

1 Introduction

2 Complexity, Games, Polytopes and Gale Strings

2.1 Some Complexity Classes

2.2 Normal Form Games and Nash Equilibria

file: background-subsection

2.3 Some Geometrical Notation

2.4 Bimatrix Games, Labels and Polytopes

file: polytopes-subsection

$$P = \{x \in \mathbb{R}^m | x \geq \mathbf{0}, B^\top x \leq \mathbf{1}\}, \quad Q = \{y \in \mathbb{R}^n | y \geq \mathbf{0}, A y \leq \mathbf{1}\}, \quad (1)$$

a point in P has label k iff $x_k = 0$ for $k \in \{1, \dots, m\}$ or $(B^\top x)_{k-m} = 0$ for $k \in \{m+1, \dots, m+n\}$; analogously, a point in Q has label k if and only if either $y_{k-m} = 0$ for $k \in \{m+1, \dots, m+n\}$ or $(A y)_k$ for $k \in \{m+1, \dots, m+n\}$.

proposition 1. Let (A, B) be a bimatrix game and (x, y) be one of its Nash equilibria. Then (z, z) , where $z = (x, y)$, is a Nash equilibrium of the symmetric game (C, C^\top) , where

$$C = \begin{pmatrix} 0 & A \\ B^\top & 0 \end{pmatrix}.$$

Theorem 1. [15] Let $l : [n] \rightarrow [m]$, and let (U, B) be the unit vector game where $U = (e_{l(1)} \cdots e_{l(n)})$. Consider the polytopes P^l and Q^l where

$$P^l = \{x \in \mathbb{R}^m | x \geq \mathbf{0}, B^\top x \leq \mathbf{1}\} \quad (2)$$

$$Q^l = \{y \in \mathbb{R}^n | y \geq \mathbf{0}, \sum_{\substack{j \in N_i \\ i \in [m]}} y_j \leq 1\} \quad (3)$$

where $N_i = \{j \in [n] | l(j) = i\}$ for $i \in [m]$.

Give a labeling l_f of the facets of P^l according to the inequality defining it, as follows:

$$x_i \geq 0 \text{ has label } i \text{ for } i \in [m] \quad (4)$$

$$B^\top x)_j \leq 1 \text{ has label } l(j) \text{ for } j \in [n] \quad (5)$$

Then $x \in P^l$ is a completely labeled point of $P^l \setminus \{\mathbf{0}\}$ if and only if there is some $y \in Q^l$ such that, after scaling, the pair (x, y) is a Nash equilibrium of (U, B)

Theorem 2. [1] Let Q be a labeled m -dimensional simplicial polytope with $\mathbf{0}$ in its interior, with vertices $e_1, \dots, e_m, c_1, \dots, c_n$, so that $F_0 = \text{conv}\{e_i \mid i \in [m]\}$ is a facet of Q .

Let $l : [n] \rightarrow [m]$, and let (U, B) be the unit vector game with $U = (e_{l(1)} \cdots e_{l(n)})$ and $B = (b_1 \cdots b_n)$, where $b_j = \frac{c_j}{1 + \mathbf{1}^\top c_j}$ for $j \in [n]$.

Label the vertices of Q as follows:

$$l_v(-e_i) = i \text{ for } i \in [m] \quad (6)$$

$$l_v(c_j) = l(j) \text{ for } j \in [n] \quad (7)$$

Then a facet $F \neq F_0$ of Q with normal vector v is completely labeled if and only if (x, y) is a Nash equilibrium of (U, B) , where $x = \frac{v+1}{1^\top(v+1)}$, and $x_i = 0$ if and only if $e_i \in F$ for $i \in [m]$. Any j so that c_j is a vertex of F represents a pure best reply to x ; the mixed strategy y is the uniform distribution on the set of the pure best replies to x .

2.5 Cyclic Polytopes and Gale Strings

2.6 Labeling and the Problem ANOTHER GALE

file: gale-def-subsection

3 Algorithmic and Complexity Results

3.1 Lemke Paths and the Lemke-Howson for Gale Algorithm

We have NEs \Leftrightarrow completely labeled things (facets, vertices, GS) We give now different versions of fundam algorithm to deal with labeling looking for compl.label. - in particular in these cases first in version on simple polytopes with labeled facets, (name: Lemke-Howson; Lemke-Howson 1964, Shapley 1974 beautiful exposition) then in its corresponding version on Gale strings (name: Lemke-Howson for Gale, where???)
We will also mention dual version on simplicial polytopes with labeled vertices (name: exchange algorithm, Edmonds - Sanità). (This case: index more problematic - see)
Sth about why (exp lemke paths; conjecture)

Consider a labeling $l : [n] \rightarrow [m]$, for a set X with $|X| = n$. Then $x = (x_1, \dots, x_m) \in X^m$ is *almost completely labeled* if $|\{j \in [n] \mid x_i = j \text{ for some } i \in [m]\}| = [m] \setminus \{k\}$ for exactly one $k \in [m]$. This mean that all labels appear once in x , except for the *missing label* k , and a *duplicate label* $\bar{k} \in [m]$ that appears twice.

now we see in poly with labeled facets, cl vertices

Let P be a simple polytope in dimension m with n facets. We define the operation of *pivoting on vertices* as moving from a vertex x of P to another vertex y such that there is an edge between x and y . Note that, since P is simple, there are exactly m possible choices for y .

