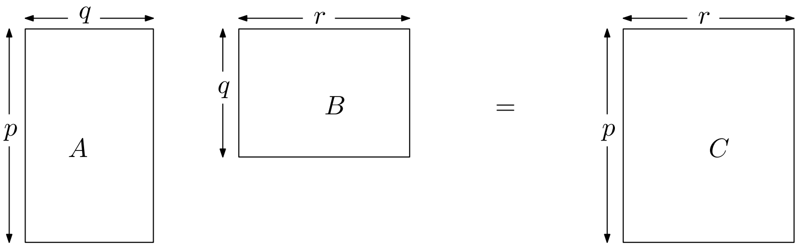
CSI - 3105 Design & Analysis of Algorithms
Course 14

Jean-Lou De Carufel

Fall 2019

Matrix Chain Multiplication

$$A : p \times q \text{ matrix}$$
$$B : q \times r \text{ matrix}$$
$$C = A\,B : p \times r \text{ matrix}$$



C has pr entries, each of which can be computed in $O(q)$ time. So C can be computed in $O(pqr)$ time. We define the *cost* of multiplying A and B to be pqr.

Consider 3 matrices

$$A_1 : 10 \times 100$$
$$A_2 : 100 \times 5$$
$$A_3 : 5 \times 50$$

How to compute $A_1 A_2 A_3$:

Consider 3 matrices

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$$A_2 : 100 \times 5$$
$$A_3 : 5 \times 50$$

How to compute $A_1 A_2 A_3$:

-
    - Compute $A_1 A_2$, $cost = 10 \times 100 \times 5 = 5000$

Consider 3 matrices

$$A_1 : 10 \times 100$$
$$A_2 : 100 \times 5$$
$$A_3 : 5 \times 50$$

How to compute $A_1 A_2 A_3$:

-
    - Compute $A_1 A_2$, $cost = 10 \times 100 \times 5 = 5000$
    - Compute $(A_1 A_2) A_3$, $cost = 10 \times 5 \times 50 = 2500$.

Consider 3 matrices

$$A_1 : 10 \times 100$$
$$A_2 : 100 \times 5$$
$$A_3 : 5 \times 50$$

How to compute $A_1 A_2 A_3$:

-
    - Compute $A_1 A_2$, $cost = 10 \times 100 \times 5 = 5000$
    - Compute $(A_1 A_2) A_3$, $cost = 10 \times 5 \times 50 = 2500$.

  For a total cost of $5000 + 2500 = 7500$.

Consider 3 matrices

$$A_1 : 10 \times 100$$
$$A_2 : 100 \times 5$$
$$A_3 : 5 \times 50$$

How to compute $A_1 A_2 A_3$:

-
  - Compute $A_1 A_2$, $cost = 10 \times 100 \times 5 = 5000$
  - Compute $(A_1 A_2) A_3$, $cost = 10 \times 5 \times 50 = 2500$.

  For a total cost of $5000 + 2500 = 7500$.
-
  - Compute $A_2 A_3$, $cost = 100 \times 5 \times 50 = 25000$

Consider 3 matrices

$$A_1 : 10 \times 100$$
$$A_2 : 100 \times 5$$
$$A_3 : 5 \times 50$$

How to compute $A_1A_2A_3$:

-
  - Compute A_1A_2, $cost = 10 \times 100 \times 5 = 5000$
  - Compute $(A_1A_2)A_3$, $cost = 10 \times 5 \times 50 = 2500$.

  For a total cost of $5000 + 2500 = 7500$.
-
  - Compute A_2A_3, $cost = 100 \times 5 \times 50 = 25000$
  - Compute $A_1(A_2A_3)$, $cost = 10 \times 100 \times 50 = 50000$.

Consider 3 matrices

$$A_1 : 10 \times 100$$
$$A_2 : 100 \times 5$$
$$A_3 : 5 \times 50$$

How to compute $A_1 A_2 A_3$:

-
  - Compute $A_1 A_2$, $cost = 10 \times 100 \times 5 = 5000$
  - Compute $(A_1 A_2) A_3$, $cost = 10 \times 5 \times 50 = 2500$.

