

## Homework 5

(1) For each of the following prime  $p$ , find a generator of  $\mathbb{Z}_p^*$ .

(a)  $p = 17$

(b)  $p = 29$

(c)  $p = 31$

(2) If you pick any integer from  $\mathbb{Z}_p^*$  randomly: what's the probability that it is a generator of  $\mathbb{Z}_p^*$ ?

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(a)  $p = 17$

$$\text{ord}(g) = 16 = 2^4$$

(b)  $p = 29$

$$\forall h \in \mathbb{Z}_p^*, \text{ord}(h) \mid |\mathbb{Z}_p^*| = 16 = 2^4$$

$$\text{ord}(h) = 2^i \text{ for some } i$$

(c)  $p = 31$

$$\text{Take } h \in \mathbb{Z}_p^*, \text{ if } h^3 \not\equiv 1 \pmod{17}$$

$$\text{then } \text{ord}(h) = 2^4$$

hence,  $h$  is a generator

(2) If you pick any integer from  $\mathbb{Z}_p^*$

randomly: What's the probability that

it is a generator of  $\mathbb{Z}_p^*$ ?

(a)  $p = 17$

$$\frac{\# \text{ of generators}}{\text{size of group}} = \frac{\phi(p-1)}{p-1}$$

(b)  $p = 29$

If  $g$  is a generator

(c)  $p = 31$

then  $g^i$  is also a generator

if  $i$  is coprime with  $p-1$ .

1 b)  
 $p = 29$

$$p-1 = 28 = 2^2 \cdot 7$$

$$\forall h \in \mathbb{Z}_p^\times, \text{ord}(h) = 2^i \cdot 7^j \quad \begin{matrix} 0 \leq i \leq 2 \\ 0 \leq j \leq 1 \end{matrix}$$

If  $h^2 \neq 1$  and  $h^{2 \cdot 7} \neq 1$ , then the  $\text{ord}(h) = 2^2 \cdot 7$ .

If  $h^2 \neq 1$ , what could be the order of  $h$ ?

$7, 2 \cdot 7, 2^2 \cdot 7$

Hence, if in addition,  $h^{2 \cdot 7} \neq 1$ , then the only

order of  $h$  is  $2^2 \cdot 7$ .

To find a generator of a group of order  $n$   
(in  $\mathbb{Z}_p^*$ ,  $n = p-1$ )

$$n = p_1^{e_1} p_2^{e_2} \dots p_n^{e_n}, \text{ } p_i \text{ distinct primes}$$

$$n_i = \frac{n}{p_i}$$

Pick  $h \in G$ . If  $h^{n_i} \neq 1$  for all  $i$

then  $h$  is a generator.

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$$h^{n_i} \neq 1 \longrightarrow p_i^{e_i} \mid \text{ord}(h)$$

Hence if  $h^{n_i} \neq 1$  for all  $i$

then  $p_i^{e_i} \mid \text{ord}(h)$  for all  $i$ .

Since  $p_i^{e_i}$ s are pairwise coprime,

$$n = p_1^{e_1} p_2^{e_2} \dots p_n^{e_n} \mid \text{ord}(h)$$

Hence,  $\text{ord}(h) = n$ .

$$2) \frac{Q(P-1)}{P-1}$$

(a)  $p=17$ , probability is  $\frac{Q(16)}{16} = \frac{Q(2^4)}{16} = \frac{1}{2}$

(b)  $p=29$ , probability is  $\frac{Q(28)}{28} = \frac{Q(2^2)Q(7)}{2^2 \cdot 7}$   
 $= \frac{2 \cdot 6}{4 \cdot 7}$   
 $= \frac{3}{7}$

generator  $g \in G$  is such that

$$\langle g \rangle = \{ g^i, i \in \mathbb{Z} \} = G$$

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$$\forall h \in G, h = g^i \text{ for some } i$$

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If  $g=3$  is a generator,  
then all elements in  $G$  is  $3^i$   
for some  $i$

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(3) Let  $g \in G$  be a group element.

prove that 
$$\text{ord}(g^i) = \frac{\text{ord}(g)}{\gcd(i, \text{ord}(g))}$$

(4)

**2.17.** Use Shanks's babystep–giantstep method to solve the following discrete logarithm problems. (For (b) and (c), you may want to write a computer program implementing Shanks's algorithm.)

(a)  $11^x = 21$  in  $\mathbb{F}_{71}$ .

(b)  $156^x = 116$  in  $\mathbb{F}_{593}$ .

(c)  $650^x = 2213$  in  $\mathbb{F}_{3571}$ .



(4)

modulo 71

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Recap:  $x = im + j$   $0 \leq i < m$   $0 \leq j < m$   $m = \lceil \sqrt{70} \rceil = 9$

L1:  $11, 11^2, 11^3, \dots, 11^8$

$u = 11^{-9} = 7$

L2:  $21, 21 \cdot 7, 21 \cdot 7^2, \dots$

$h = 21$

Find a match  $i, j$  such that  $11^i = 21 \cdot 7^j$

Verify that  $11^x = 21$  by performing square and multiply

(5)

**2.27.** Write out your own proof that the Pohlig–Hellman algorithm works in the particular case that  $p - 1 = q_1 \cdot q_2$  is a product of **two distinct primes**. This provides a good opportunity for you to understand how the proof works and to get a feel for how it was discovered.

=  $h_1$

$h$

(5)

**2.27.** Write out your own proof that the Pohlig-Hellman algorithm works in the particular case that  $p-1 = q_1 \cdot q_2$  is a product of **two distinct primes**. This provides a good opportunity for you to understand how the proof works and to get a feel for how it was discovered.

