1. Compute the tightest order, in big-O notation, of the following functions. In all cases, the variable to consider is 'n'.

A. $O(n^3)$ as $n^3 \ge n^2$ for $n \ge 1$ and $n^3 \ge -0.1n$ for $n \ge 1$. So if C = 5 and K = 1, Big O of f is $O(n^3)$.

- B. O(n) as $n \ge \sqrt{n}$ for $n \ge 1$. So if C = 3 and K = 1, Big O of f is O(n).
- C. O(1) as $1 \ge \sin(n)$ for $n \ge 1$. So if C = 1 and K = 1, Big O of f is O(1).
- D. $O(\sqrt{n})$ as $i \le \sqrt{n}$ for $n \ge 1$. So if C = 1 and K = 1, Big O of f is $O(\sqrt{n})$
- 2. If $f(n) = O(12n^2 + \log(n) + 3)$, Big O of f is $O(n^2)$. $n^2 \ge \log(n)$ for $n \ge 3$ and $n^2 \ge 3$ for $n \ge 3$. So, if C = 14 and K = 3, Big O of f is $O(n^2)$.

If $g(n) = O((2n^2 + 4n - 1))$, Big O of g is also $O(n^2)$. $4n^2 \ge 4n$ for $n \ge 1$ and $n^2 \ge -1$ for $n \ge 1$. So, if C = 7 and K = 1, Big O of g is $O(n^2)$.

Therefore, Big O of f(n) and g(n) are equal.

3.

public boolean removeStudent(Student s) {

```
L1
try {
       if(studentIndex(s) \ge 0) {
                                                L2
              int x = studentIndex(s);
                                                L3
              rawData[x] = null;
                                                L4
              for(int i = x; i < counter; i++) {
                                                L5
                     rawData[i] = rawData[i+1];L6
             }
              rawData[counter--] = null;
                                                L7
              return true;
                                                L8
```

```
}
                     return false;
                                                             L9
              } catch(NullPointerException e) {
                                                             L10
                     System.err.println(e.getMessage()); L11
                                                             L12
                     return false:
              }
       }
L1 = 1. L2 = N \rightarrow studentIndex function runs O(n) times worst case as it finds the index
by searching through rawData. L3 = N. L4 = 1. L5 = N. L(6) = N-1. L7 = 1. L8 = 1. L9 =
1. L10 = 1. L11 = 1. L12 = 1.
So, if NOT NullPointerException, 1 + N + N + 1 + N + (N-1) + 1 + 1 \rightarrow where
studentIndex(s) >= 0.
1 + N + 1 \rightarrow \text{where studentIndex}(s) < 0.
If NullPointerException, 1 + 1 + 1.
Therefore, the Big O of removeStudent for the worst case is O(n).
4.
public static double getMedian(int[] arr) {
       if(arr.length % 2 == 0) {
                                             // L1
              int a = arr.length / 2;
                                            //L2
              int b = (arr.length /2) - 1;
                                            //L3
              return (arr[a] + arr[b]) / 2.0; //L4
       } else {
              return arr[arr.length / 2];
                                           //L5
      }
}
If arr.length \% 2 = 0, then L1 = 1, L2 = 1, L3 = 1, L4 = 1. So, 1 + 1 + 1 + 1 = 4.
If arr.length \% 2 != 0, then L1 = 1, L5 = 1. So 1 + 1 = 2.
Worst case for getMedian runs with 4 constants and thus has O(1).
```