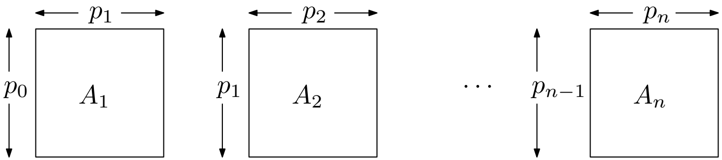
  For a total cost of $5000 + 2500 = 7500$.
-
  - Compute $A_2 A_3$, $cost = 100 \times 5 \times 50 = 25000$
  - Compute $A_1 (A_2 A_3)$, $cost = 10 \times 100 \times 50 = 50000$.

  For a total cost of $25000 + 50000 = 75000$.

Which one is better?

In general,

- $p_0, p_1, \ldots, p_n$: positive integers
- $A_1 A_2, \ldots, A_n$: matrices such that A_i has p_{i-1} rows and p_i columns.



Compute the best order to compute $A_1 A_2 \cdot \ldots \cdot A_n$, i.e., minimize the total cost.

Step 1: Structure of the Optimal Solution

Consider the best order to compute $A_i A_{i+1} \cdot \ldots \cdot A_j$. In the **last** step, we multiply

$$\underbrace{(A_i \cdot \ldots \cdot A_k)}_{\text{already computed}} \ \underbrace{(A_{k+1} \cdot \ldots \cdot A_j)}_{\text{already computed}}$$

for some k such that $i \leq k \leq j - 1$.

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How did we compute $A_i \cdot \ldots \cdot A_k$?

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How did we compute $A_i \cdot \ldots \cdot A_k$? In the best order.

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How did we compute $A_i \cdot \ldots \cdot A_k$? In the best order.
How did we compute $A_{k+1} \cdot \ldots \cdot A_j$? In the best order!

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How did we compute $A_i \cdot \ldots \cdot A_k$? In the best order.
How did we compute $A_{k+1} \cdot \ldots \cdot A_j$? In the best order!

minimum cost to compute $A_i \cdot \ldots \cdot A_j$

$=$

minimum cost to compute $A_i \cdot \ldots \cdot A_k$

$+$

minimum cost to compute $A_{k+1} \cdot \ldots \cdot A_j$

$+$

$p_{i-1} p_k p_j$

Step 2: Set Up a Recurrence for the Optimal Solution

For $1 \leq i \leq j \leq n$, define

$$m(i,j) = \text{minimum cost to compute } A_i \cdot \ldots \cdot A_j.$$

We want to compute $m(1,n)$.

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If we know k, then

$$m(i, j) = m(i, k) + m(k + 1, j) + p_{i-1} p_k p_j.$$

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But we do not know k, so try all values of k, $i \leq k \leq j-1$ and take the best one.

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But we do not know k, so try all values of k, $i \leq k \leq j-1$ and take the best one.

Recurrence:

- For $1 \leq i \leq n$: $m(i,i) = 0$.

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If we know k, then

$$m(i,j) = m(i,k) + m(k+1,j) + p_{i-1}p_k p_j.$$

But we do not know k, so try all values of k, $i \leq k \leq j-1$ and take the best one.

