

we are given a directe I graph G=(V, E)

cul capacities on the edges

Associated cul each node Ve V is a

elemand dv.

if dv > 0, no de V has a demand

of dv for flow (Sink)

if dv (o, the node v has a supply of

|dv| for flow (Source)

if dv=0 neither a sink nor a

Source

Def. A circulation with demands {dv}
is a function of that assigns
non-negative real numbers to
each edge and satisfies:

1) Cap condition
for each edge eet, off(e) (Ce

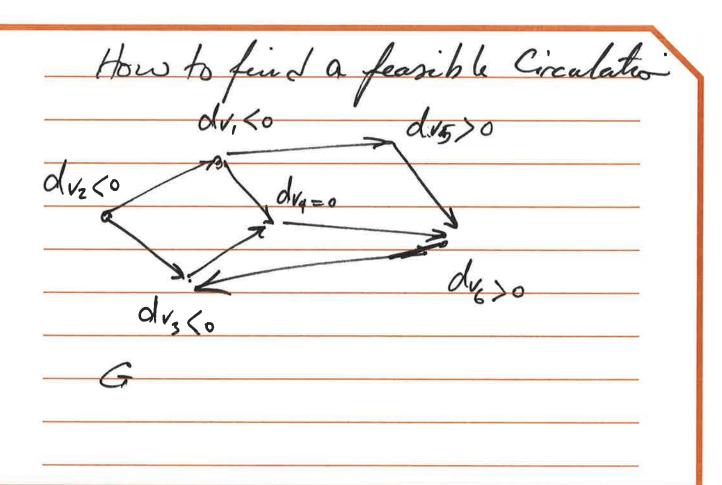
z) Demand Condition
for each veV, f'''(v)-f''(v)=dv

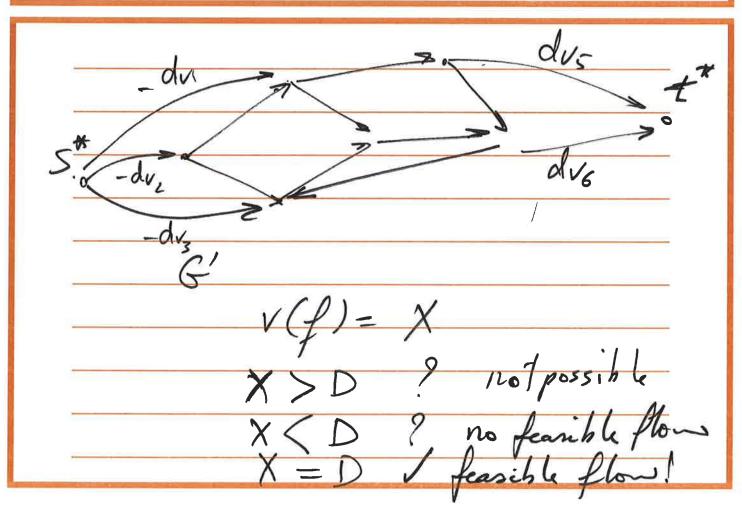
V f(e)

 $\sum_{v} dv = 0$

2 dv = 5 - dv = D V:dv>0 v:dv (o total demance

ralve





Proof. A If there is a feasible

Circulation for demand values

(d) in G, we can find

a Max flow din G' of

Value D.

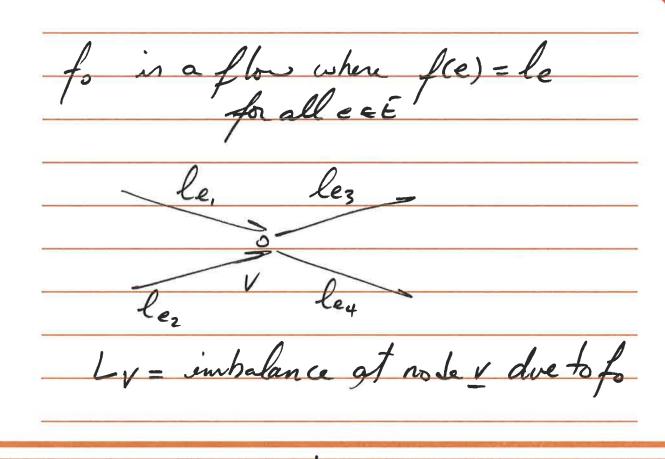
B- If there is a Max flow

in G' of value D, we can

find on circulation in G.

feasible

circulation w/ demands & lower bound
Conditions 1) Cap. Cond. for each eeE (f(e) (Ce
for each ext lefte) (Ce
2) Dema / Con /
2) Demand Cond for every v, f''(v)-fort(v)=dv



fin(v) for(v) = 5 le - 5 le
emtor eoutofr

= Lv

flow impalance
at v.

1- pash flow for three G

where f(e) = le

2. Construct G'
where Ce = Ce - le8. dv = dv - lv3- find fearible circulation in G'

Call this f,

4- if there is no fearible circ. G'

so fearible circ. G

otherwise, fearih le circ = 6=

fo + fi

