$ -code.

# Linear codes

## Gilbert-Varshamov bound

### Proposition (Gilbert-Varshamov Bound)

An  $[n, k, d]_q$  code exists if

$$q^{n-k} > \sum_{i=0}^{d-2} \binom{n-1}{i} (q-1)^i.$$

# Goppa codes



# Goppa codes

Curves and divisors

$p$  prime,  $h \in \mathbb{N}$ ,  $q = p^h$ .

$\mathcal{C}$  : non-singular plane curve over  $\mathbb{F}_q$ .

## Definition

- A divisor  $G$  is a formal power series of places of  $\mathcal{C}$ .
- The Riemann-Roch space  $\mathcal{L}(G)$  is the vector space consisting of all rational functions that are regular outside  $G$ .

## Theorem (Riemann-Roch Theorem)

$$\ell(G) = \deg(G) - g + 1 + \ell(W - G),$$

where  $g$  is the genus of the curve,  $\ell(G) = \dim(\mathcal{L}(G))$  and  $W$  is a canonical divisor. In particular, for  $\deg(G) > 2g - 2$ ,

$$\ell(G) = \deg(G) - g + 1.$$

# Goppa codes

## Construction

The functional code  $C_L(D, G)$  arises as follows: take a divisor  $G$  with support  $G \subseteq \mathcal{C}$ , and take  $P_1, \dots, P_N = D$ , and assume  $D \cap G = \emptyset$ . Then evaluating the functions  $f \in \mathcal{L}(G)$  on  $D$  produces a linear code of length  $N$  and dimension  $\ell(G)$ .

## Proposition

The minimum distance of  $C_L(D, G)$  is at least  $\delta = n - \deg(G)$ .

## Definition

The differential code  $C_\Omega(D, G)$  is the dual code  $C_L^\perp(D, G)$ .

Here  $\mathcal{C}$  is the Hermitian curve  $H(2, q^2) : Y^q + Y - X^{q+1} = 0$ ,  $G$  is an orbit of a large  $\Gamma \leq Aut(H(2, q^2)) \cong PGU(3, q)$ ,  $G \cup D = \mathcal{C}$ .

# Goppa codes

## Subgroups of $PGU(3, q)$

### Theorem

Let  $d$  be a divisor of  $q = p^k$ . The following is the list of maximal subgroups in  $PSU(3, q)$  (up to conjugacy)

- (i) The one-point stabilizer (order  $\frac{q^3(q^2-1)}{d}$ );
- (ii) The stabilizer of a non-tangent line (order  $\frac{q(q^2-1)(q+1)}{d}$ );
- (iii) the stabilizer of a self-conjugate triangle (order  $\frac{6(q+1)}{d}$ );
- (iv) the normalizer of a cyclic Singer group (order  $\frac{3(q^2-q+1)}{d}$ );

further when  $q$  is odd:

- (v) the stabilizer of a conic  $PGL(2, q)$ ;
- (vi)  $PSU(3, p^m)$ , with  $m \mid k$  and  $\frac{k}{m}$  odd;
- (vii) the subgroup containing  $PSU(3, p^m)$  as index 3 normal subgroup, with  $m \mid k$ ,  $\frac{k}{m}$  odd, and 3 divides both  $q + 1$  and  $\frac{k}{m}$ ;
- (viii) the Hessian groups of order 216 when  $9 \mid (q + 1)$  and of order 72 and 36 when  $3 \mid (q + 1)$ ;
- (ix)  $PSL(2, 7)$  when either  $p = 7$  or  $-7$  is not a square in  $\mathbb{F}_q$ ;
- (x)  $A_6$  when either  $p = 3$  and  $k$  is even, or 5 is a square in  $\mathbb{F}_q$  and  $\mathbb{F}_q$  contains no cubic roots of the unity;
- (xi)  $S_6$  when  $p = 5$  and  $k$  odd;
- (xii)  $A_7$  when  $p = 5$  and  $k$  odd...

