

MEMORY SYSTEM

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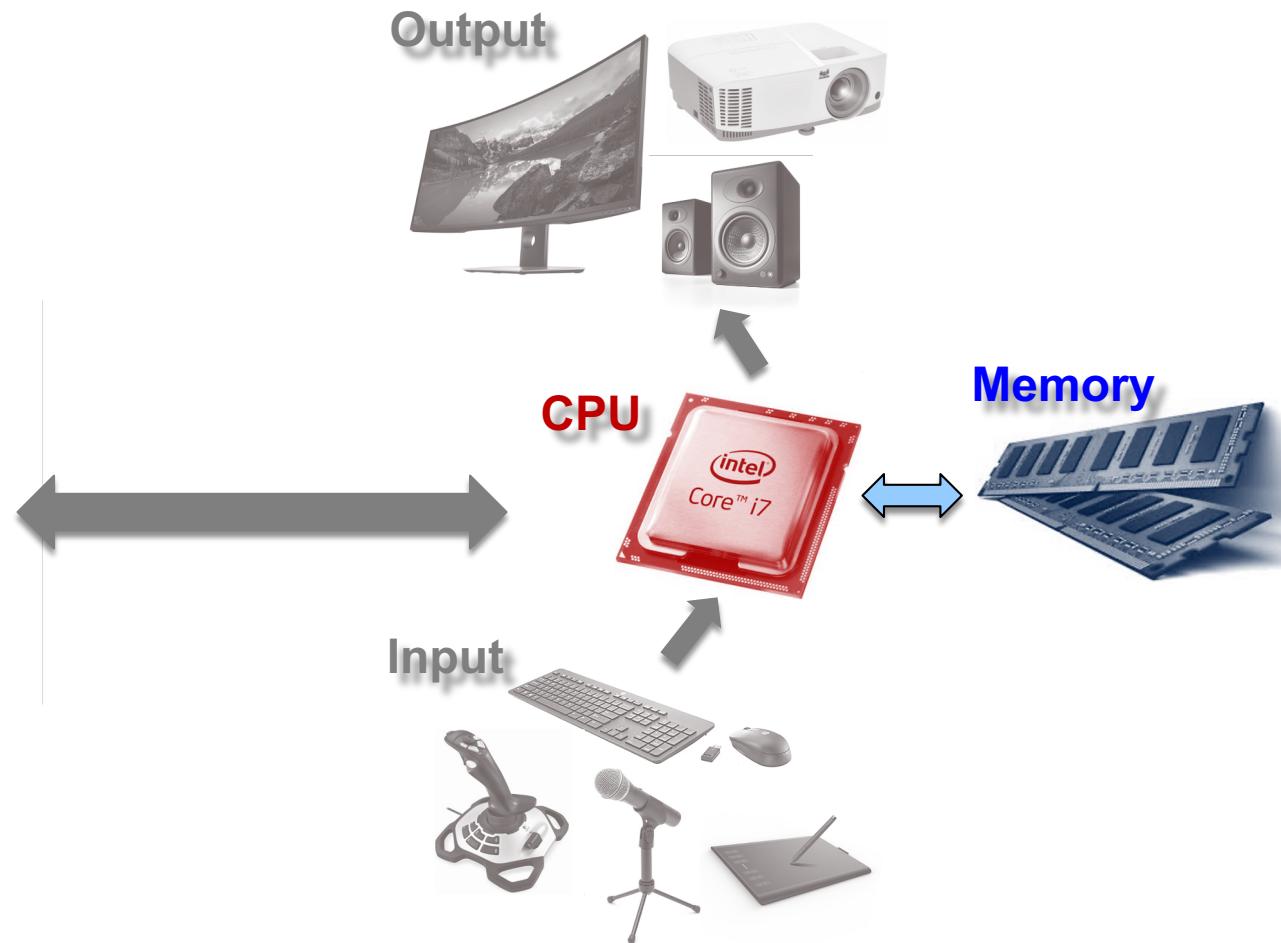
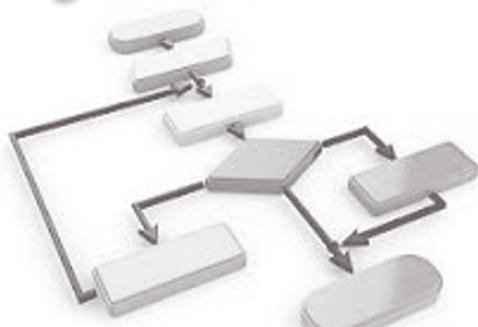
Overview

