

Goppa codes from a Singer cycle

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Linear codes

Definition

An $[n, k]_q$ -linear code \mathcal{C} is a subspace of \mathbb{F}_q^n of dimension k .

Definition

- The Hamming distance between two codewords $x = (x_1, x_2, \dots, x_n)$ and $y = (y_1, y_2, \dots, y_n)$ is the number of entries in which x and y differ:

$$d(x, y) = |\{i | x_i \neq y_i\}|.$$
- The minimum distance of a code \mathcal{C} is

$$d = d(\mathcal{C}) = \min\{d(x, y) | x, y \in \mathcal{C}, x \neq y\}.$$

In this case we say \mathcal{C} is a $[n, k, d]_q$ -linear code.

Theorem

Let \mathcal{C} be a $[n, k, d]_q$ -linear code. Then, \mathcal{C} can correct $\lfloor \frac{d-1}{2} \rfloor$ errors.
 If is used for detection, \mathcal{C} can detect $d - 1$ errors.

Linear codes

Dual codes

Definition

Let \mathcal{C} be an $[n, k]_q$ -linear code. Consider the standard inner product in \mathbb{F}_{q^n} : $x \cdot y = \sum_{i=1}^n x_i y_i$. The dual code \mathcal{C}^\perp is

$$\mathcal{C}^\perp = \{x \in \mathbb{F}_{q^n} \mid x \cdot c = 0, \forall c \in \mathcal{C}\}$$

Theorem

\mathcal{C}^\perp is a $[n, n - k]_q$ -code.

Linear codes

Gilbert-Varshamov bound

Proposition (**Gilbert-Varshamov Bound**)

An $[n, k, d]_q$ code exists if

$$q^{n-k} > \sum_{i=0}^{d-2} \binom{n-1}{i} (q-1)^i.$$

Goppa codes



Goppa codes

Curves and divisors

p prime, $h \in \mathbb{N}$, $q = p^h$.

\mathcal{C} : non-singular plane curve over \mathbb{F}_q .

Definition

- A divisor G is a formal power series of places of \mathcal{C} .
- The Riemann-Roch space $\mathcal{L}(G)$ is the vector space consisting of all rational functions that are regular outside G .

Theorem (Riemann-Roch Theorem)

$$\ell(G) = \deg(G) - g + 1 + \ell(W - G),$$

where g is the genus of the curve, $\ell(G) = \dim(\mathcal{L}(G))$ and W is a canonical divisor. In particular, for $\deg(G) > 2g - 2$,

$$\ell(G) = \deg(G) - g + 1.$$

Goppa codes

Construction

The functional code $C_L(D, G)$ arises as follows: take a divisor G with support $G \subseteq \mathcal{C}$, and take $P_1, \dots, P_N = D$, and assume $D \cap G = \emptyset$. Then evaluating the functions $f \in \mathcal{L}(G)$ on D produces a linear code of length N and dimension $\ell(G)$.

Proposition

The minimum distance of $C_L(D, G)$ is at least $\delta = n - \deg(G)$.

Definition

The differential code $C_\Omega(D, G)$ is the dual code $C_L^\perp(D, G)$.

Here \mathcal{C} is the Hermitian curve $H(2, q^2) : Y^q + Y - X^{q+1} = 0$, G is an orbit of a large $\Gamma \leq \text{Aut}(H(2, q^2)) \cong \text{PGU}(3, q)$, $G \cup D = \mathcal{C}$.

Goppa codes

Subgroups of $PGU(3, q)$

Theorem

Let d be a divisor of $q = p^k$. The following is the list of maximal subgroups in $PSU(3, q)$ (up to conjugacy)

- (i) The one-point stabilizer (order $\frac{q^3(q^2-1)}{d}$);
- (ii) The stabilizer of a non-tangent line (order $\frac{q(q^2-1)(q+1)}{d}$);
- (iii) the stabilizer of a self-conjugate triangle (order $\frac{6(q+1)}{d}$);
- (iv) the normalizer of a cyclic Singer group (order $\frac{3(q^2-q+1)}{d}$);

further when q is odd:

- (v) the stabilizer of a conic $PGL(2, q)$;
- (vi) $PSU(3, p^m)$, with $m \mid k$ and $\frac{k}{m}$ odd;
- (vii) the subgroup containing $PSU(3, p^m)$ as index 3 normal subgroup, with $m \mid k$, $\frac{k}{m}$ odd, and 3 divides both $q+1$ and $\frac{k}{m}$;
- (viii) the Hessian groups of order 216 when $9 \mid (q+1)$ and of order 72 and 36 when $3 \mid (q+1)$;
- (ix) $PSL(2, 7)$ when either $p = 7$ or -7 is not a square in \mathbb{F}_q ;
- (x) A_6 when either $p = 3$ and k is even, or 5 is a square in \mathbb{F}_q and \mathbb{F}_q contains no cubic roots of the unity;
- (xi) S_6 when $p = 5$ and k odd;
- (xii) A_7 when $p = 5$ and k odd...

