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ELCE333

Laboratory Report Experiment No.1

MICROCONTROLLER ASSEMBLY PROGRAM DEVELOPMENT

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Summary

This report demonstrates with the use of figures and results the processes taken in this laboratory experiment to edit, assemble, simulate, and then execute an assembly program written using Freescale (Code Warrior IDE). This laboratory experiment includes three tasks. The first task focuses on the idea of creating an assembly code in order to add numbers. As a matter of fact the code was given and what was required is to know how to assemble and simulate the program by single stepping through it to record the content of the registers after each step. The second task focusses on modifying a program so it would allow the subtraction of numbers. Task three shows us how to open a file called the List File which contains many different beneficial details of the program. Task three in general requires the knowledge of techniques to find the memory contents for the developed program. Finally, the report is concluded through discussing the main achievements.

1. Introduction

After being introduced to FreeScale software and CodeWarrior, this experiment encapsulates the whole introduction idea by requiring us to create a file from start to end. CodeWarrior is a program that will be used to create projects and creating programs using assembly with simulating and debugging them. The registers are described and listed below:

Accumulator A: an 8 bit accumulator that is used as a source or destination for 8bit instructions.

Accumulator B: an 8 bit accumulator that is used as a source or destination for 8bit instructions.

Accumulator D: a 16 bit accumulator that is used as a source or destination for 16bit instructions. It is the combined accumulator A and B.

Register X: 16bit register used for addressing and arithmetic operations.

Register Y: 16bit register used for addressing and arithmetic operations.

Stack Pointer: a pointer that ensures the program remains in the RAM.

Program Counter: a counter that shows how much memory can be addressed.

Condition Code Register: 5bit register in which each bit has a significant cause. The bits are:

C: when a borrow or carry occurs

V: when a 2's compliment occurs

Z: when result is zero

N: when the most significant bit is set

H: when a carry or borrow occur in bit three of the result

In assembly language, there are many important operations which are either data transfer instructions or assembly arithmetic instructions. Examples of the instructions are clear, shift, decrement and increment. For instance, we dealt with several mnemonics that represented instructions and such are DEC, INC, CLR, ABA, SBA, NEG, MUL and IDIV. In our lab we are dealing with HCS12 microcontroller; the code written will be saved in the ROM or to more precisely in the Flash EPROM. The variables of the program and the stack will be located in the RAM.

Aim:

The aim of this experiment is to introduce and familiarize students with the process of editing, assembling, simulating, downloading and subsequently executing a simple assembly program.

Objectives:

After doing this experiment, students will be able to:

- Understand how to edit an assembly program on HCS12.
- Generate a binary file and a list file by assembling the program.
- Realize how to download the program binary file to the Dragon Plus Trainer board.
- Learn how to go through the program using single-step.
- Trace a program execution and check the system registers.
- Understand a given program and modify it as required using HCS12 Instruction Set Manual.
- Learn how to realize the content of the program list file.

2. Design and Results

TASK 1: Adding Numbers

Task one summarizes the idea of creating a program from a given code. The code add the contents of the accumulator A and accumulator B after assigning the desired value in accumulator A which is \$ 55 (note that \$ is for hexadecimal) and for accumulator B which is \$ 25. Then, the result stored into memory location \$1002. The second part, the contents of accumulator A and \$10 are added and the results are stored in accumulator A. After that \$06 is added to the contents of accumulator A and the final results is stored in memory location \$1003. The code is as the following.

```
        INCLUDE 'derivative.inc'
        XDEF Entry
SUMA     EQU   $1002      ; Location of SUMA
SUMB     EQU   $1003      ; Location of SUMB

Entry:   ORG   $4000      ; Set the program counter to$4000

        CLRA             ; Clear accumulator A
        CLRB             ; Clear accumulator B
        LDAA #$55         ; the contents of accumulator A is $55
        LDAB #$25         ; the contents of accumulator B is $25
        ABA              ; Add accumulator B to accumulator A
        STAA SUMA         ; Store the contents of A to memory location $1002
        ADDA #$10         ; Add A to the immediate value $10 and store result in A
        ADDA #$06         ; Add A to the immediate value $10 and store result in A
        STAA SUMB         ; Store the contents of A to memory location $1003

HERE     JMP   HERE
```

