





## A Trick for an Accurate $e^{-|x|}$ Function in Fixed-Point Arithmetics

# Mantas Mikaitis Department of Mathematics The University of Manchester

mantas.mikaitis@manchester.ac.uk

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#### Motivation

- Improve accuracy of exponential decay
- Fixed-point arithmetic on embedded CPUs and similar
- Inputs of e<sup>x</sup> in exponential decay require both integer and fractional parts
- Taking e<sup>0</sup> as special case, outputs only fractional
- To optimize accuracy this means mixed I/O format combination
- This work: use single-format implementation of e<sup>x</sup> to get a mixed format e<sup>-|x|</sup>
- No change to the implementation of e<sup>x</sup> needed
- Beneficial when single-format ex cannot be modified

#### Fixed-Point Arithmetic

Consider two commonly used fixed-point data formats from the ISO/IEC 18037:2008 standard:

- s16.15 (accum): sign, 16 integer, 15 fractional bits
- s0.31 (fract): sign, zero integer, 31 fractional bits.

Here  $\pm 16.15$  has a range of approximately (-65536, 65536) in steps of  $\varepsilon_{s16.15}=2^{-15}$ .  $\pm 0.31$  has a range of (-1,1) in steps of  $\varepsilon_{s0.31}=2^{-31}$ .





#### Usage of two formats

Use s16.15 for range, s0.31 for accuracy.

## **Exponential in Fixed-Point Arithmetic**

Assume we have an implementation (for example, in hardware) of  $y_a = e^{x_a}$  with  $x_a$  and  $y_a$  s16.15. Here, rounding is implicit to what format is on the left-hand side.

The input domain is

$$x_a \in (\log_e(2^{-15}) = -10.397..., \log_e(2^{16} - 2^{-15}) = 11.09...)$$

since any  $x_a$  outside that would **under/overflow** output.

#### Mixed-format input/output domain

If we take input s16.15 but output s0.31, the input domain changes to  $x_a \in (\log_e(2^{-31}) = -21.487..., 0)$ . Wider input range and more accurate outputs for  $x_a < 0$ .

## Exponential Decay in Fixed-Point Arithmetic

We sometimes need only  $x \le 0$  in  $e^x$ , for example exponential decay.

We focus on improving the accuracy of  $e^{-|x|}$  using the observations above, when we have  $e^x$  in \$16.15. Take  $y_f = e^{x_a}$  to be exponential in \$0.31.

We use int(x) to get an **underlying integer representing** x **in fixed-point**.

If we first compute  $y_a$  and then do a basic shift to get  $int(y_f) = 2^{16}int(y_a)$  we would not gain any accuracy since zeros are shifted in to the bottom 16 bits.

## A Trick with Input Range Shifting

Use a property

$$2^{16}e^x = e^{\log_e 2^{16}}e^x = e^{16 \times \log_e 2 + x}$$
.

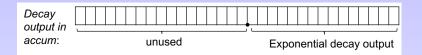
#### Input range shifting

Instead of computing  $int(y_f)$  through first computing  $y_a$  and then shifting, we can compute

$$int(y_f) = int(e^{16 \times log_e 2 + x_a}).$$

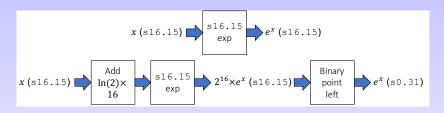
Note that this modification to the input requires only an addition of a precomputed constant and interpreting the underlying integer representation of the output from the  $\mathfrak{s}16.15$  exp as  $\mathfrak{s}0.31$ .

## A Trick with Input Range Shifting



- Most bits are unused in s16.15 exp for x < 0.
- The core idea is to shift the range of inputs from x < 0 to the original input of the s16.15 exp, compute exp using all the bits of the s16.15 and recover the answer in the original range.</p>
- The answer is placed 16 bits to the left from what it should be in \$16.15, which allows us to interpret it as \$0.31.

## A Trick with Input Range Shifting



In summary, the proposed modification includes following steps (figure, top: standard usage, bottom: proposed modification).

- Add  $16 \times \log_{e} 2 \approx 11.09$  to  $x_a$ .
- Calculate exp using the s16.15 implementation.
- Interpret the output bits as s0.31 instead of s16.15.

## Example

Consider computing  $e^{-8}$ . In basic s16.15 we get 0.000335693359375 (error  $\sim -0.00000023$ ).

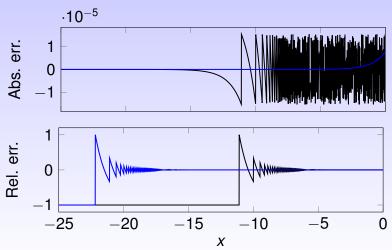
- Compute the constant  $log_e 2 \times 16 = 11.090362548828125 (s16.15).$
- Shift range:
   -8 + 11.090362548828125 = 3.090362548828125
   (s16.15).
- Compute  $e^{3.090362548828125} = 21.98504638671875$  (s16.15).
- Get the underlying int representation:  $21.98504638671875 \times 2^{15} = 720506$ .
- Interpret as  $\pm 0.31$ :  $720506/2^{31} = 0.000335511751472949981689453125$  (error  $\sim -0.000000049$ ).

## **Accuracy Comparison**

- To compare the accuracy of the two approaches, we used MATLAB's binary64 exp.
- We take  $x \in [-25, -0.01]$  (in steps of 0.01) and for each obtain  $x_a$  by rounding to the nearest s16.15 value.
- The constant  $\log_e 2 \times 16$  is represented as \$16.15 and added to  $x_a$  (note addition contains no error, only the rounding of the constant).
- We compute the binary64 exp and round it to s16.15 to simulate the s16.15 implementation.
- To implement the proposed mixed-format version, we round the output to s16.15 and then divide by 2<sup>16</sup> to obtain the s0.31 answer.

## Accuracy Comparison

Below: absolute and relative errors of the standard s16.15 exp (black) and the mixed-format version.



#### Conclusion

- We presented a basic input transformation to improve the accuracy of e<sup>x</sup> when x < 0.</li>
- Tested using the s16.15 and s0.31 fixed-point arithmetics.
- Both the **input range** where underflow does not occur and the **accuracy** of  $e^{-|x|}$  were improved.
- Future work includes generalizing the results for parameterized precisions, performing error analysis, and finding other functions that can be similarly improved.

#### Slides and code

Available https://mmikaitis.github.io/talks/

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