# Matrix Equation Solvers

Alex Harvey - mm<br/>13ah - ID: 200786528

# Contents

1	Inti	oduction	3
	1.1	Problem Statement	3
	1.2	Sylvester Equation	4
	1.3	Notation and Preliminaries	5
	1.4	Organisation	5
<b>2</b>	Sco	pe and Schedule	6
	2.1	Aim	6
	2.2	Objectives	6
	2.3	Deliverables	6
	2.4	Methodology	7
	2.5	Tasks, Milestones and Timeline	7
3	Ma	rix Solvers	8
	3.1	Direct Methods	0
		3.1.1 Kronecker Product	0
			1
		3.1.3 Similarity Transformation	13
			6
	3.2		17
			17
			18
			18
		* 0	20
4	Uno	ertainty 2	22
	4.1		22
		4.1.1 Monte Carlo	24
		4.1.2 Stochastic Galerkin	25
	4.2		29
		-	29
			29
5	Apj	lication 3	0
6	Eva	uation 3	31
	6.1	Aims	31
	6.2	Extensions	31
	6.3		21

# 1 Introduction

#### 1.1 Problem Statement

The traditional approach to solving partial differential equations (PDEs) numerically involves stacking all the unknowns of the problem into a single vector which ignores underlying structures. This prevents methods from being used which can take advantage of the problem structure to solve the problem more efficiently. This can come at a significant cost when uncertainty is introduced into the problem. An alternative approach is to formulate the problem as a matrix equation, which can be solved using a range of different methods. This project involves exploring methods that solve this alternative formulation and how these solvers compare against each other.

As an example, define spatial domain  $\mathcal{D} = \{(x,y) : 0 \le x \le 1, \ 0 \le y \le 1\}$  and let  $u : \mathcal{D} \to \mathbb{R}$  be the solution of the equation:

$$-u_{xx} - u_{yy} = f$$
 on  $\mathcal{D}$   
 $u = 0$  on  $\partial \mathcal{D}$  (1.1)

as shown in Figure 1.

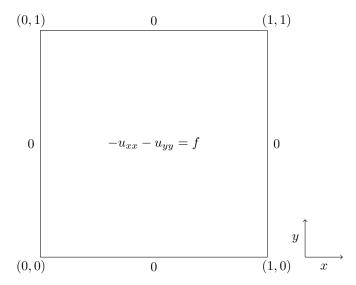


Figure 1: Domain for  $-u_{xx} - u_{yy} = f$ .

The domain of this PDE can be discretised into a mesh with uniform spacing h using the centred finite difference approximations:

$$u_{xx} \approx \frac{u_{i-1j} - 2u_{ij} + u_{i+1j}}{h^2} \tag{1.2}$$

$$u_{yy} \approx \frac{u_{ij-1} - 2u_{ij} + u_{ij+1}}{h^2} \tag{1.3}$$

where  $u_{ij} = u(x_i, y_j)$ . The mesh is shown in Figure 2.

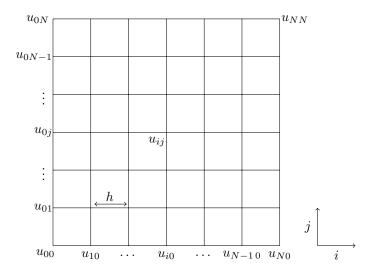


Figure 2: Discretised domain for  $-u_{xx} - u_{yy} = f$ .

The discretised form of this PDE can then be solved by computing the equation:

$$f_{ij} = -\frac{1}{h^2} \left( u_{i-1j} - 2u_{ij} + u_{i+1j} \right) - \frac{1}{h^2} \left( u_{ij-1} - 2u_{ij} + u_{ij+1} \right)$$
 (1.4)

at each internal grid point, meaning the system has  $n^2$  unknowns with n = N - 2.

The traditional approach to solving this discretised form would be to write (1.4) as:

$$f_{ij} = -\frac{1}{h^2} \left( u_{i-1j} + u_{ij-1} - 4u_{ij} + u_{i+1j} + u_{ij+1} \right)$$
 (1.5)

and then stack all unknowns  $u_{ij}$  into a single vector U, resulting in the linear system AU = F. As stated previously, this ignores the underlying structure of the problem.

## 1.2 Sylvester Equation

A Sylvester equation is a matrix equation of the form AX + XB = C, where A is a  $n \times n$  matrix, B is a  $m \times m$  matrix, and X and C are  $n \times m$  matrices. We can write (1.4) as a Sylvester equation in the form:

$$TU + UT = F (1.6)$$

where T, U and F are of size  $n \times n$ .

Define tridiag(i, j, k) as a matrix with j on the main diagonal and i and k on the left and right diagonals respectively. Then for (1.6) we have  $T = -\frac{1}{h^2} \operatorname{tridiag}(1, -2, 1)$  and  $U_{ij} = u(x_i, y_j)$ , where  $(x_i, y_j)$  are interior grid nodes for  $i, j = 1, \ldots, n$ . The system has n unknowns in each direction meaning there is a total of  $n^2$  unknowns. The matrix equation is visualised below:

$$\frac{-1}{h^2} \begin{pmatrix}
-2 & 1 & 0 & \dots & 0 & 0 \\
1 & -2 & 1 & \dots & 0 & 0 \\
0 & 1 & -2 & \dots & 0 & 0 \\
\vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\
0 & 0 & 0 & \dots & -2 & 1 \\
0 & 0 & 0 & \dots & 1 & -2
\end{pmatrix}
\begin{pmatrix}
u_{11} & u_{12} & \dots & u_{1j} & \dots & u_{1n} \\
u_{21} & u_{22} & \dots & u_{2j} & \dots & u_{2n} \\
\vdots & \vdots & \ddots & \vdots & \dots & \vdots \\
u_{n1} & u_{n2} & \dots & u_{nj} & \dots & u_{nn}
\end{pmatrix}$$

$$+ \begin{pmatrix}
u_{11} & u_{12} & \dots & u_{1j} & \dots & u_{1n} \\
u_{21} & u_{22} & \dots & u_{2j} & \dots & u_{2n} \\
\vdots & \vdots & \ddots & \vdots & \dots & \vdots \\
u_{n1} & u_{n2} & \dots & 0 & 0 \\
0 & 1 & -2 & \dots & 0 & 0 \\
0 & 1 & -2 & \dots & 0 & 0 \\
\vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\
0 & 0 & 0 & \dots & -2 & 1 \\
0 & 0 & 0 & \dots & -2 & 1 \\
0 & 0 & 0 & \dots & 1 & -2
\end{pmatrix} = F$$

There are a variety of methods that can be used to solve Sylvester equations. This project will explore and compare different methods for solving equations in this form.

## 1.3 Notation and Preliminaries

## 1.4 Organisation

# 2 Scope and Schedule

## 2.1 Aim

The aim of this project is to first study, implement and compare a range of matrix equation solvers. Following this, a specific problem will be derived with the help of my supervisor so that these solvers may be used and compared for a suitable application.

## 2.2 Objectives

The objectives of this project are as follows:

- To carry out an extensive, in-depth literature review on methods (both direct and iterative) for solving matrix equations from a wide range of sources. To decide which of these methods are appropriate to implement and to gain a solid understanding of how they work.
- To use and expand upon my programming experience to implement the chosen methods for solving matrix equations to solve the specified problem.
- To evaluate the implementation by comparing and contrasting the methods implemented to try to decide which is the best method for solving the given problem.
- To derive a suitable application equation so that the methods studied in this project can be applied to a specific problem.
- To clearly present the work carried out during the project by using and building upon my report writing skills.

### 2.3 Deliverables

The deliverables of this project include:

- The final report that will include the details of the matrix solvers that have been studied, how the solvers were implemented, an evaluation and comparison of the implemented solvers, an analysis of how these solvers were used to solve the chosen application problem, and finally an evaluation of the success of the project.
- Code that successfully implements the chosen matrix solvers so that they solve the given problem.

## 2.4 Methodology

The methodology of this project will first involve studying academic publishings to gain an understanding of various methods for solving matrix equations. Python will be used as the programming language of choice for the implementation because of my familiarity with it, the extensive amount of documentation available for it and the excellent libraries it has available (e.g. NumPy and SciPy). GitHub will be used for version control and the final report will be written using LATEX.