Recurrence:

- For $1 \leq i \leq n$: $m(i,i) = 0$.
- For $1 \leq i < j \leq n$:

$$m(i,j) = \min_{i \leq k \leq j-1} \{m(i,k) + m(k+1,j) + p_{i-1}p_k p_j\}$$

Step 3: Solve the Recurrence Bottom-Up

Compute, in this order,

$$
\begin{array}{l}
m(1,1), m(2,2), \ldots, m(n,n) \\
m(1,2), m(2,3), \ldots, m(n-1,n) \\
m(1,3), m(2,4), \ldots, m(n-2,n) \\
m(1,4), m(2,5), \ldots, m(n-3,n) \\
\qquad\quad \vdots \\
m(1,n-1), m(2,n) \\
m(1,n)
\end{array}
$$

Algorithm

Algorithm Matrix Chain Multiplication

```
for i = 1 to n do
    m(i, i) = 0
end for
for ℓ = 2 to n do
    // Compute m(1, ℓ), m(2, ℓ + 1), ..., m(n − ℓ + 1, n)
    for i = 1 to n − ℓ + 1 do
        // Compute m(i, i + ℓ − 1)
        j = i + ℓ − 1
        // Compute m(i, j) using the recurrence
        m(i, j) = ∞
        for k = i to j − 1 do
            m(i, j) = min {m(i, j), m(i, k) + m(k + 1, j) + p_{i−1} p_k p_j}
        end for
    end for
end for
return m(1, n)
```

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
Dimensions	30×35	35×15	15×5	5×10	10×20	20×25
	$p_0 \times p_1$	$p_1 \times p_2$	$p_2 \times p_3$	$p_3 \times p_4$	$p_4 \times p_5$	$p_5 \times p_6$

$$m(i,j) = \begin{cases} 0 & i = j \\ \min\limits_{i \leq k < j} \{m(i,k) + m(k+1,j) + p_{i-1}p_k p_j\} & 1 \leq i < j \leq 6 \end{cases}$$

$i \diagdown j$	1	2	3	4	5	6
1	0					
2		0				
3			0			
4				0		
5					0	
6						0

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
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$i \diagdown j$	1	2	3	4	5	6
1	0	15750				
2		0				
3			0			
4				0		
5					0	
6						0

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
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$i \diagdown j$	1	2	3	4	5	6
1	0	15750				
2		0	2625			
3			0	750		
4				0		
5					0	
6						0

Example

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3			0	750		
4				0	1000	
5					0	
6						0

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$i \diagdown j$	1	2	3	4	5	6
1	0	15750	7875			
2		0	2625			
3			0	750		
4				0	1000	
5					0	5000
6						0

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Matrices	A_1	A_2	A_3	A_4	A_5	A_6
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$i \diagdown j$	1	2	3	4	5	6
1	0	15750	7875			
2		0	2625	4375		
3			0	750		
4				0	1000	
5					0	5000
6						0

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3			0	750	2500	
4				0	1000	3500
5					0	5000
6						0

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1	0	15750	7875	9375		
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3			0	750	2500	
4				0	1000	3500
5					0	5000
6						0

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2		0	2625	4375	7125	10500
3			0	750	2500	5375
4				0	1000	3500
5					0	5000
6						0

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3			0	750	2500	5375
4				0	1000	3500
5					0	5000
6						0

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
Dimensions	30×35	35×15	15×5	5×10	10×20	20×25
	$p_0 \times p_1$	$p_1 \times p_2$	$p_2 \times p_3$	$p_3 \times p_4$	$p_4 \times p_5$	$p_5 \times p_6$

$$m(i,j) = \begin{cases} 0 & i = j \\ \min\limits_{i \leq k < j} \{m(i,k) + m(k+1,j) + p_{i-1}p_kp_j\} & 1 \leq i < j \leq 6 \end{cases}$$

$i \diagdown j$	1	2	3	4	5	6
1	0					
2		0				
3			0			
4				0		
5					0	
6						0

$i \diagdown j$	1	2	3	4	5	6
1	0					
2		0				
3			0			
4				0		
5					0	
6						0

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
Dimensions	30×35	35×15	15×5	5×10	10×20	20×25
	$p_0 \times p_1$	$p_1 \times p_2$	$p_2 \times p_3$	$p_3 \times p_4$	$p_4 \times p_5$	$p_5 \times p_6$

$$m(i,j) = \begin{cases} 0 & i = j \\ \min\limits_{i \le k < j} \{m(i,k) + m(k+1,j) + p_{i-1}p_kp_j\} & 1 \le i < j \le 6 \end{cases}$$