- This lecture
 - Memory system
 - Cache

Computer Organization

□ Classic components of a computing system

Algorithm



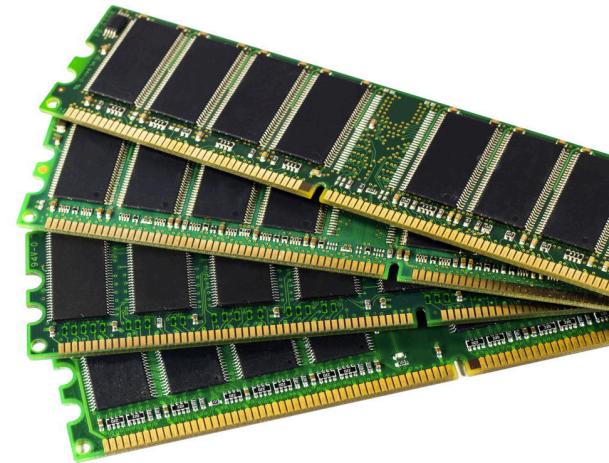
Memory System

- Data and instructions are stored on DRAM chips
 - DRAM has high bit density and low speed
 - An access DRAM may take about 300 processor cycles

Processor



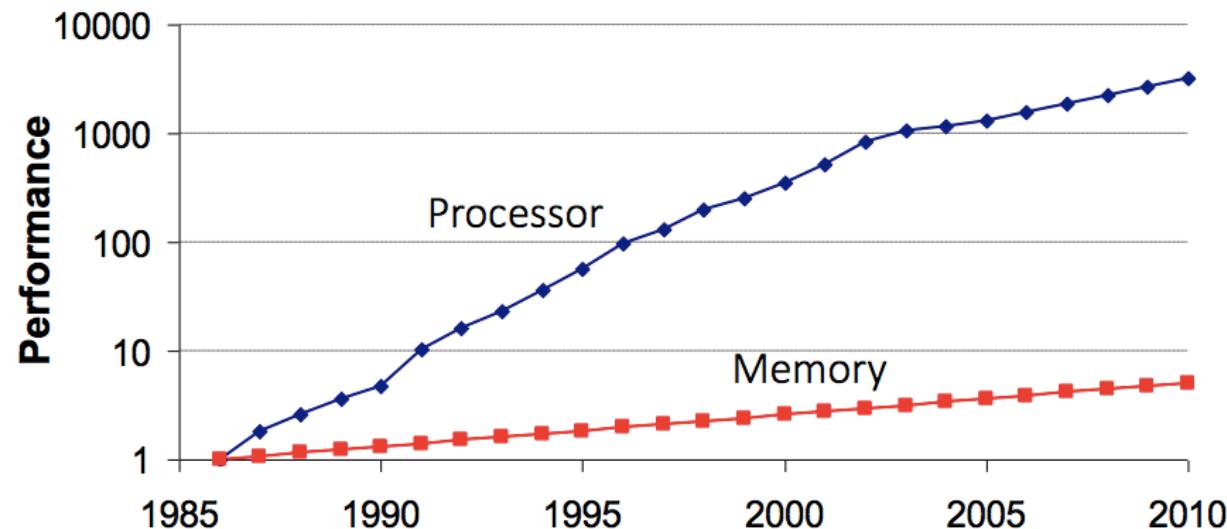
Memory



~300X

The Memory Wall

- Processor-memory performance gap increased over 50% per year
 - ▣ Processor performance historically improved ~60% per year
 - ▣ Main memory access time improves ~5% per year



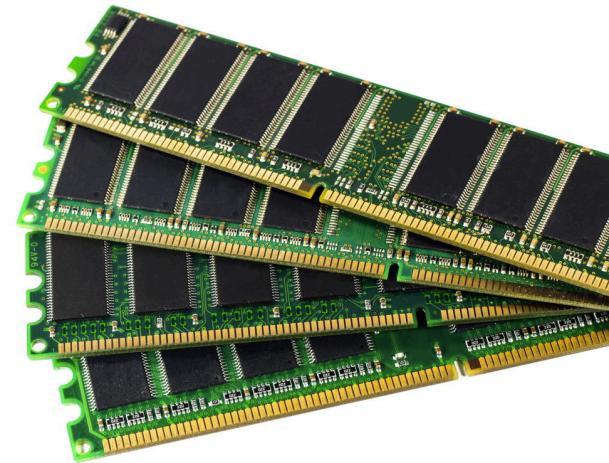
Memory System

- Data and instructions are stored on DRAM chips
 - ▣ DRAM has high bit density and low speed
 - ▣ An access DRAM may take about 300 processor cycles
- How to bridge the speed gap?