## 2.5 Tasks, Milestones and Timeline

The steps of this project will be divided into iterations, with the problem in each iteration becoming successively more complex and difficult to solve. This is because understanding is a key part of this project, and so each iteration will build on the understanding of the last. Each iteration will consist of studying and applying matrix methods to the problem, implementing them in Python to solve the problem, evaluating the results and write up. Also rough deadlines will be given for when each iteration should be completed by, to ensure the project is on track at any given stage.

The iterations are as follows:

- Introductory problem:  $-u_{xx} u_{yy} = 2\pi^2 \sin(\pi x) \sin(\pi y)$  deadline June 1st
- Problem introducing uncertainty:  $-\varepsilon u_{xx} \varepsilon u_{yy} = 2\pi^2 \sin(\pi x) \sin(\pi y)$  deadline June 22nd
- A Poisson equation on a surface defined by a height map (not yet derived) deadline July 13th

If the project deadlines are met the remaining time will be dedicated to project evaluation, write up and any possible project extensions.

## 3 Matrix Solvers

For spatial domain  $\mathcal{D} = \{(x,y) : 0 \le x \le 1, \ 0 \le y \le 1\}$  and  $u : \mathcal{D} \to \mathbb{R}$ , let  $u(x,y) = \sin(\pi x)\sin(\pi y)$  be the exact solution of the equation:

$$-u_{xx} - u_{yy} = f$$
 on  $\mathcal{D}$   
 $u = 0$  on  $\partial \mathcal{D}$  (3.1)

This gives:

$$u_{xx} = u_{yy} = -\pi^2 \sin(\pi x) \sin(\pi y)$$

$$\implies f = 2\pi^2 \sin(\pi x) \sin(\pi y)$$
(3.2)

The equation to be be solved is therefore:

$$-u_{xx} - u_{yy} = 2\pi^2 \sin(\pi x) \sin(\pi y) \quad \text{on } \mathcal{D}$$

$$u = 0 \quad \text{on } \partial \mathcal{D}$$
(3.3)

Using the centred finite difference approximations ((1.2)) and (1.3) with uniform grid spacing h to discretise the domain of this system, we have:

$$-\frac{1}{h^2}(u_{i-1j}-2u_{ij}+u_{i+1j})-\frac{1}{h^2}(u_{ij-1}-2u_{ij}+u_{ij+1})$$

$$=2\pi^2\sin(\pi x_i)\sin(\pi y_j)$$
(3.4)

to be solved at each grid point for  $i, j = 1, \dots n$ , where n is the total number of unknowns in each direction. The matrix form of this equation is therefore:

$$TU + UT = F (3.5)$$

where  $T = -\frac{1}{h^2}$  tridiag(1, -2, 1),  $U_{ij} = u(x_i, y_j)$  and  $F_{ij} = 2\pi^2 \sin(\pi x_i) \sin(\pi y_j)$  are all matrices of size  $n \times n$  and  $(x_i, y_j)$  are interior grid nodes for  $i, j = 1, \ldots, n$ .

This form allows us to explore different methods for solving this equation and compare them to the exact solution  $u = \sin(\pi x)\sin(\pi y)$ . A plot of the exact solution, with n = 1000, is given in Figure 3.

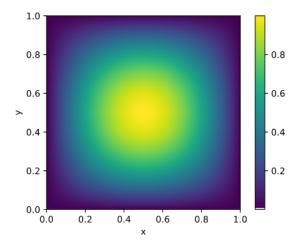


Figure 3: Plot of the solution with n = 1000.

Throughout the following section, the following measurements are given to evaluate the performance of each of the methods implemented:

- n: The number of unknowns in each direction for the system. The total number of unknowns is  $n^2$ .
- Time(s): The time taken in seconds for the method to compute the solution to the problem.
- $||u u_h||_{L^{\infty}}$ : Measures the maximum difference between the actual solution and computed solution for each u. Defined as:

$$||u - u_h||_{L^{\infty}} = \max_{i,j} |u(x_i, y_i) - u_{ij}|$$

•  $||u - u_h||_{L^2}$ : An average measure of error, defined as:

$$||u - u_h||_{L^2} = \sqrt{h^2 \sum |u(x_i, y_j) - u_{ij}|^2}$$

• Experimental order of convergence (eoc): Measures the rate of convergence of a method as the problem size is increased, which should approach 2 as the step size is increased. Defined as:

$$eoc(i) = \frac{\log(E_i/E_{i-1})}{\log(h_i/h_{i-1})}$$

where  $E_i$  is the error (chosen here as  $||u-u_h||_{L^{\infty}}$ ) and  $h_i$  is the mesh size at level i.

• Experimental order of growth (eog): Measures the order of growth of the execution time of an algorithm as the problem size is increased (that is, the approximate time complexity). Defined as:

$$eog(i) = \frac{\log(t_i/t_{i-1})}{\log(n_i/n_{i-1})}$$

where  $t_i$  is the total execution time and  $n_i$  is the problem size at level i.

• No. iters (iterative methods only): The number of iterations taken for the method to converge to the given convergence tolerance.

It is worth noting that as the total number of unknowns, and therefore the problem size, is  $n^2$ , the best time complexity that an optimal solver can achieve is  $O(n^2)$ , as it must compute a solution for each unknown.

For each of the implemented methods, a maximum running time of 6000 seconds will be given. If a method does not compute a solution in this time, this will be specified in the results. If a method does not compute a solution because it tries to use more memory than is available, this will also be specified in the results.

## 3.1 Direct Methods

#### 3.1.1 Kronecker Product

A naive approach to solving this system is to use the Kronecker product to rewrite (3.5) as a standard vector linear system. The Sylvester equation AX + XB = C can be written as the standard vector linear system:

$$\mathcal{A}x = c \tag{3.6}$$

with  $A = I \otimes A + B^* \otimes I$ , where I is the identity matrix,  $B^*$  denotes the conjugate transpose of B, x = vec(X) and c = vec(C).

The equivalent linear system for the matrix equation TU + UT = F is:

$$\mathcal{T}u = \mathcal{F} \tag{3.7}$$

where  $\mathcal{F} = \text{vec}(F)$ ,  $\mathcal{T} = I \otimes T + T \otimes I$  and u = vec(U).

This is the exact linear system that would be obtained from equation (1.5), i.e. stacking all unknowns  $u_{ij}$  into a single vector in the first place. Since the matrix  $\mathcal{T}$  is sparse, this equation can be solved using a standard direct sparse solver. This approach provides a good base case for comparison. Results solving this linear system using the direct sparse solver sparse.linalg.spsolve from the

 $<sup>^{1}\</sup>mathrm{The}$  vec operator reshapes a matrix into a vector by stacking the columns one after another.

SciPy library are shown in Table 1. As can be seen from the results, both errors decrease as the problem size n is increased. The eoc is very close to 2 for all n, demonstrating the convergence of the algorithm. As n is increased the eog grows beyond 3, which shows this algorithm has worse than cubic time complexity for large n.

n	Time(s)	$  u-u_h  _{L^{\infty}}$	$  u - u_h  _{L^2}$	eoc	eog
125	0.18141	$5.1807 \times 10^{-5}$	$2.5904 \times 10^{-5}$	-	-
250	0.60371	$1.3054 \times 10^{-5}$	$6.5275 \times 10^{-6}$	2.0001	1.7346
500	4.1663	$3.2767 \times 10^{-6}$	$1.6384 \times 10^{-6}$	1.9999	2.7868
1000	36.113	$8.2082 \times 10^{-7}$	$4.1041 \times 10^{-7}$	2.0000	3.1157
2000	404.48	$2.0539 \times 10^{-7}$	$1.0267 \times 10^{-7}$	2.0001	3.4855
4000		Me	mory Error		

Table 1: Results obtained from solving the linear system Ax = c using the direct solver sparse.linalg.spsolve from the SciPy library.

## 3.1.2 Bartels-Stewart Algorithm

The Bartels-Stewart algorithm [1] can be used to solve the Sylvester equation AX+XB=C. This algorithm works by computing the Schur decompositions of the coefficient matrices A and B, which are used to transform the equation into an equivalent form which is triangular, meaning it can be solved one element at a time. The solution of this triangular system is then used to obtain the solution to the original equation.

