Title tbd

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Chapter I

2022

1 December

1.1 (11/28) Ok, let's give this a try

I want to read *Higher Topos Theory* (HTT). That book is > 700 pages, and fairly dense ones, so it's a bit of a project. On top of that, HTT isn't really a self-contained read (and that's sort of the nature of the subject). So I'll be drawing on lots of additional material. The subject seems well worth learning, even unavoidable at times (e.g., in the areas I am interested in), so this should be a productive little activity.

...

I need to cover some ground before actually opening HTT. First I need to think about *simplicial sets*. Their role in higher category theory is ubiquitous. They give us a combinatorial model for the homotopy theory of spaces, and also a model for ∞-categories. (On the list of things to do is make precise sense of those statements.) Some references are [Rie], [Mat], [Fri08], and `kerodon.net`.

Definition 1.1. Denote by $\mathbf{\Delta}$ the *simplex category*, defined to have...

- As objects, the ordered set $[n] := \{0 < 1 < \cdots < n\}$ for each $n \geq 0$; and
- As maps, the weakly order-preserving set maps.

Definition 1.2. A *simplicial set* is a contravariant functor $\mathbf{\Delta} \to \mathsf{Set}$. The *category of simplicial sets*, denoted sSet, is the functor category $\mathrm{Fun}(\mathbf{\Delta}^{op}, \mathsf{Set})$.

Notation 1.1. Let $X : \mathbf{\Delta}^{op} \to \mathsf{Set}$ be a simplicial set. We may write it $X_\bullet$, and denote by X_n the set $X([n])$. We call the elements of X_n the *n-simplices* of X.

Notation 1.2. We may write $\langle f_0, \ldots, f_n \rangle$ to denote the function $[n] \to [m]$ given by $n \mapsto f_n$.

Simplical sets are not just simplices. They carry additional structure, that arising from morphisms in $\mathbf{\Delta}$. We can give a simple description of $\mathbf{\Delta}$. This in turn gives some intuition for what a simplicial set "is."

Proposition 1.3 (The structure of $\mathbf{\Delta}$). *For each $n \geq 0$ and $0 \leq i \leq n$, define the*

$$i\text{-th face map } d^i : [n-1] \to [n] \text{ as } \langle 0, \ldots, \hat{i}, \ldots n \rangle, \text{ and the}$$
$$i\text{-th degeneracy map } s^i : [n+1] \to [n] \text{ as } \langle 0, \ldots, i, i, \ldots, n \rangle.$$

Every morphism in $\mathbf{\Delta}$ may be written as a composition of face and degeneracy maps. (Also, the face/degeneracy maps satisfy various relations, the simplicial identities; in fact $\mathbf{\Delta}$ is the category generated by those maps, subject to these identities.)

Thus a simplicial set $X_\bullet$ can be described as a collection of sets X_n (n-cells) together with face and degeneracy maps which satisfy the "simplicial identities." I should write more about how this notion arises from topology, in particular the singular complex. That in turn would be a good time to relate all this back to topology (nerves, geometric realization, ...) which is important.

1.2 (12/1) Why simplicial sets, simplicial complexes

I had stuff written here. But it was incomplete, and the "story" here is an aside I want to write about a bit more carefully at some point. I'm leaving this day blank for the time being.

1.3 (12/4) Basic structure in sSet

We need to make some terminology regarding / record examples of simplicial sets.

Definition 1.4. The *standard n-simplex* Δ^n is the simplicial set represented by $[n]$, i.e. $\Delta^n := \mathrm{Hom}_{\mathbf{\Delta}}(-, [n])$.

Definition 1.5. Let $X_\bullet, Y_\bullet$ be simplicial sets. We say $Y_\bullet$ is a *simplicial subset* of $X_\bullet$ if $Y_n \subseteq X_n$ and $Xf|_{Y_n} = Yf$ for every $n \geq 0$ and simplicial operator f. In other words, the action of operators on Y is the restriction of their action on X. In other words, $Y_\bullet$ is a subfunctor of $X_\bullet$.

Proposition 1.6. *Let $X_\bullet$ be a simplicial set. The Yoneda lemma asserts a bijection $\mathrm{Hom}_{\mathsf{sSet}}(\Delta^n, X_\bullet) \cong X_n$. Under this bijection, each n-cell $a \in X_n$ corresponds to a map $f_a : \Delta^n \to X_\bullet$ satisfying $f_a(\mathrm{id}_{[n]}) = a$.*

Definition 1.7. Let $X_\bullet$ be a simplicial set. By the above, we may identify its n-cells with maps $\Delta^n \to X_\bullet$. Call a cell $a \in X_n$ *degenerate* if it factors as $\Delta^n \to \Delta^m \to X_\bullet$ for some $m < n$. (See [Lur22, Tag 0011] for equivalent conditions.)

Proposition 1.8. *The standard simplex Δ^n has a unique non-degenerate n-simplex, that arising from $\mathrm{id}_{[n]}$. We may call this the generator of Δ^n.*

Definition 1.9 (Boundary of Δ^n)**.** Define a simplicial subset $\partial\Delta^n$, the *boundary of Δ^n*, by

$$(\partial\Delta^n)_k := \{\text{non-surjective maps } [k] \to [n]\} \subseteq \mathrm{Hom}_{\mathbf{\Delta}^{op}}([k],[n]).$$

Proposition 1.10. *The boundary of Δ^n is the maximal proper simplicial subset of Δ^n.*

Definition 1.11 (Horns in Δ^n)**.** For $0 \leq i \leq n$, define a simplicial subset Λ^n_i, the *i-th horn in Δ^n*, by

$$(\Lambda^n_i)_k := \{f \in \mathrm{Hom}_{\mathbf{\Delta}^{op}}([k],[n])) : f([k]) \cup \{i\} \neq [n]\}.$$

In other words, its cells are those maps "missing something besides i." A horn Λ^n_i is called *outer* if $i \neq 0, n$ and *inner* otherwise.

Any simplicial operator $f : [m] \to [n]$ factors through its image, i.e. we can uniquely write $f = f^{inj} f^{surj}$, a surjection followed by an injection. Furthermore, this is unique. We get the following.

Proposition 1.12. *Let $\sigma : \Delta^n \to X_\bullet$ be an n-cell of $X_\bullet$. Then σ factors uniquely as*

$$\Delta^n \xrightarrow{\alpha} \Delta^m \xrightarrow{\tau} X_\bullet,$$

Where α represents a surjection $[n] \to [m]$ and τ is not degenerate. Call m the dimension of the cell σ. (My notation, maybe poor, not that important.)

So, degenerate n-simplices are just non-degenerate simplices in a lower dimension (their "dimension"), trying to bite off more than they can chew.

Definition 1.13 (Skeleta)**.** Let $X_\bullet$ be a simplicial set. For $k \geq -1$, define a simplicial subset $\mathrm{sk}_k(X_\bullet)$, the *k-skeleton* of $X_\bullet$, by

$$(\mathrm{sk}_k(X_\bullet))_n := \{\, n\text{-simplices of } X_\bullet \text{ with dimension at most } k \,\}.$$

Remark 1.14. The face maps $d^i : [n-1] \to [n]$ induce maps $d^i : \Delta^{n-1} \to \Delta^n$ via post-composition. Now, consider an n-cell $a \in X_n$ and its representation $a : \Delta^n \to X_\bullet$. We have that $d_i(a) \in X_{n-1}$ is represented by ad^i.

1.4 (12/6) Colimits in/functors out of sSet

Today I want to understand part of Akhil's notes, about functors out of sSet. This is closely related to understanding colimits in sSet, by general theory for presheaf categories. So we also want to understand colimits in sSet. (And we should want to understand these regardless.) Let's go over this.

Here's a standard structure result for presheaf categories.

Proposition 1.15. *If a category C is small, then every presheaf on C is canonically the colimit of representable presheaves. In particular, every simplicial set is canonically the colimit of standard simplices.*

Proof. This is written out in Akhil's notes. I'll give the idea. Also see [Lur22, Remark 00X5].
Consider a presheaf $F : \mathsf{C}^{op} \to \mathsf{Set}$. We associate to F the category D_F with

- Objects: morphisms from represented presheaves to F, i.e. arrows $[-, X] \to F$; and
- Morphisms: morphisms between represented presheaves such that the obvious triangle commutes.

There is a functor $\phi_F : \mathsf{D}_F \to \mathsf{PShv}(\mathsf{C})$ which sends objects $[-, X] \to F$ to $[-, X]$. By construction, for each object $c \in \mathsf{D}_F$, there is a morphism $\phi_F(c) \to F$, and the diagram described by ϕ_F together with these morphisms commutes. We therefore have a distinguished morphism

$$\lim_{\mathsf{D}_F} \phi_F \to F.$$

This map turns out to be an isomorphism. □

Hereafter, denote by $\overline{\mathsf{C}}$ the category of presheaves on C.

Suppose D is cocomplete. We want to understand functors $\overline{F} : \overline{\mathsf{C}} \to \mathsf{D}$. The previous proposition says that objects in $\overline{\mathsf{C}}$ are colimits of representables. So, if $\overline{F}$ preserves colimits, then $\overline{F}$ is determined by $\overline{F}|_{\mathsf{C}}$, i.e. what it does to C (embedded via Yoneda). We've described an injection of sets

$$\mathrm{Fun}'(\overline{\mathsf{C}}, \mathsf{D}) \hookrightarrow \mathrm{Fun}(\mathsf{C}, \mathsf{D}). \tag{I.16}$$

Here, Fun' denotes the set of colimit-preserving functors.

Conversely, suppose given a functor $F : \mathsf{C} \to \mathsf{D}$. Does it extend along the Yoneda embedding to a functor $\overline{F} : \overline{\mathsf{F}} \to \mathsf{D}$? We can do something here, let me write it out:

(1) As above, for each presheaf $G : \mathsf{C}^{op} \to \mathsf{Set}$, consider it as a colimit of $\phi_G : \mathsf{D}_G \to \mathsf{C}$. (We can do this because it lands in represented functors.)

(2) This is 'functorial' in the following sense: a morphism $G \to H$ induces a functor $\mathsf{D}_G \to \mathsf{D}_H$ such that the obvious triangle commutes.

(3) Define a functor $\overline{F} : \overline{\mathsf{C}} \to \mathsf{D}$ by

$$\overline{F}(G) := \varinjlim_{\mathsf{D}_G} F \circ \phi_G.$$

This is a functor because of (2).

This functor $\overline{F}$ really extends F, i.e. the obvious diagram commutes. For suppose $G = [-, c]$; then D_G has a final object $[-, c] \to [-, c]$, therefore $\overline{F}(G) = \varinjlim_{\mathsf{D}_G} F \circ \phi_G = F(G)$.

Proposition 1.17. *Suppose given a functor $F : \mathsf{C} \to \mathsf{D}$ to a cocomplete category. Then the associated functor $\overline{F} : \overline{\mathsf{C}} \to \mathsf{D}$ constructed above preserves all colimits. In fact, $\overline{F}$ is a left adjoint. The right adjoint to $\overline{F}$ is the functor defined by*

$$\mathsf{D} \ni d \mapsto (c \mapsto \mathrm{Hom}_{\mathsf{D}}(Fc, d)) \in \overline{\mathsf{C}}.$$

Proposition 1.18. *Suppose given a functor $F : \mathsf{C} \to \mathsf{D}$ to a cocomplete category. Then the mapping $F \mapsto \overline{F}$ describes a bijection of sets*

$$\mathrm{Fun}(\mathsf{C}, \mathsf{D}) \xrightarrow{\sim} \mathrm{LeftAdjoints}(\overline{\mathsf{C}}, \mathsf{D}).$$

The proofs are short and formal.

Corollary 1.19. *If a functor $\overline{F} : \overline{\mathsf{C}}^{op} \to \mathsf{Set}$ takes colimits to limits, then it is representable.*

Proof. Suppose as given $\overline{F}$. By the above, it is left adjoint to some $\overline{G} : \mathsf{Set}^{op} \to \mathsf{C}$. Define $f := \overline{G}(\{pt\})$, the image of the terminal object in Set^{op}. I claim that f represents F. (Insert short, formal proof; it's in Akhil's notes.) □

1.5 (12/23) The singular complex and geometric realization

Finals are over and I've had some time to wind down at home. Last time I worked through part of Akhil's notes about functors out of sSet. (Emily Riehl talks about something similar in her notes, but I have not gotten through those, so let me say nothing about that right now.)

Next I want to relate Top, Cat, and sSet. This is the backdrop for the idea that higher categories "bridge" topology/homotopy theory and ordinary categories. Today I'll go over the relation of sSet to Top, by which I mean the adjunction

the geometric realization functor ⊣ the total singular complex functor.

