

CSI - 3105 Design & Analysis of Algorithms

Course 3

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Fall 2019

Matrix Multiplication

$$A = \begin{pmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n,1} & a_{n2} & \cdots & a_{nn} \end{pmatrix} \quad B = \begin{pmatrix} b_{11} & b_{12} & \cdots & b_{1n} \\ b_{21} & b_{22} & \cdots & b_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ b_{n,1} & b_{n2} & \cdots & b_{nn} \end{pmatrix}$$

$$C = AB$$

$$c_{ij} = \sum_{k=1}^n a_{ik} b_{kj}$$

Matrix Multiplication

Example

$$\begin{pmatrix} \mathbf{1} & \mathbf{0} & \mathbf{-1} \\ \mathbf{0} & \mathbf{2} & \mathbf{3} \\ 1 & 5 & 0 \end{pmatrix} \begin{pmatrix} \mathbf{4} & 0 & \mathbf{7} \\ \mathbf{-2} & 1 & \mathbf{0} \\ \mathbf{6} & 1 & \mathbf{-1} \end{pmatrix} = \begin{pmatrix} -2 & -1 & \mathbf{8} \\ \mathbf{14} & 5 & -3 \\ -6 & 5 & 7 \end{pmatrix}$$

Matrix Multiplication

Input: Two square matrices $A_{n \times n}$ and $B_{n \times n}$.

Output: AB .

```
1: Initialize matrix  $C_{n \times n}$ 
2: for  $i = 1$  to  $n$  do
3:   for  $j = 1$  to  $n$  do
4:      $c_{ij} = 0$ 
5:     for  $k = 1$  to  $n$  do
6:        $c_{ij} = c_{ij} + a_{ik}b_{kj}$ 
7:     end for
8:   end for
9: end for
10: return  $C$ 
```