$i \diagdown j$	1	2	3	4	5	6
1	0	15750				
2		0				
3			0			
4				0		
5					0	
6						0

$i \diagdown j$	1	2	3	4	5	6
1	0	1				
2		0				
3			0			
4				0		
5					0	
6						0

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
Dimensions	30×35	35×15	15×5	5×10	10×20	20×25
	$p_0 \times p_1$	$p_1 \times p_2$	$p_2 \times p_3$	$p_3 \times p_4$	$p_4 \times p_5$	$p_5 \times p_6$

$$m(i,j) = \begin{cases} 0 & i = j \\ \min\limits_{i \le k < j} \{m(i,k) + m(k+1,j) + p_{i-1}p_kp_j\} & 1 \le i < j \le 6 \end{cases}$$

$i \diagdown j$	1	2	3	4	5	6
1	0	15750				
2		0	2625			
3			0			
4				0		
5					0	
6						0

$i \diagdown j$	1	2	3	4	5	6
1	0	1				
2		0	2			
3			0			
4				0		
5					0	
6						0

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
Dimensions	30×35	35×15	15×5	5×10	10×20	20×25
	$p_0 \times p_1$	$p_1 \times p_2$	$p_2 \times p_3$	$p_3 \times p_4$	$p_4 \times p_5$	$p_5 \times p_6$

$$m(i,j) = \begin{cases} 0 & i = j \\ \min\limits_{i \leq k < j} \{m(i,k) + m(k+1,j) + p_{i-1}p_kp_j\} & 1 \leq i < j \leq 6 \end{cases}$$

$i \diagdown j$	1	2	3	4	5	6
1	0	15750				
2		0	2625			
3			0	750		
4				0		
5					0	
6						0

$i \diagdown j$	1	2	3	4	5	6
1	0	1				
2		0	2			
3			0	3		
4				0		
5					0	
6						0

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
Dimensions	30×35	35×15	15×5	5×10	10×20	20×25
	$p_0 \times p_1$	$p_1 \times p_2$	$p_2 \times p_3$	$p_3 \times p_4$	$p_4 \times p_5$	$p_5 \times p_6$

$$m(i,j) = \begin{cases} 0 & i = j \\ \min_{i \le k < j} \{m(i,k) + m(k+1,j) + p_{i-1}p_kp_j\} & 1 \le i < j \le 6 \end{cases}$$

$i \diagdown j$	1	2	3	4	5	6
1	0	15750				
2		0	2625			
3			0	750		
4				0	1000	
5					0	
6						0

$i \diagdown j$	1	2	3	4	5	6
1	0	1				
2		0	2			
3			0	3		
4				0	4	
5					0	
6						0

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
Dimensions	30×35	35×15	15×5	5×10	10×20	20×25
	$p_0 \times p_1$	$p_1 \times p_2$	$p_2 \times p_3$	$p_3 \times p_4$	$p_4 \times p_5$	$p_5 \times p_6$

$$m(i,j) = \begin{cases} 0 & i = j \\ \min\limits_{i \leq k < j} \{m(i,k) + m(k+1,j) + p_{i-1}p_kp_j\} & 1 \leq i < j \leq 6 \end{cases}$$

$i \diagdown j$	1	2	3	4	5	6
1	0	15750				
2		0	2625			
3			0	750		
4				0	1000	
5					0	5000
6						0

$i \diagdown j$	1	2	3	4	5	6
1	0	1				
2		0	2			
3			0	3		
4				0	4	
5					0	5
6						0

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
Dimensions	30×35	35×15	15×5	5×10	10×20	20×25
	$p_0 \times p_1$	$p_1 \times p_2$	$p_2 \times p_3$	$p_3 \times p_4$	$p_4 \times p_5$	$p_5 \times p_6$

$$m(i,j) = \begin{cases} 0 & i = j \\ \min\limits_{i \leq k < j} \{m(i,k) + m(k+1,j) + p_{i-1}p_kp_j\} & 1 \leq i < j \leq 6 \end{cases}$$