There are a few ways to introduce this adjunction. Lurie, Charles, and Akhil each do it differently. As a matter of taste, I prefer Akhil's approach. (Possibly related: [Lur22, Tag 002D].) Lurie's approach has some important ideas behind it too, I think, but that is overruled because I am feeling sleepy today.

Definition 1.20. Define a functor $|-| : \mathbf{\Delta} \to \mathsf{Top}$ as follows.

- Each object $[n]$ is sent to the *topological n-simplex* $\Delta^n_{top} \subseteq \mathbb{R}^{n+1}$, defined as those $(t_0, \ldots, t_n)$ satisfying $t_i \geq 0$ and $\sum t_i = 1$ and given the subspace topology.
- Each morphism $f : [m] \to [n]$ is sent to the map

$$(t_0, \ldots, t_n) \mapsto (u_j), \quad u_j = \sum_{i : f(i) = j} t_i.$$

Definition 1.21 (Geometric realization)**.** Since Top is cocomplete, according to Proposition 1.17 and 1.18 the functor of Definition 1.20 extends uniquely to a left adjoint $|-| : \mathsf{sSet} \to \mathsf{Top}$. We call this *geometric realization*.

In fact, as in Proposition 1.18, we know the right adjoint to geometric realization. It sends a space X to the simplicial set $[n] \mapsto \mathrm{Hom}_{\mathsf{Top}}(\,|[n]| = \Delta^n_{top}, X)$. This is an important construction I maybe should have defined earlier.

Definition 1.22 (Singular complex)**.** Let X be a space. Denote by $\mathrm{Sing}(X)_\bullet$ the simplicial set given as follows.

- The n-cells are the continuous maps $\Delta^n_{top} \to X$, and
- Each simplicial operator $f : [m] \to [n]$ acts by precomposing with the continuous map

$$\Delta^m_{top} \to \Delta^n_{top}, \quad (t_j) \mapsto (u_j = \sum_{f(i) = j} t_i).$$

We call $\mathrm{Sing}(X)_\bullet$ the *singular complex* of X. We define a functor $\mathrm{Sing}(-)_\bullet : \mathsf{Top} \to \mathsf{sSet}$ in the obvious way.

Proposition 1.23. *Prior discussion tells us that geometric realization $|-| : \mathsf{sSet} \to \mathsf{Top}$ is left adjoint to the singular complex functor.*

Proposition 1.24. *Since geometric realization is a left adjoint, it commutes with colimits. Furthermore, geometric realization commutes with finite limits of compactly generated spaces.*

We will see later that this adjunction is homtopically well-behaved.

I next want to describe geometric realization. We already have the general construction laid out for us by Proposition 1.17 and the preceding discussion. Given $X_\bullet$, we will form the *category of simplices*, also called its *category of elements*, whose elements are the morphisms $\Delta^n \to X_\bullet$ (i.e., the cells of $X_\bullet$), and we take the colimit of $|-|$ restricted to this subcategory. (This is not circular since we are really applying the "baby" geometric realization to the simplex category, Yoneda embedded.)

Definition 1.25. Given $X_\bullet \in \mathsf{sSet}$, its *category of simplices* or *category of elements* has as objects all morphisms $\Delta^n \to X_\bullet$ for every n, and morphisms the maps $\Delta^m \to \Delta^n$ making the obvious diagram commute. We write this category $\mathsf{el}(X)$.

This category of elements/simplices $\mathsf{el}(X)$ is precisely the category D_X described on (12/6) with $\mathsf{C} = \mathbf{\Delta}$. (Lurie writes this $\mathbf{\Delta}_X$.) Also as noted there, there is a natural functor $\phi_X : \mathsf{el}(X) \to \mathsf{sSet}$. Geometric realization is by definition the colimit

$$|X| \cong \varinjlim_{\mathsf{el}(X)} |-| \circ \phi_X.$$

Here, we are thinking of the "baby" geometric realization defined only on $\mathbf{\Delta}$.

Remark 1.26. General machinery gave us geometric realization. I think there are a few things worth saying about this, but I don't totally know what. I'll leave this remark here as a "to-do." Some possibly related keywords and references: "Grothendieck construction," "Kan extension," nLab, [Rie, §4], and Subsection 01Q7.

Chapter II

2023

1 January

1.1 (1/23) Plans have changed, nerves of categories

I've gone radio silent for a month. One big reason why is that I am busy this semester. Another is that some mutuals want to organize a reading group/seminar similar, but not identical to, what I've been trying to do here, and I may join them. Maybe the biggest difference is that they want to focus on Charles' quasicategory notes (under Charles' supervision).

This will probably mean repeating myself a bit while I change tracks to Charles' notes.

In any case, I want to talk about the nerve of a category. This is part of the basic "Spaces, categories, and simplicial sets" picture. In particular, the nerve of a category is a simplicial set encoding that category.

Definition 1.1. Let C be a category. Define a simplicial set NC, the *nerve* of C to have as cells $(NC)_n := \mathrm{Hom}_{\mathsf{Cat}}([n], C)$ and so that operators $f : [m] \to [n]$ act by precomposition. This defines the *nerve functor* $N : \mathsf{Cat} \to \mathsf{sSet}$.

Here's a feel for the structure of a nerve. The n-cells of $N\mathsf{C}$ may be canonically identified with the set of length n tuples of composable arrows in C. The 0-cells in particular may be identified with objects of C. An operator $f : [m] \to [n]$, or in Charles' notation $\langle f_1, \ldots, f_m \rangle$, acts by taking n-strings of composable arrows and "collapsing edges" by composing the arrows, those collapsed edges being determined by the f_i. (At least that's how I try to think about it. I think that's correct. UPDATE: Yes this is correct, see Charles' notes, Proposition 4.4.) For instance, $\langle 0, 2 \rangle^*$ takes a pair of composable arrows $(f, g) \in (NC)_2$ and sends them to their composite $gf \in (NC)_1$. See also Charles notes, p. 13.

Now we ask a nascent question: can we characterize nerves of categories?

Proposition 1.2. *Let X be a simplicial set. For $n \geq 2$, consider the function*

$$\phi_n : X_n \to \{(g_i) \in (X_1)^n : g_j\langle 1 \rangle = g_{j+1}\langle 0 \rangle \text{ for all } i\}$$

(The latter set being the collection of "n-paths" of 1-cells in X) Which acts by $a \mapsto (a\langle 0, 1 \rangle, \ldots, a\langle n-1, n \rangle)$. These ϕ_n are bijections for all $n \geq 2$ if and only if X is the nerve of a category.

Maybe a lazy way to digest this is that "a simplicial set is the nerve of a category iff, thinking of n-cells as length n strings of arrows, their 1-*dimensional structure* exactly reflects the structure that should arise from the existence and uniqueness of composites."

Proposition 1.3. *The nerve functor $N : \mathsf{Cat} \to \mathsf{sSet}$ is fully-faithful. That is, morphisms of nerves $NC \to ND$ correspond exactly to functors $C \to D$.*

1.2 (1/26) Spines

We characterized nerves as those simplicial sets whose n-cells were exactly determined by the collection of length n strings of "composable" 1-cells, in the obvious way. This captures the existence and uniqueness of composites for morphisms in a category. We can go about this characterization a bit more systematically.

Definition 1.4. Let $n \geq 0$. The *spine* of the standard n-simplex Δ^n is the simplicial subset defined by

$$(\mathrm{Spine}^n)_k := \{\langle f_0, \ldots, f_k \rangle : [k] \to [n] : f_k \leq f_0 + 1\} \subseteq \Delta^n_k.$$

Informally, the spine is the set of vertices of Δ^n together with the arrows between adjacent vertices (considered with their total ordering).

Proposition 1.5. *Let X be a simplicial set. For every $n \geq 0$, the map*

$$\mathrm{Hom}(\mathrm{Spine}^n, X) \to \{(a_i) \in (X_1)^n : a_i\langle 1 \rangle = a_{i+1}\langle 0 \rangle\} \tag{II.6}$$

Given by $f \mapsto (f\langle 0, 1 \rangle, \ldots, f\langle n-1, n \rangle)$ is a bijection.

Pictorially, I think this is obvious. Here's a clean proof.

Proof. One point we need: we previously talked about colimits in sSet. Or at least I intended to. Here's the main fact: *a colimit of simplicial sets X_α exists and has as its n-cells the colimit of the n-cells of the X_α.* This is true for presheaves in general; we say their (co)limits are "computed objectwise."

Another point we need: here's a definition. Suppose S is a totally ordered set. We denote by Δ^S the simplicial set having $(\Delta^S)_n := \{\text{order-preserving maps } [n] \to S\}$. If S is finite and nonempty, there is a unique isomorphism $\Delta^{|S|-1} \cong \Delta^S$. In the case that $S \subseteq [n]$, this is a good way to notate subcomplexes of Δ^n.

Here's a fact I won't prove: *given a subcomplex $K \subseteq \Delta^n$, writing A for the poset of $S \subseteq [n]$ such that $\Delta^S \subseteq K$, the canonical map $\varinjlim_{S \in A} \Delta^S \to K$ is an isomorphism.*

Finally, our proposition: in the case that $K = \mathrm{Spine}^n$, the poset A consists of sets of the form $\{j\}$ and $\{j+1\}$, and we have that $\varinjlim_{S \in A} \Delta^S \cong \mathrm{Spine}^n$. Now:

$$\mathrm{Hom}(\mathrm{Spine}^n, X) \cong \mathrm{Hom}(\varinjlim_{S \in A} \Delta^S, X) \cong \lim_A \mathrm{Hom}(\Delta^S, X).$$

See that the latter set is precisely the RHS of (II.6). $\square$

Maybe the key observation is that Δ^n is "generated" precisely by the arrows of Spine^n. (Make this formal? Say this better? Well, this is how I think about it.)

Proposition 1.7. *A simplicial set X is the nerve of some category if and only if for each $n \geq 2$, every morphism $f : \mathrm{Spine}^n \to X$ extends uniquely along the inclusion $\mathrm{Spine}^n \hookrightarrow \Delta^n$.*

Proof. The unique extension condition is equivalent to bijectivity of the restriction $\mathrm{Hom}(\Delta^n, X) \to \mathrm{Hom}(\mathrm{Spine}^n, X)$. By Proposition 1.5, the latter set is isomorphic to

$$\{(a_i) \in (X_1)^n : a_i\langle 1 \rangle = a_{i+1}\langle 0 \rangle\}.$$

Then the desired result is immediate considering Proposition 2.34. $\square$

1.3 (1/30) Inner Horns

Recall that for each n and $0 \leq i \leq n$ we defined the *i-th horn* $\Lambda^n_k \subseteq \Delta^n$ to have as k-cells those cells $f : [k] \to [n]$ of Δ^n which "miss something other than i," i.e. those satisfying $f([n]) \cup \{i\} \neq [n]$. If $j \neq 0, n$ then we called Λ^n_j an *inner horn*.

Drawing some pictures of horns and thinking of 1-cells as arrows, we think of inner horns as those collections of arrows that "should be composable." Similar to how we handled spines, we may think to characterize nerves as those simplicial sets whose inner horns have unique extensions. By our analogy comparing inner horns to composable arrows, this unique extension condition is analogous the existence and unqiueness of composites.

Theorem 1.8. *A simplicial set X is the nerve of some category if and only if for every $n \geq 2$ and inner horn Λ^n_j, every morphism $\Lambda^n_j \to X$ extends uniquely along $\Lambda^n_j \hookrightarrow \Delta^n$.*

Proof. Charles gives a full proof on p. 21 of his notes. The "only if" direction is not complicated. For the "if" direction:

- We can construct a category C whose nerve realizes X explicitly. The objects and morphisms are specified by X_0 and X_1.
- Existence *and* uniqueness of fillers are necessary for the existence and uniqueness of composites.
- Existence of a filler for Λ^3_1 or Λ^3_2 is necessary for the associative law to hold.
- Then one exhibits an isomorphism $X \to NC$.

$\square$

2 February

2.1 (2/4) Quasicategories

A *quasicategory* or *∞-category* is a simplicial set X such that every inner horn $\Lambda^n_j \to X$ has a filler (i.e. an extension along $\Lambda^n_j \hookrightarrow \Delta^n$.) We have shown that ordinary categories are precisely those quasicategories with *unique* horn extensions.

Definition 2.1. Some terminology. Let X, Y be quasicategories.

- *Objects* of $X := X_0$.
- *Morphisms* of $X := X_1$.
- *Identity morphism* of $x \in X_0 := x\langle 0, 0\rangle$.
- *Products of quasicategories* are just their products as simplicial sets.
- *Coroducts of quasicategories* are just their products as simplicial sets.
- A *morphism of quasicategories* $X \to Y$ is a map of simplicial sets.
- A *Natural transformation* $f_0 \Longrightarrow f_1$ of functors $f_0, f_1 : X \to Y$ is a map of simplicial sets $\phi : X \times \Delta^1 \to Y$ such that $\phi|_{X\times\{i\}} = f_i$.