Processor



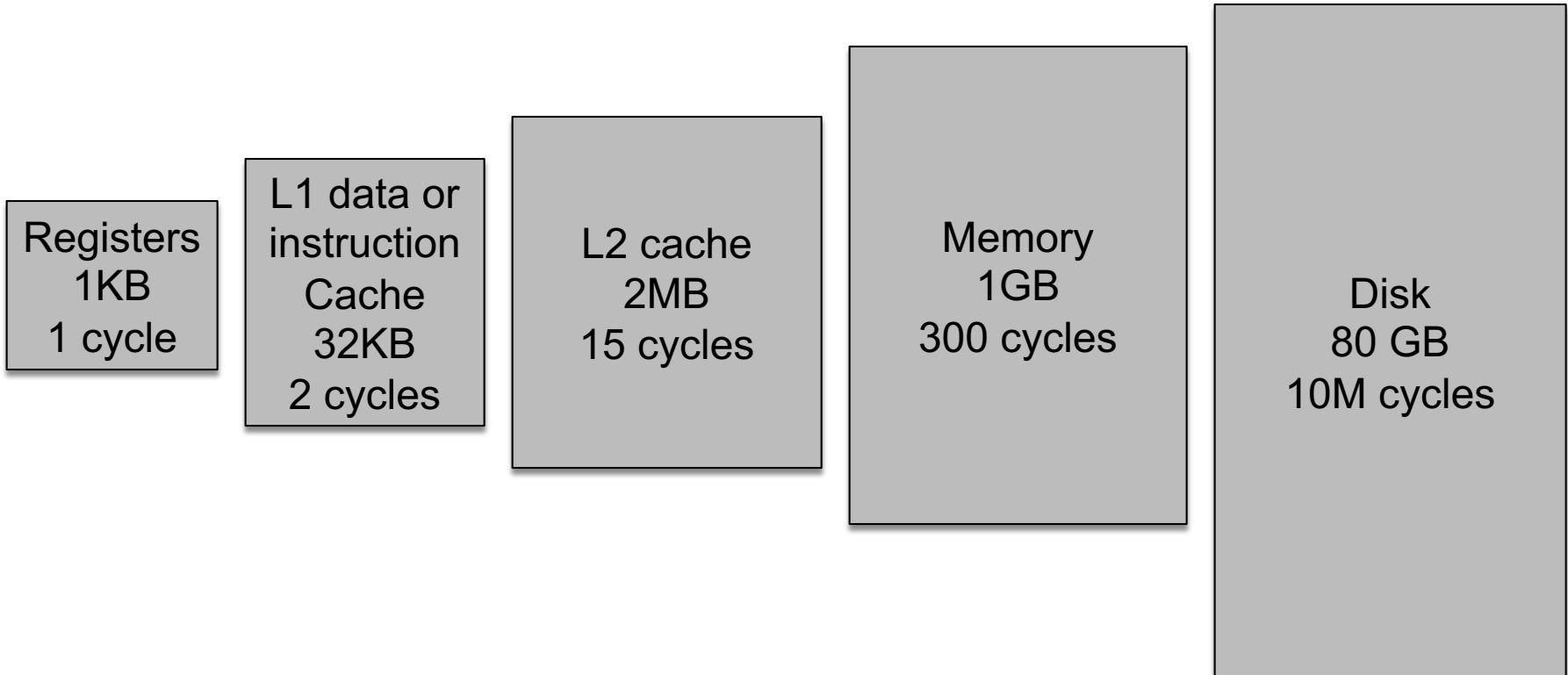
Memory



~300X

Memory Hierarchy

- The basic structure of a memory hierarchy.



