Garbage Collector in Java

how your VM does the "chores" instead of you

Questions to answer

- What is a garbage collector?
- What approaches a garbage collector uses?
- How is the JVM's garbage collector working?
- How can we improve GC's performance?

What is a garbage collector?

- Memory management handles where and how objects are represented in memory
- Problem: objects not used anymore occupy space
- Solution: remove them !!!
- Ways: explicit (C++ style) vs automatic (Java style)

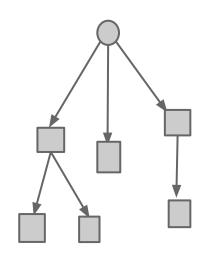
What is a garbage collector?

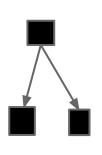
- A garbage collector is an automatic memory management system
- Represent a set of data structures and algorithms hidden from the developer
- Main task is to find the objects that are no longer used and remove them

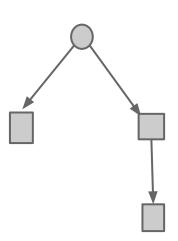
Terminology

- Root local or global variable
- Live object object referenced either by a root or another live object
- Garbage object object created by the program which is not live
- Reachability an object is reachable from another one if there is a chain of references that links them

Graph analogy







- Root
- Live Object
- Garbage

Design choices

- Serial vs Parallel
- Concurrent vs Stop-the-world
- Compacting vs Non-compacting vs Copying

Performance metrics

- Throughput
- Overhead
- Pause time
- Frequency of collection
- Footprint
- Promptness

Basic algorithms: Mark-and-Sweep

- A 2-phase algorithm
- Mark phase a DFS search from roots to find all the live objects and mark them
- Sweep phase visiting all objects and deleting those unmarked
- Compacting option moving live objects left most
- Performance: linear in terms of no of objects
- Problems: fragmentation*

Basic algorithms: Copying

- A 1-phase algorithm
- There are 2 equal memory allocation spaces: old space and new space
- During the mark phase live objects are copied from old space to new space
- At the end algorithm the roles of the physical spaces are swapped
- Fast but cuts the available space by half !!!

Basic algorithms: Generational

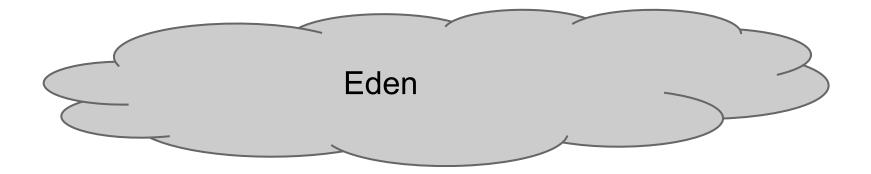
- Idea: the longer one object lives the more probable is the object would live longer
- At least 2 generations of objects can be observed
- Young generation objects that have a short life (usually small)
- Old generation objects that have a long life (usually larger)
- Algorithms tuned for each generation can be used

- Uses 3 generations: young, old and permanent
- Young generation contains Eden and 2 survivor spaces
- Almost all the objects are allocated in Eden and then promote to survivor spaces
- Old generation contains objects that survived several collections
- Permanent generation info about classes and methods

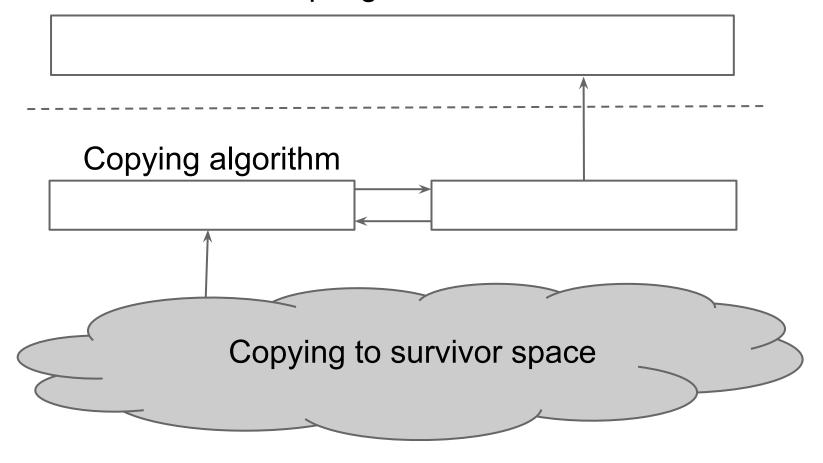
Old generation

New survivor space

Old survivor space



Mark-and-Sweep algorithm



- The objects are allocated in Eden
- When the Eden is collected the live objects are copied in the current new survivor space and the Eden is erased
- In the survivor space the copying algorithm is used
- After an object survived several collections in the survivor space is moved in old generation space
- In the old generation space the Mark-and-Sweep algorithm is used

