Zadanie 3. Podaj rozmiar w bajtach poniższych struktur przyjmując, że wskaźnik jest 64-bitowy (architektura x86–64). Pod jakim przesunięciem, względem początku struktury, znajdują się poszczególne pola? Jak zreorganizować pola struktury, by zajmowała mniej miejsca? Z czego wynika takie zachowanie kompilatora?

```
1 struct A {
2   int8_t a;
3   void *b;
4   int8_t c;
5   int16_t d;
6 }:
1 struct B {
2   uint16_t a;
4   double b;
4   void *c;
5   int16_t d;
6 }:
```

Wskazówka: Użyj kompilatora, aby się dowiedzieć jaki jest rozmiar powyższych struktur - przyda się słowo kluczowe «sizeof».

Data structure padding [edit]

Although the compiler (or interpreter) normally allocates individual data items on aligned boundaries, data structures often have members with different alignment requirements. To maintain proper alignment the translator normally inserts additional unnamed data members so that each member is properly aligned. In addition, the data structure as a whole may be padded with a final unnamed member. This allows each member of an array of structures to be properly aligned.

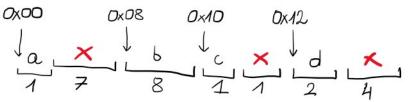
Catoix ma by will but of no sugar and proper alignment the translator normally inserts additional unnamed data members so that each member is properly aligned. In addition, the data structure as a whole may be padded with a final unnamed member. This allows each member of an array of structures to be properly aligned.

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Padding is only inserted when a structure member is followed by a member with a larger alignment requirement or at the end of the structure. By changing the ordering of members in a structure, it is possible to change the amount of padding required to maintain alignment.

| C Data Type | Typical 32-bit | Typical 64-bit | x86-64 |
|-------------|----------------|----------------|--------|
| char | 1 | 1 | 1 |
| short | 2 | 2 | 2 |
| int | 4 | 4 | 4 |
| long | 4 | 8 | 8 |
| float | 4 | 4 | 4 |
| double | 8 | 8 | 8 |
| pointer | 4 | 8 | 8 |

| member | offset | |
|--------|--------------|--|
| a | <i>0</i> ×00 | |
| Ь | 0×08 | |
| C | 0x 10 | |
| d | 0×12 | |



lepszy będzie taki

```
b d a c ×
1 struct A {
   void *b;
2
   int16_t d;
3
                       sizeof(A)=16
4
   int8_t a;
   int8_t c;
5
6 };
  struct B {
                                  offset
                          member
    uint16_t a; 2
2
                                   0×00
    double b; 8
                                   0x08
   void *c; 8
5 };
                                sizenf(B)=24
    2 6 8 B
          nie da się poprawio
      2+8+8=18 = zawsze co najmniej 24, 60 ma byco wielohrotnościa 8
```

The CPU accesses memory by a single <u>memory word</u> at a time. As long as the memory word size is at least as large as the largest <u>primitive data type</u> supported by the computer, aligned accesses will always access a single memory word. This may not be true for misaligned data accesses.

If the highest and lowest bytes in a datum are not within the same memory word the computer must split the datum access into multiple memory accesses. This requires a lot of complex circuitry to generate the memory accesses and coordinate them.