Processor Cache

- Occupies a large fraction of die area in modern microprocessors

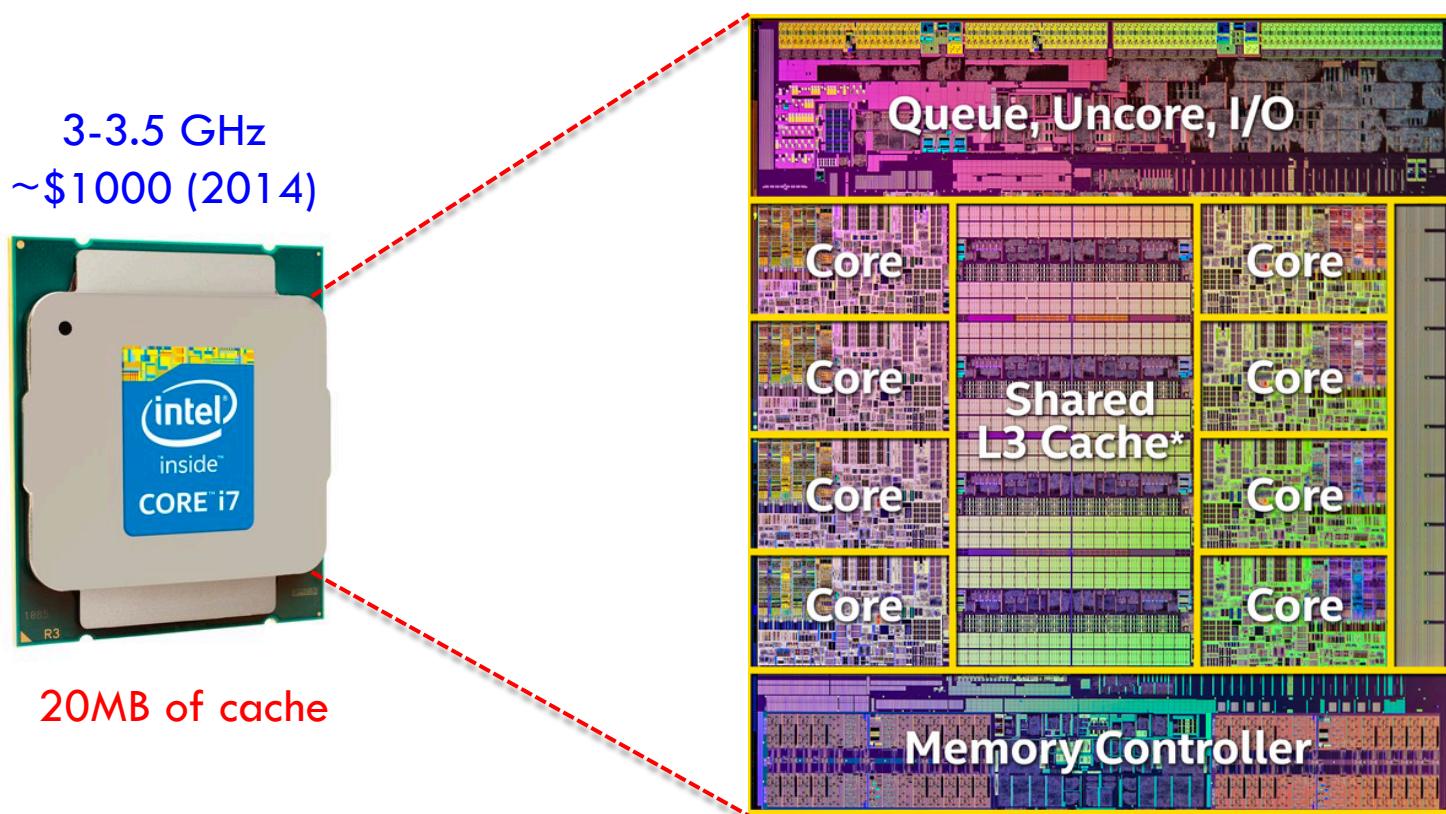
3-3.5 GHz
~\$1000 (2014)



