

# Beginner's Practical: Hypergraph Conductance-based Clustering

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Hypergraph clustering is a problem of partitioning the nodes of a hypergraph into "true clusters" which are densely connected internally and sparsely connected externally. Conductance is a metric measuring the quality of a cut which is defined as the ratio of the weight of the cut to the minimal volume of one side of the cut. In this project, we implement a hypergraph clustering algorithm that heuristically tries to minimize the maximum value of the conductance over all cuts between clusters to find a high-quality hypergraph clustering. The algorithm is based on the existing shared-memory parallel algorithm for hypergraph partitioning Mt-KaHyPar.

## 1 INTRODUCTION

Hypergraph clustering is a problem of partitioning the nodes of a hypergraph into "true clusters" which are densely connected internally and sparsely connected externally. It is useful in many applications, such as semi-supervised learning [6], gene expression analysis [5] and image segmentation [3]. Conductance metric, which is defined as the ratio of the weight of a given cut to the minimal volume of one side of the cut, is widely used in graph clustering [2, 4] but to the best of our knowledge, it has only been applied as a metric for evaluating the quality of a hypergraph clustering [1]. Therefore in this project, we implement a heuristic hypergraph clustering algorithm with the objective of minimizing the maximum value of the conductance over all cuts between the clusters while partitioning the hypergraph in at most  $k$  clusters.

Our algorithm is based on the existing shared-memory parallel hypergraph  $k$ -way partitioner Mt-KaHyPar and uses the same multilevel approach as the original algorithm. First, the hypergraph is coarsened in stages by detecting and merging node communities, then the coarsest hypergraph is partitioned into  $k$  clusters, and finally the initial partitioning is refined by uncoarsening the hypergraph nodes and applying a local search algorithm to improve the given objective. We implement two objective functions for the local search algorithm - *conductance\_local* and *conductance\_global* - and compare their performance on a benchmark set `ibm01...18` against a state of the art hypergraph partitioner HyperModularity.

In following, we define the used concepts in Section 2, then we describe related work on hypergraph partitioning and provide an overview of the existing conductance-based clustering algorithms in Section 3. In Section 4, we describe the implementation of the new objectives for optimizing the conductance. Afterwards, we present the experimental results of our implementation in Section 5. Finally, we conclude the paper and discuss future work in Section 6.

## 2 PRELIMINARIES

In this section, we give an overview of the used concepts and notations and define conductance of a hypergraph.

**Hypergraph** is a generalization of a graph, where a hyperedge connects one, two or more nodes. Formally, a hypergraph  $H = (V, E, c, \omega)$  consists of a non-empty set of nodes  $V$ , a set of hyperedges (also called *nets* and simply *edges*)  $E \subseteq 2^V$  and weight functions for hypernodes and hyperedges  $c : V \rightarrow \mathbb{N}$ ,  $\omega : E \rightarrow \mathbb{N}$ . If case of unweighted nodes or hyperedges, we set  $c(v) = 1$  for all  $v \in V$  or  $\omega(e) = 1$  for all  $e \in E$ , respectively.

**$k$ -way clustering** of a hypergraph  $H = (V, E, c, \omega)$  is a partitioning of the set of nodes  $V$  into  $k$  disjoint subsets  $V_1, \dots, V_k$  called *clusters*. Empty clusters are allowed, so  $k$  is only an upper bound

on the number of clusters. The algorithm is allowed to eventually move all the nodes of a cluster to other clusters if it would be beneficial for the objective function.

**Cut**  $\partial S$  of a hypergraph  $H = (V, E, c, \omega)$  is partitioning of the set of nodes  $V$  into two disjoint non-empty subsets  $S$  and  $V \setminus S$ . A hyperedge  $e \in E$  is called a **cutting edge** of the cut  $\partial S$  if it has at least one pin in  $S$  and at least one pin in  $V \setminus S$ . The weight of the cut is defined as the sum of the weights of the cutting edges of the cut  $\partial S$ :

$$\omega(\partial S) = \sum_{e \in \partial S} \omega(e)$$

Therefore In case of unweighted hyperedges, the weight of the cut is equal to the number of cutting edges of the cut.

**Pin** of a hyperedge  $e \in E$  is a node  $v \in V$  such that  $v \in e$ . A hyperedge with only one pin is called a **single-pin net**. In the implementation, we use number of pins  $PinNum_i(e)$  of a hyperedge  $e \in E$  in a cluster  $V_i$  to decide whether to move a node  $v$  from one cluster  $V_i$  to another cluster  $V_j$ .

**Weighted degree** of a node  $v \in V$  is defined the same way in hypergraphs as in graphs. It is the sum of the weights of all hyperedges  $e \in E$  that contain the node  $v$ :

$$\deg_{\omega}(v) = \sum_{e \in E: v \in e} \omega(e)$$

As during the coarsening phase of the algorithm we merge nodes, we also use **original weighted degree** of a node  $v \in V$   $origDeg_{\omega}(v)$  which is defined as the sum of the weighted degrees of all nodes  $v' \in V$  that were merged into  $v$  during the coarsening phase of the algorithm.

**Volume** of a hypergraph  $H = (V, E, c, \omega)$  is defined as the sum of the weighted degrees of all nodes  $v \in V$ . Analogously, we define the **volume of a cluster**  $V_i$  as the sum of the weighted degrees of all nodes  $v \in V_i$ :

$$\text{vol}(H) = \sum_{v \in V} \deg_{\omega}(v) \quad \text{vol}(V_i) = \sum_{v \in V_i} \deg_{\omega}(v)$$

And for the same reason as for the original weighted degree, we also define the **original volume** of a hypergraph and a cluster:

$$\text{origVol}(H) = \sum_{v \in V} \text{origDeg}_{\omega}(v) \quad \text{origVol}(V_i) = \sum_{v \in V_i} \text{origDeg}_{\omega}(v)$$

This way, the original volume  $origVol$  of a cluster or the whole hypergraph in the coarsened hypergraph is equal to the corresponding volume in the initial - *original* - hypergraph.

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**Conductance of a cut**  $S \subseteq V$ :

$$\varphi(S) = \frac{\sum_{e \in \partial S} \omega(e)}{\min\{\text{vol}(S), \text{vol}(V \setminus S)\}}$$

$k$ -way partition of  $H$ :  $V = \sqcup_{i=1}^k V_k$

**Conductance of a  $k$ -way partition:**

$$\varphi(V_1, \dots, V_k) = \max_{\emptyset \neq I \subset \{1, \dots, k\}} \{\varphi(\sqcup_{i \in I} V_i)\}$$

99 3 RELATED WORK  
100 4 IMPLEMENTATION  
101 5 EXPERIMENTAL RESULTS  
102 6 CONCLUSION  
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104 REFERENCES

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