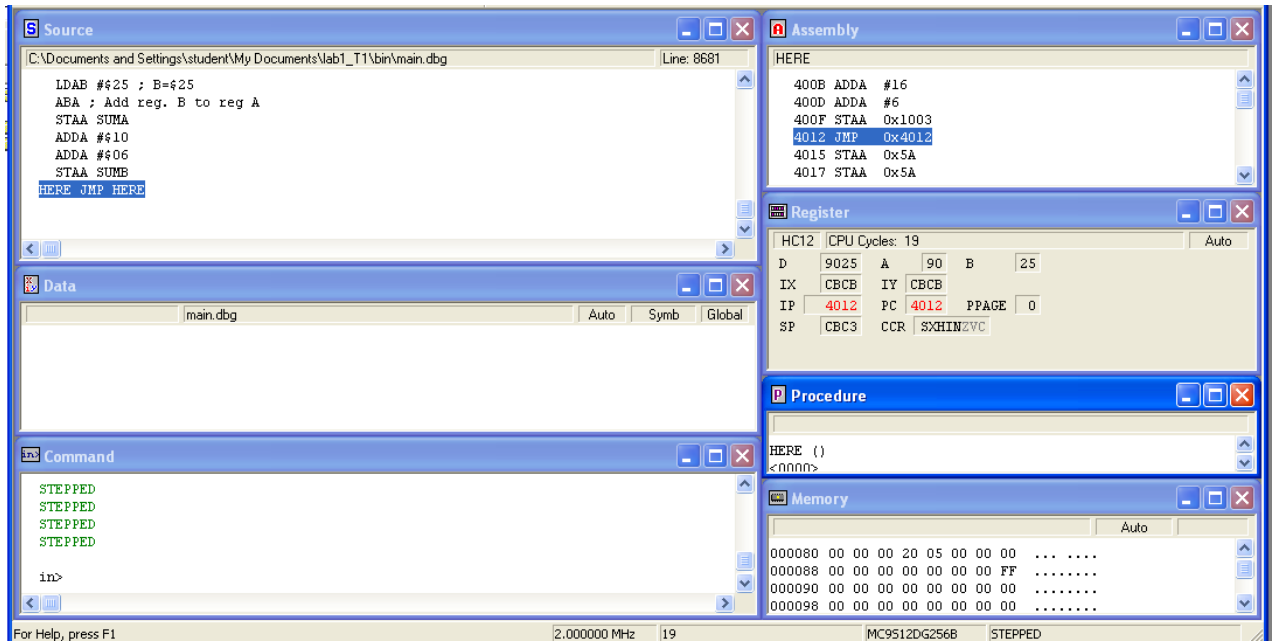


Figure 1: Registers at instruction STAA SUMB

We recorded the details required about each of the instructions shown in the code in Table 1.

Table 1: Registers contents while single stepping in the program

Instruction	A	B	D	Comments
CLRA	0	25	25	D= 0+25= 0025, Clear acc. A
CLRB	0	0	00	D=0000, Clear acc. B
LDAA #\$55	55	0	5500	D=55+0=5500, Load acc. A with the value \$55
LDAB #\$25	55	25	5525	D=55+25=5525, Load acc. B with the value \$25
ABA	7A	25	7A25	D=7A+25=7A25, Add A to B and save the result in A
STAA SUMA	7A	25	7A25	D=7A+25=7A25, Store the value of A at SUMA
ADDA #\$10	8A	25	8A25	D=8A+25=8A25, Add the value \$10 to A
ADDA #\$06	90	25	9025	D=90+25=9025, Add the value \$06 to A
STAA SUMB	90	25	9025	D=90+25=9025, Store the value of A at SUMB

Finally, after the program is finished we examined both memory locations \$1002 and \$1003 and recorded the values before starting and after finishing as shown below in the table and in Figure 2.