Source: Intel Core i7

Processor Cache

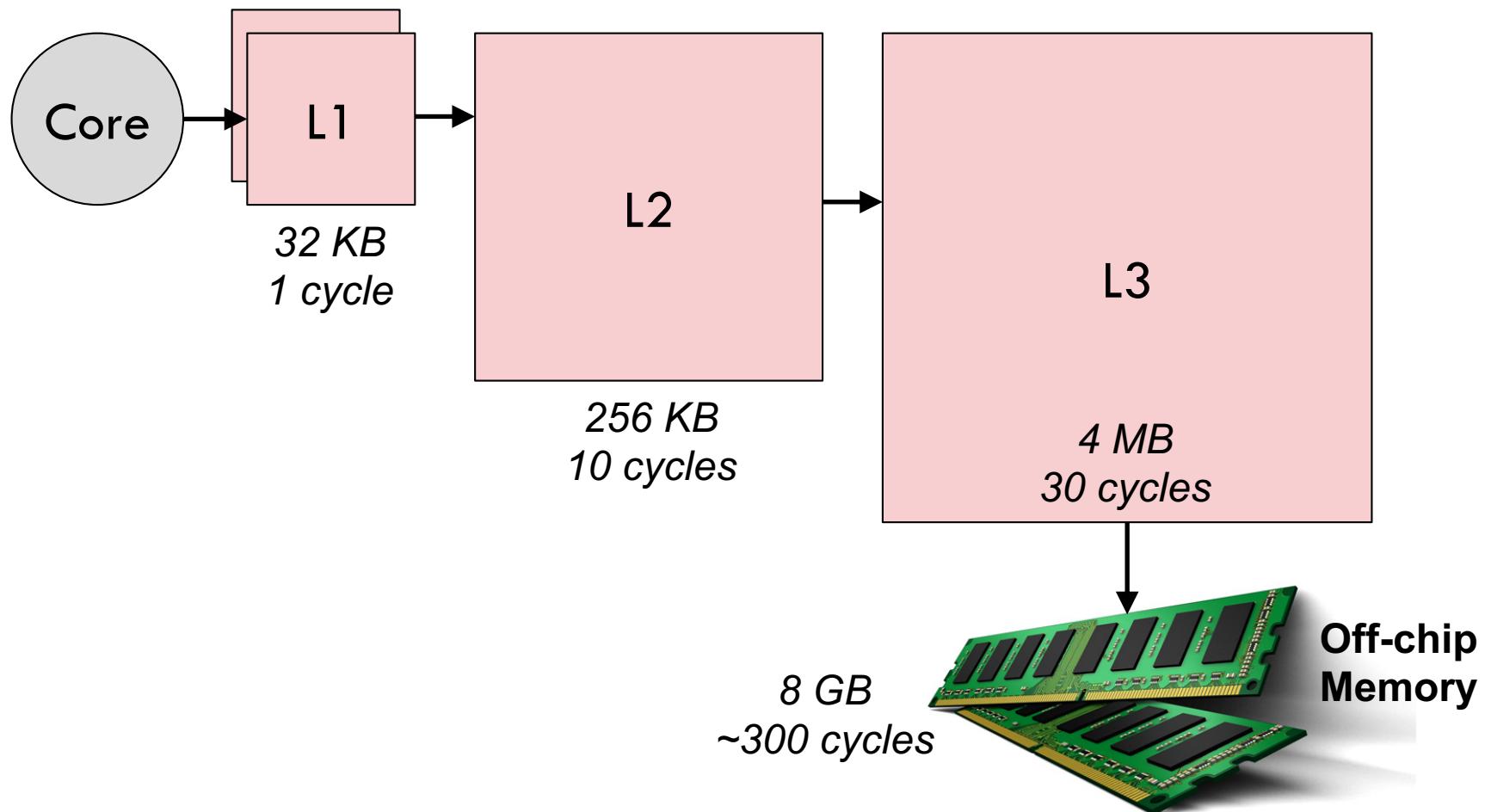
