

CSCI 570 Spring 2015 Discussion 2

1. Suppose we have an array A of n integers and we wish to compute an $n \times n$ array B where $B[i,j]$ holds the sum of $A[i]$ through $A[j]$. The following code fragment is one way to solve this problem:

```
for i = 1 to n
  for j = i to n
    sum = 0;
    for k = i to j
      sum += A[k];
    B[i, j] = sum;
```

Handwritten annotations: $O(n)$ (with an arrow pointing to the outer loop), $O(n)$ (next to the inner loop), $O(n)$ (next to the innermost loop), and $O(n^3)$ (to the right of the code).

- What is the running time of the code fragment? Hint: it is $\Omega(n^2)$ but not $O(n^2)$.
- Do you think this *problem* can be solved in time better than $O(n^2)$? Why or why not?
- Give an $O(n^2)$ time solution to this problem.

2. Mathematicians often keep track of a statistic called their Erdős Number, after the great 20th century mathematician. Paul Erdős himself has a number of zero. Anyone who wrote a mathematical paper with him has a number of one, anyone who wrote a paper with someone who wrote a paper with him has a number of two, and so forth and so on. Supposing that we have a database of all mathematical papers ever written along with their authors:

- Explain how to represent this data as a graph.
- Explain how we would compute the Erdős number for a particular researcher.
- Explain how we would determine all researchers with Erdős number at most two.

3. In class, we discussed finding the shortest path between two vertices in a graph. Suppose instead we are interested in finding the *longest* simple path in a directed acyclic graph. In particular, I am interested in finding a path (if there is one) that visits all vertices. Given a DAG, give a linear-time algorithm to determine if there is a simple path that visits all vertices.

Solution Outlines:

1.
 - a. This is $O(n^3)$.
 - b. No; there are $O(n^2)$ output items.
 - c. As follows:

```
for i = 1 to n
    B[i,i] = A[i]
    for j = i+1 to n
        B[i,j] = B[i,j-1] + A[j]
```
2.
 - a. Set up a graph with vertices as researchers and (undirected) edges between co-authors.
 - b. BFS from Erdos or from the researcher; the layer the other is in is the number.
 - c. BFS from Erdos, stopping after layer two.
3. Do a topological sort. If that's a path, it's what we're looking for. If it isn't, no such path exists: it would need to have an inversion relative to that path, and by definition of topological sort, nothing that appears later in the output can have a path to anything former, so *that* wouldn't be a full path either.
(I'm debating whether or not to explicitly call this a Hamiltonian Path problem).

B

j →

i ↓	1	2	3					<u>n</u>
		1	2					n-1
			1	2				n-2
				1				
					1			
						1		
							1	
								1

$$1 \cdot n + 2 \cdot (n-1) + 3 \cdot (n-2) + \dots + n \cdot 1$$

$$\frac{n}{2} \cdot \frac{n}{2} + (\frac{n}{2}-1)(\frac{n}{2}+1) + (\frac{n}{2}-2)(\frac{n}{2}+2) + \dots$$

$$\frac{n^2}{4} + \frac{n^2}{4} - 1 + \frac{n^2}{4} - 4 + \frac{n^2}{4} - 9 \dots$$

$$\frac{n^2}{4} \cdot \frac{n}{2}$$

$$\frac{n^3}{8}$$

$$1 + 4 + 9 + \dots + \left(\frac{n}{2}\right)^2$$

$$\sum = \frac{n^3}{24}$$

$$\frac{n^3}{8} - \frac{n^3}{24} = \underline{\underline{C n^3}} = O(n^3)$$

b - No!

c - for $i=1$ to n

$$B[i, i] = A[i]$$

for $j = i+1$ to n

$$B[i, j] = B[i, j-1] + A[j]$$

- Nodes are researchers
- edges represent co-authorship