It's a fact to be proven that the (co)product of quasicategories is again a quasicategory. Here's more terminology.

Definition 2.2. Let X be a simplicial set. Let $\sim$ denote the equivalence relation on the set $\coprod X_n$ of cells of X generated by the relation which identifies a cell a with any other cell of the form af for some simplicial operator f. A *connected component* of X is an equivalence class of $\sim$. We write $\pi_0 X$ for the set of equivalence classes. A simplicial set is called *connected* if $\pi_0 X$ is a singleton.

Proposition 2.3. *Let X be a simplicial set and suppose x, y are cells in the same connected component of X, i.e. $x = yf$ for some $f : [m] \to [n]$. If $F : X \to Y$ is a map of simplicial sets, then $F(yf) = F(y)Y(Xf)$. The latter is $F(x)$ by hypothesis, thus $F(x) \sim F(y)$. So morphisms $F : X \to Y$ induce maps $\pi_0 X \to \pi_0 Y$ on connected components.*

Proposition 2.4. *The induced map $\pi_0(X \times Y) \to \pi_0(X) \times \pi_0(Y)$ is a bijection.*

2.2 (2/6) Sub, opposite quasicategories

Definition 2.5. Let C be a quasicategory. A subcomplex $C' \subseteq C$ is called a *subcategory* of C if for all $n \geq 2$ and $0 < k < n$, every $f : \Lambda^n_k \to C$ such that $f(\Lambda^n_k) \subseteq C'$ extends "into C'," i.e. extends to a map $f : \Delta^n \to C'$.

It is clear that subcategories of quasicategories are quasicategories.

Definition 2.6. Let C be a simplicial set. A simplicial subset $C' \subseteq C$ is called *full* if

- For every cell $\sigma : \Delta^n \to C$ such that for every $0 \leq i \leq n$ the vertex $\sigma(i) \in C$ belongs in C', the cell σ belongs in C'.

If C is a quasicategory and C' is a full subcomplex, then it is a quasicategory. In this case we say C' is a *full subcategory*.

Next we can define the opposite of a quasicategory. In ordinary categories, we do this by reversing composition. We can do something similar once we identify an involution on the simplex category $\mathbf{\Delta}$.

Definition 2.7. Define an involution functor $\mathrm{op} : \mathbf{\Delta} \to \mathbf{\Delta}$ as follows.

- It acts as the identity on objects.
- It sends the morphism $\langle f_0, \dots, f_m\rangle : [m] \to [n]$ to its "reverse" $\langle n - f_m, \dots, n - f_0\rangle : [m] \to [n]$.

Definition 2.8. Let $X : \mathbf{\Delta}^{op} \to \mathrm{Set}$ be a simplicial set. We define the *opposite simplicial set* as $X^{op} := X \circ \mathrm{op}$.

One sees that $(\Delta^n_j)^{op} \cong \Delta^n_{n-j}$ and that $(NC)^{op} = N(C^{op})$. The former fact ensures that opposites of quasicategories are quasicategories. The latter ensures that this notion of opposites restricts to the usual 1-categorical notion.

2.3 (2/7) Examples of ∞-categories

Example 2.1. The nerve NC of a category C is a quasicategory. This is immediate by our characterization of nerves (Theorem 1.8).

Example 2.2. The singular complex $\mathrm{Sing}(X)$ of a space X is a quasicategory. In fact, we can say a little bit more. Denote by $(\Lambda^n_j)_{top}$ the *topological horn*, defined as you might expect:

$$(\Lambda^n_j)_{top} := \{t \in \Delta^n : t_i = 0 \text{ for some } i \in [n]/\{j\}\}.$$

It is clear that for any j, the simplex Δ^n_{top} retracts onto $(\Lambda^n_j)_{top}$. Thus by precomposing with the retract we get, *for every* j, an inverse to the restriction

$$\mathrm{Hom}(\Delta^n, \mathrm{Sing}X) \to \mathrm{Hom}(\Lambda^n_j, \mathrm{Sing}X).$$

In other words, we can fill every $\Lambda^n_j \to \mathrm{Sing}X$ via the retract. This shows that $\mathrm{Sing}(X)$ is a quasicategory. In fact, we've shown all horns fill, not just inner horns. We call such simplicial sets *Kan complexes*.

Example 2.3. Let A be an abelian group and $d \geq 0$ an integer. In spaces, the *Eilenberg-Maclane spaces* $K(A,d)$ represent $H^d(-;A)$. We will define an analogous simplicial set $K = K(A,d)$, the *Eilenberg-Maclane objects* in sSet, like so.

- An element of K_n is a collection $(a_\delta \in A)_\delta$, where the index δ occurs over all operators $\delta : [d] \to [n]$, so that
  - If δ is not injective, $a_\delta = 0$; and
  - For each operator $\gamma : [d+1] \to [n]$ we have $\sum_{j=0}^{d+1} (-1)^j a_{\delta d^j} = 0$.
- For each operator $f : [m] \to [n]$ and $a \in K_n$, we define $(af)_\delta := a_{f\delta}$.

These $K(A,d)$'s are ∞-categories. In fact, they are simplicial abelian groups, which are always Kan complexes. In sSet, they represent *normalized d-cocycles with values in A*. (See Charles' notes, p. 29.)

Example 2.4. There is a simplicial set of ordinary categories, denoted Cat_1. We define it like so.

- Each n-cell $(\mathsf{Cat}_1)_n$ is the data of $(C_i, F_{ij}, \zeta_{ijk})$ where
  - For each $i \in [n]$, C_i is a category,
  - For each $i \leq j$ in $[n]$, $F_{ij} : C_i \to C_j$ is a functor, and
  - For each $i \leq j \leq k$ in $[n]$, $\zeta_{ijk} : F_{ik} \to F_{jk}F_{ij}$ is a natural isomorphism,
  - And furthermore, these data are subject to certain basic properties (e.g. $F_{ii} = \mathrm{id}_{C_i}$).
- Each operator $f : [m] \to [n]$ acts on an n-cell $(C_i, F_{ij}, \zeta_{ijk})$ by composing with the indices.

The simplicial set Cat_1 is an ∞-category . Let's discuss fillers.

- A 2-horn $\Lambda^2_1 \to \mathsf{Cat}_1$ is the data of functors $C_0 \xrightarrow{F_{01}} C_1 \xrightarrow{F_{12}} C_2$. An extension is the data of a functor $F_{02} : C_0 \to C_2$ and a natural isomorphism $\zeta_{012} : F_{12}F_{01} \Longrightarrow F_{02}$. An obvious *but not necessarily unique* candidate is $F_{02} := F_{12}F_{01}$.
- The data of a 3-horn $\Lambda^3_1 \to \mathsf{Cat}_1$ is a bit of a picture. A filler amounts to finding a natural isomorphism to fill the following diagram.

$$\begin{tikzcd} F_{03} \arrow[r, "\zeta_{013}"] \arrow[d, dashed] & F_{13}F_{01} \arrow[d, "\zeta_{123}F_{01}"] \\ F_{23}F_{02} \arrow[r, "F_{23}\zeta_{012}"'] & F_{23}F_{12}F_{01} \end{tikzcd}$$

  We can always find this and it is unique, since we required the ζ's to be natural *isomorphisms*.

2.4 (2/8) The fundamental category of a simplicial set

The fundamental groupoid $\pi_{\leq 1}X$ of a space X can be recovered from its singular complex $\mathrm{Sing}(X)$. We will recast this construction $\pi_{\leq 1}X$ for any ∞-category. The result will no longer be a groupoid in general (it will only be so for Kan complexes, I think). Let's see how far we get.

First we will look at a certain construction for all simplicial sets. By its definition, it's essentially a left adjoint to the nerve functor.

Definition 2.9. Let X be a simplicial set. A *fundamental category of X* is a category hX and a map $\alpha : X \to N(hX)$ such that for every nerve NC, the restriction

$$\alpha^* : \mathrm{Hom}(N(hX), NC) \to \mathrm{Hom}(X, NC)$$

Is a bijection. This characterizes hX up to unique isomorphism, if it exists. (It always does.)

Proposition 2.10. *Every simplicial set has a fundamental category.*

Proof. Charles sketches this on p. 30. The objects of our category are X_0. The morphisms are (those generated by) the edges X_1, where we identify composites according to the 2-cells of X. So we turn X into a category in the most obvious way, flattening the higher-categorical structure in the process. The map $\alpha : X \to N(hX)$ is the one you'd expect. □

Proposition 2.11. *The fundamental category describes a functor h : $\mathsf{sSet} \to \mathsf{Cat}$, and this functor is left adjoint to the nerve functor N.*

2.5 (2/9) Homotopy for ∞-categories

Now we start down a long, dark path. Neither of these adjectives matter up-to-homotopy, however.

Definition 2.12. Let C denote an ∞-category and let $f, g : x \to y$ be two morphisms between objects x, y in C. A *homotopy from f to g* is a 2-cell $a \in C_2$ such that $a_{01} = f$, $a_{12} = \mathrm{id}_y$, and $a_{02} = g$.

Proposition 2.13. *The homotopy relation is is an equivalence relation on $\mathrm{Hom}_C(x, y)$, i.e. the set of edges i with $i_0 = x$ and $i_1 = y$. Thus we may unamibguously say maps are (or are not) homotopic and speak of homotopy classes.*

Remark 2.14. The existence of inner horn extensions is necessary for this relation to be symmetric and transitive. So quasicategories stand out amongst simplicial sets as those having a good notion of homotopy.

Proposition 2.15. *Maps f, g are homotopic in C if and only if they are homotopic in C^{op}.*

It's maybe a little weird that "f homotopic to g" is a slightly asymmetric definition, in that even if f is homotopic to g, the data of a homotopy does not suffice to get a homotopy from f to g in C^{op}. I don't think this matters much, in light of the previous proposition. Lurie also gives an alternate, symmetric notion of homotopy to address this point [Lur22, Tag 00V0].

Suppose as given $f \in \mathrm{Hom}_C(x, y)$ and $g \in \mathrm{Hom}_C(y, z)$. We say an edge $h \in \mathrm{Hom}_C(x, z)$ is a *composite* of (g, f) if there exists a 2-cell a such that (what you expect).

Proposition 2.16. *Composition respects the homotopy relation on morphisms. Thus, composites are unique up to homotopy.*

Definition 2.17. Let C be an ∞-category . Its *homotopy category hC* is the category having as objects C_0 and as morphisms the homotopy classes of morphisms of C.

Now we have defined the *fundamental category of a simplicial set* and the *homotopy category of an ∞-category* . The fundamental category is supposed to be the homotopy category to some extent, so we should compare these two notions where they are both defined (∞-categories).

Definition 2.18. Let C be an ∞-category . There is a natural map $\pi : C \to N(hC)$ that "passes to homotopy." It acts like so.

- An object is sent to itself (note $C_0 = (hC)_0 = N(hC)_0$).
- A morphism f is sent to its homotopy class.

- An n-cell $a \in C_n$ is sent to the unique $\pi(a) \in N(hC)_n$ that satisfies $\pi(a)_{i-1,i} = \pi(a_{i-1,i})$ for all i. (See also [Lur22, Construction 004G].)

This map $\pi : C \to N(hC)$ is compatible with simplicial operators, in the sense that given a cell $a \in C_n$, one has $[a_{01}] \circ \cdots \circ [a_{n-1,n}] = [a_{0n}]$.

Proposition 2.19. *If C is an ordinary category, then $f \simeq g$ iff $f = g$. Thus if an ∞-category C is isomorphic to a nerve of a 1-category, then $\pi : C \to N(hC)$ is an isomorphism, so it must be isomorphic to the nerve of its homotopy category.*

Proposition 2.20 (Universal property of homotopy category)**.** *Let C be an ∞-category and D a small category. If $f = C \to ND$ is a map of simplicial sets, then there exists a unique map $g : N(hC) \to ND$ such that $f = g \circ \pi$.*

Proof. We will construct g. We will do so by constructing a functor $g : hC \to D$. On objects $c \in ob(hC) = C_0$, we define $g(c) := f(c) \in (ND)_0 = ob(D)$. On morphisms, we define $g([h]) := f(h)$. This is well-defined, for if $h \simeq h'$ exhibited by some $a \in C_2$, then $\phi(a) \in (ND)_2$ exhibits the identity $f(h') = \mathrm{id} \circ h$. $\square$

Corollary 2.21. *The homotopy category construction is left adjoint to the nerve functor:*

$$h : \mathsf{qCat} \rightleftarrows \mathsf{Cat} : N$$

(Easy to-do: homotopy category of products.)

