

## HW08

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### Task 1: Breadth-First Search Running Time

```

Set<Integer> nbrsExcluding(
    UndirectedGraph<Integer> G,
    Set<Integer> vtxes,
    Set<Integer> excl
){
    Set<Integer> union = new TreeSet<>(); // not HashMap //O(1)
    for (Integer src : vtxes) { //It takes n times to run so O(n)*O(logn) = O(nlogn)
        for (Integer dst : G.adj(src)) //O(logn)
            if (!excl.contains(dst)) { union.add(dst); } // O(logn)
        }
    return union; //O(1)
}

Set<Integer> bfs(UndirectedGraph<Integer> G, int s) {
    Set<Integer> frontier = new TreeSet<>(Arrays.asList(s)); //O(1)
    Set<Integer> visited = new TreeSet<>(Arrays.asList(s)); //O(1)
    while (!frontier.isEmpty()) {
        frontier = nbrsExcluding(G, frontier, visited); //O(nlogn)
        visited.addAll(frontier); // the i-th position is what's reached at i hops //O(nlogn)
    }
    return visited; //O(1)
}

```

- The total time it takes is  $O(n \log n)$ , where  $n$  is the number of vertices.

## Task 2: Mathematical Facts

i)

```
int minSoFar = Integer.MAX_VALUE;
int numUpdate = 0;
for (int i=1;i<=n;i++) {
    if (p(i) < minSoFar) {
        minSoFar = p(i);
        numUpdate++;
    }
}
```

Prove that at the end of the for-loop,  $\ln(n+1) \leq E[\text{numUpdate}] \leq 1 + \ln n$ .

- The start is when  $p(1)$  and the end of the loop is when  $p(n)$ .

(ii) Let  $G=(V,E,w)$  be an undirected connected weighted graph with distinct edge weights. Show that  $G$  has a unique minimum spanning tree.

- If each edge has a distinct edge weight, then each weight will appear in the graph only once, meaning there will be a pair of edges  $(V,E)$  and  $(E,V)$  that have the same weight  $(w)$ .

## Task3: HackerRank Problems

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