$i \diagdown j$	1	2	3	4	5	6
1	0	15750	7875			
2		0	2625			
3			0	750		
4				0	1000	
5					0	5000
6						0

$i \diagdown j$	1	2	3	4	5	6
1	0	1	1			
2		0	2			
3			0	3		
4				0	4	
5					0	5
6						0

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
Dimensions	30×35	35×15	15×5	5×10	10×20	20×25
	$p_0 \times p_1$	$p_1 \times p_2$	$p_2 \times p_3$	$p_3 \times p_4$	$p_4 \times p_5$	$p_5 \times p_6$

$$m(i,j) = \begin{cases} 0 & i = j \\ \min\limits_{i \le k < j} \{m(i,k) + m(k+1,j) + p_{i-1}p_k p_j\} & 1 \le i < j \le 6 \end{cases}$$

$i \diagdown j$	1	2	3	4	5	6
1	0	15750	7875			
2		0	2625	4375		
3			0	750		
4				0	1000	
5					0	5000
6						0

$i \diagdown j$	1	2	3	4	5	6
1	0	1	1			
2		0	2	3		
3			0	3		
4				0	4	
5					0	5
6						0

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
Dimensions	30×35	35×15	15×5	5×10	10×20	20×25
	$p_0 \times p_1$	$p_1 \times p_2$	$p_2 \times p_3$	$p_3 \times p_4$	$p_4 \times p_5$	$p_5 \times p_6$

$$m(i,j) = \begin{cases} 0 & i = j \\ \min\limits_{i \leq k < j} \{m(i,k) + m(k+1,j) + p_{i-1}p_k p_j\} & 1 \leq i < j \leq 6 \end{cases}$$

$i \diagdown j$	1	2	3	4	5	6
1	0	15750	7875			
2		0	2625	4375		
3			0	750	2500	
4				0	1000	
5					0	5000
6						0

$i \diagdown j$	1	2	3	4	5	6
1	0	1	1			
2		0	2	3		
3			0	3	3	
4				0	4	
5					0	5
6						0

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
Dimensions	30×35	35×15	15×5	5×10	10×20	20×25
	$p_0 \times p_1$	$p_1 \times p_2$	$p_2 \times p_3$	$p_3 \times p_4$	$p_4 \times p_5$	$p_5 \times p_6$

$$m(i,j) = \begin{cases} 0 & i = j \\ \min\limits_{i \le k < j} \{m(i,k) + m(k+1,j) + p_{i-1}p_kp_j\} & 1 \le i < j \le 6 \end{cases}$$

$i \setminus j$	1	2	3	4	5	6
1	0	15750	7875			
2		0	2625	4375		
3			0	750	2500	
4				0	1000	3500
5					0	5000
6						0

$i \setminus j$	1	2	3	4	5	6
1	0	1	1			
2		0	2	3		
3			0	3	3	
4				0	4	5
5					0	5
6						0

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
Dimensions	30×35	35×15	15×5	5×10	10×20	20×25
	$p_0 \times p_1$	$p_1 \times p_2$	$p_2 \times p_3$	$p_3 \times p_4$	$p_4 \times p_5$	$p_5 \times p_6$

$$m(i,j) = \begin{cases} 0 & i = j \\ \min\limits_{i \leq k < j} \{m(i,k) + m(k+1,j) + p_{i-1}p_kp_j\} & 1 \leq i < j \leq 6 \end{cases}$$

$i \diagdown j$	1	2	3	4	5	6
1	0	15750	7875	9375		
2		0	2625	4375		
3			0	750	2500	
4				0	1000	3500
5					0	5000
6						0

$i \diagdown j$	1	2	3	4	5	6
1	0	1	1	3		
2		0	2	3		
3			0	3	3	
4				0	4	5
5					0	5
6						0