2.6 (2/12) About composition in ∞-categories

Let C be an ∞-category . Let $x, y, z \in C_0$ be objects and let $f \in \mathrm{Hom}_C(x, y)$ and $g \in \mathrm{Hom}_C(y, z)$ be morphisms (i.e. 1-cells starting/ending at their domains/targets.) Last time we defined a *composite of morphisms f, g in C* to be any $h \in \mathrm{Hom}_C(x, z)$ such that there exists a 2-cell $a \in C_2$ such that $a_{01} = f$, $a_{12} = g$, and $a_{02} = h$. Composites exist and are unique up-to-homotopy (thus we can compose homotopy classes), but are not uniquely determined in general. We may ask whether every representative of a homotopy class of a composite can be realized on-the-nose as the extension of its (compositees?) The answer is yes.

Proposition 2.22. *If $f : x \to y, g : y \to z$, and $h : x \to z$ are morphisms in an ∞-category C, then $h \in [g] \circ [f]$ if and only if h is a composite of f with g, i.e. there exists $u \in C_2$ satisfying*

$$u|_{\Delta^{0,1}} = f, \quad u|_{\Delta^{1,2}} = g, \quad u|_{\Delta^{0,2}} = h.$$

The proof is nice. I'd reproduce it here but I don't feel like making that diagram right now.

2.7 (2/13) Isomorphisms and inverses in ∞-categories

Denote by C an ∞-category and $f : x \to y$ a morphism in C. We say f is an *isomorphism* or an *equivalence* if $[f]$ is an isomorphism in hC. Unwinding a bit, this is equivalent to the existence of a $g : y \to x$ such that $[f] \circ [g] = [1_y]$ and $[g] \circ [f] = [1_x]$. The property of being an isomorphism is related to inverses; the following is elementary.

Proposition 2.23. *Let $f : x \to y$ be a morphism in an ∞-category C. A morphism $g : y \to x$ is called a preinverse to f if $[f] \circ [g] = [\mathrm{id}_y]$, and a postinverse if $[g] \circ [f] = [\mathrm{id}_x]$. If g is both, we call it an inverse. TFAE.*

1. *f is an isomorphism.*
2. *f has an inverse.*
3. *f has a preinverse and a postinverse.*
4. *f has a preinverse with a preinverse.*
5. *f has a postinverse with a postinverse.*

As for composites, inverses are not generally unique, but are so up-to-homotopy.

Proposition 2.24. *If $F : C \to D$ is a map of quasicategories, then F sends isomorphisms to isomorphisms.*

Proof. Suppose that $f : x \to y$ is an isomorphism in C. By Proposition 2.23, f admits an inverse $g : y \to x$. By definition, it satisfies $[f] \circ [g] = [\mathrm{id}_y]$, so by Proposition 2.22 there exists $u \in C_2$ witnessing $\mathrm{id}_y = f \circ g$. By this, I mean the obvious identities with simplicial operators hold. Morphisms of simplicial sets commute with operators, hence $F(u)$ witnesses $\mathrm{id}_{F(y)} = F(f) \circ F(g)$. This shows that $F(g)$ is a preinverse to $F(f)$. By an identical argument, one sees that $F(g)$ is a postinverse. So $F(g)$ is an inverse and we are done. $\square$

2.8 (2/13) ∞-groupoids, cores, and Kan complexes

Here are some definitions.

- An ∞-category C is called a *quasigroupid* or *∞-groupoid* if hC is a groupoid, i.e. if every morphism is an isomorphism.
- For an ordinary category C, its *core* C^{core} is the subcategory with the same objects and only the isomorphisms of C.
- For an ∞-category C, its *core* C^{core} is the simplicial subset consisting of all cells of C whose edges are all isomorphisms.
- Recall that a simplicial set is called a *Kan complex* if all horns (not necessarily inner) have extensions.

Note that $N(C^{core}) = (NC)^{core}$, so our terminology is justified. Now here are some facts.

Proposition 2.25. *If C is an ∞-category , then $\pi_0 C^{core} = \{$objects of $C\}/\cong$.*

Proposition 2.26. *If C is an ∞-category , then C^{core} is a subcategory and an ∞-groupoid . Furthermore, every sub-∞-groupoid is contained in C^{core}. In other words, C^{core} is the maximal subcategory which is also an ∞-groupoid .*

Proposition 2.27. *Every Kan complex is an ∞-groupoid .*

Proof. Suppose that K is a Kan complex and $f : x \to y$ is a morphism in K. Consider the horn $u : \Lambda^2_0 \to K$ with $u_{01} = f$ and $u_{02} = id_x$. Extending this horn gives us $g := u_{12}$ with $[g] \circ [f] = [\mathrm{id}_x]$. Thus, every morphism admits a preinverse, and this turns out to be sufficient for an ∞-category to be an ∞-groupoid . $\square$

Proposition 2.28 (Joyal's theorem; harder)**.** *Every ∞-groupoid is a Kan complex.*

Definition 2.29. Since $\mathrm{Sing}\, X$ is a Kan complex, the proposition allows us to define the *fundamental ∞-groupoid* of a space X as $\mathrm{Sing}\, X$.

2.9 (2/15) The functor quasicategory

The nerve functor $N : \mathsf{Cat} \to \mathsf{sSet}$ is fully faithful, so functors of categories correspond to morphisms of their nerves. Maybe this suggests that if we want a "mapping space," or really a "mapping ∞-category ," it's 0- and 1-categorical structure should consist of morphisms $X \to Y$ of simplicial sets and natural transformations between them. Morphisms are the same thing as maps $X \times \Delta^0 \to Y$. A natural transformation is a suitable map $X \times \Delta^1 \to Y$. This suggests the following.

Definition 2.30. Let X, Y be simplicial sets. Their *function complex* is the simplicial set $\mathrm{Fun}(X, Y)$ with

$$(\mathrm{Fun}(X, Y))_n := \mathrm{Hom}_{\mathsf{sSet}}(\Delta^n \times X, Y).$$

Proposition 2.31. *There is a natural bijection*

$$\mathrm{Hom}(X \times Y, Z) \xrightarrow{\sim} \mathrm{Hom}(X, \mathrm{Fun}(Y, Z)).$$

Proposition 2.32. *For ordinary C and D, one has $N(\mathrm{Fun}(C, D)) \cong \mathrm{Fun}(NC, ND)$.*

Eventually we will see that *function complexes to an ∞-category are ∞-categories* . This is what we want. We can't prove this yet.

2.10 (2/16) Lifting properties time — weakly saturated classes

My assessment to become a personal fitness trainer is approaching, so I'm shifting focus to that for a little while. I'll probably be leaving more to-do's than I should be.

Quasicategories are defined in terms of lifting properties. Now we will take some time to generally study lifting properties, which will be useful for studying quasicategories.

Definition 2.33. Let C be a category admitting small colimits. A class of morphisms A is called a *weakly saturated class* if it

1. Contains all isomorphisms,
2. Is closed under cobase change (also called pushouts), composition, transfinite composition, coproducts, and retracts. (See Charles p. 38 for the definitions.)

Given any class of morphisms S, its *weak saturation* $\bar{S}$ is the smallest weakly saturated class containing S.

Example 2.5. Take $C =$ Set. The weak saturation of $\{\{0, 1\} \to \{1\}\}$ is the class of surjective maps. The weak saturation of $\{\varnothing \to \{1\}\}$ is the class of injective maps.

Example 2.6. Take $C =$ Set. The surjections/injections also arise as weak saturations. Of what?

Proposition 2.34. *Let S be a category with small colimits and let $\mathcal{C}$ be a class of objects. Let $\mathcal{A}$ be the class of maps with the following lifting property: if $i : A \to B$ is in $\mathcal{A}$, then for every $f : A \to C$ to an object of $\mathcal{C}$, we can fill the following diagram:*

$$\begin{tikzcd} A \arrow[r, "f"] \arrow[d, "i"'] & C \\ B \arrow[ur, dashed] & \end{tikzcd}$$

Then $\mathcal{A}$ is a weakly saturated class. (Example: $S =$ sSet, $\mathcal{C} = \{\infty$-categories $\}$.

Proof. To-do. (Worked out in meeting.) □

2.11 (2/16) Classes of horns, anodyne morphisms

As indicated, we are interested in ∞-categories , so we ought to study lifting properties of horns in particular. We make some definitions for this.

Definition 2.35. We define the following sets of horns.

$$\begin{aligned} \mathrm{InnHorn} :=& \{\Lambda^n_k \hookrightarrow \Delta^n : 0 < k < n, n \geq 2\}, \\ \mathrm{Horn} :=& \{\Lambda^n_k \hookrightarrow \Delta^n : 0 \leq k \leq n, n \geq 1\}, \\ \mathrm{RHorn} :=& \{\Lambda^n_k \hookrightarrow \Delta^n : 0 < k \leq n, n \geq 1\}, \\ \mathrm{LHorn} :=& \{\Lambda^n_k \hookrightarrow \Delta^n : 0 \leq k < n, n \geq 1\}. \end{aligned}$$

We call their weak saturations in sSet the *(inner, right, left) anodyne morphisms*.

Proposition 2.36. *Monomorphisms of simplicial sets form a weakly saturated class. Therefore, since* Horn *consists of monomorphisms, its weak saturation must too. So (inner, right, left) anodyne maps are always monomorphisms.*

Proposition 2.37. *Let C be an ∞-category . If $A \hookrightarrow B$ is an inner anodyne inclusion, then every $f : A \to C$ extends to B.*

Proof. Let $\mathcal{A}$ denote the set of maps of simplicial sets $X \to Y$ which extend along every map $X \to C$. (C is fixed here.) Since C is a quasicategory, we have $\mathrm{InnHorn} \subseteq \mathcal{A}$. By Proposition 2.34, the class $\mathcal{A}$ is weakly saturated, so $\overline{\mathrm{InnHorn}} \subseteq \mathcal{A}$. That is what we wanted to show. □

Proposition 2.38. *(To-do: Prop 16.10.)*

Example 2.7. Here are some examples of inner anodyne morphisms. (Important to-do: finish.)

- The inclusions of spines $I^n \hookrightarrow \Delta^n$ are inner anodyne for every n. In particular, if C is an ∞-category , every map $I^n \to C$ extends to a $\Delta^n \to C$.

- $\vdots$

2.12 (2/16) Lifting calculus

Here's a definition I expected a bit earlier in Charles' notes.

Definition 2.39. Say an object X *satisfies the extension property* for $f : A \to B$ if for every $u : A \to X$ we can find an extension $B \to X$.

Definition 2.40. Suppose as given maps $f : A \to B$ and $g : X \to Y$. A *lifting problem* for (f, g) is a pair of maps $u : A \to X$ and $v : B \to Y$ making a commutative square. A *lift* for the lifting problem is a fill s to the obvious diagram:

$$\begin{tikzcd} A \arrow[r, "u"] \arrow[d, "f"'] & X \arrow[d, "g"] \\ B \arrow[r, "v"'] \arrow[ur, dashed, "s"] & Y \end{tikzcd}$$

Definition 2.41. Let f, g be morphisms in a category. We write *$f \boxslash g$* if every lifting problem for (f, g) admits a lift. We call this the *lifting relation* on morphisms. If $f \boxslash g$, we say:

- f has the left lifting property rel. to g, or
- g has the right lifting property rel. to f, or
- f lifts against g.

Definition 2.42. Let A be a class of morphisms. We define the *right complement* $A^{\boxslash} := \{g : a \boxslash g, \forall a \in A\}$. We define the *left complement* ${}^{\boxslash}A$ similarly.

Proposition 2.43. *Let $\mathcal{A}$ be any class of morphisms in a category with small colimits. The left complement ${}^{\boxslash}\mathcal{A}$ is weakly saturated, and the right complement $A^{\boxslash}$ is weakly cosaturated.*

Proof. (To-do: prove. Also, Charles' related exercises.) $\square$

Example 2.8. Let C be an abelian category, let $\mathcal{P} = \{0 \to P : P \text{ projective}\}$, and let $\mathcal{B}$ be the class of epimorphisms in C. By the definition of projective objects, we have $\mathcal{P} \boxslash \mathcal{B}$. Thus $\mathcal{B} \subseteq \mathcal{P}^{\boxslash}$.

(To-do: show converse?)

2.13 (2/17) Inner fibrations

A map p of simplicial sets is called an *inner fibration* if $\mathrm{InnHorn} \boxslash p$. Thus, $\mathrm{InnFib} = \mathrm{InnHorn}^{\boxslash}$.

Proposition 2.44. *A simplicial set C is an ∞-category iff $C \to *$ is an inner fibration.*

Proposition 2.45. InnFib *is defined as a right complement, thus* InnFib *is weakly cosaturated. This implies, for instance, that if C is an ∞-category and $D \to C$ is an inner fibration, then D is an ∞-category .*

Proposition 2.46 (. *kerodon] If X is a simplicial set, then a morphism $X \to ND$ is an inner fibration iff X is an ∞-category .*

Proof. (To-do.) $\square$

2.14 (2/20) Factorizations

Recall that we defined *inner fibrations* InnFib as the right complement of $\mathrm{InnHorn}$. This tells us something—maybe this is why we call it a "complement."