In the general case the algorithm is as follows:

- 1. Compute the Schur forms  $A^* = PRP^*$  and  $B = QSQ^*$
- 2. Solve  $R^*V + VS = P^*CQ$  for V
- 3. Compute  $X = PVQ^*$

where  $A^*$  denotes the conjugate transpose of A.

For the equation TU + UT = F, the algorithm is as follows:

- 1. Compute the Schur form  $T = PRP^*$
- 2. Solve  $R^*V + VR = P^*FP$  for V
- 3. Compute  $U = PVP^*$

Results using this algorithm are shown in Table 2. These results show that both errors decrease as n is increased. Also the eoc is close enough to 2 to conclude that this algorithm converges. The eog demonstrates that this algorithm has an approximate cubic time complexity.

n	Time(s)	$  u-u_h  _{L^{\infty}}$	$  u - u_h  _{L^2}$	eoc	eog
125	1.2486	$5.1807 \times 10^{-5}$	$2.5904 \times 10^{-5}$	-	-
250	9.6712	$1.3054 \times 10^{-5}$	$6.5275 \times 10^{-6}$	2.0001	2.9124
500	77.888	$3.2767 \times 10^{-6}$	$1.6384 \times 10^{-6}$	1.9999	3.0096
1000	640.22	$8.2088 \times 10^{-7}$	$4.1044 \times 10^{-7}$	2.0000	3.0391
2000	5093.1	$2.0467 \times 10^{-7}$	$1.0237 \times 10^{-7}$	2.0052	2.9919
4000		Tiı	ming Error		

Table 2: Results using the Bartels-Stewart algorithm.

By breaking this algorithm down into its component parts, we can time each step to see which is the most costly. The results of doing so are given in Table 3, where the steps are:

- 1. Compute the Schur form  $T = PRP^*$  (using linalg.schur from the SciPy library)
- 2. Solve the  $R^*V + VR = P^*FP$  for V, which is a triangular system
- 3. Compute the solution  $U = PVP^*$

As can be seen from the results in Table 3, the most costly part of the algorithm is by far the second step, which is due to the fact that solving the triangular system uses back substitution which requires multiple nested for loops.

n	1	2	3	Total
125	0.012667	1.2357	0.00023031	1.2486
250	0.048434	9.6218	0.0010102	9.6712
500	0.26262	77.619	0.0064600	77.888
1000	0.083478	639.34	0.043818	640.22
2000	5.1290	5087.5	0.41275	5093.1

Table 3: Timings for each step using the Bartels-Stewart algorithm.

The SciPy library has a built in solver for solving Sylvester equations, linalg.solve  $\_$ sylvester, which uses the Bartels-Stewart algorithm. Results using this solver are given in Table 4. The errors here are similar to the errors given in Table 2. Again the eoc is close enough to 2 to conclude that this algorithm converges. However, the eog seems to vary depending on the problem size n.

n	Time(s)	$  u-u_h  _{L^{\infty}}$	$  u-u_h  _{L^2}$	eoc	eog
125	0.059959	$5.1807 \times 10^{-5}$	$2.5904 \times 10^{-5}$	-	-
250	0.11421	$1.3054 \times 10^{-5}$	$6.5275 \times 10^{-6}$	2.0001	0.92964
500	1.9383	$3.2767 \times 10^{-6}$	$1.6384 \times 10^{-6}$	1.9999	4.0850
1000	4.0920	$8.2084 \times 10^{-7}$	$4.1042 \times 10^{-7}$	2.0000	1.0780
2000	32.029	$2.0503 \times 10^{-7}$	$1.0259 \times 10^{-7}$	2.0027	2.9685
4000	279.95	$4.9949 \times 10^{-8}$	$2.4876 \times 10^{-8}$	2.0377	3.1277
8000	2812.2	$8.6380 \times 10^{-9}$	$4.1216 \times 10^{-9}$		
16000		Ti	ming error		

Table 4: Results using SciPy's linalg.solve\_sylvester.

As can be seen from the results above, using the built-in SciPy solver results in a significant speed-up in time as n is increased. This is likely because it makes use of of LAPACK, which is an optimised software library for solving linear algebra problems.

#### 3.1.3 Similarity Transformation

A similarity transformation [5] can be used to solve the Sylvester equation AX + XB = C. This method uses an eigendecomposition to rewrite the equation into a form that the solution can be easily obtained from. The method is as follows.

Assuming matrices A and B can be diagonalised, let  $P^{-1}AP = \operatorname{diag}(\lambda_1, \ldots, \lambda_n)$  and  $Q^{-1}BQ = \operatorname{diag}(\mu_1, \ldots, \mu_m)$ , where  $\lambda_1, \ldots, \lambda_n$  are the eigenvalues of A,  $\mu_1, \ldots, \mu_m$  are the eigenvalues values of B and the columns of P and Q are the eigenvectors of A and B, respectively. Letting  $\widetilde{C} = P^{-1}CQ$ , the solution is then:

$$X = P\widetilde{X}Q^{-1}$$
, with  $\widetilde{X}_{ij} = \frac{\widetilde{C}_{ij}}{\lambda_i + \mu_j}$ 

For TU + UT = F, we have P = Q and  $P^{-1}TP = \operatorname{diag}(\lambda_1, \ldots, \lambda_n)$ , where  $\lambda_1, \ldots, \lambda_n$  are the eigenvalues of T and the columns of P are the eigenvectors of T. Letting  $\widetilde{F} = P^{-1}FP$ , the solution is then:

$$U = P\widetilde{U}P^{-1}$$
, with  $\widetilde{U}_{ij} = \frac{\widetilde{F}_{ij}}{\lambda_i + \lambda_j}$ 

## Using numpy.linalg.eig

Using the method above, the eigenvalues and eigenvectors can be computed using NumPy's linalg.eig function. The results doing this are given in Table 5. The errors are again similar to the previous methods. However, here the eoc

moves away from 2 as n is increased. This is likely due to the fact that there is no general formula for calculating eigenvalues and eigenvectors for an arbitrary matrix, meaning the eigenvalues and eigenvectors are most likely approximated by NumPy's linalg.eig. It is likely that the approximations become less accurate as n is increased. The eog here is however slightly better than the previous methods.

n	Time(s)	$  u-u_h  _{L^{\infty}}$	$  u - u_h  _{L^2}$	eoc	eog
125	0.034534	$5.1802 \times 10^{-5}$	$2.5904 \times 10^{-5}$	-	-
250	0.13690	$1.3054 \times 10^{-5}$	$6.5275 \times 10^{-6}$	2.0000	1.9870
500	0.59718	$3.2767 \times 10^{-6}$	$1.6384 \times 10^{-6}$	1.9999	2.1250
1000	2.2211	$8.2086 \times 10^{-7}$	$4.1043 \times 10^{-7}$	1.9999	1.8950
2000	11.571	$2.0467 \times 10^{-7}$	$1.0237 \times 10^{-7}$	2.0053	2.3812
4000	70.984	$4.9966 \times 10^{-8}$	$2.4876 \times 10^{-8}$	2.0350	2.6170
8000	466.83	$8.6360 \times 10^{-9}$	$4.1223 \times 10^{-9}$	2.5330	2.7173
16000		Mer	nory Error		

Table 5: Results using a similarity transformation, calculating the eigenvalues and eigenvectors using numpy.linalg.eig.

