THE1 - CENG223 - Fall 2021

Student Information

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Q. 1

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 \begin{array}{l} \neg (p \wedge q) \leftrightarrow (\neg q \rightarrow p) \\ \equiv (\neg p \vee \neg q) \leftrightarrow (\neg q \rightarrow p) \quad De \ Morgan's \ law \\ \equiv (\neg p \vee \neg q) \leftrightarrow (\neg (\neg q) \vee p)) \quad conditional \ statement \ logical \ equivalence(table \ 7 \ row \ 1) \\ \equiv (\neg p \vee \neg q) \leftrightarrow (q \vee p) \quad double \ negation \ law \\ \equiv [(\neg p \vee \neg q) \rightarrow (q \vee p)] \wedge [(q \vee p) \rightarrow (\neg p \vee \neg q)] \quad biconditional \ statement \ logical \ equivalence(table \ 8 \ row \ 1) \\ \equiv [\neg (\neg p \vee \neg q) \vee (q \vee p)] \wedge [\neg (q \vee p) \vee (\neg p \vee \neg q)] \quad conditional \ statement \ logical \ equivalence(table \ 7 \ row \ 1) \\ \equiv [(p \wedge q) \vee (q \vee p)] \wedge [(\neg q \wedge \neg p) \vee (\neg p \vee \neg q)] \quad De \ Morgan's \ law \\ \equiv [(p \vee q \vee p) \wedge (q \vee q \vee p)] \wedge [(\neg q \vee \neg p \vee \neg q) \wedge (\neg p \vee \neg p \vee \neg q)] \quad distributive \ laws \\ \equiv [(p \vee q) \wedge (p \vee q)] \wedge [(\neg p \vee \neg q) \wedge (\neg p \vee \neg q)] \quad idempotent \ law \\ \equiv (p \wedge q) \wedge (\neg p \vee \neg q) \quad idempotent \ law \end{array}
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Q. 2

a.

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\forall_x \forall_a \forall_y (I(x,y) \land I(a,y) \land x \neq a \rightarrow \exists_z \exists_b (E(x,z) \land E(a,b) \land b \neq z))
where domain of x and a is all people, domain of y is all faculties and domain of b,z is all employee IDs
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b.

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\forall_y \exists_x (I(x,y) \land S(x,x)) where domain of y is all faculties and domain of x is all people
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c.

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\forall_x \big( J(x, medicine) \rightarrow \forall_a \forall_b \forall_c (A(a, x) \land I(a, medicine) \land A(b, x) \land I(b, medicine) \land A(c, x) \land I(c, medicine)) \rightarrow (a = b \lor a = c \lor b = c) \big)
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where domain of x is all job positions, domain of a,b,c is all people

Q. 3

a.

 $p \vee \neg q, p \vee r \vdash (r \to q) \to p$

b.

$$\vdash ((q \to p) \to q) \to q$$

	((1	• /	-/	•
				assumption
				assumption
				assumption
$ \begin{array}{ c c c c } & 4. & \neg p \\ 5. & \neg q \\ 6. & q \\ 7. & \bot \\ \end{array} $				assumption copy 2 copy 3 $\neg e 5,6$
$ \begin{array}{ c c c c c } \hline & 8. & \neg \neg p \\ 9. & p \\ \hline \end{array} $				$\neg i \ 4-7$ $\neg \neg e \ 8$
$\begin{array}{ c c c c }\hline 10. & q \to p \\ 11. & q \\ 12. & \bot \\ \hline \end{array}$				$\begin{array}{l} \rightarrow i \ 3.9 \\ \rightarrow e \ 1.10 \\ \neg e \ 2.11 \end{array}$
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$				$\neg i \ 2\text{-}12$ $\neg \neg e \ 13$
$ 15. ((q \rightarrow p) \rightarrow q) $	$\rightarrow q$			$\rightarrow i$ 1-14

^{*} proved in last page

Q. 4

a.

^{*} proved in last page

b.

1.
$$\forall_x \forall_y (P(x,y) \to \neg P(y,x))$$
 premise $2. \forall_x \exists_y P(x,y)$ premise $2. \forall_x \exists_y P(x,y)$ premise $3. a$

4. $\forall_y (P(a,y) \to \neg P(y,a))$
5. $P(a,a) \to \neg P(a,a)$ $\forall_y e \ 5$
6. $\exists_y P(a,y)$ $\forall_x e \ 2$

7. $P(a,a)$ assumption $egline egline eg$

 $\forall_x \forall_y (P(x,y) \to \neg P(y,x)), \forall_x \exists_y P(x,y) \vdash \neg \exists_v \forall_z P(z,v)$

Proofs for used derived rules

a. Disjunctive Syllogism

b. De Morgan's law

$$\neg (P \land Q) \vdash \neg P \lor \neg Q$$
1. $\neg (P \land Q)$ premise

$$\begin{vmatrix}
2. \neg (\neg P \lor \neg Q) & \text{assumption} \\
4. \neg P \lor \neg Q & & \forall i \ 3 \\
5. \bot & & \neg e4, 2
\end{vmatrix}$$
6. $\neg \neg P$ $\neg i \ 3-5$

$$\begin{vmatrix}
7. \neg Q & \text{assumption} \\
8. \neg P \lor \neg Q & & \forall i \ 7 \\
9. \bot & & \neg e \ 8, 2
\end{vmatrix}$$
10. $\neg \neg Q$ $\Rightarrow e \ 8, 2$

10. $\neg \neg Q$ $\Rightarrow e \ 8, 2$

11. $\Rightarrow e \ 10$ $\Rightarrow e \ 11, 14$

16. $\Rightarrow \neg (\neg P \lor \neg Q)$ $\Rightarrow e \ 16$