

- Define the problem

- Present a solution

- Prove the solution is correct

- Determine the complexity of the sol.

## Stable Matching

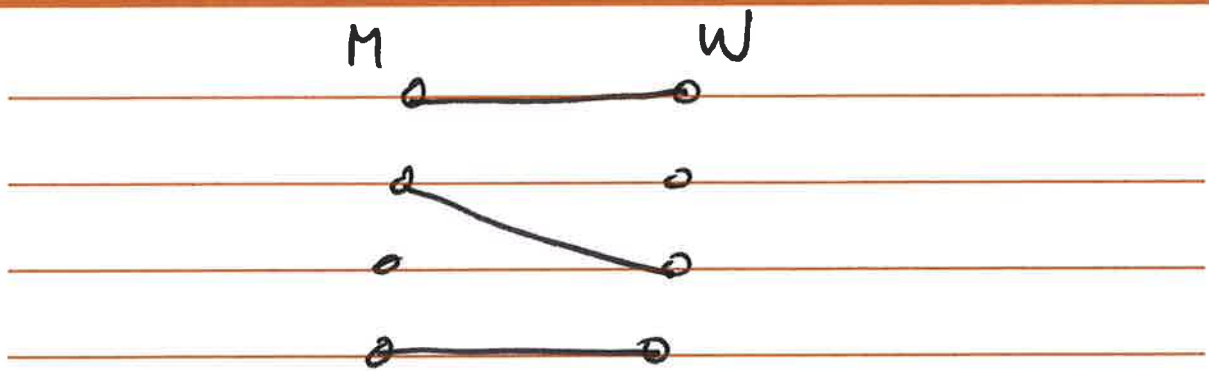
Problem we are interested in matching  $n$  men with  $n$  women so that they could end up happily married ever after.

we have a set of  $n$  men  $M = \{m_1, \dots, m_n\}$

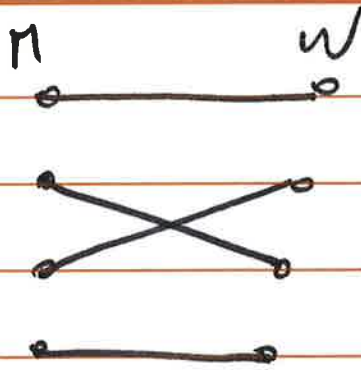
" " " " " women  $W = \{w_1, \dots, w_n\}$

Notation: an ordered pair  $m_i, w_j$  is shown as  $(m_i, w_j)$

Def. A matching  $S$  is a set of ordered pairs



Def. A perfect matching  $S'$  is a matching with the property that each member of  $M$  and each member of  $W$  appears in exactly one pair in  $S'$ .



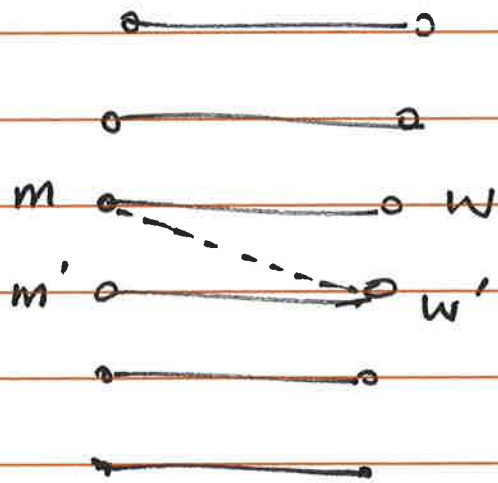
Add notion of preference

Each man  $m \in M$  ranks all women  
 $m$  prefers  $w$  to  $w'$  if  $m$   
 ranks  $w$  higher than  $w'$

o ordered ranking of  $m$  is his  
 preference list

$$P_{m_i} = \{w_{i_1}, w_{i_2}, \dots, w_{i_n}\}$$

Same for women



Def. Such a pair  $(m, w')$  is an instability wRT  $S$ .

( $m$  prefers  $w'$  to  $w$  &  $w'$  prefers  $m$  to  $m'$ )

Def. Matching  $S$  is stable if

- 1 - It is perfect
- 2 - There are no instabilities wRT  $S$ .



Input: preference lists for a set of  $n$  men &  $n$  women.

Output: Set of marriages ~~set~~ ~~no~~ ~~instabilities~~ representing a stable matching

Gale-shapley alg.

Proof of correctness:

- ① From the w's perspective, she starts single and once she gets engaged, she can only get into better engagements

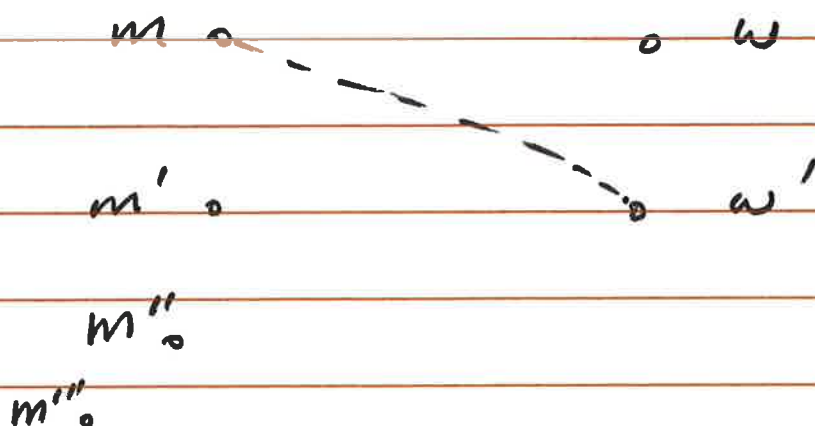
② From the  $m$ 's perspective, he starts single and gets engaged, and might be dropped repeatedly only to ~~see~~ settle for a lower ranking woman.

③ Alg terminates after  $n^2$  iterations

④ Solution is a perfect matching

⑤ Solution is stable

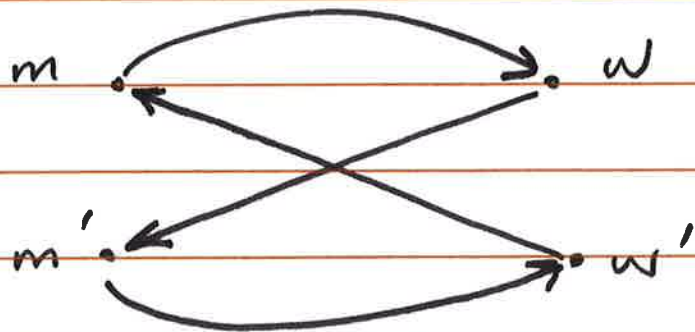
Assume instability with two pairs  
 $(m, w)$  and  $(m', w')$



Question: Did  $m$  propose to  $w'$   
at some point in the execution?

If no, then  $w$  must be higher than  
 $w'$  on his list  $\Rightarrow$  Contradiction!

If yes, he must have been rejected  
in favor of  $m''$  and due to ①  
either  $m'' = m'$  or  $m'$  is better  
than  $m'' \Rightarrow$  Contradiction!



$\rightarrow$  Men proposing:  $(m, w), (m', w')$

$\rightarrow$  Women proposing:  $(m', w), (m, w')$



## Complexity of Gale-Shapley

- 1 - Identify a free man
- 2 - For a man  $m$ , identify the highest ranked woman to whom he has not yet proposed.
- 3 - For a woman  $w$ , decide if  $w$  is engaged, if so to whom
- 4 - For a woman  $w$  and two men  $m$  &  $m'$  decide which man is preferred by  $w$ .
- 5 - place a man back in the list of free men