Proposition 2.47 (Small object argument). *Let S be a set of morphisms in* sSet. *Then every map f of simplicial sets admits a factorization $f = p \circ j$ with $p \in S^{\boxslash}$ and $j \in \overline{S}$.*

Corollary 2.48. *If S is any set of morphisms in* sSet, *then $\overline{S} = {}^{\boxslash}(S^{\boxslash})$.*

Proof. Since ${}^{\boxslash}(S^{\boxslash})$ is a left complement, it is weakly saturated (Proposition 2.43), thus $\overline{S} \subseteq {}^{\boxslash}(S^{\boxslash})$.

Now suppose that $f \boxslash S^{\boxslash}$. By the previous proposition, we may write $f = pj$ for $p \in S^{\boxslash}$ and $j \in \overline{S}$, and by assumption f admits a lift in the following diagram.

$$\begin{tikzcd}
\bullet \arrow[r, "j"] \arrow[d, "f"'] & \bullet \arrow[d, "p"] \\
\bullet \arrow[r, "\mathrm{id}"'] \arrow[ur, dashed, "s"] & \bullet
\end{tikzcd}$$

Thus, we get the following commutative diagram.

$$\begin{tikzcd}
\bullet \arrow[r, "\mathrm{id}"] \arrow[d, "f"'] & \bullet \arrow[r, "\mathrm{id}"] \arrow[d, "j"] & \bullet \arrow[d, "f"] \\
\bullet \arrow[r, "s"'] & \bullet \arrow[r, "p"'] & \bullet
\end{tikzcd}$$

This exhibits f as a retract of j. Since $j \in \overline{S}$ and weak saturations are closed under retracts, we have $f \in \overline{S}$. $\square$

Corollary 2.49. *Every map f of simplicial sets can be factored $f = pj$ with p an inner fibration and j inner anodyne.*

2.15 (2/21) Factorization systems and unique lifts

Last time, we proved that for a class S of maps in sSet, we have $\overline{S} = {}^{\boxslash}(S^{\boxslash})$ and we can factor every map as a composite of one map from $\overline{S}$ and $S^{\boxslash}$. Starting with $S = \mathrm{InnHorn}$, we got a factorization of an arbitrary map as an inner anodyne followed by an inner fibration. We study such systems in general.

Definition 2.50. A *weak factorization system* in a category is a pair of classes of maps $(\mathcal{L}, \mathcal{R})$ with the following properties.

1. Every morphism factors as rl for $l \in \mathcal{L}$ and $r \in \mathcal{R}$; and
2. $\mathcal{L} = {}^{\boxslash}\mathcal{R}$ and $\mathcal{R} = \mathcal{L}^{\boxslash}$.

Example 2.9. The pair $(\overline{\mathrm{InnHorn}}, \mathrm{InnHorn}^{\boxslash})$ is a weak factorization system.

We would like to understand lifting problems with unique solutions.

Definition 2.51. In a category with coproducts, let $f : A \to B$ be a morphism. We define the *fold* of f, denoted $f^{\wedge}$: $B \coprod_A B \to B$, as the unique map making the following diagram commute.

$$\begin{tikzcd}
B & & \\
& B \coprod_A B \arrow[ul, dashed, "f^{\vee}"] & B \arrow[l] \arrow[ull, "\mathrm{id}"'] \\
& B \arrow[u] \arrow[uul, "\mathrm{id}"] & A \arrow[l, "f"'] \arrow[u, "f"']
\end{tikzcd}$$

Proposition 2.52. *Let f, g be morphisms. The following are equivalent.*

1. *We have $f \boxslash g$ and $f^\vee \boxslash g$.*
2. *The solution to any lifting problem for (f, g) exists and is unique.*

Proof. $(2) \Longrightarrow (1)$ is obvious. For $(1) \Longrightarrow (2)$, existence is assumed, so we need only show uniqueness. Wait, what does uniqueness mean in an arbitrary category? □

Definition 2.53. A weak factorization system $(\mathcal{L}, \mathcal{R})$ is called an *orthogonal factorization system* if $\mathcal{L} = {}^{\boxslash}\mathcal{R}$ and $\mathcal{R} = \mathcal{L}^{\boxslash}$ are realized by *unique* lifts.

Proposition 2.54. *The factorization $f = rl$ is unique up to unique isomorphism in an orthogonal factorization system.*

Proposition 2.55. *$(\{\text{surjections}\}, \{\text{injections}\})$ form an orthogonal factorization system in* Set. *(Proof is obvious.)*

Proposition 2.56. *For any class of simplicial maps S, the pair $(\overline{S \cup S^\vee}, (S \cup S^\vee)^{\boxslash})$ is an orthogonal system.*

2.16 (2/23) Degenerate cells

We want to concretely understand monomorphisms of simplicial sets. For this, recall that we defined the *boundary of* Δ^n as the subcomplex of Δ^n whose k-cells are the non-surjective maps $[k] \to [n]$. Write Cell for the class of inclusions $\partial\Delta^n \hookrightarrow \Delta^n$ and $\mathrm{InnFib} := \mathrm{Cell}^{\boxslash}$. Since Cell consists of monos, we know $\overline{\mathrm{Cell}}$ contains all monomorphisms. Our main theorem is the converse.

Proposition 2.57. *The class $\overline{\mathrm{Cell}}$ is exactly the class of monomorphisms of simplicial sets.*

We'll prove this (Proof 2.17) once we've set some stuff up.

Toward proving this, recall the notion of degenerate cells: a cell $\sigma : \Delta^n \to X$ is called *degenerate* if there exists a non-injective operator $f : [m] \to [n]$ such that $\sigma = \tau f$. Since every simplicial operator factors uniquely as $f = f^{inj} f^{surj}$, we see that if σ is degenerate if and only if there is some non-identity *surjective* f such that $a = bf$. A cell which is not degenerate is called *non-degenerate*. We write $X_n = X_n^{deg} \coprod X_n^{nd}$ for the decomposition of X_n into (non)-degenerate cells. Neither assemble to a subcomplex.

Proposition 2.58. *Here are some straightforward facts about degenerate cells.*

1. *If $f : X \to Y$ is a map of simplicial sets, then $f(X_n^{deg}) \subseteq Y_n^{deg}$.*
2. *If $f : X \to Y$ is a map of simplicial sets, then $f^{-1}(Y_n^{nd}) \subseteq X_n^{nd}$.*
3. *If $A \hookrightarrow X$ is a subcomplex, then*
$$A_n^{nd} = X_n^{nd} \cap A_n, \text{ and}$$
$$A_n^{deg} = X_n^{deg} \cap A_n.$$
4. *The elements of $(\Delta^n)_k^{nd}$ are in bijection with the subsets of $[n]$ of size k.*
5. *The simplicial n-sphere $\Delta^n/\partial\Delta^n$, defined as the pushout of $\Delta^n \hookleftarrow \partial\Delta^n \to \Delta^0$, has exactly two nondegenerate cells: its unique vertex and the generator $\langle 0, 1, \ldots, n\rangle$. In other words, the image of $\Delta^0 \to \Delta^n/\partial\Delta^n$ and of the generator in $\Delta^n \to \Delta^n/\partial\Delta^n$ in the pushout square:*
$$\begin{tikzcd} \partial\Delta^n \arrow[r, tail] \arrow[d] & \Delta^n \arrow[d] \\ \Delta^0 \arrow[r] & \Delta^n/\partial\Delta^n \end{tikzcd}$$

The *Eilenberg-Zilber lemma* says that every cell a of X occurs *uniquely* as $a = b\sigma$ for a nondegenerate b and surjective operator σ. (This is not too complicated; the nontrivial part is uniqueness.) Let's state this in a slightly different, stronger form.

Proposition 2.59. *If X is a simplicial set, then for every n the map*

$$\coprod_{j\geq 0} X_j^{nd} \times \mathrm{Hom}_{\Delta^{surj}}([n],[j]) \to X_n$$

Given by $(j,a,\sigma) \mapsto a\sigma$ is a bijection. Furthermore, this map is natural with respect to surjective operators $[n'] \to [n]$ and with respect to monomorphisms of simplicial sets $X \to X'$.

There are various things to be said now. I think I will move on, and refer back to these things when they are needed.

2.17 (2/23) The skeletal filtration

If $\sigma : \Delta^n \to X$ is an n-cell, it uniquely factors as $\Delta^n \to \Delta^m \to X$ where the first map is surjective and the second is nondegenerate. So σ is "really" an m-cell, for some $m \leq n$. Now, we want a notion of the k-skeleton of X. Its n-cells should be the n-cells of X which are "really" j-cells for some $j \leq k$.

Definition 2.60. Let X be a simplicial set. The *k-skeleton* of X, written $Sk_k X$, is the smallest subcomplex containing all cells of dimension $\leq k$. Thus, we have

$$(Sk_k X)_n = \bigcup_{0\leq j\leq k} \{yf : y \in X_j \text{ and } f : [n] \to [j]\}.$$

A nondegenerate cell $\Delta^k \to X$ determines a cell $\Delta^k \to Sk_k X$. This map carries $\partial\Delta^{k-1}$ to $Sk_{k-1}X$.

Proposition 2.61. *The evident square*

$$\begin{tikzcd}
\coprod_{i\in X_k^{nd}} \partial\Delta^k \arrow[r] \arrow[d, tail] & Sk_{k-1}X \arrow[d, tail] \\
\coprod_{i\in X_k^{nd}} \Delta^k \arrow[r] & Sk_k X
\end{tikzcd}$$

Is a pushout square. More generally, if $A \subseteq X$ is a subcomplex, the following is a pushout square.

$$\begin{tikzcd}
\coprod_{i\in X_k^{nd}/A_k^{nd}} \partial\Delta^k \arrow[r] \arrow[d, tail] & A \cup Sk_{k-1}X \arrow[d, tail] \\
\coprod_{i\in X_k^{nd}/A_k^{nd}} \Delta^k \arrow[r] & A \cup Sk_k X
\end{tikzcd}$$

It is in this sense that simplicial sets are built out of standard simplices: a simplicial set X is filtered by $X_0 = Sk_0 X \subseteq Sk_1 X \subseteq Sk_2 X \subseteq \cdots$, and each Sk_n is obtained from Sk_{n-1} by attaching copies of Δ^n as in Proposition 2.61.

Now we are ready to prove our characterization of monomorphisms (Proposition 2.57).

Proof. A monomorphism of simplicial sets is isomorphic to an inclusion $A \hookrightarrow X$. It is clear that $X \cong \varinjlim_k A \cup Sk_k X$. But see that, by the above proposition, the maps $A \cup Sk_{k-1}X \to A \cup Sk_k X$ arise via cobase change from coproducts of maps in Cell. Then the inclusion is exhibited as a countable composition(?) of maps in $\overline{\text{Cell}}$, thus is in $\overline{\text{Cell}}$. □

And so we have some handle on monomorphisms in sSet now.

Corollary 2.62 (Geometric realizations are CW)**.** *Recall that we constructed the geometric realization* *functor* $|-|$: sSet $\to$ Top *as a left adjoint. Left adjoints preserve colimits, hence we have a pushout diagram*

$$\begin{tikzcd}
\coprod_{a\in X_k^{nd}} \partial\Delta^k_{top} \arrow[r] \arrow[d, tail] & {|Sk_{k-1}X|} \arrow[d, tail] \\
\coprod_{a\in X_k^{nd}} \Delta^k_{top} \arrow[r] & {|Sk_k X|}
\end{tikzcd}$$

Additionally, we have $|X| = \lim_{\to} |Sk_k X|$. *This describes a canonical CW structure on the geometric realization* $|X|$ *of a simplicial set. Evidently, cells of* $|X|$ *correspond to nondegenerate simplices of* X.