HotSpot VM Algorithms

- Serial Collector
 - -XX:+UseSerialGC
- Parallel Collector
 - -XX:+UseParallelGC
- Parallel Compacting Collector
 - -XX:+UseParallelOldGC
- Concurrent Mark-and-Sweep Collector
 - -XX:+UseConcMarkSweepGC

HotSpot VM: Serial Collector

- The collections for both new and old generations are done serially with the execution of program
- Stop-the-world happens when the collector is running
- Useful for most applications run on a clientstyle machine (1 CPU) without constraints on pause time
- Enabled by default on client machines

HotSpot VM: Parallel Collector

- Collection of young generation is done using multiple threads in parallel
- Collection of old generation is done serial
- Still a Stop-the-world algorithm but the pause time is decreased
- Decreases overhead, increases throughput
- Useful on multi-core machines for applications with medium to large data sets

HotSpot VM: Parallel Compacting

- Same as parallel collector, just uses multiple threads for old generation collection too
- The Mark-and-Sweep algorithm is used with 3 phases
- Mark phase parallelized
- Summary phase serial
- Sweep phase parallelized
- Better parameters for throughput, overhead and pause time than parallel collector

HotSpot VM: Concurrent Collector

- Use the same algorithm for young generation as parallel collector
- Collection for old generation is done concurrently with the program execution
- The concurrent algorithm has 4 phases
- Initial mark stop-the-world, serial
- Concurrent marking
- Remark stop-the-world, parallel
- Concurrent sweep

HotSpot VM: Concurrent Collector

- Is a non-compacting algorithm (leads to fragmentation - bad!!!)
- Overhead when allocating objects in old generation
- The collections starts before the heap is full
- Floating garbage can survive between collections
- Significantly decreases pause time

HotSpot VM Selecting collector

- Use serial algorithm when:
 - app has small data set
- 1 CPU available and no constraint on pause time
- Use parallel algorithm when:
- performance is most important and pause time constraint is not very "heavy"

HotSpot VM Selecting collector

- Use parallel compacting algorithm when:
- the same conditions as for parallel algorithm where met and the machine has more than 2 CPUs

- Use concurrent mark-and-sweep when:
- response time more important than throughput and pause times have to be <1s

HotSpot VM Ergonomics

- Parameter tuning of algorithms is done using automatic selection and behavioural tuning
- Automatic selection selects the type of algorithm and parameters best for the hardware configuration
- A machine is considered server if it has 2 or more CPUs and 2 GB RAM or more

HotSpot VM Ergonomics

- Behavioural tuning refers to setting a goal for one parameter and let the runtime tune its parameters to achieve it
- Maximum Pause Time Goal
 - -XX:MaxGCPauseMillis = n
- Throughput Goal
 - -XX:GCTimeRatio = n ratio of garbage time = 1/ (1+n)

HotSpot VM Extra-tuning

- The best approach for extra-tuning is the donothing approach !!!
- Change the default collector or the behavioural parameters (if appropriate)
- Use debugging tools available in the JDK
- May want to modify sizes used in the algorithms: -Xmx,-Xms, NewRatio, MinHeapFreeRatio, MaxHeapFreeRatio, SurvivorRatio

Code considerations

- The Java API has 2 functions: System.gc(), Runtime.gc()
- These functions do not force execution of GC, but merely gives a hint to the VM
- Methods to leverage the GC: object pooling and use of weak references

Code consideration: Object Pooling

- Design pattern consisting of creating objects beforehand and keeping them alive during the execution of the program
- A pool is a set of objects ready to use without needing allocation
- Useful just when: initialization cost is high, rate of instantiation is high and instances in use at a time is low
- Controversial as modern GC are powerful

Code considerations: Weak references

- Different types of references: strong, soft, weak and phantom
- Weak references do not count when considering reachability
- Soft references same as weak references but are not collected if can be saved in memory
- Weak and soft references are useful in situations like caching or dynamic relations between instances

Sources

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Questions?