# Goppa codes

## Constructions of Goppa codes

- Group (i) *On Goppa codes and Weierstrass gaps at several points*,  
C. Carvalho, F. Torres, Designs, Codes and Cryptography, 2005, 35, pp. 211-225;
- Group (v) *Hermitian curves with automorphism group isomorphic to  $PGL(2, q)$  with  $q$  odd*, G. Korchmáros, P. Speziali,  
Finite Fields and their Applications, 2017, 44, pp. 1-17;
- Group (vi) *Codes and gap sequences of Hermitian curves*,  
G. Korchmáros, G. P. Nagy, M. Timpanella, IEEE Transactions  
of Information Theory, 2019, 66(6), pp. 3547-3554.

# Singer cycles on $PG(2, q^6)$

The group (iv)

The group of size  $3(q^2 - q + 1)$  is the normalizer of a Singer cycle. The Singer cycle acts on the Hermitian curve  $H(2, q^2)$  regularly on a point-orbit of length  $q^2 - q + 1$ . The matrices representing such a subgroup of  $PGU(3, q)$  may be represented by the  $3 \times 3$  matrices of the shape

$$\begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ a & b & c \end{pmatrix},$$

where  $X^3 + aX^2 + bX + c \in \mathbb{F}_{q^2}[X]$  is an irreducible polynomial.

# Singer cycles on $PG(2, q^6)$

Cubic extension

$$PG(2, q^2) \subseteq PG(2, q^6).$$

$a$  primitive  $(q^4 + q^2 + 1)$ -th root of the unity in  $\mathbb{F}_{q^6}$ .

$$M = \begin{pmatrix} a & 1 & a^{q^2+1} \\ a^{q^2+1} & a & 1 \\ 1 & a^{q^2+1} & a \end{pmatrix}$$

maps the canonical subplane  $PG(2, q^2)$  onto

$$\Pi = \{(a^i : a^{i(q^2+1)} : 1) \mid i = 0, 1, \dots, q^4 + q^2\}.$$

$$A_1(1 : 0 : 0) \mapsto A'_1(a : a^{q^2+1} : 1)$$

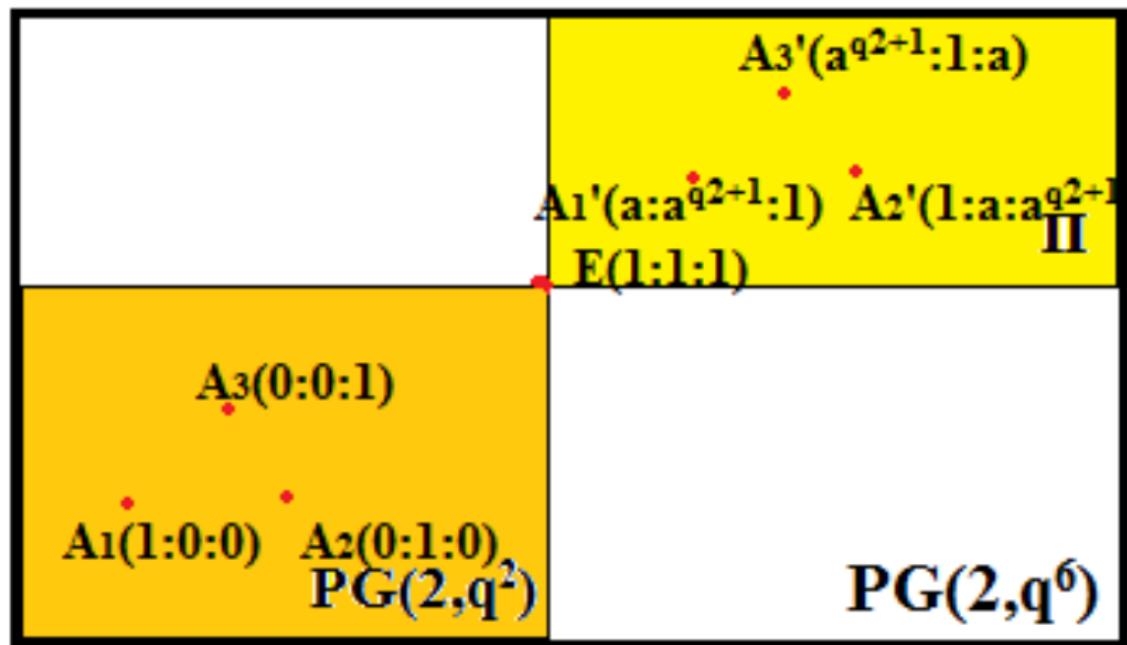
$$A_2(0 : 1 : 0) \mapsto A'_2(1 : a : a^{q^2+1})$$

$$A_3(0 : 0 : 1) \mapsto A'_3(a^{q^2+1} : 1 : a)$$

$E(1 : 1 : 1)$  is fixed.