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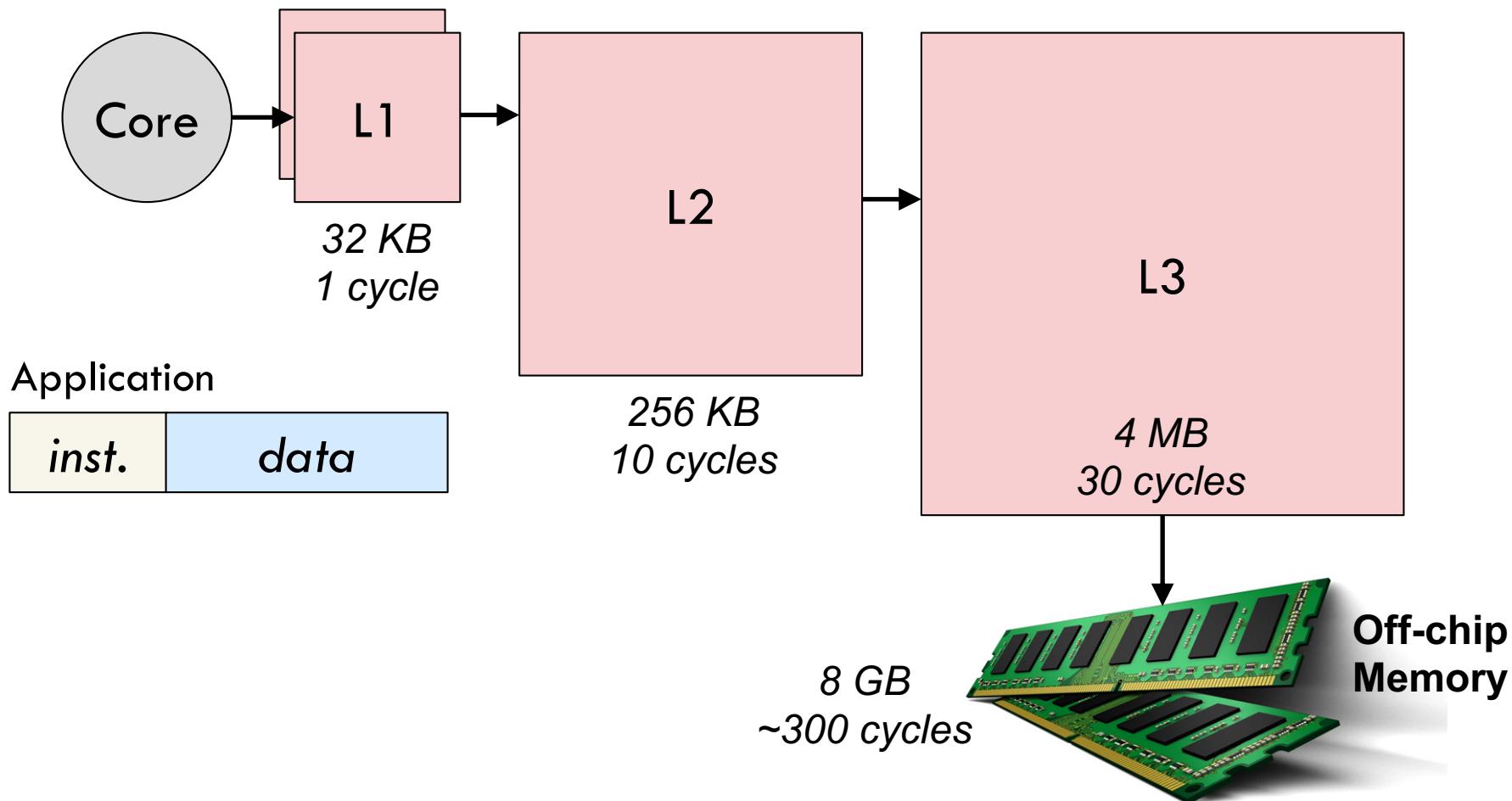
Cache Hierarchy