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7:     end for
8:   end for
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10: return  $C$ 

```

$$T(n) = \Theta(n^3)$$

How can we Divide and Conquer?

Assume $n = 2^k$.

$$\begin{pmatrix} 1 & 0 & -1 & 2 \\ 3 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 2 & 4 & 7 & 1 \end{pmatrix} \begin{pmatrix} 0 & 1 & 2 & 3 \\ -2 & 1 & -1 & 1 \\ 1 & 0 & 1 & 0 \\ 5 & 1 & -2 & -1 \end{pmatrix} = \begin{pmatrix} 9 & 3 & -3 & 1 \\ 4 & 5 & 4 & 9 \\ 5 & 1 & -2 & -1 \\ 4 & 7 & 5 & 9 \end{pmatrix}$$

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$$\left(\begin{pmatrix} 1 & 0 \\ 3 & 1 \\ 0 & 0 \\ 2 & 4 \end{pmatrix} \begin{pmatrix} -1 & 2 \\ 1 & 1 \\ 0 & 1 \\ 7 & 1 \end{pmatrix} \right) \left(\begin{pmatrix} 0 & 1 \\ -2 & 1 \\ 1 & 0 \\ 5 & 1 \end{pmatrix} \begin{pmatrix} 2 & 3 \\ -1 & 1 \\ 1 & 0 \\ -2 & -1 \end{pmatrix} \right) = \left(\begin{pmatrix} 9 & 3 \\ 4 & 5 \\ 5 & 1 \\ 4 & 7 \end{pmatrix} \begin{pmatrix} -3 & 1 \\ 4 & 9 \\ -2 & -1 \\ 5 & 9 \end{pmatrix} \right)$$

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$$A = \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \quad B = \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}$$

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$$A = \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \quad B = \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}$$

$$C = \begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix} = \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}$$

$$C_{11} = A_{11}B_{11} + A_{12}B_{21}$$

$$C_{12} = A_{11}B_{12} + A_{12}B_{22}$$

$$C_{21} = A_{21}B_{11} + A_{22}B_{21}$$

$$C_{22} = A_{21}B_{12} + A_{22}B_{22}$$

How can we Divide and Conquer?

Assume $n = 2^k$.

$$\begin{pmatrix} 1 & 0 & -1 & 2 \\ 3 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 2 & 4 & 7 & 1 \end{pmatrix} \begin{pmatrix} 0 & 1 & 2 & 3 \\ -2 & 1 & -1 & 1 \\ 1 & 0 & 1 & 0 \\ 5 & 1 & -2 & -1 \end{pmatrix} = \begin{pmatrix} 9 & 3 & -3 & 1 \\ 4 & 5 & 4 & 9 \\ 5 & 1 & -2 & -1 \\ 4 & 7 & 5 & 9 \end{pmatrix}$$

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$$A = \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \quad B = \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}$$

$$C = \begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix} = \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix} \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix}$$

$$C_{11} = A_{11}B_{11} + A_{12}B_{21}$$

$$C_{12} = A_{11}B_{12} + A_{12}B_{22}$$

$$C_{21} = A_{21}B_{11} + A_{22}B_{21}$$

$$C_{22} = A_{21}B_{12} + A_{22}B_{22}$$

$$T(n) = 8 T(n/2) + 4 (n/2)^2 = 8 T(n/2) + n^2$$

$$T(n) = \begin{cases} 1 & \text{if } n=1 \\ 8T(n/2) + n^2 & \text{if } n \geq 2 \end{cases}$$

Assume $n = 2^k$.

$$T(n) = 8T(n/2) + n^2$$

$$= 8 \left(8T(n/2^2) + \left(\frac{n}{2}\right)^2 \right) + n^2$$

$$= 8^2 T(n/2^2) + 2n^2 + n^2$$