# Singer cycles on $PG(2, q^6)$

Cubic extension



# Singer cycles on $PG(2, q^6)$

Cubic extension

## Construction

*In  $\Pi$ , the Singer cycle is represented by*

$$B = \begin{pmatrix} \beta & 0 & 0 \\ 0 & \beta^{q^2+1} & 0 \\ 0 & 0 & 1 \end{pmatrix},$$

*where  $\beta$  is a primitive  $(q^2 - q + 1)$ -th root of the unity.*

# Singer cycles on $PG(2, q^6)$

Cubic extension

$\mathcal{C} = v(G(\overline{X_0}, \overline{X_1}, \overline{X_2}))$  is the zero locus of the polynomial

$$\begin{aligned} G(\overline{X_0}, \overline{X_1}, \overline{X_2}) = & \overline{X_1}^{2q} \overline{X_2}^{2q} + \overline{X_0}^{2q} \overline{X_1}^{2q} + \overline{X_0}^{2q} \overline{X_2}^{2q} + \\ & -2(\overline{X_0}^{q+1} \overline{X_1}^q \overline{X_2} + \overline{X_0}^q \overline{X_1} \overline{X_2}^{q+1} + \overline{X_0} \overline{X_1}^{q+1} \overline{X_2}^q). \end{aligned}$$

$$g(\mathcal{C}) = \frac{q^2 - q}{2}, \quad |\mathcal{C}(\mathbb{F}_{q^6})| = q^6 + q^5 - q^4 + 1.$$

The singular points of  $\mathcal{C}$  have coordinates  $(1 : 0 : 0)$ ,  $(0 : 1 : 0)$ ,  $(0 : 0 : 1)$  in the system of coordinates  $(\overline{X_0}, \overline{X_1}, \overline{X_2})$ .

# Singer cycles on $PG(2, q^6)$

Cubic extension

$$M^{-1} = \frac{1}{|M|} \begin{pmatrix} a^2 - a^{q^2+1} & 1 - a^{q^2+2} & a^{2q^2+2} - a \\ a^{2q^2+2} - a & a^2 - a^{q^2+1} & 1 - a^{q^2+2} \\ 1 - a^{q^2+2} & a^{2q^2+2} - a & a^2 - a^{q^2+1} \end{pmatrix}.$$

$\mathcal{D} = v(H(X_0, X_1, X_2))$  is a plane model  $H(2, q^2)$ , where

$$H(X_0, X_1, X_2) = G(aX_0 + X_1 + a^{q^2+1}X_2,$$

$$a^{q^2+1}X_0 + aX_1 + X_2, X_0 + a^{q^2+1}X_1 + aX_2).$$

The singular points of  $\mathcal{D}$  have coordinates defined by the three columns of  $M^{-1}$ .

# The new codes

The functional code  $C_L(D, G)$

$P_1, \dots, P_{q^2-q+1}$  orbit of a Singer cycle.

$$G = P_1 + \dots + P_{q^2-q+1}.$$

$D$  divisor whose support is  $H(2, q^2) \setminus \{P_1, \dots, P_{q^2-q+1}\}$ .