Table 2: Memory Locations Contents

	\$1002	\$1003	Comments
Before	--	--	The addresses were empty
After	7A	90	The content of A is shown at the address 1002 after adding A&B while the final content of A is shown in 1003, after doing the second set of addition operation

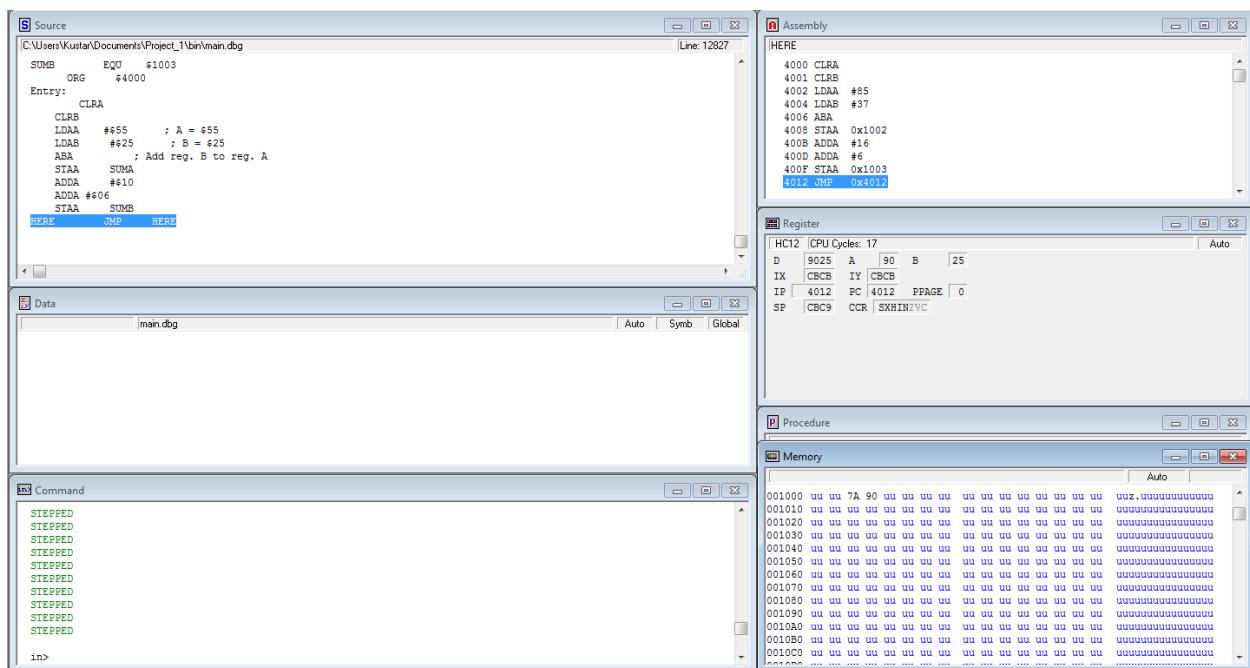
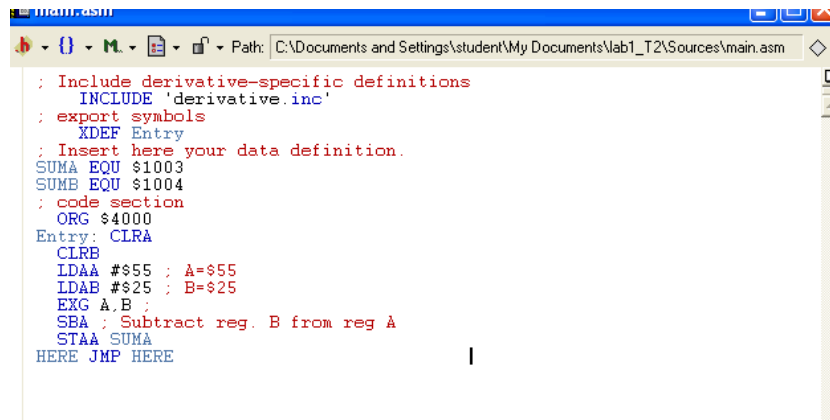


Figure 2: The memory contents for \$1002 and \$1003 memory address

Comment: The addresses \$1002 and \$1003 were assigned for the labels SUMA and SUMB. In 1002, the addition of \$ 55 (decimal value: 85) and \$ 25 (decimal value: 37) is 7A (decimal value: 122). Then, the result stored into memory location \$1002, as shown above. For the second set, the address \$1003, \$ 55+ \$ 25=\$ 80 and then this is added to \$ 10 which will result in \$ 90, as shown above.