$$= 8^2 \left(8T(n/2^3) + \left(\frac{n}{2^2}\right)^2 \right) + 2n^2 + n^2$$

$$= 8^3 T(n/2^3) + 4n^2 + 2n^2 + n^2$$

$$= 8^3 \left(8T(n/2^4) + \left(\frac{n}{2^3}\right)^2 \right) + 4n^2 + 2n^2 + n^2$$

$$= 8^4 \cdot T(n/2^4) + 8n^2 + 4n^2 + 2n^2 + n^2$$

$$= 8^4 \cdot T(n/2^4) + (2^0 + 2^1 + 2^2 + 2^3) n^2$$

\vdots

$$= 8^k \cdot \overbrace{T(n/2^k)}^{T(1)=1} + \underbrace{(2^0 + 2^1 + \dots + 2^{k-1})}_{2^k - 1 = n - 1} n^2$$

$(2^3)^k = (2^k)^3 = n^3$

$$= n^3 + (n-1)n^2$$

$$= O(n^3), \quad \text{no improvement!}$$

Strassen Algorithm (1969)

$$\begin{array}{lllll} S_1 = B_{12} - B_{22} & S_2 = A_{11} + A_{12} & S_3 = A_{21} + A_{22} & S_4 = B_{21} - B_{11} & S_5 = A_{11} + A_{22} \\ S_6 = B_{11} + B_{22} & S_7 = A_{12} - A_{22} & S_8 = B_{21} + B_{22} & S_9 = A_{11} - A_{21} & S_{10} = B_{11} + B_{12} \end{array}$$

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$$P_1 = A_{11}S_1 = A_{11}B_{12} - A_{11}B_{22}$$

$$P_2 = S_2B_{22} = A_{11}B_{22} + A_{12}B_{22}$$

$$P_3 = S_3B_{11} = A_{21}B_{11} + A_{22}B_{11}$$

$$P_4 = A_{22}S_4 = A_{22}B_{21} - A_{22}B_{11}$$

$$P_5 = S_5S_6 = A_{11}B_{11} + A_{11}B_{22} + A_{22}B_{11} + A_{22}B_{22}$$

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$$P_5 + P_4 - P_2 + P_6$$

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$$\begin{array}{rcll} & P_5 + P_4 - P_2 + P_6 & & \\ = & A_{11}B_{11} & + & A_{11}B_{22} & + & A_{22}B_{11} & + & A_{22}B_{22} \\ & & & & & - & A_{22}B_{11} & & + & A_{22}B_{21} \\ & & & & & - & A_{11}B_{22} & & & - & A_{12}B_{22} \\ & & & & & & & & & - & A_{22}B_{22} & - & A_{22}B_{21} & + & A_{12}B_{22} & + & A_{12}B_{21} \end{array}$$

Strassen Algorithm (1969)

$$\begin{array}{lllll} S_1 = B_{12} - B_{22} & S_2 = A_{11} + A_{12} & S_3 = A_{21} + A_{22} & S_4 = B_{21} - B_{11} & S_5 = A_{11} + A_{22} \\ S_6 = B_{11} + B_{22} & S_7 = A_{12} - A_{22} & S_8 = B_{21} + B_{22} & S_9 = A_{11} - A_{21} & S_{10} = B_{11} + B_{12} \end{array}$$

$$P_1 = A_{11}S_1 = A_{11}B_{12} - A_{11}B_{22}$$

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$$P_5 = S_5S_6 = A_{11}B_{11} + A_{11}B_{22} + A_{22}B_{11} + A_{22}B_{22}$$

$$P_6 = S_7S_8 = A_{12}B_{21} + A_{12}B_{22} - A_{22}B_{21} - A_{22}B_{22}$$

$$P_7 = S_9S_{10} = A_{11}B_{11} + A_{11}B_{12} - A_{21}B_{11} - A_{21}B_{12}$$

$$\begin{aligned} & P_5 + P_4 - P_2 + P_6 \\ = & \begin{array}{ccccccc} A_{11}B_{11} & + & A_{11}B_{22} & + & A_{22}B_{11} & + & A_{22}B_{22} \\ & & & & - & A_{22}B_{11} & \\ & & & & & & + A_{22}B_{21} \\ & & - & A_{11}B_{22} & & & - A_{12}B_{22} \\ & & & & - & A_{22}B_{22} & - A_{22}B_{21} & + A_{12}B_{22} & + A_{12}B_{21} \end{array} \\ = & \frac{A_{11}B_{11}}{A_{11}B_{11}} \end{aligned}$$

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$$\begin{array}{lllll} S_1 = B_{12} - B_{22} & S_2 = A_{11} + A_{12} & S_3 = A_{21} + A_{22} & S_4 = B_{21} - B_{11} & S_5 = A_{11} + A_{22} \\ S_6 = B_{11} + B_{22} & S_7 = A_{12} - A_{22} & S_8 = B_{21} + B_{22} & S_9 = A_{11} - A_{21} & S_{10} = B_{11} + B_{12} \end{array}$$

