

1a.

Each account a has:

stack of amounts

$\text{count}[0..T]$

Insert(a, t):

push t onto $\text{account}[a].\text{stack}$

$\text{count}[a][t] = \text{count}[a][t] + 1$

Search(a, t):

if $\text{count}[a][t] > 0$ then

return true

else

return false

Most_recent(a):

if stack is empty then

return NULL

else

return top value of stack

Delete_most_recent(a):

if stack is empty then

return

$t = \text{pop from stack}$

$\text{count}[a][t] = \text{count}[a][t] - 1$

Insert(a, t):

When we add a transaction, we push t onto the stack and increment $\text{count}[a][t]$. Both steps take constant time.

Time complexity: $O(1)$

Search(a, t):

To check if an amount t exists, we simply test if $\text{count}[a][t] > 0$. Array access is constant time.

Time complexity: $O(1)$

Most_recent(a):

To find the most recent transaction, we return the top element of the stack. Accessing the top of a stack is constant time.

Time complexity: $O(1)$

Delete_most_recent(a):

To delete the most recent transaction, we pop the top of the stack and decrement count[a][t]. Both operations are constant time.

Time complexity: $O(1)$

1b.

Each account a has:

- stack of amounts

- balanced BST (amount \rightarrow frequency)

Insert(a, t):

- push t onto account[a].stack

- if t exists in BST then

 - freq[t] = freq[t] + 1

- else

 - insert t into BST with freq = 1

Search(a, t):

- if t exists in BST then

 - return true

- else

 - return false

Most_recent(a):

- if stack is empty then

 - return NULL

- else

 - return top value of stack

Delete_most_recent(a):

- if stack is empty then

 - return

- t = pop from stack

- freq[t] = freq[t] - 1

- if freq[t] == 0 then

 - remove t from BST

INSERT(a, t):

Push t onto the stack ($O(1)$), then search for t in the BST. If found, increment its frequency; otherwise, insert a new node. Searching and inserting in a BST both take $O(\log n)$.

Time complexity: $O(\log n)$

Search(a, t):

Look up t in the BST. Searching in a balanced BST takes $O(\log n)$.

Time complexity: $O(\log n)$

Most_recent(a):

Return the top element of the stack. Stack top lookup takes constant time.

Time complexity: $O(1)$

Delete_most_recent(a):

Pop the top of the stack ($O(1)$), then search for that amount t in the BST and decrement its frequency. If the frequency reaches zero, remove the node from the BST. Both searching and deleting in a BST take $O(\log n)$.

Time complexity: $O(\log n)$

2a.

Build_bst(arr, start, end):

if start > end then
return NULL

mid = (start + end) / 2
root = new Node(arr[mid])

root.left = Build_bst(arr, start, mid - 1)
root.right = Build_bst(arr, mid + 1, end)

return root

2b.

Merge_trees(T1, T2):

list1 = Inorder(T1)
list2 = Inorder(T2)

merged = Merge_sorted(list1, list2)

root = Build_bst(merged, 0, length(merged) - 1)

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return root
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Inorder(root):  
    if root == NULL:  
        return []  
    return Inorder(root.left) + [root.key] + Inorder(root.right)
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Merge_sorted(list1, list2):  
    i = 0, j = 0  
    merged = []  
    while i < length(list1) and j < length(list2):  
        if list1[i] < list2[j]:  
            merged.append(list1[i])  
            i = i + 1  
        else:  
            merged.append(list2[j])  
            j = j + 1  
    append remaining elements of list1 or list2  
    return merged
```