TASK 2: Subtraction Numbers

Task two is about modifying the code given in task one –which adds numbers- to create a program that subtracts numbers instead. We used the SBA instruction to subtracted and final results were saved in \$1003. We did not use the memory location \$1004 because we exchanged the contents of A and B. Figure 3 below shows the modified code.



```
; Include derivative-specific definitions
INCLUDE 'derivative.inc'
; export symbols
XDEF Entry
; Insert here your data definition.
SUMA EQU $1003
SUMB EQU $1004
; code section
ORG $4000
Entry: CLRA
      CLRB
      LDAA #$55 ; A=$55
      LDAB #$25 ; B=$25
      EXG A,B ;
      SBA ; Subtract reg. B from reg A
      STAA SUMA
      HERE JMP HERE
```

Figure 3: Subtraction code

Figure 4 shows the results we obtained from this task. Using the single step facility, we were able to view the registers at each instruction and when tested, the output we got was D0 which gives a hexadecimal value in correspondence to the subtraction value. As seen in the memory window, the \$1000 row and fourth column shows the location of \$1003 which as previewed is D0.

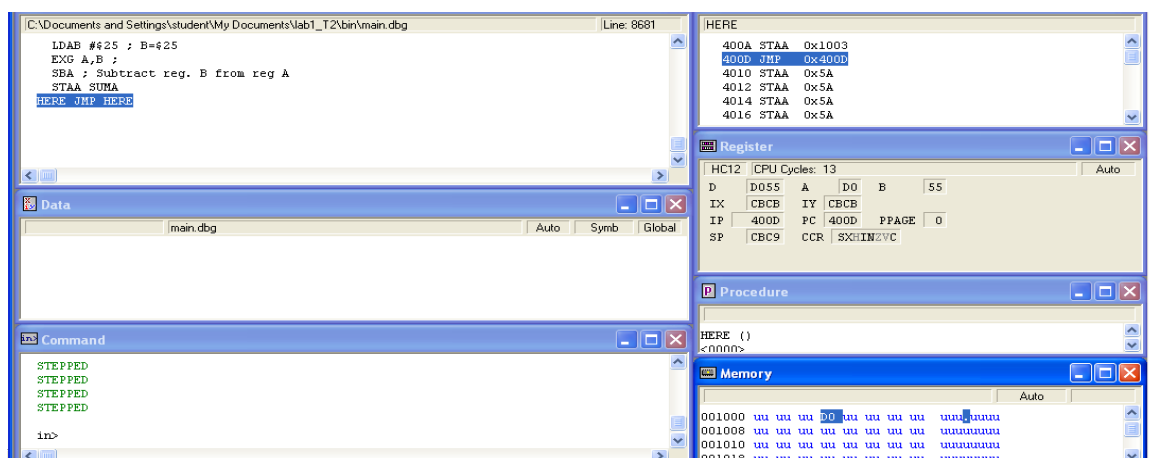


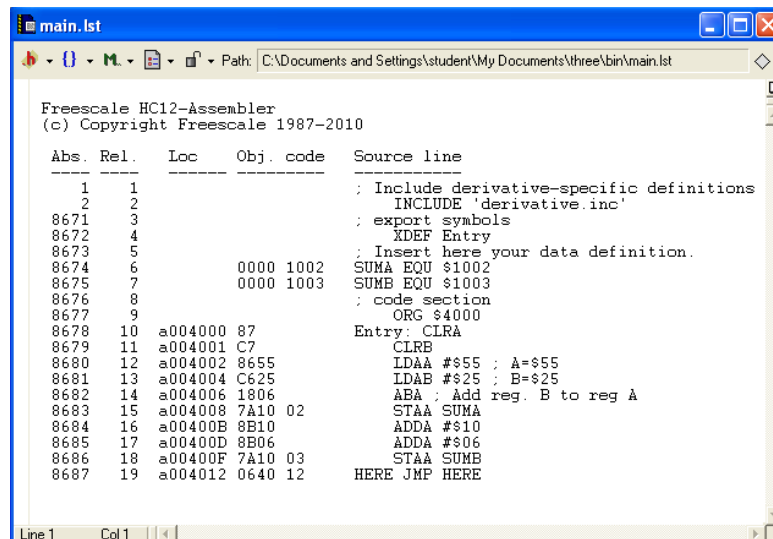
Figure 4: Result of subtraction program

Comment: The addresses \$1003 has the contents of D0 which is the result of the subtraction of accumulator A and accumulator B.