2.18 (2/25) Pushout-products, pullback-homs

Let $f : A \to B$ and $g : K \to L$ be morphisms in sSet. We define the *pushout-product* of f and g, denoted fg, as the unique dotted map making the following pushout square diagram commute.

$$\begin{tikzcd}
A \times K \arrow[r, "\mathrm{id}\times g"] \arrow[d, "f\times\mathrm{id}"'] & A \times L \arrow[d] \arrow[ddr, "f\times\mathrm{id}"] & \\
B \times K \arrow[r] \arrow[drr, "\mathrm{id}\times g"'] & (B\times K)\coprod_{A\times K}(A\times L) \arrow[dr, dashed, "f\Box g"] & \\
& & B \times L
\end{tikzcd}$$

Dually, we define the *pullback-hom* of f and g, denoted $f^{\Box g}$, to be the unique dotted map making the following pullback square diagram commute.

$$\begin{tikzcd}
\mathrm{Fun}(L,A) \arrow[dr, dashed, "f^{\Box g}"] \arrow[drr, "\mathrm{Fun}(g,\mathrm{id})"] \arrow[ddr, "\mathrm{Fun}(\mathrm{id},f)"'] & & \\
& \mathrm{Fun}(K,A)\times_{\mathrm{Fun}(K,B)}\mathrm{Fun}(L,B) \arrow[r] \arrow[d] & \mathrm{Fun}(K,A) \arrow[d, "\mathrm{Fun}(\mathrm{id},f)"] \\
& \mathrm{Fun}(L,B) \arrow[r, "\mathrm{Fun}(g,\mathrm{id})"'] & \mathrm{Fun}(K,B)
\end{tikzcd}$$

3 March

3.1 (3/5) Pullback-hom as an enriched lifting problem

Suppose given $g : K \to L$ and $h : X \to Y$. We have the pullback-hom

$$h^{\square g} : \mathrm{Fun}(L, X) \to \mathrm{Fun}(K, X) \times_{\mathrm{Fun}(K,Y)} \mathrm{Fun}(L, Y).$$

The vertices of $\mathrm{Fun}(L, X)$ are morphisms $L \to X$ in sSet. The vertices of $\mathrm{Fun}(K, X) \times_{\mathrm{Fun}(K,Y)} \mathrm{Fun}(L, Y)$ are those pairs of morphisms $(s : K \to X, t : L \to Y)$ such that $hs = tg$, i.e. lifting problems for (g, h). The pullback-hom $h^{\square g}$ takes a morphism $w : L \to X$ and composes it to (wg, gh). This gives us a lifting problem for (g, h) that is solvable. Then the following is clear.

Proposition 3.1. *The pullback-hom $h^{\square g}$ is surjective on vertices iff $g \boxslash h$.*

In this sense, $h^{\square g}$ encodes an "enriched" lifting problem for (g, h). The target $\mathrm{Fun}(K, X) \times_{\mathrm{Fun}(K,Y)} \mathrm{Fun}(L, Y)$ parametrizes lifting problems for (g, h) while the source $\mathrm{Fun}(L, X)$ parametrizes families of lifting problems together with a chosen lift.

Also, let's talk about so-called *adjunctions of lifting problems*. The product and function complex constructions are adjoint. Ultimately, this leads to the following.

Proposition 3.2. *One has $(f \square g) \boxslash h$ if and only if $f \boxslash h^{\square g}$.*

Here's a special case. Take $K = \varnothing$ and $Y = *$. Then the proposition gives us that

$$(f \times \mathrm{id}_L) \boxslash (X \to *) \Longleftrightarrow f \boxslash (\mathrm{Fun}(L, X) \to *).$$

4 April

5 May

5.1 (5/3) Operads for (Peter) May

For May, I will learn about operads and monads. Here are my motives.

(1) Peter **May** coined the term *operad*.

(2) They're interesting and fit into the ∞-categorical framework, eventually.

(3) I read some of Moerdijk-Weiss's *Dendroidal sets* for Charles' Kan Seminar and thought it was exicting.

(4) Peter May tipped me off that monads and operads would make a big appearance in his talks for the 2023 UChicago REU. (Probably related to his recent work with Ruoqi Zhang and Foling Zou.)

Here are some potential references.

(1) May, *The Geometry of Iterated Loop Spaces* (1971)

(2) Markl-Shnider-Stasheff, *Operads in Algebra, Topology, and Physics* (2000)

(3) Heuts, *Simplicial and Dendroidal Homotopy Theory* (2022)

(4) Markl, *Operads and PROPs* (2006)

(5) Lawson, *E_n-ring spectra and Dyer-Lashof operations*

Let me say what I *think* operads are supposed to do/be before I dive into it:
An operad abstracts away the structure of "an operation on a structure its identities/coherences." We'll see this worked out later, for the first time in Example 5.4. Here's a vague indication as to why this is useful: say an algebraic thing X has operations which "play nice" with its algebraic structure. If this occurs, it can have useful consequences. It's natural that we then (1) find and study objects for which this occurs, and (2) study them in concert. But this has problems: (A) it may require *lots* of data to verify or realize that X's operations "play nice" with its structure (*coherence data*), especially for complicated X, and/or especially if we're thinking up-to-homotopy, and (B) if we want to study such objects relative to each other, we'll have to compare several of these huge packages of data. Operads do the work for us: the idea is to say, "let $\mathcal{C}$ be the operad codifying the possession of coherent operations." Then, given an objekt X, a choice (if one exists) of such structure on X amounts to a morphism $\mathcal{C} \to \mathcal{E}$

nd$_X$, the latter being a canonical "endomorphism operad" associated to X. That morphism essentially says, "we can interpret the structure within $\mathcal{C}$ as some class of operations on X." In other words, for an algebraic structure X, an operad classifies "coherent" operations on X, the details (e.g., how coherent?) dependent on which operad you're considering.

Maybe another way of putting it is that operads represent the formal algebraic theory, while we're interested in "instantiations" of these theories, i.e. their representations—I think we call these "algebras over the operads."

OK, let me actually learn what these are now.

Definition 5.1. An *operad*, *symmetric operad*, or *classical operad* $\mathcal{C}$ is a collection of sets $(\mathcal{C}(i))_{i\geq 0}$ with a distinguished operation $1_{\mathcal{C}} \in \mathcal{C}(1)$ and functions $\gamma : \mathcal{C}(n) \times \mathcal{C}(k_1) \times \cdots \times \mathcal{C}(k_n) \to \mathcal{C}(\sum k_s)$, which we regard as *operations*, the *unit operation*, and as *composition*, respectively, and in this regard we require that these structures are suitably associative, unital, and equivariant up to reordering of inputs.

Concretely, $\mathcal{C}$ is the data of

- (Operations) For each $i \geq 0$, a set $\mathcal{C}(i)$ called *i-ary operations*; and
- (Composites) For each $n \geq 0$ and $k_1, \ldots, k_n \geq 0$, a *composition* map $\gamma : \mathcal{C}(n) \times \mathcal{C}(k_1) \times \cdots \times \mathcal{C}(k_n) \to \mathcal{C}(\sum k_s)$; and
- (Unit) A distinguished *identity* operation $1_{\mathcal{C}} \in \mathcal{C}(1)$; and

- (Symmetries) For each $i \geq 0$, an action $\Sigma_i \to \mathrm{Aut}\mathcal{C}(i)$

satisfying the following conditions.

- (Unitality) For every operation $d \in \mathcal{C}(j)$, we have $\gamma(1; d) = d$ and $\gamma(d; 1^{\times j}) = d$.
- (Associativity) Blah blah blah
- (Σ_n-Equivariance) For each $\Sigma_1^k j_s = j, c \in \mathcal{C}(k)$, $d_s \in \mathcal{C}(j_s), \sigma \in \Sigma_k$, and $\tau_s \in \Sigma_{j_s}$, we have

$$\gamma(c\sigma; (d_i)) = \gamma(c; (d_{\sigma^{-1}i})) \cdot \sigma(j_1, \ldots, j_s), \quad \text{and}$$
$$\gamma(c; (d_i\tau_i)) = \gamma(c; (d_i))(\tau_1 \oplus \cdots \oplus \tau_k).$$

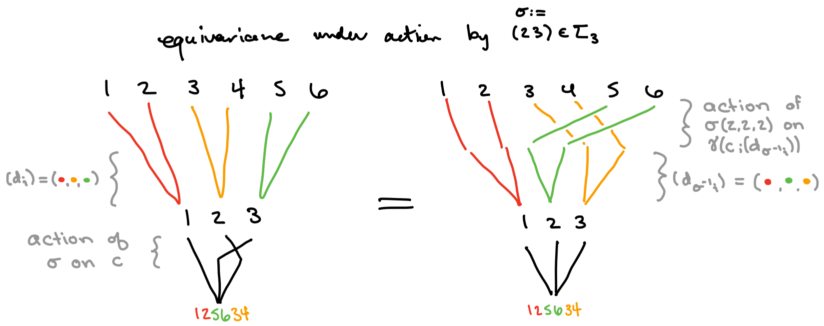
Here, $\sigma(j_1, \ldots, j_s) :=$ the permutation of j letters given by permuting the k blocks of letters determined by the partition $j = \Sigma j_s$ according to σ.

Definition 5.2. An operad with no Σ_i-actions or equivariance is called *plain* or *non-Σ* or *non-symmetric*.

Definition 5.3. Above, we defined an operad in Set. We can make an analogous definition in any bicomplete symmetric monoidal category $(\mathsf{C}, \otimes, \mathbf{1})$. In this case, the unit/identity is a distinguished morphism $\mathbf{1} \to \mathcal{C}(1)$. The symmetries become maps $\Sigma_i \to \mathrm{Iso}(\mathcal{C}(i), \mathcal{C}(i))$. Such a thing is called an *operad in* C. This is like "enriching" an operad over a category.

Definition 5.4. A *morphism of operads* $\mathcal{C} \to \mathcal{C}'$ is a collection of maps $f_i : \mathcal{C}(i) \to \mathcal{C}'(i)$ such that $f_1(1) = 1$ and (equivariance, compatibility with composition).

Remark 5.5. As indicated, we think of $\mathcal{C}(i)$ as a set of i-ary operations, and the functions $\gamma : \mathcal{C}(k) \times \mathcal{C}(n_1) \times \cdots \times \mathcal{C}(n_k) \to \mathcal{C}(\sum n_s)$ as taking a k-ary operation c and plugging in k other operations (d_i). The Σ_n-equivariance demands that if we tamper with the inputs for c then plug in the (d_i), that is the same as plugging in the (d_i) in a different order then tampering with the order of their inputs. See the little picture.



An operad's purpose in life is to help define *algebras over an operad*. Such a thing establishes an "algebraic structure representing the operad" upon an object. Here is the definition.

Definition 5.6. Let $\mathcal{C}$ be an operad in a symmetrical monoidal C. A *$\mathcal{C}$-algebra* A is an object A and maps $\mathcal{C}(i) \otimes A^{\otimes i} \to A$ that are suitably associative, unital, and equivariant. (We take $A^{\otimes 0} = \mathbf{1}_{\mathsf{C}}$.)

I'll go over many examples soon. Some of these will let us reinterpret some of the above structures. But for the rest of today, I'll just make a little remark.

Remark 5.7 (Classical operads generalize monoids)**.** Let $\mathcal{C}$ be a classical operad. (A non-Σ operad works too.) There is an associated category $j_!\mathcal{C}$ with one object and morphisms given by $\mathcal{C}(1)$, the unary operations. Given $f, g \in \mathcal{C}(1)$, their composite $g \circ f :=$ their image in $\mathcal{C}(1) \times \mathcal{C}(1) \to \mathcal{C}(1)$, and the unit is the identity operation $1_{\mathcal{C}} \in \mathcal{C}(1)$. This checks out thanks to the unitality and associativity axioms.

Conversely, given a one-object category M, i.e. a monoid, we may form an operad $j^*\mathsf{M}$ with solely unitary operations, given by $(j^*\mathsf{M})(1) := \mathrm{Hom}_{\mathsf{M}}(*, *)$. The unit and composition functions are obvious, and the Σ_i-actions are trivial.

Altogether we get an *adjunction*

$$\begin{tikzcd} \text{Monoids} \arrow[r, shift left=1.5, "j_!"] \arrow[r, phantom, "\perp"] & \text{Operads} \arrow[l, shift left=1.5, "j^*"] \end{tikzcd}$$

5.2 (5/5) Basic examples of operads

In what follows, let $(\mathsf{C}, \otimes, \mathbf{1})$ denote a symmetric monoidal category.

Example 5.1. Let A be an object of C. If C is closed,[1] we denote by $\mathcal{E}\mathrm{nd}_A$ the *endomorphism operad* of A, defined by $(\mathcal{E}\mathrm{nd}_A)(i) := \mathrm{Hom}(A^{\otimes i}, A)$. The unit is $\mathrm{id}_A \in \mathcal{E}\mathrm{nd}_A(1)$ and the compositions are given by composing tensor product maps. The right Σ_i-action is given by the left Σ_i-action on tensor powers.

Proposition 5.8. *Let $A \in \mathsf{C}$ and let $\mathcal{C}$ be an operad in C. Via the tensor-Hom adjunction, a $\mathcal{C}$-algebra structure on A is "the same thing as" a morphism $\mathcal{C} \to \mathcal{E}\mathrm{nd}_A$.*

Example 5.2. We denote by $\mathcal{C}\mathrm{omm}$ the *commutative operad* in Set. It is defined to have a single operation $\mathcal{C}\mathrm{omm}(i) := \{*\}$ for every i.