Similarly to the Bartels-Stewart algorithm, this method can be split into component parts and each part can be timed, to see which part is the most costly. The steps of the method are:

- 1. Calculate eigenvalues and eigenvectors of T
- 2. Diagonalise T (i.e. calculate P and  $P^{-1}$ )
- 3. Calculate  $\widetilde{F}=P^{-1}FP$
- 4. Calculate  $\widetilde{U}$ , where  $\widetilde{u}_{ij} = \frac{\widetilde{f}_{ij}}{\lambda_i + \lambda_i}$
- 5. Calculate solution  $U = P\widetilde{U}P^{-1}$

The results of doing so are shown in Table 6.

	n	1	2	3	4	5	Total
ſ	125	0.013860	0.00042105	0.00075817	0.018881	0.00061440	0.034534
	250	0.052402	0.00068998	0.0026329	0.078654	0.0025253	0.13690
	500	0.25900	0.0020170	0.013870	0.30815	0.014147	0.59718
	1000	0.91819	0.0011129	0.084650	1.1328	0.084323	2.2211
	2000	5.8236	0.0035102	0.61534	4.5150	0.61360	11.571
	4000	43.651	0.0074658	4.6290	18.071	4.6257	70.984
	8000	325.30	0.027416	35.464	70.695	35.345	466.83

Table 6: Timing results for each step of the similarity transformation method, calculating the eigenvalues and eigenvectors using numpy.linalg.eig.

## Calculating eigenvalues and eigenvectors explicitly

As can be seen from the results in Table 6, the most costly part of this method is calculating the eigenvalues and eigenvectors. As T is a matrix in Toeplitz form, we can instead compute the eigenvalues and eigenvectors directly as:

$$\lambda_i = \frac{2}{h^2} \left( \cos \left( \frac{i\pi}{n+1} \right) - 1 \right) \tag{3.8}$$

and:

$$t_{ij} = \sqrt{\frac{2}{n+1}} \sin\left(\frac{ij\pi}{n+1}\right) \tag{3.9}$$

Also the matrix P is Hermitian, which means that  $P^{-1} = P^T$ . We can use both of these properties to improve the timings of this algorithm. Results of doing so are given in Table 7, and timings for each step are given in Table 8.

n	Time(s)	$  u-u_h  _{L^{\infty}}$	$  u - u_h  _{L^2}$	eoc	eog
125	0.028324	$5.1807 \times 10^{-5}$	$2.5904 \times 10^{-5}$	-	-
250	0.080759	$1.3054 \times 10^{-5}$	$6.5275 \times 10^{-6}$		
500	0.22574	$3.2767 \times 10^{-6}$	$1.6384 \times 10^{-6}$		
1000	0.87985	$8.2082 \times 10^{-7}$	$4.1041 \times 10^{-7}$		
2000	4.0533	$2.0541 \times 10^{-7}$	$1.0270 \times 10^{-7}$		
4000	15.696	$5.1313 \times 10^{-8}$	$2.5656 \times 10^{-8}$		
8000	73.321	$1.2813 \times 10^{-8}$	$6.4065 \times 10^{-9}$		
16000	428.46	$7.7527 \times 10^{-10}$	$3.8764 \times 10^{-10}$		
32000		Memory	Error??		

Table 7: Results using a similarity transformation, calculating the eigenvalues and eigenvectors explicity.

n	1	2	3	4	5	Total
125	0.057252	0.00025916	0.0045545	0.018010	0.00059748	0.080673
250	0.21954	0.00047016	0.0081279	0.073295	0.0025220	0.30396
500	0.88125	0.0012608	0.018555	0.28724	0.013819	1.2021
1000	3.7073	0.00096869	0.67056	1.1304	0.081200	5.5904
2000	13.812	0.0092294	1.3690	4.4495	0.60793	20.247
4000	54.907	0.0068946	4.9050	17.381	4.5936	82.243
8000	217.26	0.026015	35.357	71.581	36.129	360.35
16000	892.51	168.51	217.68	288.92	204.25	1771.9

Table 8: Timing results for each step of the similarity transformation method, calculating the eigenvalues and eigenvectors explicitly.

As can be seen from the results in Table 8, calculating the eigenvalues and

eigenvectors explicitly outperforms calculating them using the NumPy library when n is large.

## 3.1.4 Shifted System

A projection method [5] can be used to solve AX + XB = C. This method works by computing the eigendecomposition of B to reform the problem as solving n independent linear systems. Each of these linear systems can be solved simultaneously to obtain the solution X. The steps of this method are as follows:

- 1. Compute  $B = WSW^{-1}$ , where  $S = \operatorname{diag}(s_1, \ldots, s_n)$  are the eigenvalues of B and the columns of W are the eigenvectors of B
- 2. Compute  $\hat{C} = CW$
- 3. For i = 1 to n, solve the system  $(A + s_i I)(\hat{X})_i = (\hat{C})_i$
- 4. Compute solution  $X = \hat{X}W$

where  $(\hat{X})_i$  denotes the  $i^{\text{th}}$  column of  $\hat{X}$ .

In the case of TU + UT = F, the steps are as follows:

- 1. Compute  $T = WSW^{-1}$ , where  $S = \text{diag}(s_1, \dots, s_n)$  are the eigenvalues of T and the columns of W are the eigenvectors of T
- 2. Compute  $\hat{F} = FW$
- 3. For i = 1 to n, solve the system  $(T + s_i I)(\hat{U})_i = (\hat{F})_i$
- 4. Compute solution  $U = \hat{U}W$

Results using this method are given in Table 9. The errors are approximately the same as the errors given previously. The eoc is approximately 2 and the eog is > 3 for large n demonstrating that this algorithm converges with a worse than cubic time complexity.

		,			
n	Time(s)	$  u-u_h  _{L^{\infty}}$	$  u-u_h  _{L^2}$	eoc	eog
125	0.092085	$5.1807 \times 10^{-5}$	$2.5904 \times 10^{-5}$	-	-
250	0.43609	$1.3054 \times 10^{-5}$	$6.5274 \times 10^{-6}$	2.0001	2.2436
500	2.1851	$3.2767 \times 10^{-6}$	$1.6384 \times 10^{-6}$	1.9999	2.3250
1000	19.498	$8.2082 \times 10^{-7}$	$4.1041 \times 10^{-7}$	2.0000	3.1576
2000	194.13	$2.0541 \times 10^{-7}$	$1.0271 \times 10^{-7}$	2.0000	3.3156
4000	2452.5	$5.1347 \times 10^{-8}$	$2.5674 \times 10^{-8}$	2.0009	3.6592
8000		Tir	ning error		

Table 9: Results using a projection method to solve n linear systems.

The steps of this method can also be broken down into component parts and each part timed. The results of doing so are given in Table 10. Unsuprisingly, the most costly part of the algorithm is solving the n linear systems.

n	1	2	3	4	Total
125	0.057421	0.0043032	0.3028	$7.1764 \times 10^{-7}$	0.092085
250	0.23967	0.0041912	0.19177	0.00046206	0.43609
500	0.85408	0.0079994	1.3200	0.0030336	2.1851
1000	3.3709	0.026276	16.079	0.021675	19.498
2000	13.278	0.50065	180.81	0.16621	194.13
4000	53.300	1.8301	2396.1	1.2866	2452.5

Table 10: Timing results of each step using a projection method to solve n linear systems.

#### 3.2 Iterative Methods

#### 3.2.1 Kronecker Product

Similarly to Section 3.1.1, the Kronecker product can be used to write the matrix equation as a standard vector linear system. A a standard iterative solver can then be used to solve the system, which can provide a base case for comparison. Results using scipy.sparse.linalg.cg, which is a sparse solver that uses the conjugate gradient iterative method, are given in Table 11, using a convergence tolerance of  $10^{-9}$ . The results show errors similar to all the direct methods. The eoc demonstrates that the algorithm converges very well. Although this algorithm is much faster than all the previous direct methods, the eog demonstrates that when n is significantly large, the algorithm exhibits a worse than cubic time complexity, which is shown here when n = 8000.

n	Time(s)	no. iters	$  u-u_h  _{L^{\infty}}$	$  u-u_h  _{L^2}$	eoc	eog			
125	0.014182	1	$5.1807 \times 10^{-5}$	$2.5904 \times 10^{-5}$	-	-			
250	0.042605	1	$1.3054 \times 10^{-5}$	$6.5275 \times 10^{-6}$	2.0001	1.5870			
500	0.14136	1	$3.2767 \times 10^{-6}$	$1.6384 \times 10^{-6}$	1.9999	1.7303			
1000	0.54308	1	$8.2082 \times 10^{-7}$	$4.1041 \times 10^{-7}$	2.0000	1.9418			
2000	2.3162	1	$2.0541 \times 10^{-7}$	$1.0271 \times 10^{-7}$	2.0000	2.0925			
4000	9.8487	1	$5.1379 \times 10^{-8}$	$2.5689 \times 10^{-8}$	2.0000	2.0882			
8000	103.34	2	$1.2848 \times 10^{-8}$	$6.4239 \times 10^{-9}$	2.0000	3.3913			
16000		Memory Error							

Table 11: Results obtained from solving the linear system Ax = c using the iterative solver sparse.linalg.cg from the SciPy library.