TASK 3: Simple Program

Task three concepts about dealing with list files with extension “.lst”. This file provides the following with the source-code program, The assembled code in hexadecimal format (including memory locations), any error messages and a list of user-defined symbols used and their values.

We were asked to copy the same given code in task 1 and use the “create a listing file” checkbox to create the file we need. It is saved in the “bin” file and opened from that source. When we opened the file, several columns appeared. A screenshot of the file was taken and shown in Figure 4 in the next page. We were required to note down several details about every instruction and the information are jotted in Table 3.



Abs.	Rel.	Loc.	Obj.	code	Source line
1	1				; Include derivative-specific definitions
2	2				INCLUDE 'derivative.inc'
8671	3				; export symbols
8672	4				XDEF Entry
8673	5				; Insert here your data definition.
8674	6		0000	1002	SUMA EQU \$1002
8675	7		0000	1003	SUMB EQU \$1003
8676	8				; code section
8677	9				ORG \$4000
8678	10	a004000	87		Entry: CLRA
8679	11	a004001	C7		CLRB
8680	12	a004002	8655		LDAA #\$55 ; A=\$55
8681	13	a004004	C625		LDAB #\$25 ; B=\$25
8682	14	a004006	1806		ABA ; Add reg. B to reg A
8683	15	a004008	7A10	02	STAA SUMA
8684	16	a00400B	8B10		ADDA #\$10
8685	17	a00400D	8B06		ADDA #\$06
8686	18	a00400F	7A10	03	STAA SUMB
8687	19	a004012	0640	12	HERE JMP HERE

Figure 5: list file

The table shows that the ORG instruction forces the program to start from address 4000. The difference between the start address of the required instruction and the next instruction shows the number of bytes.

Table 3: The memory contents for the developed program:

Instruction	Start Address	No. of Bytes	Instruction Code
CLRA	a004000	1	87
CLRB	a004001	1	C7
LDAA #\$55	a004002	2	8655
LDAB #\$25	a004004	2	C625
ABA	a004006	2	1806
STAA SUMA	a004008	3	7A1002
ADDA #\$10	a00400B	2	8B06
ADDA #\$06	a00400D	3	7A1003
STAA SUMB	a00400F	3	064012

Comment: The list file shows the operation that must be done and the operands that are being operated upon or each instruction in the assembly language. If we take the third instruction in the table which is LDAA #\$55, the first byte which is 86 is a unique code for each operation to be executed by the microcontroller; this is known as the opcode byte. The second byte is for the operands of the operation.

3. Assignment Questions

1. Find the status of CCR flags - C, V, H, N, Z and Acc A after executing A, B, C and D individually and comment on the results.

Table 4- Results of simulation for question 1

	Instructions	C	V	H	N	Z	Acc A	Acc B	Comment
A	LDAA #\$A5	-	0	0	1	0	A5		<ul style="list-style-type: none"> ➤ <i>Addressing mode:</i> Immediate ➤ <i>Operation:</i> (\$A5) → A ➤ <i>C:</i> Executing LDAA instructions doesn't change the carry bit ➤ <i>V:</i> 2's complement overflow is set to zero ➤ <i>H:</i> no half carry occurred ➤ <i>N:</i> Because the msb is 1 $(A5)_{16} = (10100101)_2$ ➤ <i>Z:</i> The value contents of acc. A isn't zero
	ADDA #\$89	1	1	0	0	0	12E		<ul style="list-style-type: none"> ➤ <i>Addressing mode:</i> Immediate ➤ <i>Operation:</i> (\$A5) + (\$89) → A ➤ <i>C:</i> the result of the instruction is $(12E)_{16} = (100101110)_2$ which is 9 bits so we have a carry ➤ <i>V:</i> the binary representation shows a 2's complement overflow ➤ <i>H:</i> no half carry occurred ➤ <i>N:</i> the msb (the 8th bit) is zero ➤ <i>Z:</i> The value contents of acc. A isn't zero
B	LDAA #\$80	-	0	0	1	0	80		<ul style="list-style-type: none"> ➤ <i>Addressing mode:</i> Immediate ➤ <i>Operation:</i> (\$80) → A ➤ <i>C:</i> Executing LDAA instructions doesn't change the carry bit

