

1. (30 points) For each of the following two functions $f(n)$ and $g(n)$, indicate whether $f = O(g)$, or $f = \Omega(g)$ or both (in which case $f = \Theta(g)$).

| | | |
|----|--|-----------------|
| a. | $f(n) = n - 100$ and $g(n) = n - 200$ | $f = \Theta(g)$ |
| b. | $f(n) = n^{1/2}$ and $g(n) = n^{2/3}$ | $f = O(g)$ |
| c. | $f(n) = 100n + \lg n$ and $g(n) = n + (\lg n)^2$ | $f = O(g)$ |
| d. | $f(n) = n \lg n$ and $g(n) = 10n \lg(10n)$ | $f = O(g)$ |
| e. | $f(n) = 10 \lg n$ and $g(n) = \lg(n^2)$ | $f = \Omega(g)$ |
| f. | $f(n) = n^2 / \lg n$ and $g(n) = n (\lg n)^2$ | $f = \Omega(g)$ |
| g. | $f(n) = n^{0.1}$ and $g(n) = (\lg n)^{10}$ | $f = \Omega(g)$ |
| h. | $f(n) = \sqrt{n}$ and $g(n) = (\lg n)^3$ | $f = O(g)$ |
| i. | $f(n) = n^{2^n}$ and $g(n) = 3^n$ | $f = O(g)$ |
| j. | $f(n) = 2^n$ and $g(n) = 2^{n+1}$ | $f = O(g)$ |

2. Given a collection of n nuts and a collection of n bolts, arranged in an increasing order of size, give an $O(n)$ time algorithm to check if there is a nut and a bolt that have the same size. The sizes of the nuts and bolts are stored in the sorted arrays NUTS [1... n] and BOLTS [1... n], respectively. Your algorithm can stop as soon as it finds a single match (i.e., you do not need to report all matches).

```

i = 0, j = 0;
While (i < n and j < n)
Do
    if (NUTS[i] == BOLTS[j]) then
        //Match found
        return (True);
    else if (NUTS[i] < BOLTS[j]) then
        i = i + 1;
    else j = j + 1;
end
end
return (False);

```

3. Let $A[1..n]$ be an array of distinct positive integers, and let t be a positive integer
 - a. Assuming that A is sorted, show that in $O(n)$ time it can be decided if A contains two distinct elements x and y such that $x + y = t$.

two_sum

```

var ht = []; //Hash Table
for i = 0 to n do
  if (A[i] > t) then continue; end
  diff = t - A[i];
  if (ht[diff] != undefined) then
    return [diff, A[i]];
  end
  ht[A[i]] = A[i];
end
return ('Not Found');
```

- b. Use part (a) to show that the following problem, referred to as the 3-Sum problem, can be solved in $O(n^2)$ time: 3-Sum Given an array $A[1..n]$ of distinct positive integers that is not (necessarily) sorted, and a positive integer t , determine whether or not there are three distinct elements x, y, z in A such that $x + y + z = t$

Without using two-sum

```

second=0; third=0;

sort A

for i = 0 to n-2 do
  second=i+1;
  third=n-1;
  while(second<third) do
    if(A[i]+A[second]+A[third] == t) then
      return [ A[i], A[second], A[third]];
    else if (A[i]+A[second]+A[third] < t) then
      second=second+1;
    else third= third-1;
  end
```

```

    end
end
return ('Not Found');

```

With using two-sum

```

For i = 0 to n-2
    if (A[i] > t) then continue; end
    diff = t - A[i];
    match = two_sum (A. slice(i+1), diff); // send array from i+1 element and diff
    //to the 2-sum function to find the next two elements using the logic from (a)
    if (match) then
        return [A[i]] + match;
    end
end
return ("Not Found");

```

4. Let $A[1..n]$ be an array of positive integers (A is not sorted). Pinocchio claims that there exists an $O(n)$ -time algorithm that decides if there are two integers in A whose sum is 1000. Is Pinocchio right, or will his nose grow? If you say Pinocchio is right, explain how it can be done in $O(n)$ time; otherwise, argue why it is impossible.

Pinocchio is right, it can be done in $O(n)$ time using the two-sum method using Hash table discussed in 3(a), this solution works for unsorted arrays as well

```

var ht = []; //Hash Table
for i = 0 to n do
    if (A[i] > 1000) then continue; end
    diff = 1000 - A[i];
    if (ht[diff] != undefined) then
        return [diff, A[i]];
    end
    ht[A[i]] = A[i];
end

```

```
return ('Not Found');
```