- Example three-level cache organization



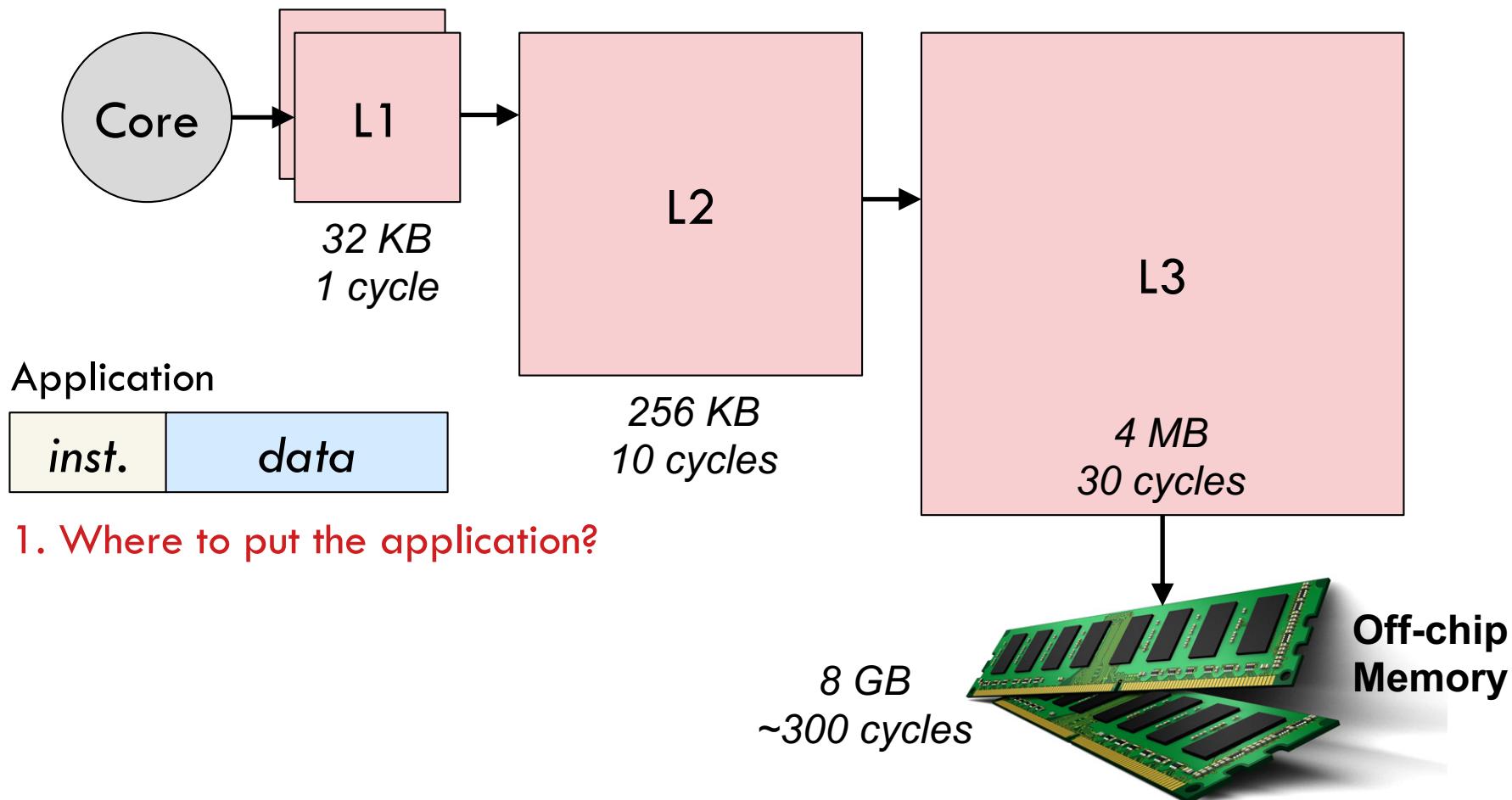
Cache Hierarchy

□ Example three-level cache organization



Cache Hierarchy

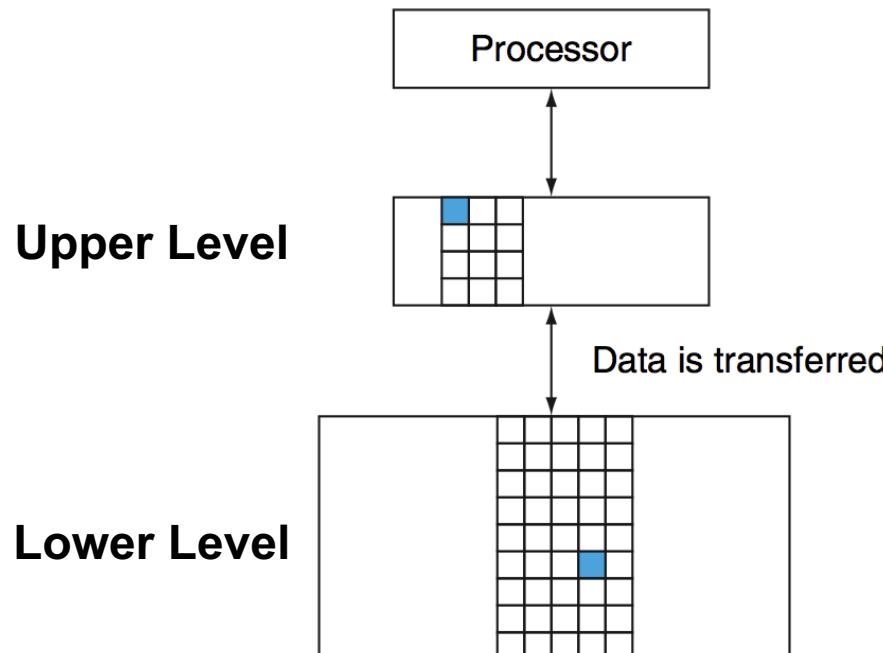
□ Example three-level cache organization



Memory Hierarchy

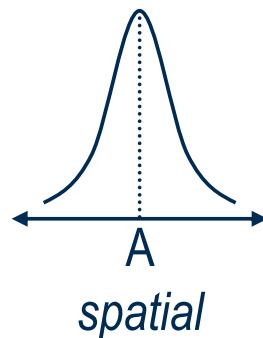
- The basic structure of a memory hierarchy.
- Multiple levels of the memory

Idea: keep important data closer to processor.

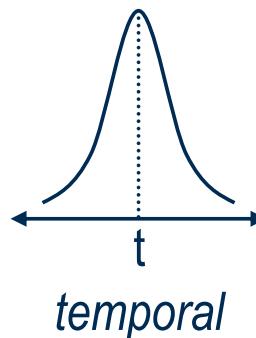


Principle of Locality

- Memory references exhibit localized accesses
- Types of locality
 - *spatial*: probability of access to $A+\delta$ at time $t+\varepsilon$ highest when $\delta \rightarrow 0$
 - *temporal*: probability of accessing $A+\varepsilon$ at time $t+\delta$ highest when $\delta \rightarrow 0$



spatial



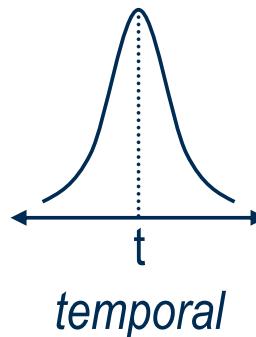
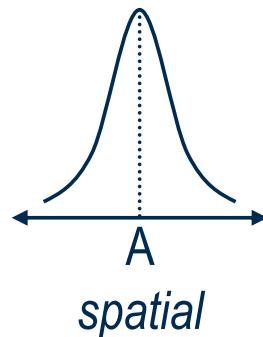
temporal

```
for (i=0; i<1000; ++i) {  
    sum = sum + a[i];  
}
```