Example 5.3. We denote by $\mathcal{A}\mathrm{ssoc}$ the *associative operad* in Set. It is defined to have $\mathcal{A}\mathrm{ssoc}(i) := \Sigma_i$ for every i. The unit and Σ_i-action are obvious. The maps $\gamma : \Sigma_n \times \Sigma_{k_1} \times \cdots \times \Sigma_{k_n} \to \Sigma_{k_1+\cdots+k_n}$ are defined as follows: given $\sigma \in \Sigma_n$ and $\tau_j \in \Sigma_{k_j}$ regarded as matrices, one inserts τ_j in place of the 1 in the j-th column of σ, for each $1 \le j \le n$.

Proposition 5.9. *In Set, $\mathcal{A}\mathrm{ssoc}$-algebras (resp. $\mathcal{C}\mathrm{omm}$-algebras) are precisely monoids (resp. commutative monoids).*

In fact, we can encode monoids in the arbitrary C with operad actions. If $(\mathsf{C}, \otimes, \mathbf{1})$ has finite coproducts, for a finite set S let $\mathbf{1}[S]$ denote the coproduct $\coprod_S \mathbf{1}$.

Definition 5.10. We denote by $\mathcal{C}\mathrm{omm}$ the *commutative operad* in C. It is defined to have $\mathcal{C}\mathrm{omm}(i) := \mathbf{1}$.

Definition 5.11. If C has finite coproducts, we denote by $\mathcal{A}\mathrm{ssoc}$ the *associative operad* in C. It is defined to have $\mathcal{A}\mathrm{ssoc}(i) := \mathbf{1}[\Sigma_i]$. The rest of its structure is mostly obvious.

Proposition 5.12. *In C, the $\mathcal{C}\mathrm{omm}$-algebras are precisely the monoids in C. If C has finite coproducts, then the $\mathcal{A}\mathrm{ssoc}$-algebras are precisely the commutative monoids in C.*

Remark 5.13 (Algebras over symmetric vs. plain operads)**.** Above, we are regarding $\mathcal{A}\mathrm{ssoc}$ and $\mathcal{C}\mathrm{omm}$ as *symmetric* operads, and this is manifest in the structure of algebras over them. We can instead skip any mention of Σ_n's and consider $\mathcal{A}\mathrm{ssoc}$, $\mathcal{C}\mathrm{omm}$ as *plain* operads. If we do, then $\mathcal{C}\mathrm{omm}$-algebras become precisely monoids in C. And $\mathcal{A}\mathrm{ssoc}$-algebras become...something? Maybe this indicates we should avoid plain operads if possible.

5.3 (5/6) Warm-up: monoids are $\mathcal{A}\mathrm{ssoc}$-algebras

Let me work out a concrete example of how monoids are "the same thing as" $\mathcal{A}\mathrm{ssoc}$-algebras.

Example 5.4 (How do $\mathcal{A}\mathrm{ssoc}$-algebras encode monoids?)**.** Let X be a set. Suppose it is a monoid, i.e. that we have chosen a unital, associative product $\mu : X \times X \to X$. As I said yesterday, the monoid (X, μ) is "the same thing as" an $\mathcal{A}\mathrm{ssoc}$-algebra which is "the same thing as" a choice of morphism $f : \mathcal{A}\mathrm{ssoc} \to \mathcal{E}\mathrm{nd}_X$.

I'll describe the morphism f. Let me write e_n for the identity in Σ_n.

- f sends $\mathcal{A}\mathrm{ssoc}(1) = \{*\}$ to the identity $\mathrm{id}_X \in \mathcal{E}\mathrm{nd}_X(1)$.
- f sends $e_2 \in \mathcal{A}\mathrm{ssoc}(2) = \Sigma_2$ to $\mu \in \mathcal{E}\mathrm{nd}_X(2)$.
- Where does f send $\sigma \in \mathcal{A}\mathrm{ssoc}(2) = \Sigma_2$? Equivariance demands that $f(\sigma) = f(e_2)\sigma =$"swap inputs then do $f(e_2) = \mu$."

[1] I.e., if it has an internal Hom

- Similarly, the value of f on $\mathcal{A}\text{ssoc}(3)$ is already totally determined by its value $f(e_2) = \mu$. Here's why. In $\mathcal{A}\text{ssoc}$, the composition $\gamma : \mathcal{A}\text{ssoc}(2) \times \mathcal{A}\text{ssoc}(2) \times \mathcal{A}\text{ssoc}(1) \to \mathcal{A}\text{ssoc}(3)$ satisfies $\gamma(e_2, e_2, e_1) = e_3$. Since f respects composition, it is determined at e_3:

  $$f(e_3) = \gamma(f(e_2), f(e_2), f(e_1)) = \gamma(\mu, \mu, e) = \big((a, b, c) \mapsto \mu(\mu(a, b), c)\big).$$

  Moreover, notice that $\gamma(e_2, e_1, e_2) = e_3$. Then by the same argument, we find that

  $$f(e_3) = \big((a, b, c) \mapsto \mu(a, \mu(b, c))\big).$$

  So, $(ab)c = a(bc)$, where we're suppressing μ from notation. Associativity! Similarly, we could take into account the Σ_3-action and show the other associativity identities hold, e.g. $a(cb) = (ac)b$. So, the composition and equivariance conditions gave us associativity (with three inputs)!

- We could repeat the above to see that f's values on $\mathcal{A}\text{ssoc}(n)$ are forced by $f(e_2) = \mu$. Again, the various equalities $\gamma(e_k, e_{n_1}, \dots, e_{n_k}) = e_{n_1+\cdots+n_k}$, the equivariance equalities, and the fact that f must respect these force all the associativity laws.

Remark 5.14. The above example was nice in that after we specified $f(e_2 \in \mathcal{A}\text{ssoc}(2)) = \mu$, we could determine the rest of f based on composition laws, equivariance, and the fact that f must respect those. This is because in some sense, $\mathcal{A}\text{ssoc}$ is "generated by" the element $e_2 \in \mathcal{A}\text{ssoc}(2)$.

So there's an example of how we can use an operad to describe an algebraic operation with coherence. In this case, coherence was strict associativity: we had equalities such as

$$(ab)c = a(bc),\ a(cb) = (ac)b, \text{ and } a((bc)d) = ((ab)c)d.$$

These equalities were present "formally" in $\mathcal{A}\text{ssoc}$, and since $\mu : X \times X \to X$ was associative, we could find a corresponding $f : \mathcal{A}\text{ssoc} \to \mathcal{E}$
nd$_X$ implementing μ.

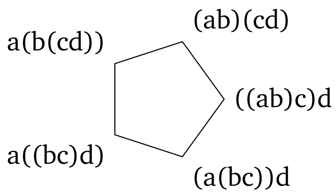
That example was overkill. We "just know" what the "coherence" is—it's associativity. The utility of operads arises when the coherences are more complicated. Next time I will look at an example wherein we'll replace the set X with a space, μ with a continuous map, and strict associativity with "associativity up to specified homotopy." In this case, we'll have e.g. $(ab)c \simeq a(bc)$ *and the data of a homotopy realizing this equivalence.*

5.4 (5/9) Associativity up to homotopy, Stasheff associahedra, and A_∞-operads

Let Y denote a based space. If $X = \Omega Y$, then X has a multiplication (loop concatenation). Let's parameterize it like so: for $x, y \in X$, define $xy : [0, 1]/\sim \to Y$ as the loop "do y over the first half of the interval, then x over the other."

Generally, we have $(ab)c \neq a(bc)$, thus X is not an $\mathcal{A}\text{ssoc}$-algebra. Loop concatenation is not "strictly associative." But clearly $(ab)c \simeq a(bc)$, realized by (say) a linear reparametrization of $[0, 1]$. That would be a map $[0, 1] \times I \to X$ starting at $a(bc)$ and landing at $(ab)c$. In $\mathrm{Fun}(X^3, X)$, in which $(ab)c$ and $a(bc)$ reside,[2] such a homotopy is a path from $(ab)c$ to $a(bc)$.

Now say we're concatenating four loops. There are five ways to do this (by reordering parentheses). Say we want a *choice* of homotopy between each reordering. This is the same as taking the pentagon below and sending it to $\mathrm{Fun}(X^4, X)$, each vertex going to the labeled 4-ary operation. **This pentagon is not filled in.**



Again, we think of e.g. $((ab)c)d$ and $(ab)(cd)$ as points in $\mathrm{Fun}(X^4, X)$ and the image in $\mathrm{Fun}(X^4, X)$ of the edge between them as a choice of homotopy equivalence $((ab)c)d \simeq (ab)(cd)$.

[2]That is, the maps $X^3 \to X$ given by "concatenate (a, b, c) to $(ab)c$, or to $a(bc)$."

Actually, if we chose a path in $\mathrm{Fun}(X^3, X)$ from $a(bc)$ to $(ab)c$ (i.e., if we've made a choice of homotopy equivalence realizing $(ab)c \simeq a(bc)$), that already gives us the five homotopy equivalences above—that is, where to send the pentagon above in $\mathrm{Fun}(X^4, X)$.

Notice that there are determined TWO homotopy equivalences between e.g. $a(b(cd))$ and $(ab)(cd)$. One follows the path in $\mathrm{Fun}(X^4, X)$ from $a(b(cd))$ to $(ab)(cd)$ given by traversing the pentagon clockwise, the other counterclockwise. **Since X is a loop space, it turns out that there is a higher homotopy equivalence between these two homotopies.** (In general, this need not be the case, since $\pi_1 \mathrm{Fun}(X^4, X)$ is not generally trivial.) Therefore, **there is determined a continuous map from the *solid* pentagon to** $\mathrm{Fun}(X^4, X)$. This is like a "higher" level of associativity, a "higher" level of coherence.

In fact, since X is a loop space, its multiplication is associative up to *all* higher homotopy coherences. Let's define an operad whose algebras are spaces with a multiplication that is associative up to all higher homotopy coherences.

Definition 5.15. Denote by $\mathcal{K}$ the non-Σ *Stasheff operad*. It is defined to have $\mathcal{K}(n) :=$ the convex $(n-2)$-dimensional polygon with a vertex for each parenthetization of n ordered letters. (The composition maps can be defined if we use a more explicit description; I won't give that.)

It is somewhat clear (to me, maybe everyone) that the Stasheff operad $\mathcal{K}$ works largely because $\mathcal{K}(n)$ is contractible for each n.

Definition 5.16. Let $\mathcal{C}$ be a non-Σ operad in Top. Say it is an *A_∞-operad* if each $\mathcal{C}(n)$ is contractible. Say that a space is an *A_∞-space* if it an algebra over an A_∞-operad.

Here's the main thing.

Theorem 5.17. *Up to weak equivalence,*

1. *A_∞-spaces are precisely the $\mathcal{K}$-algebras, and*
2. *Loop spaces are precisely the grouplike $\mathcal{K}$-algebras.*

5.5 (5/11) A_n-operad stuff

Let X be a space. Given an operation $\mu : X^2 \to X$, we may ask if it is...

(1) Associative *up to first homotopy*, i.e. we can choose an equivalence $\mu(\mu(-,-),-) \simeq \mu(-,\mu(-,-))$. A choice is the same data as a path $\mathcal{K}(3) = I \to \mathrm{Fun}(X^3, X)$ from $(ab)c$ to $a(bc)$.

(2) (Harder) Associative *up to second homotopy*, i.e. not only can we choose an equivalence $\mu(\mu(-,-),-) \simeq \mu(-,\mu(-,-))$, but can do so in such a way that the resulting homotopy equivalences between e.g. $\mu(\mu(\mu(-,-),-),-)$ and $\mu(-,\mu(\mu(-,-),-))$ are themselves realted by a higher homotopy equivalence. (A "second order" homotopy equivalence.) Such a choice amounts to (A) the structure described in (2), plus (B) a "compatible" map from the solid pentagon $\mathcal{K}(4)$ to $\mathrm{Fun}(X^4, X)$ which sends the vertices to the various parenthetizations of $abcd$. See (5/9).

$\vdots$

(∞) (Even harder) Associative *up to all higher homotopies*, i.e. we can choose a morphism $\mathcal{K} \to \mathcal{E}$

nd$_X$ such that the path $\mathcal{K}(2) = I \to \mathcal{E}$

nd$_X$ connects $(ab)c$ and $a(bc)$.

(?) (Too hard) *Strictly* associative, i.e. $\mu(\mu(-,-),-) = \mu(-,\mu(-,-))$.

The structure of (∞) on X is intuitively captured by a choice of $\mathcal{K}$-algebra structure on X ($\Longleftrightarrow$ a choice of A_∞-algebra structure). We saw that if $X \simeq \Omega Y$ and we take $\mu =$ loop concatenation, then X has a $\mathcal{K}$-algebra structure. Moreover, up to (some notion of equivalence between $\mathcal{K}$-algebras?), all grouplike $\mathcal{K}$-algebras arise from a loop space and loop concatenation. That's the $n = 1$ case of May's recognition principle.