#### 3.2.2 Gradient Based Method

In [7] a gradient based method for solving Sylvester equations is given. The equation TU + UT = F can be written as two recursive sequences:

$$U_k^{(1)} = U_{k-1}^{(1)} + \kappa T(F - TU_{k-1}^{(1)} - U_{k-1}^{(1)}T)$$
(3.10)

$$U_k^{(2)} = U_{k-1}^{(2)} + \kappa (F - TU_{k-1}^{(2)} - U_{k-1}^{(2)} T)T$$
(3.11)

where  $\kappa$  represents the relative step size. The approximate solution  $U_k$  is taken as the average of these two sequences:

$$U_k = \frac{U_k^{(1)} + U_k^{(2)}}{2} \tag{3.12}$$

This solution only converges if:

$$0 < \kappa < \frac{1}{\lambda_{\max}(T^2)} \tag{3.13}$$

where  $\lambda_{\max}(T^2)$  denotes the maximum eigenvalue of  $T^2$ . We can use the formula (3.8) to calculate the maximum eigenvalue of  $T^2$  as:

$$\lambda_{\max}(T^2) = \frac{4}{h^4} \max\left(\left(\cos\left(\frac{i\pi}{n+1}\right) - 1\right)^2\right) \tag{3.14}$$

 $\lambda_{\max}(T^2)$  therefore scales with  $\frac{1}{h^4}$  meaning its reciprocal scales with  $h^4$ . This implies that  $\kappa$  will need to be significantly small as n is increased for the solution to converge. Even for small values of n, this is impractical and therefore this method is not appropriate for solving this equation. This can be seen from the fact that, using this method with n=125 and a step size of  $\kappa=1/2\lambda_{\max}(T^2)$ , the method took longer than the maximum time of 6000 seconds allowed to compute a solution.

#### 3.2.3 Modified Conjugate Gradient

In [4] a modified conjugate gradient (MCG) algorithm is proposed, which is adapted for solving Sylvester Equations. The general algorithm for solving AX + XB = C, where A, B, C and X are all  $n \times n$  matrices, is as follows:

- 1. Choose initial guess  $X^{(0)}$  and calculate:
  - $R^{(0)} = C AX^{(0)} X^{(0)}B$
  - $Q^{(0)} = A^T R^{(0)} + R^{(0)} + R^{(0)} B^T$
  - $Z^{(0)} = Q^{(0)}$
- 2. If  $R^{(k)} < \text{tol or } Z^{(k)} < \text{tol. stop. Else, go to 3.}$

3. Calculate:

•  $\gamma_k = \frac{[R^{(k)}, R^{(k)}]}{[Z^{(k)}, Z^{(k)}]}$ , where  $[R, R] = \text{tr}(R^T R)$  is the trace of the matrix  $R^T R$ 

• 
$$X^{(k+1)} = X^{(k)} + \gamma_k Z^{(k)}$$

• 
$$R^{(k+1)} = C - AX^{(k+1)} - X^{(k+1)}B$$

• 
$$Q^{(k+1)} = A^T R^{(k+1)} + R^{(k+1)} + R^{(k+1)} B^T$$

and return to Step 2.

In the case of the matrix equation TU + UT = F, the algorithm is as fol-

1. Choose initial guess  $U^{(0)}$  and calculate:

• 
$$R^{(0)} = F - TU^{(0)} - U^{(0)}T$$

$$Q^{(0)} = TR^{(0)} + R^{(0)} + R^{(0)}T$$

• 
$$Z^{(0)} = Q^{(0)}$$

2. If  $R^{(k)} < \text{tol or } Z^{(k)} < \text{tol, stop. Else, go to 3.}$ 

3. Calculate:

• 
$$\gamma_k = \frac{[R^{(k)}, R^{(k)}]}{[Z^{(k)}, Z^{(k)}]}$$

• 
$$U^{(k+1)} = U^{(k)} + \gamma_k Z^{(k)}$$

• 
$$R^{(k+1)} = F - TU^{(k+1)} - U^{(k+1)}T$$

• 
$$Q^{(k+1)} = TR^{(k+1)} + R^{(k+1)} + R^{(k+1)}T$$

and return to Step 2.

Results using this algorithm with a convergence tolerance of  $10^{-3}$  are given in Table 12. These results show that, even when using this relatively large convergence tolerance, this algorithm is slow. Also the errors does not always decrease as n is increased, which is reflected in the eoc. The errors here are also bigger than those for previous methods for large n, which is due to the large convergence tolerance. The eog demonstrates that this algorithm has a worse than cubic time complexity for large n.

n	Time(s)	$  u-u_h  _{L^{\infty}}$	$  u - u_h  _{L^2}$	eoc	eog
125	11.006	$1.1451 \times 10^{-6}$	$5.7254 \times 10^{-7}$	-	-
250	53.234	$3.7599 \times 10^{-5}$	$1.8800 \times 10^{-5}$	-5.0662	2.2741
500	397.11	$4.6673 \times 10^{-5}$	$2.3337 \times 10^{-5}$	-0.31294	2.8991
1000	3865.5	$3.3353 \times 10^{-5}$	$1.6677 \times 10^{-5}$	0.48547	3.2830

Table 12: Results obtained using the MCG algorithm.

#### 3.2.4 Preconditioned MCG

[4] also proposes a preconditioned version of the MCG algorithm which accelerates the speed of convergence. This version of the algorithm solves AX + XB = C by solving the equation  $X^{(k+1)} = \widetilde{A}X^{(k)}\widetilde{B} + \widetilde{C}$ , with:

$$\widetilde{A} = I + 2(\alpha A - I)^{-1}$$
 
$$\widetilde{B} = I + 2(\alpha B - I)^{-1}$$
 
$$\widetilde{C} = -2\alpha(\alpha A - I)^{-1}C(\alpha B - I)^{-1}$$

where  $\alpha$  is a parameter that needs to be chosen and I is the identity matrix. [4] suggests choosing  $\alpha = \pm \sqrt{n}/||A||_F$  or  $\alpha = \pm \sqrt{n}/||B||_F$ , where  $||A||_F = \sqrt{[A,A]} = \sqrt{\operatorname{tr}(A^TA)}$ , for good convergence. Once  $\alpha$  has been chosen and  $\widetilde{A}, \widetilde{B}$  and  $\widetilde{C}$  have been calculated, the algorithm is as follows:

- 1. Choose initial guess  $X^{(0)}$  and calculate:
  - $\bullet \ \ R^{(0)} = -\widetilde{C} + X^{(0)} \widetilde{A}X^{(0)}\widetilde{B}$
  - $Q^{(0)} = \widetilde{A}^T R^{(0)} \widetilde{B}^T + R^{(0)}$
  - $Z^{(0)} = Q^{(0)}$
- 2. If  $R^{(k)} < \text{tol or } Z^{(k)} < \text{tol, stop. Else, go to 3.}$
- 3. Calculate:
  - $\gamma_k = \frac{[R^{(k)}, R^{(k)}]}{[Z^{(k)}, Z^{(k)}]}$
  - $X^{(k+1)} = X^{(k)} + \gamma_k Z^{(k)}$
  - $\bullet \ R^{(k+1)} = -\widetilde{C} + X^{(k+1)} \widetilde{A}X^{(k+1)}\widetilde{B}$
  - $Q^{(k+1)} = \widetilde{A}^T R^{(k+1)} \widetilde{B}^T R^{(k+1)}$
  - $\eta_k = \frac{[R^{(k+1)}, R^{(k+1)}]}{[R^{(k)}, R^{(k)}]}$
  - $Z^{(k+1)} = Q^{(k+1)} + \eta_k Z^{(k)}$

and return to Step 2.