Goppa codes

Constructions of Goppa codes

- Group (i) *On Goppa codes and Weierstrass gaps at several points*,
C. Carvalho, F. Torres, Designs, Codes and Cryptography,
2005, 35, pp. 211-225;
- Group (v) *Hermitian curves with automorphism group isomorphic to*
 $PGL(2, q)$ *with q odd*, G. Korchmáros, P. Speziali,
Finite Fields and their Applications, 2017, 44, pp. 1-17;
- Group (vi) *Codes and gap sequences of Hermitian curves*,
G. Korchmáros, G. P. Nagy, M. Timpanella, IEEE Transactions
of Information Theory, 2019, 66(6), pp. 3547-3554.

Singer cycles on $PG(2, q^6)$

The group (iv)

The group of size $3(q^2 - q + 1)$ is the normalizer of a Singer cycle. The Singer cycle acts on the Hermitian curve $H(2, q^2)$ regularly on a point-orbit of length $q^2 - q + 1$. The matrices representing such a subgroup of $PGU(3, q)$ may be represented by the 3×3 matrices of the shape

$$\begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ a & b & c \end{pmatrix},$$

where $X^3 + aX^2 + bX + c \in \mathbb{F}_{q^2}[X]$ is an irreducible polynomial.

Singer cycles on $PG(2, q^6)$

Cubic extension

$$PG(2, q^2) \subseteq PG(2, q^6).$$

a primitive $(q^4 + q^2 + 1)$ -th root of the unity in \mathbb{F}_{q^6} .

$$M = \begin{pmatrix} a & 1 & a^{q^2+1} \\ a^{q^2+1} & a & 1 \\ 1 & a^{q^2+1} & a \end{pmatrix}$$

maps the canonical subplane $PG(2, q^2)$ onto

$$\Pi = \{(a^i : a^{i(q^2+1)} : 1) \mid i = 0, 1, \dots, q^4 + q^2\}.$$

$$A_1(1 : 0 : 0) \mapsto A'_1(a : a^{q^2+1} : 1)$$

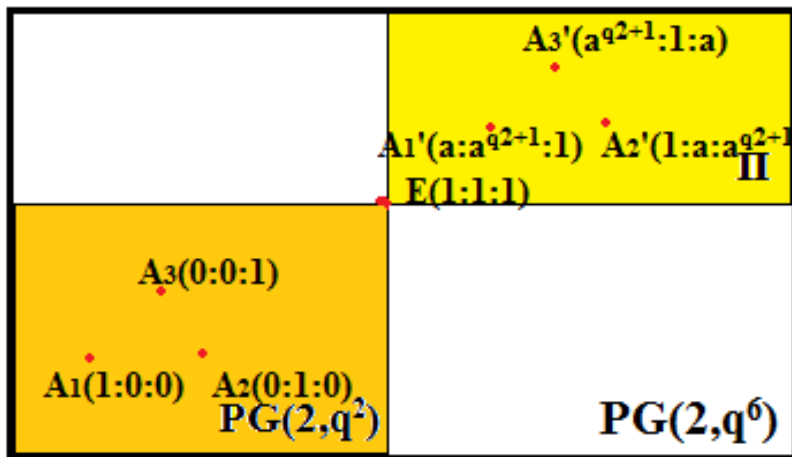
$$A_2(0 : 1 : 0) \mapsto A'_2(1 : a : a^{q^2+1})$$

$$A_3(0 : 0 : 1) \mapsto A'_3(a^{q^2+1} : 1 : a)$$

$E(1 : 1 : 1)$ is fixed.

Singer cycles on $PG(2, q^6)$

Cubic extension



Singer cycles on $PG(2, q^6)$

Cubic extension

Construction

In Π , the Singer cycle is represented by

$$B = \begin{pmatrix} \beta & 0 & 0 \\ 0 & \beta^{q^2+1} & 0 \\ 0 & 0 & 1 \end{pmatrix},$$

where β is a primitive $(q^2 - q + 1)$ -th root of the unity.

Singer cycles on $PG(2, q^6)$

Cubic extension

$\mathcal{C} = v(G(\overline{X}_0, \overline{X}_1, \overline{X}_2))$ is the zero locus of the polynomial

$$G(\overline{X}_0, \overline{X}_1, \overline{X}_2) = \overline{X}_1^2 \overline{X}_2^{2q} + \overline{X}_0^2 \overline{X}_1^{2q} + \overline{X}_0^{2q} \overline{X}_2^2 + \\ - 2(\overline{X}_0^{q+1} \overline{X}_1^q \overline{X}_2 + \overline{X}_0^q \overline{X}_1 \overline{X}_2^{q+1} + \overline{X}_0 \overline{X}_1^{q+1} \overline{X}_2^q).$$

$$g(\mathcal{C}) = \frac{q^2 - q}{2}, |\mathcal{C}(\mathbb{F}_{q^6})| = q^6 + q^5 - q^4 + 1.$$

The singular points of \mathcal{C} have coordinates $(1 : 0 : 0)$, $(0 : 1 : 0)$, $(0 : 0 : 1)$ in the system of coordinates $(\overline{X}_0, \overline{X}_1, \overline{X}_2)$.

