

Outline

Discrete Logarithm Problem

Brute force : $O(N)$ steps and $O(1)$ space

Baby-step-Giant-step : $O(\sqrt{N} \log N)$ steps
 $O(\sqrt{N})$ space

Discrete Logarithm Problem (DLP)

Given a group G with operation \cdot

(written multiplicatively), and its identity element is 1 .

Let $g, h \in G$. Find an integer x such that $g^x = h$.

In \mathbb{Z}_n^* , the operation is multiplication

given g, h , want to compute x such that

$$g^x = h \pmod n$$

Hard

Hardness of DLP depends on group

In \mathbb{Z}_n with operation $+$, written additively
Given $g, h \in \mathbb{Z}_n$, want to x .

$$x \cdot g = h \pmod{n}$$

To find x ,

$$x = h g^{-1} \pmod{n}$$

Easy problem

Brute - Force Method

Given g, h find x s.t. $g^x = h \in G$

Successive multiplication of g until
we reach h .

Running time: $O(\text{ord}(g))$

$O(|G|)$ exponential
in # of bits
to store $|G|$

$$|G| = 2^{\log |G|}$$

Trivial Space and Running Time

Given a group G ,

$g, h \in G$.

g has order N .

Then there exists an algorithm to solve DLP in $O(N)$ steps and in $O(1)$ space.

each step is group multiplication

Baby-step Giant-Step

Trade-off time with space

$$g^x = h, \quad N = \text{ord}(g)$$

$$x = im + j \quad m = \lceil \sqrt{N} \rceil$$

$$0 \leq i < m$$

$$0 \leq j < m$$

Create two lists:

$$L_1: g^0, g^1, g^2, \dots, g^{m-1}$$

$$L_2: hu^0, hu^1, hu^2, \dots, hu^{m-1}$$

$$u = g^{-m}$$

Find the matched value g^j, hu^i

$$x = im + j.$$

$$x = im + \bar{j}$$

$$g^x = g^{im + \bar{j}} = h$$

$$\begin{aligned} g^j &= hg^{-mi} \\ &= hu^i \end{aligned}$$

Running time :

$O(m)$ multiplications

$O(m \log m)$ sorting & finding
match

Total $O(m \log m)$

$$= O(\sqrt{N} \log \sqrt{N})$$

$$= O(\sqrt{N} \log N) \text{ steps}$$

$$\text{Space} = O(m) = O(\sqrt{N})$$

Proposition 2.21 (Shanks's Babystep–Giantstep Algorithm). *Let G be a group and let $g \in G$ be an element of order $N \geq 2$. The following algorithm solves the discrete logarithm problem $g^x = h$ in $\mathcal{O}(\sqrt{N} \cdot \log N)$ steps using $\mathcal{O}(\sqrt{N})$ storage.*

(1) Let $n = 1 + \lfloor \sqrt{N} \rfloor$, so in particular, $n > \sqrt{N}$.

(2) Create two lists,

List 1: $e, g, g^2, g^3, \dots, g^n$,

List 2: $h, h \cdot g^{-n}, h \cdot g^{-2n}, h \cdot g^{-3n}, \dots, h \cdot g^{-n^2}$.

(3) Find a match between the two lists, say $g^i = hg^{-jn}$.

(4) Then $x = i + jn$ is a solution to $g^x = h$.