Solve  $g^x = h$  where  $\text{ord}(g) = p-1 = q_1 \cdot q_2$  in  $\mathbb{Z}_p^*$

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$$g_1 = g^{q_2}, \quad h_1 = h^{q_2}, \quad \text{ord}(g_1) = \text{ord}(g^{q_2}) = q_1$$

$$g_2 = g^{q_1}, \quad h_2 = h^{q_1}, \quad \text{ord}(g_2) = q_2$$

Solve  $x_1$  and  $x_2$  such that

$$g_1^{x_1} = h_1$$

$$g_2^{x_2} = h_2$$

Solve  $x$  such that

$$x \equiv x_1 \pmod{q_1}$$
$$x \equiv x_2 \pmod{q_2}$$

Since  $\text{gcd}(q_1, q_2) = 1$ , there exists integers  $u, v$  such that

$$q_1 u + q_2 v = 1$$

$$x = x_1 q_2 v + x_2 q_1 u$$

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Verify that  $g^x = h$

$$\begin{aligned} g^x &= g^{x_1 q_2 v + x_2 q_1 u} = g^{x_1 q_2 v} \cdot g^{x_2 q_1 u} = g^{q_2 x_1 v} \cdot g^{q_1 x_2 u} \\ &= g_1^{x_1 v} \cdot g_2^{x_2 u} \\ &= h_1^v \cdot h_2^u \\ &= h^{q_2 v} \cdot h^{q_1 u} = h^{q_2 v + q_1 u} = h^1 = h \end{aligned}$$

**3.14.** We stated that the number 561 is a Carmichael number, but we never checked that  $a^{561} \equiv a \pmod{561}$  for every value of  $a$ .

(a) The number 561 factors as  $3 \cdot 11 \cdot 17$ . First use Fermat's little theorem to prove that

$$a^{561} \equiv a \pmod{3}, \quad a^{561} \equiv a \pmod{11}, \quad \text{and} \quad a^{561} \equiv a \pmod{17}$$

for every value of  $a$ . Then explain why these three congruences imply that  $a^{561} \equiv a \pmod{561}$  for every value of  $a$ .

(b) Mimic the idea used in (a) to prove that each of the following numbers is a

The next six Carmichael numbers are (sequence [A002997](#) in the [OEIS](#)):

$$\begin{array}{ll} 1105 = 5 \cdot 13 \cdot 17 & (4 \mid 1104; \quad 12 \mid 1104; \quad 16 \mid 1104) \\ 1729 = 7 \cdot 13 \cdot 19 & (6 \mid 1728; \quad 12 \mid 1728; \quad 18 \mid 1728) \\ 2465 = 5 \cdot 17 \cdot 29 & (4 \mid 2464; \quad 16 \mid 2464; \quad 28 \mid 2464) \\ 2821 = 7 \cdot 13 \cdot 31 & (6 \mid 2820; \quad 12 \mid 2820; \quad 30 \mid 2820) \\ 6601 = 7 \cdot 23 \cdot 41 & (6 \mid 6600; \quad 22 \mid 6600; \quad 40 \mid 6600) \\ 8911 = 7 \cdot 19 \cdot 67 & (6 \mid 8910; \quad 18 \mid 8910; \quad 66 \mid 8910). \end{array}$$

If  $n$  is a Carmichael number then  $n$  is a product of distinct primes.

$$n = p_1^{e_1} p_2^{e_2} \dots p_n^{e_n} \quad p_i \text{ are distinct primes}$$

$$\mathbb{Z}_n^\times \cong \mathbb{Z}_{p_1^{e_1}}^\times \times \mathbb{Z}_{p_2^{e_2}}^\times \times \dots \times \mathbb{Z}_{p_n^{e_n}}^\times$$

(6)

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(b) Mimic the idea used in (a) to prove that each of the following numbers is a

$n$  is carmichael number if  $\forall a \in [1, \dots, n-1]$

$$a^{n-1} \equiv 1 \pmod{n} \quad \text{and} \quad n \text{ is composite}$$

$$\approx a^n \equiv a \pmod{n}$$

Fermat little theorem (2nd version)

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If  $p$  is prime, then for all integers  $a$ ,

$$a^p \equiv a \pmod{p}$$

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Fermat little theorem (1st version)

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If  $p$  is prime, then for all integers  $a$  coprime

to  $p$ ,  $a^{p-1} \equiv 1 \pmod{p}$

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(b) Mimic the idea used in (a) to prove that each of the following numbers is a

(a) By Fermat's little theorem

$$a^2 \equiv 1 \pmod{3}$$

$$a^{10} \equiv 1 \pmod{11}$$

$$a^{16} \equiv 1 \pmod{17}$$

$$a^{561} \equiv a^{561 \bmod 2} \equiv a \pmod{3} \rightarrow 3 \mid a^{561} - a$$

$$a^{561} \equiv a^{561 \bmod 10} \equiv a \pmod{11} \rightarrow 11 \mid a^{561} - a$$

$$a^{561} \equiv a^{561 \bmod 16} \equiv a \pmod{17} \rightarrow 17 \mid a^{561} - a$$

Using the fact that if  $a \mid c$  and  $b \mid c$  and  $\gcd(a, b) = 1$ , then  $ab \mid c$ .

$$\text{We have } 3 \cdot 11 \cdot 17 \mid a^{561} - a$$

$$\text{so, } a^{561} \equiv a \pmod{3 \cdot 11 \cdot 17}$$