## Theorem

*The code  $C_L(D, G)$  is a*

*$[q(q^2 - q + 1), \frac{q^2 - q}{2} + 2, (q - 1)(q^2 - q + 1)]_{q^2}$ -linear code.*

$$n = (q^3 + 1) - (q^2 - q + 1) = q^3 - q^2 + q = q(q^2 - q + 1).$$

$$k = \deg(G) - g + 1 = q^2 - q + 1 - \frac{q^2 - q}{2} + 1 = \frac{q^2 - q}{2} + 2.$$

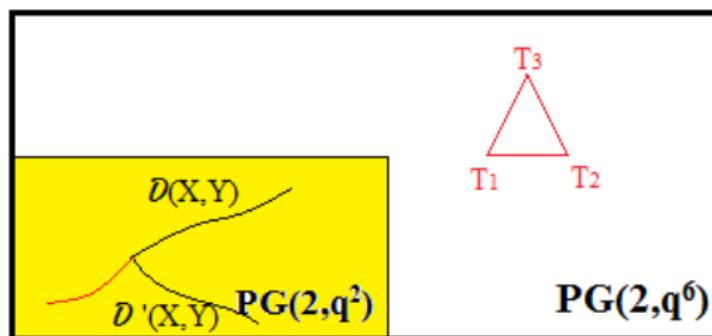
# The new codes

The minimum distance of  $C_L(D, G)$

$$\delta = n - \deg(G) = q(q^2 - q + 1) - (q^2 - q + 1) = (q - 1)(q^2 - q + 1).$$

## Construction

Take the codeword given by a further Hermitian curve  $\mathcal{D}'(X, Y)$ , intersecting  $\mathcal{D}(X, Y)$  at  $q^2 - q + 1$  points, while

$$q(q^2 - q + 1) - (q^2 - q + 1) = \delta.$$


# The new codes

The differential code  $C_{\Omega}(D, G)$

## Result

*There exists a canonical divisor W such that  $C_{\Omega}(D, G) \cong C_L(D, W + D - G)$ .*

$$W = \frac{F^2}{L} dx$$

$\mathcal{C}$  is another Hermitian curve  $F_q$  of equation  $F(x, y) = 0$  through the support of  $G$ , and  $L$  is the product of  $q^2 - q$  lines through an external point  $R$  to  $H_q$  together with the polar line of  $R$ .

$$W + D - G \equiv (q^3 - q^2 - 5q - 3)Y_{\infty} + 2qT$$

where  $T = T_1 + T_2 + T_3$ , the common points of  $H_q$  and  $F_q$  in  $PG(2, q^6)$ . Since  $D + qT \equiv (q + 1)^2 Y_{\infty}$ , this can also be written as

$$(q^2 - 1)(q + 1)Y_{\infty} - 2D.$$

# The new codes

The differential code  $C_{\Omega}(D, G)$

## Theorem

The code  $C_{\Omega}(D, G)$  is a  $[q(q^2 - q + 1), q^3 - \frac{3}{2}q^2 + \frac{3}{2}q - 2, \frac{1}{2}(q^2 - q + 4)]_{q^2}$ -linear code.

$$k = \deg(W + D - G) - g(H_q) + 1 = q^3 - \frac{3}{2}q^2 + \frac{3}{2}q - 2.$$

$$\delta = q(q^2 - q + 1) - \deg(W + D - G) = 3.$$

The minimum distance is  $d = \frac{1}{2}(q^2 - q + 4) > 3$ .

# The new codes

The minimum distance of  $C_\Omega(D, G)$

## Construction

*Take a chord  $\ell$  of  $D$  not passing through  $Y_\infty$*

*$\Lambda$  is the orbit of  $\ell$  under the action of the Singer cycle and consists of  $q^2 - q + 1$  pairwise distinct chords of  $D$  not through  $Y_\infty$ .*

*$\Lambda$  together with a further curve  $C$  of degree  $q - 2$  define a reducible curve  $L$  of degree  $q^2 - 1$ .*

$$\text{div}_0(L) - 2D = A_1 + A_2$$

*where  $A_1 = A_1 + \dots + A_N$  with  $N = (q - 1)(q^2 - q + 1)$  and  $A_2$  is the intersection divisor  $H_q \circ C$ .*

$$\deg(A_1) + \deg(A_2) = q^3 - 2q^2 + 2q - 1 + \frac{1}{2}(q^2 - q - 2).$$

*Therefore, the weight of the codeword  $A_1 + A_2$  equals  $d = \frac{1}{2}(q^2 - q + 4)$ .*