To solve the matrix equation TU+UT=F, we instead solve the equation  $U^{(k+1)}=\widetilde{T}U^{(k)}\widetilde{T}+\widetilde{F}$ , with:

$$\widetilde{T} = I + 2(\alpha T - I)^{-1}$$
 
$$\widetilde{F} = -2\alpha(\alpha A - I)^{-1}C(\alpha B - I)^{-1}$$

Choosing  $\alpha = \pm \sqrt{n}/||T||_F$  and calculating  $\widetilde{T}$  and  $\widetilde{F}$ , the algorithm proceeds as follows:

1. Choose initial guess  $U^{(0)}$  and calculate:

$$\bullet \ R^{(0)} = -\widetilde{F} + U^{(0)} - \widetilde{T}U^{(0)}\widetilde{T}$$

• 
$$Q^{(0)} = \widetilde{T}R^{(0)}\widetilde{T} + R^{(0)}$$

• 
$$Z^{(0)} = Q^{(0)}$$

- 2. If  $R^{(k)} < \text{tol or } Z^{(k)} < \text{tol, stop. Else, go to 3.}$
- 3. Calculate:

$$\bullet \ \gamma_k = \frac{[R^{(k)}, R^{(k)}]}{[Z^{(k)}, Z^{(k)}]}$$

• 
$$U^{(k+1)} = U^{(k)} + \gamma_k Z^{(k)}$$

$$\bullet \ \ R^{(k+1)} = -\widetilde{F} + U^{(k+1)} - \widetilde{T}U^{(k+1)}\widetilde{T}$$

• 
$$Q^{(k+1)} = \widetilde{T}R^{(k+1)}\widetilde{T} - R^{(k+1)}$$

• 
$$\eta_k = \frac{[R^{(k+1)}, R^{(k+1)}]}{[R^{(k)}, R^{(k)}]}$$

• 
$$Z^{(k+1)} = Q^{(k+1)} + \eta_k Z^{(k)}$$

and return to Step 2.

Results using the preconditioned MCG algorithm with a tolerance of  $10^{-9}$  are given in Table 13. These results demonstrate that this method vastly outperforms the standard MCG algorithm. The errors are similar to previous methods (excluding the MCG algorithm) and the eoc is close to 2 in most cases demonstrating that this algorithm converges well apart from in the case where n=8000, which could be because the error is approaching the convergence tolerance. The eog shows that although this algorithm is fast, for large n it has close to cubic time complexity.

n	Time(s)	$  u-u_h  _{L^{\infty}}$	$  u-u_h  _{L^2}$	eoc	eog
125	0.0010328	$5.1807 \times 10^{-5}$	$2.5904 \times 10^{-5}$	-	-
250	0.0040712	$1.3054 \times 10^{-5}$	$6.5275 \times 10^{-6}$	2.0001	1.9789
500	0.027955	$3.2767 \times 10^{-6}$	$1.6384 \times 10^{-6}$	1.9999	2.7796
1000	0.58508	$8.2078 \times 10^{-7}$	$4.1038 \times 10^{-7}$	2.0001	4.3875
2000	3.5886	$2.0585 \times 10^{-7}$	$1.0275 \times 10^{-7}$	1.9968	2.6167
4000	39.681	$5.3078 \times 10^{-8}$	$2.6012 \times 10^{-8}$	1.9561	3.4670
8000	336.49	$1.7720 \times 10^{-8}$	$6.6084 \times 10^{-9}$	1.5830	3.0840

Table 13: Results obtained using the preconditioned MCG algorithm.

[4] also proposes a parallel implementation of the preconditioned MCG algorithm, which would result in better performance than given above. However the parallel version of this algorithm goes beyond the scope of this project.

# 4 Uncertainty

Uncertainty can arise in PDEs if, for example, coefficients, boundary conditions or initial conditions are unknown. A solution to dealing with uncertainty is to model unknown parameters as random variables, and then appropriate methods can be used to solve the system.

## 4.1 Single Parameter

The simplest way to introduce uncertainty into a PDE is by introducing a single unknown parameter. Define the spatial domain  $\mathcal{D}$  as  $\mathcal{D} = \{(x,y) : 0 \le x \le 1, \ 0 \le y \le 1\}$ , and let  $\Omega$  denote the sample space. We can introduce an unknown coefficient  $\varepsilon(\omega)$ , where  $\omega \in \Omega$  denotes dependence on a random variable and  $\varepsilon : \Omega \to \mathbb{R}$  is uniformly distributed over the interval [1,3]. Let  $u(x,y,\varepsilon) = \sin(\pi x)\sin(\pi y) + \varepsilon\sin(3\pi x)\sin(5\pi y)$  be the exact solution of the equation:

$$-\varepsilon u_{xx} - \varepsilon u_{yy} = f \quad \text{ on } \mathcal{D} \times \Omega$$
$$u = 0 \quad \text{ on } \partial \mathcal{D} \times \Omega$$
(4.1)

which gives:

$$u_{xx} = -\pi^2 \sin(\pi x) \sin(\pi y) - 9\pi^2 \varepsilon \sin(3\pi x) \sin(5\pi y)$$

$$u_{yy} = -\pi^2 \sin(\pi x) \sin(\pi y) - 25\pi^2 \varepsilon \sin(3\pi x) \sin(5\pi y)$$

$$\implies f = 2\pi^2 \varepsilon \sin(\pi x) \sin(\pi y) + 34\pi^2 \varepsilon^2 \sin(3\pi x) \sin(5\pi y)$$
(4.2)

The equation to be solved is therefore:

$$-\varepsilon u_{xx} - \varepsilon u_{yy} = 2\pi^2 \varepsilon \sin(\pi x) \sin(\pi y) + 34\pi^2 \varepsilon^2 \sin(3\pi x) \sin(5\pi y) \quad \text{on } \mathcal{D} \times \Omega$$

$$u = 0 \quad \text{on } \partial \mathcal{D} \times \Omega$$
(4.3)

By discretising the domain using the centred finite difference approximations, (4.3) can be formed as a matrix equation:

$$TU + UT = F (4.4)$$

with  $T = -\frac{\varepsilon}{h^2} \operatorname{tridiag}(1, -2, 1)$ ,  $U_{ij} = u(x_i, y_j)$  and  $F_{ij} = 2\pi^2 \varepsilon \sin(\pi x_i) \sin(\pi y_j) + 34\pi^2 \varepsilon^2 \sin(3\pi x_i) \sin(5\pi y_j)$ . This problem is significantly harder to solve accurately and efficiently than the problem in Section 3 due to the fact that  $\varepsilon$  is an unknown, random coefficient.

As  $\varepsilon$  is unknown, the exact solution to (4.3) cannot be computed and therefore solutions to (4.4) cannot be compared to the true solution. Instead, certain

quantities of interest can be computed using the probability density function  $\rho(\omega)$  of  $\varepsilon$ . Firstly, the expectation  $\mathbb{E}[u]$  can be computed as:

$$\mathbb{E}[u] = \int_{\Omega} u(x, y, \varepsilon) \rho(\omega) d\omega \tag{4.5}$$

The variance Var[u] can also be computed as:

$$Var[u] = \int_{\Omega} (u(x, y, \varepsilon) - \mathbb{E}[u])^{2} \rho(\omega) d\omega$$
 (4.6)

and the standard deviation as  $Std[u] = \sqrt{Var[u]}$ .

