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# Outline

- What is cryptography
- Ciphers
- Cryptanalysis
- Shift cipher
- Modular arithmetic
- Group

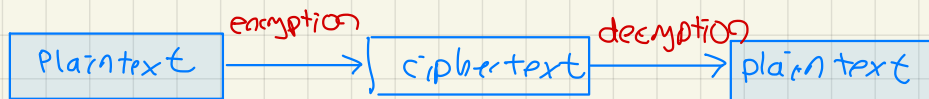
# What is cryptography?

- Protocols for secret communication, under the presence of eavesdropper



- Idea: Transform original message before sending it

- Basic terms:



- secure key:
- ① mutual (shared key) (symmetric)
  - ② (public key, private key) (asymmetric)  
                encryption                  decryption

cryptanalysis: Attack on the cryptography scheme

Read: even without secure key

change: change the original msg

Fog: send msg to Bob in the name of Alice

Encryption and decryption are known to be public

Security of cryptosystem is based on secure key.

# Shift Cipher

Alphabets : A B C D E F . . . X Y Z

secret key :  $K$  is an integer from 0 to 25

Encryption : Each letter in the plaintext, replaced it with the letter  $K$  position to the right in the alphabet list.

$k=1$   
A  $\rightarrow$  B  $\rightarrow$  C  
B  $\rightarrow$  C  $\rightarrow$  D  
C  $\rightarrow$  D  
D  $\rightarrow$  E

Plaintext : HELLO

$k=1$ , ciphertext : IFMMO

$k=2$ , ciphertext : JGNNP

X  $\rightarrow$  Y  
Y  $\rightarrow$  Z  
Z  $\rightarrow$  A

Decryption : "shift to left by  $K$  position"

Cryptanalysis : Brute Force attack

Try all possible keys 26

## Formal Definitions

$e_k(m)$  is an encryption function that maps plaintext  $m$  to ciphertext  $C$  with the key  $k$   
 $d_k(m)$  is a decryption function maps ciphertext  $\rightarrow$  plaintext with the key  $k$ .

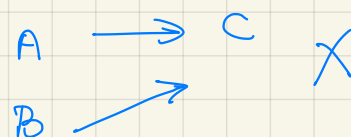
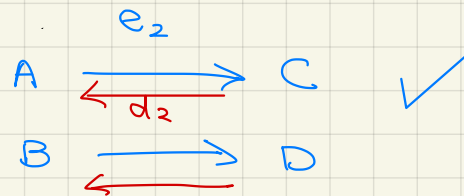
inverse of  $e_k(m)$

$$m = d_k(e_k(m))$$

$e_k$  is a one-to-one function

If  $e_k(m) = e_k(m')$  then

$$m = d_k(e_k(m)) = d_k(e_k(m')) = m'$$



$K$  : space of keys

$M$  : space of plaintexts

$C$  : space of ciphertexts

$e$  : encryption scheme

$d$  : decryption scheme

A  $(K, M, C, e, d)$  cipher is "successful"

- (1) easy to compute ciphertext  $e_k(m)$ ,  $k \in K$ ,  $m \in M$  Alice
- (2) easy to decrypt ciphertext if key  $k$  is known Bob
- (3) difficult to decrypt ciphertext without key  $k$ . Eve
- (4) secure against known plaintext attack
- (5) secure against chosen plaintext attack

easy vs difficult

# Understand Integers

Several "hard" problems provide crypto keys

Notation:  $\mathbb{Z}$

## Divisibility

- $a$  divides  $b$  ( $a|b$ ),  $a$  does not divide  $b$  ( $a \nmid b$ )  
→  $b = ac$  for some integer  $c$
- statement not an operation
- $a$  is a divisor of  $b$

## Remainder

- $n = qd + r$

$d$  is divisor (modulus,  $m$ )

$q$  is quotient (N/A)

$r$  is remainder (residue)

$$\begin{array}{r} 2 \\ 13 \overline{) 37} \\ \underline{26} \\ 11 \end{array} \leftarrow \text{remainder}$$

- Given  $n, d$ , there are infinite possibilities for  $q, r$

$$37 = 2 \cdot 13 + 11$$

$$37 = 3 \cdot 13 - 2$$

- By convention,  $0 \leq r < d$

# Modular Arithmetic

modulus

$$a, b \in \mathbb{Z}, m \in \mathbb{Z}$$

$$a \equiv b \pmod{m} \quad (a \text{ congruent to } b \text{ modulo } m)$$

$$\text{if } m \mid a - b$$

Example:

$$\begin{array}{lcl} 17 \equiv 7 \pmod{5} & (5 \mid 17 - 7) \\ 19 \not\equiv 6 \pmod{11} & (11 \nmid 19 - 6) \end{array}$$

## Reduce modulo m

a reduce modulo m: take the remainder of division of a by m

$$37 \pmod{13} = 11$$

$$64 \pmod{26} = ? (12)$$

$$50 \pmod{26} = 24$$

$$-1 \pmod{26} = ? \text{ 25}$$

$$\begin{array}{r} 64 \\ -26 \\ \hline 38 \\ -26 \\ \hline 12 \end{array} \quad \begin{array}{r} -1 \\ +26 \\ \hline 25 \end{array}$$

## Shift ciphers using modular arithmetic

a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25

Table 1.7: Assigning numbers to letters

Encryption :  $\text{plaintext} + k \pmod{26}$

Decryption :  $\text{ciphertext} - k \pmod{26}$

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$$\mathbb{Z}/26\mathbb{Z} = \{0, 1, 2, \dots, 25\}$$

Set of remainders modulo 26