

# Homework 11

## Question 1

**5.39.** Solve the discrete logarithm problem  $10^x = 106$  in the finite field  $\mathbb{F}_{811}$  by finding a collision among the random powers  $10^i$  and  $106 \cdot 10^i$  that are listed in Table 5.17.

$i$	$g^i$	$h \cdot g^i$	$i$	$g^i$	$h \cdot g^i$	$i$	$g^i$	$h \cdot g^i$	$i$	$g^i$	$h \cdot g^i$
116	96	444	519	291	28	791	496	672	406	801	562
497	326	494	286	239	193	385	437	95	745	194	289
225	757	764	298	358	642	178	527	714	234	304	595
233	517	465	500	789	101	471	117	237	556	252	760
677	787	700	272	24	111	42	448	450	326	649	670
622	523	290	307	748	621	258	413	795	399	263	304

Table 5.17: Data for Exercise 5.39,  $g = 10$ ,  $h = 106$ ,  $p = 811$

## Question 2

**5.40.** Table 5.18 gives some of the computations for the solution of the discrete logarithm problem

$$11^t = 41387 \quad \text{in } \mathbb{F}_{81799} \quad (5.62)$$

using Pollard's  $\rho$  method. (It is similar to Table 5.11 in Example 5.52.) Use the data in Table 5.18 to solve (5.62).

$i$	$x_i$	$y_i$	$\alpha_i$	$\beta_i$	$\gamma_i$	$\delta_i$
0	1	1	0	0	0	0
1	11	121	1	0	2	0
2	121	14641	2	0	4	0
3	1331	42876	3	0	12	2
4	14641	7150	4	0	25	4
$\vdots$						
151	4862	33573	40876	45662	29798	73363
152	23112	53431	81754	9527	37394	48058
153	8835	23112	81755	9527	67780	28637
154	15386	15386	81756	9527	67782	28637

Table 5.18: Computations to solve  $11^t = 41387$  in  $\mathbb{F}_{81799}$  for Exercise 5.40

### Question 3

This exercise describe Pollard-p-factorization algorithm.

$N = pq$  where  $p$  and  $q$  are odd primes

$$\text{Let } f(x) = x^2 + 1 \pmod{N}.$$

$$\text{Let } x_0 = y_0 = 2$$

For  $i = 1, 2, \dots$

(a) compute  $x_i = f(x_{i-1})$

(b) compute  $y_i = f(f(y_{i-1}))$

(c) Compute  $g_i = \gcd(|x_i - y_i|, N)$ .

If  $g_i \neq 1$ , returned  $g_i$  as the prime divisor of  $N$ .

(1) For each of the following cases, compute the smallest  $k$  such that  $g_k \neq 1$  and the ratio  $\frac{k}{\sqrt{N}}$ .

(A)  $N = 8051$

(B)  $N = 10403$

(C)  $N = 9409613$

(2) Let  $p$  be the smallest prime divisor of  $N$ .

Suppose that the function  $f$  is random, show that the algorithm factors  $N$  in  $O(\sqrt{p})$  steps.