As  $\varepsilon$  is uniformly distributed over the interval [1, 3], it has a probability distribution function:

$$\rho(\omega) = \frac{1}{2} \tag{4.7}$$

The expectation is therefore:

$$\mathbb{E}[u] = \frac{1}{2} \int_{1}^{3} \sin(\pi x) \sin(\pi y) + \varepsilon \sin(3\pi x) \sin(5\pi y) d\omega$$

$$= \frac{1}{2} \left[ \varepsilon \sin(\pi x) \sin(\pi y) + \frac{\varepsilon^{2}}{2} \sin(3\pi x) \sin(5\pi y) \right]_{1}^{3}$$

$$= \sin(\pi x) \sin(\pi y) + 2 \sin(3\pi x) \sin(5\pi y)$$
(4.8)

and the variance is:

$$\operatorname{Var}[u] = \frac{1}{2} \int_{1}^{3} \left( \left( \sin(\pi x) \sin(\pi y) + \varepsilon \sin(3\pi x) \sin(5\pi y) \right) - \left( \sin(\pi x) \sin(\pi y) + 2 \sin(3\pi x) \sin(5\pi y) \right) \right)^{2} d\omega$$

$$= \frac{1}{2} \int_{1}^{3} (\varepsilon - 2)^{2} \sin^{2}(3\pi x) \sin^{2}(5\pi y) d\omega$$

$$= \frac{1}{2} \left[ \left( \frac{\varepsilon^{3}}{3} - 2\varepsilon^{2} + 4\varepsilon \right) \sin^{2}(3\pi x) \sin^{2}(5\pi y) \right]_{1}^{3}$$

$$= \frac{1}{3} \sin^{2}(3\pi x) \sin^{2}(5\pi y)$$

$$(4.9)$$

which gives the standard deviation as  $Std[u] = \frac{1}{\sqrt{3}}\sin(3\pi x)\sin(5\pi y)$ .

Methods can now be used which compute approximate solutions to these quantities of interest, which can then be compared to the known quantities above to give an estimation of the error.

#### 4.1.1 Monte Carlo

A simple way of solving (4.4) is to use the Monte Carlo method. The Monte Carlo method works by taking M random samples of  $\varepsilon$  over its range, and then solves (4.4) for each of these samples. This allows the most effective method from Section 3 to be used as a 'black box' strategy. Once solved, the mean solution can be calculated as:

$$\mu = \frac{1}{M} \sum_{i=1}^{M} u_h(x, y, \varepsilon)$$
 (4.10)

where M is the number of samples used in the Monte Carlo method and  $u_h$  is the solution obtained from the finite difference method using a mesh of size h. Similarly, the variance can also be computed as:

$$\sigma^2 = \frac{1}{M-1} \sum_{i=1}^{M} (u_h - \mu)^2 \tag{4.11}$$

and the standard deviation is  $\sigma$ . Any of these quantites can be calculated and compared to the known values as a measure of error for the problem. By choosing the mean we can measure the error as:

$$Error = \left| \left| \mathbb{E}[u] - \mu \right| \right| \tag{4.12}$$

which can be evaluated using the  $L^{\infty}$  and  $L^2$  error measures as used in the previous section. As noted in [2], for the error to converge to 0 the mesh size must be decreased as the number of samples is increased. This is because the error depends both on the mesh size and the number of samples, so by changing only one of these factors the method will be converging to the error of the factor that is not being changed, rather than to 0.

It is also worth noting that as the solution and therefore the error is dependent on a random variable, the error is then itself a random variable and therefore not guaranteed to decrease at each level, as the samples are increased. However, the law of large numbers tells us that by taking more and more samples the mean solution should approximate the expected value and hence the error should tend towards 0.

Results using the Monte Carlo method, with a single uncertain parameter, are given in Table 14. As can be seen by these results, the error here is much bigger than the PDE without uncertainty.

n	M	Time(s)	Error
125	10	0.15511	0.071286
250	40	2.2698	0.19918
500	160	39.878	0.074265
1000	640	641.37	0.0050958
2000	2560	11957	0.0045691

Table 14: Results using the Monte Carlo method with a single uncertain parameter.

#### 4.1.2 Stochastic Galerkin

The stochastic Galerkin method uses the properties of orthogonal polynomials to produce a functional approximation to the solution of a stochastic PDE (a PDE with some kind of random input). It does this by using a Galerkin projection to "discretise the random dimensions to allow computation" [3].

A set of polynomials  $\{\Psi_n(x), n \in \mathbb{N}\}$ , where  $\Psi_n(x)$  is a polynomial of exact degree n, is orthogonal if it satisfies the orthogonality condition:

$$\int_{\Omega} \Psi_n(x)\Psi_m(x)w(x)dx = h_n\delta_{nm}, \quad n, m \in \mathbb{N}$$
 (4.13)

where  $\Omega$  represents the support of  $\{\Psi_n\}$  (the subset of the domain containing elements not mapped to zero), w(x) is a weight function,  $h_n$  are non-zero constants and  $\delta_{nm}$  is the Kronecker delta ( $\delta_{nm}=1$  for n=m and 0 otherwise). This condition is used to define the inner product of two polynomial functions, f(x) and g(x), with respect to a weight function w(x), as:

$$\langle f, g \rangle = \int_{\Omega} f(x)g(x)w(x)dx$$
 (4.14)

As noted in [6], certain classes of orthogonal polynomials have the exact same weight function as the probability distributions of certain random variables. This property allows the solution u of a stochastic PDE to be approximated via a truncated polynomial chaos expansion using the distribution of the random variable, as:

$$u(x, y, \varepsilon) = \sum_{k=0}^{P} u_k(x, y) \Psi_k(\varepsilon)$$
 (4.15)

and can then be substituted into the PDE. Here  $P+1=\frac{(K+N)!}{K!N!}$ , where N is the number of random variables and K is a convergence parameter to be chosen. Any random quantities can be represented via a Karhunen-Loeve expansion as:

$$\alpha = \bar{\alpha} + \sum_{k=1}^{Q} \sqrt{\lambda_k} \phi_k \varepsilon_k \tag{4.16}$$

for a spatially varying random field  $\alpha$  with mean  $\bar{\alpha}$ , where  $\lambda_k$  and  $\phi_k$  are the eigenvalues and eigenfunctions of the covariance function  $C_{\alpha}$ .

A Galerkin projection is then performed on the PDE by multiplying it by  $\Psi_k$  for  $k=0,\ldots,P$  and taking the inner product, which gives a system of coupled deterministic differential equations, which can then be discretised and solved via, for example, the finite difference method. The resulting solution matrix will be of size  $n^2\times(P+1)$ , where n is the number of unknowns in each direction of the mesh and P is defined as above. The mean and variance are then computed as:

$$\mu = u_0 \tag{4.17}$$

$$\sigma^2 = \sum_{k=1}^{P} u_k^2 \mathbb{E}[\Psi_k^2]$$
 (4.18)

For equation (4.4),  $\varepsilon$  is uniformly distributed over the interval [1, 3]. We introduce the substitution  $\varepsilon = \varepsilon_0 + 2$  into this equation, with  $\varepsilon_0 \sim u(-1, 1)$ , so that the random variable in the equation has zero mean. We then use the Legendre polynomials, which have the same weighting function as the probability density function of the uniform distribution, to solve the equation. They are defined as the polynomials which solve the equation:

$$\frac{d}{d\varepsilon} \left[ (1 - \varepsilon^2) \frac{d\Psi_n(\varepsilon)}{d\varepsilon} \right] + n(n+1)\Psi_n(\varepsilon) = 0$$
(4.19)

The SciPy library has a built in function, special.legendre, for calculating Legendre polynomials, which will be used along with the numerical integration function integrate.quad, to evaluate the inner product.

From equation (4.3), we now have:

$$(\varepsilon_0 + 2)(-u_{xx} - u_{yy}) = 2\pi^2(\varepsilon_0 + 2)\sin(\pi x)\sin(\pi y) + 34\pi^2(\varepsilon_0 + 2)^2\sin(3\pi x)\sin(5\pi y)$$
(4.20)

Approximating the solution u using a polynomial chaos expansion and subtituting gives:

$$(\varepsilon_0 + 2) \sum_{i=0}^{P} (-(u_i)_{xx} - (u_i)_{yy}) \Psi_i = 2\pi^2 (\varepsilon_0 + 2) \sin(\pi x) \sin(\pi y)$$
$$+ 34\pi^2 (\varepsilon_0 + 2)^2 \sin(3\pi x) \sin(5\pi y) \quad (4.21)$$

Finally, we multiply both sides by  $\Psi_k$  and take the inner product, giving:

$$\sum_{i=0}^{P} (-(u_i)_{xx} - (u_i)_{yy}) \langle \Psi_k(\varepsilon_0 + 2)\Psi_i \rangle = 2\pi^2 \sin(\pi x) \sin(\pi y) \langle \Psi_k(\varepsilon_0 + 2) \rangle + 34\pi^2 \sin(3\pi x) \sin(5\pi y) \langle \Psi_k(\varepsilon_0 + 2)^2 \rangle$$
(4.22)

for k = 0, ..., P. The uncertainty has been removed from this PDE and we now have a system of coupled deterministic differential equations, which can be solved for u.