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$$P_3 = S_3B_{11} = A_{21}B_{11} + A_{22}B_{11}$$

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$$\begin{aligned} & P_5 + P_4 - P_2 + P_6 \\ = & \begin{array}{ccccccc} A_{11}B_{11} & + & A_{11}B_{22} & + & A_{22}B_{11} & + & A_{22}B_{22} \\ & & & & - & A_{22}B_{11} & \\ & & & & & & + A_{22}B_{21} \\ & & - & A_{11}B_{22} & & & - A_{12}B_{22} \\ & & & & - & A_{22}B_{22} & - A_{22}B_{21} & + A_{12}B_{22} & + A_{12}B_{21} \end{array} \\ = & \frac{A_{11}B_{11}}{+ A_{12}B_{21}} \\ = & C_{11} \end{aligned}$$

Strassen Algorithm (1969)

$$\begin{array}{lllll} S_1 = B_{12} - B_{22} & S_2 = A_{11} + A_{12} & S_3 = A_{21} + A_{22} & S_4 = B_{21} - B_{11} & S_5 = A_{11} + A_{22} \\ S_6 = B_{11} + B_{22} & S_7 = A_{12} - A_{22} & S_8 = B_{21} + B_{22} & S_9 = A_{11} - A_{21} & S_{10} = B_{11} + B_{12} \end{array}$$

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$$P_2 = S_2 B_{22} = A_{11} B_{22} + A_{12} B_{22}$$

$$P_3 = S_3 B_{11} = A_{21} B_{11} + A_{22} B_{11}$$

$$P_4 = A_{22}S_4 = A_{22}B_{21} - A_{22}B_{11}$$

$$P_5 = S_5 S_6 = A_{11} B_{11} + A_{11} B_{22} + A_{22} B_{11} + A_{22} B_{22}$$

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$$\begin{aligned}
&= \frac{P_5 + P_4 - P_2 + P_6}{A_{11}B_{11} + A_{11}B_{22} - A_{11}B_{22} - A_{22}B_{11} + A_{22}B_{22} - A_{22}B_{21} - A_{12}B_{22} + A_{12}B_{22} + A_{12}B_{21}} \\
&= \frac{A_{11}B_{11}}{A_{11}B_{11} + A_{11}B_{22} - A_{11}B_{22} - A_{22}B_{11} + A_{22}B_{22} - A_{22}B_{21} - A_{12}B_{22} + A_{12}B_{22} + A_{12}B_{21}} \\
&= C_{11} \quad C_{11} = P_5 + P_4 - P_2 + P_6
\end{aligned}$$

Strassen Algorithm (1969)

$$\begin{array}{lllll} S_1 = B_{12} - B_{22} & S_2 = A_{11} + A_{12} & S_3 = A_{21} + A_{22} & S_4 = B_{21} - B_{11} & S_5 = A_{11} + A_{22} \\ S_6 = B_{11} + B_{22} & S_7 = A_{12} - A_{22} & S_8 = B_{21} + B_{22} & S_9 = A_{11} - A_{21} & S_{10} = B_{11} + B_{12} \end{array}$$

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$$\begin{aligned} & P_5 + P_4 - P_2 + P_6 \\ = & \begin{array}{ccccccc} A_{11}B_{11} & + & A_{11}B_{22} & + & A_{22}B_{11} & + & A_{22}B_{22} \\ & & & & - & A_{22}B_{11} & \\ & & & & & & + A_{22}B_{21} \\ & & - & A_{11}B_{22} & & & - A_{12}B_{22} \\ & & & & - & A_{22}B_{22} & - A_{22}B_{21} & + & A_{12}B_{22} & + & A_{12}B_{21} \end{array} \\ = & \frac{A_{11}B_{11}}{+ A_{12}B_{21}} \\ = & C_{11} \end{aligned}$$

$$C_{11} = P_5 + P_4 - P_2 + P_6$$

$$C_{12} = P_1 + P_2$$

$$C_{21} = P_3 + P_4$$

$$C_{22} = P_5 + P_1 - P_3 - P_7$$

Strassen Algorithm (1969)

$$\begin{array}{lllll} S_1 = B_{12} - B_{22} & S_2 = A_{11} + A_{12} & S_3 = A_{21} + A_{22} & S_4 = B_{21} - B_{11} & S_5 = A_{11} + A_{22} \\ S_6 = B_{11} + B_{22} & S_7 = A_{12} - A_{22} & S_8 = B_{21} + B_{22} & S_9 = A_{11} - A_{21} & S_{10} = B_{11} + B_{12} \end{array}$$