Singer cycles on $PG(2, q^6)$

Cubic extension

$$M^{-1} = \frac{1}{|M|} \begin{pmatrix} a^2 - a^{q^2+1} & 1 - a^{q^2+2} & a^{2q^2+2} - a \\ a^{2q^2+2} - a & a^2 - a^{q^2+1} & 1 - a^{q^2+2} \\ 1 - a^{q^2+2} & a^{2q^2+2} - a & a^2 - a^{q^2+1} \end{pmatrix}.$$

$\mathcal{D} = v(H(X_0, X_1, X_2))$ is a plane model $H(2, q^2)$, where

$$H(X_0, X_1, X_2) = G(aX_0 + X_1 + a^{q^2+1}X_2,$$

$$a^{q^2+1}X_0 + aX_1 + X_2, X_0 + a^{q^2+1}X_1 + aX_2).$$

The singular points of \mathcal{D} have coordinates defined by the three columns of M^{-1} .

The new codes

The functional code $C_L(\mathbb{D}, \mathbb{G})$

P_1, \dots, P_{q^2-q+1} orbit of a Singer cycle.

$$\mathbb{G} = P_1 + \dots + P_{q^2-q+1}.$$

\mathbb{D} divisor whose support is $H(2, q^2) \setminus \{P_1, \dots, P_{q^2-q+1}\}$.

Theorem

The code $C_L(\mathbb{D}, \mathbb{G})$ is a

$[q(q^2 - q + 1), \frac{q^2-q}{2} + 2, (q-1)(q^2 - q + 1)]_{q^2}$ -linear code.

$$n = (q^3 + 1) - (q^2 - q + 1) = q^3 - q^2 + q = q(q^2 - q + 1).$$

$$k = \deg(\mathbb{G}) - g + 1 = q^2 - q + 1 - \frac{q^2-q}{2} + 1 = \frac{q^2-q}{2} + 2.$$

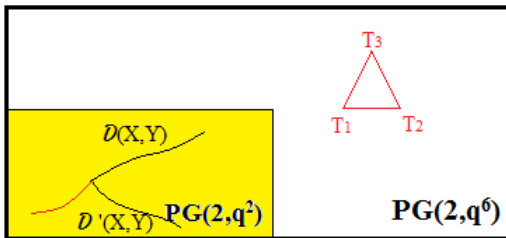
The new codes

The minimum distance of $C_L(\mathbb{D}, \mathbb{G})$

$$\delta = n - \deg(\mathbb{G}) = q(q^2 - q + 1) - (q^2 - q + 1) = (q - 1)(q^2 - q + 1).$$

Construction

Take the codeword given by a further Hermitian curve $\mathcal{D}'(X, Y)$, intersecting $\mathcal{D}(X, Y)$ at $q^2 - q + 1$ points, while $q(q^2 - q + 1) - (q^2 - q + 1) = \delta$.



The new codes

The differential code $C_\Omega(D, G)$

Result

There exists a canonical divisor W such that $C_\Omega(D, G) \cong C_L(D, W + D - G)$.

$$W = \frac{F^2}{L} dx$$

C is another Hermitian curve F_q of equation $F(x, y) = 0$ through the support of G , and L is the product of $q^2 - q$ lines through an external point R to H_q together with the polar line of R .

$$W + D - G \equiv (q^3 - q^2 - 5q - 3)Y_\infty + 2qT$$

where $T = T_1 + T_2 + T_3$, the common points of H_q and F_q in $PG(2, q^6)$. Since $D + qT \equiv (q + 1)^2 Y_\infty$, this can also be written as

$$(q^2 - 1)(q + 1)Y_\infty - 2D.$$

The new codes

The differential code $C_\Omega(D, G)$

Theorem

The code $C_\Omega(D, G)$ is a $[q(q^2 - q + 1), q^3 - \frac{3}{2}q^2 + \frac{3}{2}q - 2, \frac{1}{2}(q^2 - q + 4)]_{q^2}$ -linear code.

$$k = \deg(W + D - G) - g(H_q) + 1 = q^3 - \frac{3}{2}q^2 + \frac{3}{2}q - 2.$$

$$\delta = q(q^2 - q + 1) - \deg(W + D - G) = 3.$$

The minimum distance is $d = \frac{1}{2}(q^2 - q + 4) > 3$.

The new codes

The minimum distance of $C_\Omega(D, G)$

Construction

Take a chord ℓ of D not passing through Y_∞

Λ is the orbit of ℓ under the action of the Singer cycle and consists of $q^2 - q + 1$ pairwise distinct chords of D not through Y_∞ .

Λ together with a further curve C of degree $q - 2$ define a reducible curve L of degree $q^2 - 1$.

$$\operatorname{div}_0(L) - 2D = A_1 + A_2$$

where $A_1 = A_1 + \dots + A_N$ with $N = (q - 1)(q^2 - q + 1)$ and A_2 is the intersection divisor $H_q \circ C$.

$$\deg(A_1) + \deg(A_2) = q^3 - 2q^2 + 2q - 1 + \frac{1}{2}(q^2 - q - 2).$$

Therefore, the weight of the codeword $A_1 + A_2$ equals $d = \frac{1}{2}(q^2 - q + 4)$.

