

1a.

Each account a has:

stack of amounts
count[0..T]

Insert(a, t):

push t onto account[a].stack
count[a][t] = count[a][t] + 1

Search(a, t):

if count[a][t] > 0 then
 return true
else
 return false

Most_recent(a):

if stack is empty then
 return NULL
else
 return top value of stack

Delete_most_recent(a):

if stack is empty then
 return
t = pop from stack
count[a][t] = count[a][t] - 1

Insert(a, t):

When we add a transaction, we push t onto the stack and increment count[a][t]. Both steps take constant time.

Time complexity: O(1)

Search(a, t):

To check if an amount t exists, we simply test if count[a][t] > 0. Array access is constant time.

Time complexity: O(1)

Most_recent(a):

To find the most recent transaction, we return the top element of the stack. Accessing the top of a stack is constant time.

Time complexity: O(1)

Delete_most_recent(a):

To delete the most recent transaction, we pop the top of the stack and decrement count[a][t]. Both operations are constant time.

Time complexity: O(1)**1b.**

Each account a has:

stack of amounts
balanced BST (amount → frequency)

Insert(a, t):

```
push t onto account[a].stack
if t exists in BST then
    freq[t] = freq[t] + 1
else
    insert t into BST with freq = 1
```

Search(a, t):

```
if t exists in BST then
    return true
else
    return false
```

Most_recent(a):

```
if stack is empty then
    return NULL
else
    return top value of stack
```

Delete_most_recent(a):

```
if stack is empty then
    return
t = pop from stack
freq[t] = freq[t] - 1
if freq[t] == 0 then
    remove t from BST
```

INSERT(a, t):

Push t onto the stack ($O(1)$), then search for t in the BST. If found, increment its frequency; otherwise, insert a new node. Searching and inserting in a BST both take $O(\log n)$.

Time complexity: $O(\log n)$

Search(a, t):

Look up t in the BST. Searching in a balanced BST takes $O(\log n)$.

Time complexity: $O(\log n)$

Most_recent(a):

Return the top element of the stack. Stack top lookup takes constant time.

Time complexity: $O(1)$

Delete_most_recent(a):

Pop the top of the stack ($O(1)$), then search for that amount t in the BST and decrement its frequency. If the frequency reaches zero, remove the node from the BST. Both searching and deleting in a BST take $O(\log n)$.

Time complexity: $O(\log n)$

2a.**Build_bst(arr, start, end):**

```
if start > end then  
    return NULL
```

```
mid = (start + end) / 2  
root = new Node(arr[mid])
```

```
root.left = Build_bst(arr, start, mid - 1)  
root.right = Build_bst(arr, mid + 1, end)
```

```
return root
```

2b.**Merge_trees(T1, T2):**

```
list1 = Inorder(T1)  
list2 = Inorder(T2)
```

```
merged = Merge_sorted(list1, list2)
```

```
root = Build_bst(merged, 0, length(merged) - 1)
```

```
return root
```

```
Inorder(root):
    if root = NULL:
        return []
    return Inorder(root.left) + [root.key] + Inorder(root.right)
```

```
Merge_sorted(list1, list2):
    i = 0, j = 0
    merged = []
    while i < length(list1) and j < length(list2):
        if list1[i] < list2[j]:
            merged.append(list1[i])
            i = i + 1
        else:
            merged.append(list2[j])
            j = j + 1
    append remaining elements of list1 or list2
    return merged
```