$$P_1 = A_{11}S_1 = A_{11}B_{12} - A_{11}B_{22}$$

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$$P_5 = S_5S_6 = A_{11}B_{11} + A_{11}B_{22} + A_{22}B_{11} + A_{22}B_{22}$$

$$P_6 = S_7S_8 = A_{12}B_{21} + A_{12}B_{22} - A_{22}B_{21} - A_{22}B_{22}$$

$$P_7 = S_9S_{10} = A_{11}B_{11} + A_{11}B_{12} - A_{21}B_{11} - A_{21}B_{12}$$

$$\begin{aligned} & P_5 + P_4 - P_2 + P_6 \\ = & \begin{array}{ccccccc} A_{11}B_{11} & + & A_{11}B_{22} & + & A_{22}B_{11} & + & A_{22}B_{22} \\ & & & & - & A_{22}B_{11} & \\ & & & & & & + A_{22}B_{21} \\ & & - & A_{11}B_{22} & & & - A_{12}B_{22} \\ & & & & - & A_{22}B_{22} & - A_{22}B_{21} & + & A_{12}B_{22} & + & A_{12}B_{21} \end{array} \\ = & \frac{A_{11}B_{11}}{+ A_{12}B_{21}} \\ = & C_{11} \end{aligned}$$

$$C_{11} = P_5 + P_4 - P_2 + P_6$$

$$C_{12} = P_1 + P_2$$

$$C_{21} = P_3 + P_4$$

$$C_{22} = P_5 + P_1 - P_3 - P_7$$

$$T(n) = 7T(n/2) + 18(n/2)^2 = 7T(n/2) + (9/2)n^2$$

$$T(n) = \begin{cases} 1 & \text{if } n=1 \\ 7 T(n/2) + \frac{9}{2} n^2 & \text{if } n \geq 2 \end{cases}$$

Assume $n = 2^k$.

$$\begin{aligned} T(n) &= 7 T(n/2) + \frac{9}{2} n^2 \\ &= 7 \left(7 T(n/2^2) + \frac{9}{2} (n/2)^2 \right) + \frac{9}{2} n^2 \\ &= 7^2 T(n/2^2) + \frac{7}{4} \cdot \frac{9}{2} n^2 + \frac{9}{2} n^2 \\ &= 7^2 \left(7 T(n/2^3) + \frac{9}{2} (n/2^2)^2 \right) + \frac{7}{4} \cdot \frac{9}{2} n^2 + \frac{9}{2} n^2 \\ &= 7^3 T(n/2^3) + \left(\frac{7}{4} \right)^2 \cdot \frac{9}{2} n^2 + \frac{7}{4} \cdot \frac{9}{2} n^2 + \frac{9}{2} n^2 \\ &= 7^3 \left(7 T(n/2^4) + \frac{9}{2} (n/2^3)^2 \right) + \left(\frac{7}{4} \right)^2 \frac{9}{2} n^2 + \frac{7}{4} \cdot \frac{9}{2} n^2 + \frac{9}{2} n^2 \\ &= 7^4 T(n/2^4) + \left(\frac{7}{4} \right)^3 \frac{9}{2} n^2 + \left(\frac{7}{4} \right)^2 \frac{9}{2} n^2 + \frac{7}{4} \cdot \frac{9}{2} n^2 + \frac{9}{2} n^2 \\ &= 7^4 T(n/2^k) + \frac{9}{2} \left(\left(\frac{7}{4} \right)^0 + \left(\frac{7}{4} \right)^1 + \left(\frac{7}{4} \right)^2 + \left(\frac{7}{4} \right)^3 \right) n^2 \\ &\vdots \\ &= 7^k \left(T(n/2^k) \right) + \frac{9}{2} \left(\left(\frac{7}{4} \right)^0 + \left(\frac{7}{4} \right)^1 + \dots + \left(\frac{7}{4} \right)^{k-1} \right) n^2 \end{aligned}$$

$T(1) = 1$

Recall:

$$\sum_{i=0}^{k-1} x^i = 1 + x + x^2 + \dots + x^{k-1} = \frac{x^k - 1}{x - 1} \quad (x \neq 1)$$

$$= 7^k + \frac{9}{2} \cdot \frac{\left(\frac{7}{4}\right)^k - 1}{\frac{7}{4} - 1} n^2$$

$$= 7^k + 6 \left(\left(\frac{7}{4}\right)^k - 1 \right) n^2$$

$$= \left(2^{\log_2(7)}\right)^k + 6 \left(\left(2^{\log_2(7/4)}\right)^k - 1 \right) n^2$$

$$= \left(2^k\right)^{\log_2(7)} + 6 \left(\left(2^k\right)^{\log_2(7/4)} - 1 \right) n^2$$

$$= n^{\log_2(7)} + 6 \left(n^{\log_2(7/4)} - 1 \right) n^2$$

$$= 7 n^{\log_2(7)} - 6 n^2$$

$$= O(n^{\log_2(7)}),$$

where $\log_2(7) \approx 2.81$

Strassen Algorithm (1969)

$$T(n) = 7 T(n/2) + (9/2)n^2$$

$$T(n) = O\left(n^{\log_2(7)}\right),$$

where $\log_2(7) \approx 2.81$.

Strassen Algorithm (1969)