									<ul style="list-style-type: none"> ➤ <i>V</i>: 2's complement overflow is set to zero ➤ <i>H</i>: no half carry occurred ➤ <i>N</i>:msb is zero $(80)_{16}=(10000000)_2$ ➤ <i>Z</i>: The value contents of acc. A isn't zero
	ADDA #\$80	1	1	0	0	1	0		<ul style="list-style-type: none"> ➤ <i>Addressing mode</i>: Immediate ➤ <i>Operation</i>: $(\\$80) + (80) \rightarrow A$ ➤ <i>C</i>:the result of the instruction is $(100)_{16}=(100000000)_2$ which is 9 bits so we have a carry ➤ <i>V</i>: the binary representation shows a 2's complement overflow ➤ <i>H</i>: no half carry occurred ➤ <i>N</i>:the msb (the 8th bit) is zero ➤ <i>Z</i>: the bits of acc. A are zeros and the 9th overflow bit is stored in next memory address
C	LDAB #\$A5	-	0	0	1	0		A5	<ul style="list-style-type: none"> ➤ <i>Addressing mode</i>: Immediate ➤ <i>Operation</i>:$(\\$A5) \rightarrow B$ ➤ <i>C</i>:Executing LDAB instructions doesn't change the carry bit ➤ <i>V</i>: 2's complement overflow is set to zero ➤ <i>H</i>: no half carry occurred ➤ <i>N</i>:Because the msb is 1 $(A5)_{16}=(10100101)_2$ ➤ <i>Z</i>: The value contents of acc. A isn't zero
	SUBB #\$68	0	1	0	0	0		3D	<ul style="list-style-type: none"> ➤ <i>Addressing mode</i>: Immediate ➤ <i>Operation</i>: $(\\$A5) - (68) \rightarrow B$ ➤ <i>C</i>:no carry $(3D)_{16}=(00111101)_2$ ➤ <i>V</i>: 2's complement is used for this instruction ➤ <i>H</i>: no half carry occurred

									➤ <i>N</i> : the msb (the 8 th bit) is zero Z: the result isn't zero
D	LDAB #\$65	-	0	0	0	0		65	➤ <i>Addressing mode</i> : Immediate ➤ <i>Operation</i> : (\$65) → B ➤ <i>C</i> : Executing LDAB instructions doesn't change the carry bit ➤ <i>V</i> : 2's complement overflow is set to zero ➤ <i>H</i> : no half carry occurred ➤ <i>N</i> : Because the msb is 1 (65) ₁₆ = (01100101) ₂ ➤ <i>Z</i> : The value contents of acc. A isn't zero
	SUBB #\$65	0	0	0	0	1		0	➤ <i>Addressing mode</i> : Immediate ➤ <i>Operation</i> : (\$65) – (\$65) → B ➤ <i>C</i> : no carry occurred ➤ <i>V</i> : no 2's complement overflow occurred ➤ <i>H</i> : no half carry occurred ➤ <i>N</i> : msb is 0 (0) ₁₆ = (00000000) ₂ ➤ <i>Z</i> : The value contents of acc. A is zero

2. Write an instruction sequence to multiply two 8-bit numbers using MUL instruction (unsigned). Initialize Accumulator A and Accumulator B with \$F4 and \$F2 respectively using the load instructions. Highlight the result in the D register and the CCR flags.

The code:

```
; Include derivative-specific definitions
INCLUDE 'mc9s12dg256.inc'
; export symbols
XDEF Entry
; Insert here your data definition.

Entry:

CLRA
CLRB
LDAB    #$F4
LDAA    #$F2
MUL     |
```

Table 5- Results of simulation for question 2

Instruction	C	V	H	N	Z	Acc A	Acc B	Acc D
CLRA	0	0	0	0	1	0	CB	CB
CLRB	0	0	0	0	1	0	0	0
LDAB #\$F4	0	0	0	1	0	0	F4	F4
LDAA #\$F2	0	0	0	1	0	F2	F4	F2F4
MUL	1	0	0	1	0	E6	A8	E6A8

4. Conclusion

In the first task it was required to assemble and simulate a code for addition using ABA. Same program of task one was modified in the second task in order to do the subtraction operation and the results are stored in specific memory location.

The purpose of the last task was to learn some further details and features about the assembled program. Start Address, number of bytes, Instruction Code are essential tools in every program for debugging and finding the errors in the program. Actually, a computer instruction is the combination of an operation (what the computer is to do) and zero, or more operands (what the computer is going to do to it).