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
Dimensions	30×35	35×15	15×5	5×10	10×20	20×25
	$p_0 \times p_1$	$p_1 \times p_2$	$p_2 \times p_3$	$p_3 \times p_4$	$p_4 \times p_5$	$p_5 \times p_6$

$$m(i,j) = \begin{cases} 0 & i = j \\ \min\limits_{i \le k < j} \{m(i,k) + m(k+1,j) + p_{i-1}p_k p_j\} & 1 \le i < j \le 6 \end{cases}$$

$i \setminus j$	1	2	3	4	5	6
1	0	15750	7875	9375		
2		0	2625	4375	7125	
3			0	750	2500	
4				0	1000	3500
5					0	5000
6						0

$i \setminus j$	1	2	3	4	5	6
1	0	1	1	3		
2		0	2	3	3	
3			0	3	3	
4				0	4	5
5					0	5
6						0

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
Dimensions	30×35	35×15	15×5	5×10	10×20	20×25
	$p_0 \times p_1$	$p_1 \times p_2$	$p_2 \times p_3$	$p_3 \times p_4$	$p_4 \times p_5$	$p_5 \times p_6$

$$m(i,j) = \begin{cases} 0 & i = j \\ \min\limits_{i \le k < j} \{m(i,k) + m(k+1,j) + p_{i-1}p_k p_j\} & 1 \le i < j \le 6 \end{cases}$$

$i \diagdown j$	1	2	3	4	5	6
1	0	15750	7875	9375		
2		0	2625	4375	7125	
3			0	750	2500	5375
4				0	1000	3500
5					0	5000
6						0

$i \diagdown j$	1	2	3	4	5	6
1	0	1	1	3		
2		0	2	3	3	
3			0	3	3	3
4				0	4	5
5					0	5
6						0

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
Dimensions	30×35	35×15	15×5	5×10	10×20	20×25
	$p_0 \times p_1$	$p_1 \times p_2$	$p_2 \times p_3$	$p_3 \times p_4$	$p_4 \times p_5$	$p_5 \times p_6$

$$m(i,j) = \begin{cases} 0 & i = j \\ \min\limits_{i \le k < j} \{m(i,k) + m(k+1,j) + p_{i-1}p_kp_j\} & 1 \le i < j \le 6 \end{cases}$$

$i \diagdown j$	1	2	3	4	5	6
1	0	15750	7875	9375	11875	
2		0	2625	4375	7125	
3			0	750	2500	5375
4				0	1000	3500
5					0	5000
6						0

$i \diagdown j$	1	2	3	4	5	6
1	0	1	1	3	3	
2		0	2	3	3	
3			0	3	3	3
4				0	4	5
5					0	5
6						0

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
Dimensions	30×35	35×15	15×5	5×10	10×20	20×25
	$p_0 \times p_1$	$p_1 \times p_2$	$p_2 \times p_3$	$p_3 \times p_4$	$p_4 \times p_5$	$p_5 \times p_6$

$$m(i,j) = \begin{cases} 0 & i = j \\ \min\limits_{i \le k < j} \{m(i,k) + m(k+1,j) + p_{i-1}p_kp_j\} & 1 \le i < j \le 6 \end{cases}$$

$i \diagdown j$	1	2	3	4	5	6
1	0	15750	7875	9375	11875	
2		0	2625	4375	7125	10500
3			0	750	2500	5375
4				0	1000	3500
5					0	5000
6						0

$i \diagdown j$	1	2	3	4	5	6
1	0	1	1	3	3	
2		0	2	3	3	3
3			0	3	3	3
4				0	4	5
5					0	5
6						0