$$T(n) = 7 T(n/2) + (9/2)n^2$$

$$T(n) = O\left(n^{\log_2(7)}\right),$$

where $\log_2(7) \approx 2.81$.

Coppersmith & Winograd (1990): $O(n^{2.38})$

Theorem (Master Theorem)

Let $T : \mathbb{N} \rightarrow \mathbb{R}^+$ be a function. Suppose that

$$T(n) = \begin{cases} 1 & \text{if } n = 1, \\ a \cdot T(n/b) + n^d & \text{if } n > 1, \end{cases}$$

where $a \geq 1$, $b > 1$ and $d \geq 0$.

- If $d > \log_b(a)$, then $T(n) = O(n^d)$.
- If $d = \log_b(a)$, then $T(n) = O(n^d \log(n))$.
- If $d < \log_b(a)$, then $T(n) = O(n^{\log_b(a)})$.

Merge sort:

$$T(n) = 2T(n/2) + n^1$$

Merge sort:

$$T(n) = 2T(n/2) + n^1$$

$$a = 2, b = 2, d = 1$$

Merge sort:

$$T(n) = 2T(n/2) + n^1$$

$$a = 2, b = 2, d = 1$$

$$d = \log_b(a)$$

Merge sort:

$$T(n) = 2T(n/2) + n^1$$

$$a = 2, b = 2, d = 1$$

$$d = \log_b(a)$$

$$T(n) = O(n^1 \log(n)) = O(n \log(n))$$

Merge sort:

$$T(n) = 2T(n/2) + n^1$$

$$a = 2, b = 2, d = 1$$

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Strassen:

$$T(n) = 7T(n/2) + n^2$$

Merge sort:

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$$T(n) = O(n^1 \log(n)) = O(n \log(n))$$

Strassen:

$$T(n) = 7T(n/2) + n^2$$

$$a = 7, b = 2, d = 2$$

Merge sort:

$$T(n) = 2T(n/2) + n^1$$

$$a = 2, b = 2, d = 1$$

$$d = \log_b(a)$$

$$T(n) = O(n^1 \log(n)) = O(n \log(n))$$

Strassen:

$$T(n) = 7T(n/2) + n^2$$

$$a = 7, b = 2, d = 2$$

$$d < \log_b(a)$$

Merge sort:

$$T(n) = 2T(n/2) + n^1$$

$$a = 2, b = 2, d = 1$$

$$d = \log_b(a)$$

$$T(n) = O(n^1 \log(n)) = O(n \log(n))$$

Strassen:

$$T(n) = 7T(n/2) + n^2$$

$$a = 7, b = 2, d = 2$$

$$d < \log_b(a)$$

$$T(n) = O\left(n^{\log_2(7)}\right)$$

Merge sort:

$$T(n) = 2T(n/2) + n^1$$

$$a = 2, b = 2, d = 1$$

$$d = \log_b(a)$$

$$T(n) = O(n^1 \log(n)) = O(n \log(n))$$

Strassen:

$$T(n) = 7T(n/2) + n^2$$

$$a = 7, b = 2, d = 2$$

$$d < \log_b(a)$$

$$T(n) = O\left(n^{\log_2(7)}\right)$$

Binary Search:

$$T(n) = T(n/2) + n^0$$

Merge sort:

$$T(n) = 2T(n/2) + n^1$$

$$a = 2, b = 2, d = 1$$

$$d = \log_b(a)$$

$$T(n) = O(n^1 \log(n)) = O(n \log(n))$$

Strassen:

$$T(n) = 7T(n/2) + n^2$$

$$a = 7, b = 2, d = 2$$

$$d < \log_b(a)$$

$$T(n) = O\left(n^{\log_2(7)}\right)$$

Binary Search:

$$T(n) = T(n/2) + n^0$$

$$a = 1, b = 2, d = 0$$

Merge sort:

$$T(n) = 2T(n/2) + n^1$$

$$a = 2, b = 2, d = 1$$

$$d = \log_b(a)$$

$$T(n) = O(n^1 \log(n)) = O(n \log(n))$$

Strassen:

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$$d < \log_b(a)$$

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Binary Search:

$$T(n) = T(n/2) + n^0$$

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$$T(n) = 2T(n/2) + n^1$$

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Strassen:

$$T(n) = 7T(n/2) + n^2$$

$$a = 7, b = 2, d = 2$$

$$d < \log_b(a)$$

$$T(n) = O\left(n^{\log_2(7)}\right)$$

Binary Search:

$$T(n) = T(n/2) + n^0$$

$$a = 1, b = 2, d = 0$$

$$d = \log_b(a)$$

$$T(n) = O(n^0 \log(n)) = O(\log(n))$$