Now let $l : [n] \rightarrow [m]$ be a labeling of the facets of P such that there is at least one completely labeled vertex x_0 of P . Note that if we pivot from a vertex we “leave behind” a facet, that has label k ; we call this *dropping label* k . We will then reach a vertex that lies on a new facet, that has label j ; we will call this *picking up label* j We give the *Lemke-Howson algorithm* as follows:

reference!

The steps of the Lemke-Howson algorithm result in a *Lemke path* that connects two completely labeled vertices through *k -almost complementary*

Algorithm 1: Lemke-Howson algorithm

input : A simple m -polytope P with n facets. A labeling
 $l : [n] \rightarrow [m]$ of the facets of P . A vertex x of P , completely labeled for l .

output: A completely labeled vertex $y \neq x_0$ of P .

- 1 choose a label $k \in [n]$
- 2 pivot from x to y dropping label k
- 3 **while** y is not completely labeled **do**
- 4 let j be the duplicate label of y
- 5 pivot from y to y' dropping label j on the facet shared with x
- 6 rename $y = y'$
- 7 **return** y

vertices and edges, that is, almost completely labeled vertices and edges where the only missing label is k . It remains to show that $y \neq x$. This comes from the fact that the Lemke paths are *simple paths*, that is, there are no “loops” where a vertex is visited more than once. This is not possible because at each vertex there are only two edges corresponding to the missing label k , since P is not degenerate; one is the edge that is traversed to get to the vertex, one is the one that is traversed to leave it in the next step. This proves the following.

proposition 2. *The Lemke-Howson algorithm returns a solution to the problem ANOTHER COMPLETELY LABELED VERTEX.*

In the context of finding the Nash equilibrium of a bimatrix game (A, B) , there are two equivalent implementations of the Lemke-Howson algorithm.

We can consider the game C as in proposition 1, and the associated polytope $S = \{z \in \mathbb{R}^{m+n} \mid z \geq \mathbf{0}, Cz \leq \mathbf{1}\}$, labeling the $2(m+n)$ inequalities defining the facets of S as $1, \dots, m+n, 1, \dots, m+n$. Then applying the Lemke-Howson algorithm starting from vertex $\mathbf{0}$ returns a Nash equilibrium

(z, z) of C and a corresponding $(x, y) = z$ a Nash equilibrium of (A, B) .

We can also follow the “traditional” version of the Lemke-Howson algorithm; a very clear exposition of this can be found in Shapley [16]. Let P and Q be the best response polytopes of (A, B) as in 1. We then move alternately on P and Q , starting from the couple of vertices $(\mathbf{0}, \mathbf{0})$. Since we move in \mathbb{R}^m and \mathbb{R}^n instead of \mathbb{R}^{m+n} , this version is more practical to visualize, as shown in the following example.

ex Savani - von Stengel, pag. 11; fig 8 are Schegel diagrams of BR polytopes.

Example 3.1.

(?) thm Savani '06 - pg 12 SvS-15: not only LH paths on $P \times Q$ simple path, but projection to simple paths on both P and Q .

In the case of a unit vector game (U, B) , where $U = (e_{l(1)} \cdots e_{l(n)})$ for a labeling $l : [n] \rightarrow [m]$, let as in 2; theorem 1 then shows that P^l encodes all the Nash equilibria of (U, B) as completely labeled vertices. We can apply the Lemke-Howson algorithm to both P^l and to the product of the best response polytopes P and Q . The projection of a Lemke path for a missing label $i \in [m]$ on $P \times Q$ to P defines a Lemke path in P^l . However, $P \times Q$ has $m + n$ labels, therefore there could be Lemke paths for a missing label $m + j$ with $j \in [n]$ on $P \times Q$ that get lost in the projection on P^l . The following theorem, proved in Savani and von Stengel [15], shows that there is no loss of generality in studying Lemke paths on P^l .

Theorem 3. *Let (U, B) be a unit vector game, with $U = (e_{l(1)} \cdots e_{l(n)})$ for a labeling $l : [n] \rightarrow [m]$; let $P = \{x \in \mathbb{R}^m | x \geq \mathbf{0}, B^\top x \leq \mathbf{1}\}$ and $Q = \{y \in \mathbb{R}^n | y \geq \mathbf{0}, A y \leq \mathbf{1}\}$, as in 1; and let $P^l = \{x \in \mathbb{R}^m | x \geq \mathbf{0}, B^\top x \leq \mathbf{1}\}$ as in 2. Then the Lemke path on $P \times Q$ for the missing label k projects to a path on P that is the Lemke path on P^l for missing label k if $k \in [m]$, and for missing label $l(j)$ if $k = m + j$ with $j \in [n]$.*

if poly dual cyclic, when cons from pow gale strings we could take the str of labels $l(n+j) \cdots l(n+m)$ instead of $l(1) \cdots l(n+m)$, that is, we could cut the “artificial” first labels $12\dots n$.

Do we go straight for this, or? YES!

After all, in main we’re studying ANOTHER GALE in general, not nec starting from $12\dots n$; and we’re interested in finding *one* eq that’s not the one we started from (and is at other end of LPath, since index and so on), *not all equilibria*; but the eq we started from is not nec the artificial one - actually, if we go with this we can take any NE to start looking for another, and we’re sure to find a “non-artificial” one. Note: if we were looking for all NE, LH doesn’t work anyway - see ex by Wilson in Shapley, where “disconnected” paths between equilibria.

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complexity considerations, PPAD

Needed for why result interesting; it can be done at the end, with Morris paths

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extend term “Lemke paths” to paths of every LH-style algorithms we see. ... In this case, Lemke paths “translate” as / correspond to...

3.2 The Complexity of GALE and ANOTHER GALE

file: main-result-subsection

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