Example

Matrices	A_1	A_2	A_3	A_4	A_5	A_6
Dimensions	30×35	35×15	15×5	5×10	10×20	20×25
	$p_0 \times p_1$	$p_1 \times p_2$	$p_2 \times p_3$	$p_3 \times p_4$	$p_4 \times p_5$	$p_5 \times p_6$

$$m(i,j) = \begin{cases} 0 & i = j \\ \min\limits_{i \le k < j} \{m(i,k) + m(k+1,j) + p_{i-1}p_kp_j\} & 1 \le i < j \le 6 \end{cases}$$

$i \diagdown j$	1	2	3	4	5	6
1	0	15750	7875	9375	11875	15125
2		0	2625	4375	7125	10500
3			0	750	2500	5375
4				0	1000	3500
5					0	5000
6						0

$i \diagdown j$	1	2	3	4	5	6
1	0	1	1	3	3	3
2		0	2	3	3	3
3			0	3	3	3
4				0	4	5
5					0	5
6						0

Running Time

There are 3 nested loops, so this algorithm takes $O(n^3)$ time.

More careful counting:

- ℓ: 2 to n
- For each ℓ, we have i: 1 to $n - \ell + 1$
- For each i, we have k: i to $i + \ell - 2$

Algorithm

Algorithm Matrix Chain Multiplication

for $i = 1$ to n **do**
    $m(i, i) = 0$
end for
for $\ell = 2$ to n **do**
    // Compute $m(1, \ell), m(2, \ell + 1), \ldots, m(n - \ell + 1, n)$
    for $i = 1$ to $n - \ell + 1$ **do**
        // Compute $m(i, i + \ell - 1)$
        $j = i + \ell - 1$
        // Compute $m(i, j)$ using the recurrence
        $m(i, j) = \infty$
        for $k = i$ to $j - 1$ **do**
            $m(i, j) = \min\{m(i, j), m(i, k) + m(k + 1, j) + p_{i-1} p_k p_j\}$
        end for
    end for
end for
return $m(1, n)$

Running Time

Total time:

$$\sum_{\ell=2}^{n} \sum_{i=1}^{n-\ell+1} \sum_{k=i}^{i+\ell-2} 1$$

Running Time

Total time:

$$\sum_{\ell=2}^{n} \sum_{i=1}^{n-\ell+1} \sum_{k=i}^{i+\ell-2} 1$$

$$= \sum_{\ell=2}^{n} \sum_{i=1}^{n-\ell+1} (\ell - 1)$$

Running Time

Total time:

$$
\begin{aligned}
&\sum_{\ell=2}^{n} \sum_{i=1}^{n-\ell+1} \sum_{k=i}^{i+\ell-2} 1 \\
&= \sum_{\ell=2}^{n} \sum_{i=1}^{n-\ell+1} (\ell - 1) \\
&= \sum_{\ell=2}^{n} (n - \ell + 1)(\ell - 1)
\end{aligned}
$$

Running Time

Total time:

$$
\begin{aligned}
& \sum_{\ell=2}^{n} \sum_{i=1}^{n-\ell+1} \sum_{k=i}^{i+\ell-2} 1 \\
&= \sum_{\ell=2}^{n} \sum_{i=1}^{n-\ell+1} (\ell - 1) \\
&= \sum_{\ell=2}^{n} (n - \ell + 1)(\ell - 1) \\
&= \sum_{\ell=1}^{n} (n - \ell + 1)(\ell - 1) \qquad \text{since the summand is 0 when } \ell = 0
\end{aligned}
$$

Running Time

Total time:

$$\begin{aligned}
&\sum_{\ell=2}^{n} \sum_{i=1}^{n-\ell+1} \sum_{k=i}^{i+\ell-2} 1 \\
&= \sum_{\ell=2}^{n} \sum_{i=1}^{n-\ell+1} (\ell - 1) \\
&= \sum_{\ell=2}^{n} (n - \ell + 1)(\ell - 1) \\
&= \sum_{\ell=1}^{n} (n - \ell + 1)(\ell - 1) \qquad \text{since the summand is 0 when } \ell = 0 \\
&= \frac{n^3 - n}{6} \\
&= \Theta(n^3).
\end{aligned}$$