We can now use the finite difference method to discretise the system in space and represent the unknowns  $u_i$  as a single matrix<sup>2</sup>:

$$Lu_i = -(u_i)_{xx} - (u_i)_{yy} (4.23)$$

where L has dimension  $n^2 \times n^2$  (the total number of unknowns) and  $u_i$  is a vector of length  $n^2$ . We can also represent the inner product on the left-hand side as a  $(P+1) \times (P+1)$  matrix:

$$P_{ij} = \langle \Psi_i(\varepsilon_0 + 2)\Psi_j \rangle \tag{4.24}$$

Finally, we represent the right-hand side of the system as a  $(P+1) \times n^2$  matrix:

$$F_k = 2\pi^2 \sin(\pi x_i) \sin(\pi y_j) \langle \Psi_k(\varepsilon_0 + 2) \rangle + 34\pi^2 \sin(3\pi x_i) \sin(5\pi y_j) \langle \Psi_k(\varepsilon_0 + 2)^2 \rangle$$
(4.25)

where  $k=0,\ldots,P$  represent the random discretisation (the rows) and  $i,j=1,\ldots,n$  represent the spatial discretisation (the columns). By reshaping F into a vector we can represent the problem as a linear system:

$$Au = f (4.26)$$

where  $A = P \otimes L$ , f = vec(F) and u is a vector of length  $(P+1)n^2$ . This system can now be solved in a variety of different ways.

## Linear System

The simplest way to solve (4.26) is to solve it as a linear system using a sparse solver. Results using sparse.linalg.spsolve from the SciPy library are given in Table 15.

n	K	Time(s)	Error
125	1	0.20171	0.0021941
250	1	1.3178	0.00055258
500	1	10.284	0.00013872
1000	1	84.773	$3.4752 \times 10^{-5}$
$\geq 2000$	Memory Error		

Table 15: Results using the stochastic Galerkin method with a single uncertain parameter by solving the linear system using sparse.linalg.spsolve.

<sup>&</sup>lt;sup>2</sup>Note that here  $L = I \otimes T + T \otimes I$ , where  $T = -\frac{1}{h^2} \operatorname{tridiag}(1, -2, 1)$  and I is the identity matrix, because we are representing all unknowns in a single matrix, rather than using two matrices as before.

## **Matrix Equation**

An alternative approach to solving (4.26) is to rewrite it as a matrix equation. (4.26) can be written as the matrix equation:

$$LXP = F (4.27)$$

where u = vec(X) and L, P and F are defined as before. This can then be solved using a variation of the Bartels-Stewart algorithm. The basic idea is the same as the standard Bartels-Stewart algorithm. Firstly, the Schur decomposition of L and P is computed, so that the equation can be transformed into an equivalent upper-triangular system that can be solved element by element. The solution to the original equation is then easily obtained using the solution of the triangular system. The steps of the algorithm are as follows:

- 1. Compute the Schur forms  $L = U\hat{L}U^*$  and  $P = V\hat{P}V^*$ , and let  $\hat{F} = U^*FV$ .
- 2. Solve  $\hat{L}Y\hat{P} = \hat{F}$  for Y, where  $\hat{L}$  and  $\hat{P}$  are upper triangular.
- 3. Compute solution  $X = UYV^*$ .

The problem with this method is that the Schur forms in the first step are computed using the QR algorithm, which only works on dense matrices. Therefore a huge amount of memory is required to compute the Schur form of L (of size  $n^2 \times n^2$ ), which is unpractical even for a relatively small problem size. Therefore solving the stochastic PDE in this form is not appropriate.

## Sylvester Equation

We can rewrite (4.26) as an alternative matrix equation in the form:

$$2LXG_0 + LXG_1 = F (4.28)$$

where u = vec(X),  $(G_0)_{ij} = \langle \Psi_i \Psi_j \rangle$  and  $(G_1)_{ij} = \langle \Psi_i \varepsilon_0 \Psi_j \rangle$ . By left and right multiplying by  $L^{-1}$  and  $G_0^{-1}$  respectively, we obtain the equation in Sylvester form:

$$AX + XB = C (4.29)$$

where  $A = L^{-1}2L$ ,  $B = G_1G_0^{-1}$  and  $C = L^{-1}FG_0^{-1}$ . This can now be solved using any of the methods from Section 3, assuming that they can be adapted for sparse matrices. Results using the similarity transformation method, using sparse.linalg.eigs from the SciPy library to calculate the eigenpairs, are given in Table 15.

## 4.2 Multiple Parameters

Multiple unknown parameters can be introduced into the PDE, which makes it significantly harder to solve. These are again modelled as random variables. Let  $\varepsilon$  now be defined as:

$$\varepsilon = (\varepsilon_0 + 2) + \sum_{p,q=1}^{N} \varepsilon_{pq} \sin(p\pi x) \sin(q\pi y)$$
 (4.30)

with  $\varepsilon_0 \sim u(-1,1)$  and  $\varepsilon_{pq} \sim \left(-2^{-\sqrt{p^2+q^2}}, 2^{-\sqrt{p^2+q^2}}\right)$ . Then let  $u = \sin(\pi x)\sin(\pi y) + \varepsilon$  be the solution of the PDE:

$$-\varepsilon u_{xx} - \varepsilon u_{yy} = f \quad \text{on } \mathcal{D} \times \Omega$$
$$u = 0 \quad \text{on } \partial \mathcal{D} \times \Omega$$
(4.31)

which gives:

$$u_{xx} = -\pi^{2} \sin(\pi x) \sin(\pi y) - \sum_{p,q=1}^{N} p^{2} \pi^{2} \varepsilon_{pq} \sin(p\pi x) \sin(q\pi y)$$

$$u_{yy} = -\pi^{2} \sin(\pi x) \sin(\pi y) - \sum_{p,q=1}^{N} q^{2} \pi^{2} \varepsilon_{pq} \sin(p\pi x) \sin(q\pi y)$$

$$\implies f = \varepsilon \left( 2\pi^{2} \sin(\pi x) \sin(\pi y) + \pi^{2} \sum_{p,q=1}^{N} (p^{2} + q^{2}) \varepsilon_{pq} \sin(p\pi x) \sin(q\pi y) \right)$$

$$= \left( (\varepsilon_{0} + 2) + \sum_{p,q=1}^{N} \varepsilon_{pq} \sin(p\pi x) \sin(q\pi y) \right)$$

$$\times \left( 2\pi^{2} \sin(\pi x) \sin(\pi y) + \pi^{2} \sum_{p,q=1}^{N} (p^{2} + q^{2}) \varepsilon_{pq} \sin(p\pi x) \sin(q\pi y) \right)$$

$$(4.32)$$

## 4.2.1 Monte Carlo

#### 4.2.2 Stochastic Galerkin

# 5 Application

- 6 Evaluation
- 6.1 Aims
- 6.2 Extensions
- 6.3 Conclusion

# References

- [1] R. H. Bartels and G. W. Stewart. Solution of the Matrix Equation AX + XB = C. Commun. ACM, 15(9):820–826, Sept. 1972.
- J. Bishop and O. E. Strack. A statistical method for verifying mesh convergence in Monte Carlo simulations with application to fragmentation. 88:279

   306, 10 2011.
- [3] P. Constantine. A Primer on Stochastic Galerkin Methods. 03 2007.
- [4] J. Hou, Q. Lv, and M. Xiao. A Parallel Preconditioned Modified Conjugate Gradient Method for Large Sylvester Matrix Equation. 2014:1–7, 03 2014.
- [5] V. Simoncini. Computational Methods for Linear Matrix Equations. 58:377–441, 01 2016.
- [6] D. Xiu and G. Karniadakis. The Wiener-Askey Polynomial Chaos for Stochastic Differential Equations. SIAM Journal on Scientific Computing, 24(2):619–644, 2002.
- [7] J. Zhou, W. Ruirui, and Q. Niu. A Preconditioned Iteration Method for Solving Sylvester Equations. 2012, 07 2012.