Key idea: store local data in fast cache levels

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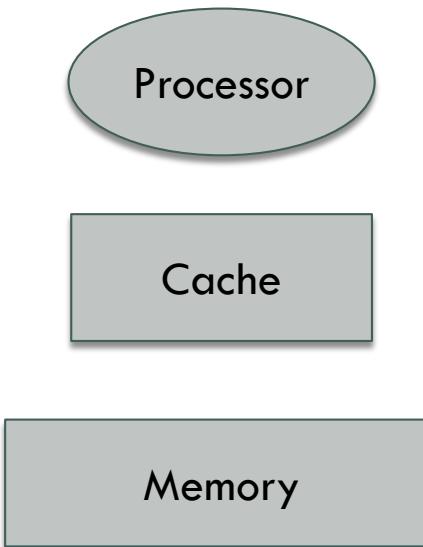


```
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temporal   spatial
```

Key idea: store local data in fast cache levels

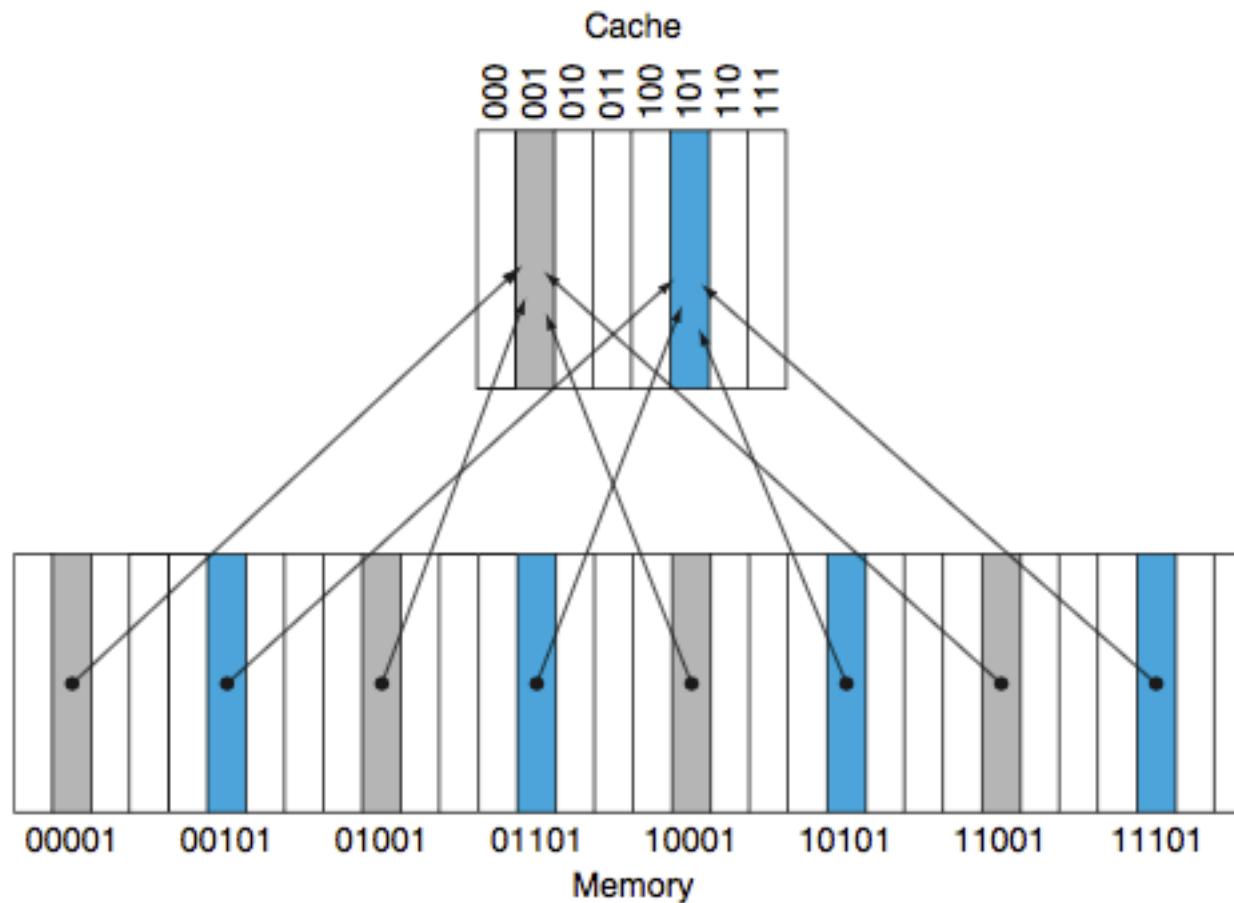
Cache Architecture

- Design principles
 - ▣ Temporal locality: if you used some data recently, you will likely use it again
 - ▣ Spatial locality: if you used some data recently, you will likely access its neighbors
- Cache terminology
 - ▣ Access time
 - ▣ Hit vs. miss
 - ▣ Miss penalty



Direct-Mapped Cache

Cache address



Direct-Mapped Cache

□ Cache lookup

