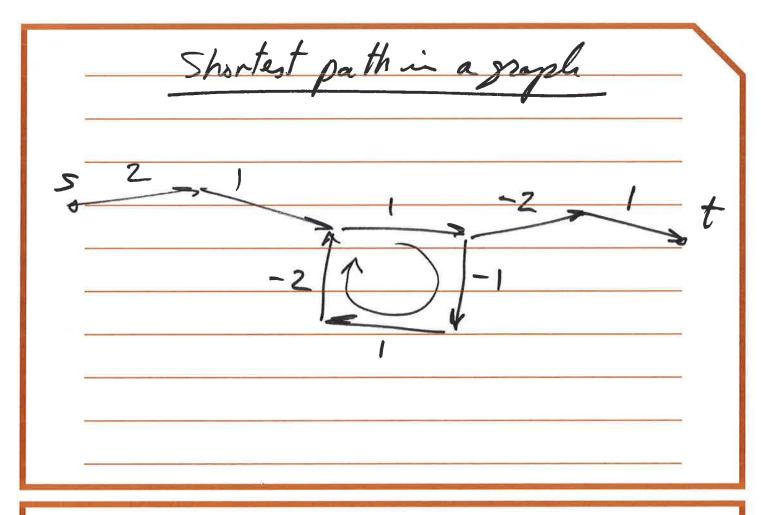
Shortest Path Problem
Dynamic Programming

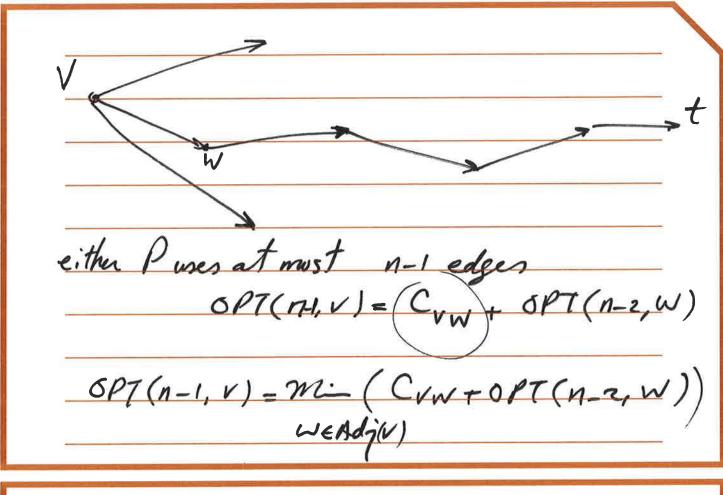


If G has no negative cycles, then
there is a shortest path from s to t
that is simple and hence has
at most n-1 edges.

OPT(i,v) denotes the win cost of a

Vet path using at most i edges

we want to compute of T(N-1, s)



$$\delta PT(\stackrel{\leftarrow}{=} n-1, v) = \delta PT(n-2, v)$$

$$\delta PT(n-1, v) = Mi \quad (\delta PT(n-2, v),$$

$$Mi \quad (C_{vw} + \delta PT(n-2, w))$$

$$w \in Adj(v)$$