We're often thinking about the case where μ is only associative up to n-th homotopy for some $n < \infty$. For example, *homotopy associative H-spaces* are precisely those with an operation that is associative up to first homotopy.[3] (And, in my notation, there is something which might be called "associative up to no homotopies," whose algebras are precisely not-necessarily-homotopy-associative H-spaces.)

We want to characterize these structures with operads too. This means taking the data of an A_∞-space and truncating its coherence data at the n-th level. So for example, an "A_3-space" should be a based space X, a homotopy monoidal structure $\mu : X^2 \to X$, and a homotopy equivalence $\mathcal{K}(3) = I \to \mathrm{Fun}(X^3, X)$ from $(ab)c$ to $a(bc)$. Then an "A_n-operad" should be an operad whose algebras are precisely A_n spaces, modulo weak homotopy equivalence.

Our flagship A_∞-operad is the Stasheff operad $\mathcal{K}$. Each associahedron $\mathcal{K}(i)$ is a CW complex (in fact a simplicial complex) in an obvious way. My first thought was, "maybe taking the $(n-2)$-skeleton of every $\mathcal{K}(i)$ will produce an A_n-operad." After all, taking e.g. $n = 3$, if we take the 1-skeleton of $\mathcal{K}$, call it $\mathcal{K}_3$, then an algebra over this operad will have a homotopy monoidal structure (specified by where we send $\mathcal{K}_3(2) = *$) that is homotopy associative (specified by where we send $\mathcal{K}_3(3) = I$), but it will NOT be "associative up to second homotopy" since $\mathcal{K}_3(4)$ is the *hollow* pentagon. (The image of its interior is what *would have* specified the "second-order homotopy coherence" for associativity.)

BUT, this thing $\mathcal{K}_3$ is NOT an operad. For a simple reason: if $\mathcal{K}_3(4)$ contains intervals, then composition necessitates that $\mathcal{K}_3(5)$ contains products of intervals. (We haven't precisely defined composition, but no matter.) But this proposal for $\mathcal{K}_3$ is such that $\mathcal{K}_3(5)$ has no products of intervals, so it doesn't work.

I don't think we're far off, though. There's a *free operad* construction on "collections of objects" (yet-undefined) that I suspect will recover an operad from an appropriate n-truncation of $\mathcal{K}$ which should be rightfully called an A_n-operad.

5.6 (5/13) Rings via operads?

Tangent today. So far, I've thought of operads as devices for studying the structure of "an operation with coherences" on a given object. We also care about *rings* and *ringlike* structures. These have TWO operations, plus coherence data for identities that may involve BOTH operations at once. (Distributivity.)

Question 5.18. Can we use operads to capture the structure of rings?

The answer depends on the base category! We cannot use operads to characterize rings "all at once," i.e. in Set. We can instead try changing our base category to "handle" some of one operation first, though—we'll find that we can recover rings as algebras over operads in Ab.

Let's start in Set. We want an operad $\mathcal{R}$ such that $\mathcal{R}$-algebras are precisely rings. Or commutative rings—it won't matter, since $\mathcal{R}$ does not exist in either case. I'll give two proofs why.

The first proof is moral, not a real proof. Suppose $\mathcal{R}$ is such that $\mathcal{R}$-algebras are precisely rings. Then there should be "addition" and "multiplication" elements $A, M \in \mathcal{R}(2)$ such that $\mathcal{R}$'s compositions realize the identity $a(b + c) = ab + ac$. The issue is that the 3-ary function $ab + ac$ calls a in more than one input spot. **This is NOT expressible using an operad's notion of composition.** The closest we can get is $\gamma(A, M, M) \in \mathcal{R}(4)$, but this is $ab + cd$, not really what we wanted. Thus, we cannot impose distributivity.

Here's a formal proof. First note: there's an operad $\mathcal{S}\mathrm{et}^\times$ whose i-th object is a set of size i and whose composition arises from the Cartesian product. For an operad $\mathcal{C}$, a $\mathcal{C}$-algebra is precisely a morphism $\mathcal{C} \to \mathcal{S}\mathrm{et}^\times$. The category Operad has a terminal object whose algebras are monoids,[4] so there's a functor $S : \mathsf{Monoid} \to \mathsf{Alg}_\mathcal{C}$. It takes each map $* \to \mathcal{S}\mathrm{et}^\times$ and postcomposes it with the unique map $\mathcal{C} \to *$. **I THINK** the functor $S : \mathsf{Monoid} \to \mathsf{Alg}_\mathcal{C}$ just takes a monoid, forgets its structure, and endows it with the structure of a $\mathcal{C}$-algebra; THUS, this functor fits into a commutative triangle with the forgetful functor. BUT, the following proposition says that there are NO functors $\mathsf{Monoid} \to \mathsf{Ring}$ commuting with the forgetful functor, so there must not exist an operad $\mathcal{C}$ in Set such that $\mathsf{Alg}_\mathcal{C} \cong \mathsf{Ring}$.

Proposition 5.19. *There are no functors* $\mathsf{Monoid} \to \mathsf{Ring}$ *commuting with the forgetful functors to* Set.

Zhen Lin proves this in his MSE answer here.

So there are no operads in Set whose algebras are rings. Here's a workaround: we can replace Set with a category of objects which have addition already built in, that would be Ab. Then we can finish the job with an operad in Ab and its algebras will recover Ring.

[3]Actually, this is wrong as I've described "associativity up to n-th homotopy," but I'm just trying to give some intuition for A_n-algebra structures, so I'll gloss over this point as it'll all work when I actually get to A_n-operads.

[4]If we're thinking of symmetric operads, the terminal operad is $\mathcal{C}\mathrm{omm}$. If we're thinking of plain operads, it is $\mathcal{A}\mathrm{ssoc}$.

Let $\mathcal{C}\mathrm{omm}$ denote the commutative monoid operad in Ab. Its objects are trivial groups and its actions are trivial also.

Suppose G is a $\mathcal{C}\mathrm{omm}$-algebra, i.e. an abelian group with a map $f : \mathcal{C}\mathrm{omm} \to \mathcal{E}$

nd$_G$. This data distinguishes two elements $1_G := f_0(*) \in G$ and $\times := f_2(*) \in \mathrm{Hom}(G^2, G)$. Since $\mathcal{C}(1) = \{*\}$, we have that

$$\gamma(\times, \mathrm{id}, 1_G) = \gamma(\times, 1_G, \mathrm{id}) \in \mathcal{C}(1).$$

Since f must preserve composition, we get that $g \times 1_G = 1_G \times g = g$ for every $g \in G$. (How to finish???)

OK, in the above I started with and tried to show its algebras in Ab are commutative rings, I should've started with $\mathcal{A}\mathrm{ssoc}$ and tried to show its algebras in Ab are rings, but I'm getting tired. (I couldn't figure out how to get the distributive property anyway...)

Anyway, I'll remark that we do something similar in homotopy theory. We want a good notion of "spectra that are ringlike up-to-homotopy," and we must play the same game: find a category of objects with addition built-in (that would be Sp) and then take algebras over an operad for commutativity "up to homotopy." (I think that's little disks?)

We are foreshadowing!

5.7 The rest of May

There are some notes I have not texed, and I also spent some time working on a condensed math seminar I'll be organizing at the University of Chicago in June. (I will also be there, probably studying more operad stuff with Peter May.) Maybe I will upload the rest of my May notes eventually?

6 June

6.1 (6/13) June activities, monoidal categories

There are a few things going on.

1. Peter is giving lectures on operads and algebraic K-theory.
2. I'll probably be reading some of these lecture notes about algebraic K-theory.
3. I'm organizing a seminar on elementary condensed math. (And teaching quite a bit of basic category for that, as well as for other REU participants who just want to learn basic category theory.)

All this will be taking up most of my time. And it all somehow relates somehow to my goal of understanding higher category theory, especially (2). So I'll be sporadically writing here my inner monologue as I learn/do this stuff.

Today Peter spoke about A_∞-spaces. I already wrote about those. But Peter also mentioned monoidal categories, and this led me to a little question we were not sure about.

Let me get to explaining my thought.

Definition 6.1. A *monoidal category* is the data of a category C together with

- A functor $\otimes : \mathsf{C} \times \mathsf{C} \to \mathsf{C}$ called the *tensor product*;
- A distinguished object $\mathbf{1} \in \mathsf{C}$ called the *unit*;
- Isomorphisms $\lambda : \mathbf{1} \otimes - \to -$ and $\rho : - \otimes \mathbf{1} \to -$, we call the *left/right unitor* and whose components we denote λ_X, ρ_X; and
- An isomorphism
$$A : (- \otimes -) \otimes - \cong - \otimes (- \otimes -)$$
which we call the *associator* and whose components we write $A_{X,Y,Z}$.

And these data must have the following properties.

- The *triangle identity* holds:
$$\begin{tikzcd}
(X \otimes \mathbf{1}) \otimes Y \arrow[rr, "A_{X,\mathbf{1},Y}"] \arrow[dr, "\rho_X \otimes \mathbf{1}"'] & & X \otimes (\mathbf{1} \otimes Y) \arrow[dl, "\mathbf{1} \otimes \lambda_Y"] \\
& X \otimes Y &
\end{tikzcd}$$
- The *pentagon identity* holds:

Put it in?

Follow-up after discussion

6.2 (6/15) Initial, final objects

I'm thinking about initial/final objects so that I can define a stable ∞-category .

Topological categories are one model for ∞-categories. In a topological category, one potential definition of a final object is obvious: it is a final object in the underlying ordinary category. Call this a "strict final object." Consider (say) the topological category $CGHaus$. There, the point $*$ is a strict final object. However, there are homotopy-equivalent spaces (i.e., contractible spaces) which are not isomorphic to $*$ in $CGHaus$—that's not good since an ∞-categorical definition (i.e. of final objects) should be homotopy-invariant.

But this kind of immediately indicates the (first) correct definition.

Definition 6.2. An object of a topological category C is called *final* if it is final in hoC (regarded as enriched over hoCW). Thus, an object X is final iff $\mathrm{Hom}_{\mathsf{hoC}}(Y, X) \in \mathsf{hoCW}$ is contractible for every Y.

Why weakly contractible in hoCW? Why not final as an object in the category? Is there a different

How to port this idea to quasicategories (which we must do in a way that results in the same notion upon passage to the homotopy category)? A geometric definition: if $x \in \mathsf{C}$ is terminal, then everything "has an arrow to x," so the "collection of arrows to x" should be "like a deformation retract of C." Maybe that made sense, here's the definition.

Definition 6.3. Let x be a vertex in a quasicategory C. We say x is *intial* if $\mathsf{C}_{x/} \to \mathsf{C}$ is a trivial fibration. Likewise, we say x is *terminal* if $\mathsf{C}_{/x} \to \mathsf{C}$ is a trivial fibration.

Eventually define over and under-categories.

Remark 6.4. By definition, a vertex $x \in \mathsf{C}$ is terminal $\iff$ every map $f : \partial\Delta^n \to \mathsf{C}$ such that $f(n) = x$ extends to a map $f : \Delta^n \to \mathsf{C}$. Dually for terminal objects.

Remark 6.5. If C is a nerve, restriction $\mathrm{Hom}(\Delta^n, \mathsf{C}) \to \mathrm{Hom}(\partial\Delta^n, \mathsf{C})$ is an equivalence for $n \geq 3$. Therefore the terminal objects of nerves are the terminal objects of their underlying categories.

Proposition 6.6. *Let $\mathrm{Hom}^R_{\mathsf{C}}(x, y)$ denote the ∞-category of right-fibrations $x \to y$ in an ∞-category C. It is a Kan complex. Furthermore, an object y is terminal $\iff$ $\mathrm{Hom}^R_{\mathsf{C}}(x, y)$ is contractible for every object x. (Dually for initial objects.)*

Proposition 6.7. *An object y is terminal $\iff$ $\mathrm{Hom}_{\mathsf{C}}(x, y)$ is contractible for every object x. (Dually for initial objects.)*

Proposition 6.8. *Let C denote a quasicategory and C' the full subcategory of final objects. Then C' is a contractible Kan complex. (Likewise for the full subcategory of initial objects.)*

6.3 (6/18) Stable ∞-categories

Definition 6.9. Let C denote an ∞-category. We say an object in C is a *zero object* if it is initial and terminal. If C has a zero object, we call C *pointed*.

If C is pointed, its full subcategory of zero objects is a contractible Kan complex, so zero objects are unique up to equivalence.

Remark 6.10. As in the ordinary case, an ∞-category C is pointed if and only if it has an initial object $\varnothing$, a terminal object 1, and a morphism $1 \to \varnothing$.

Remark 6.11. If C is pointed, then $\mathrm{Hom}_{\mathsf{C}}(X, 0) \times \mathrm{Hom}_{\mathsf{C}}(0, Y)$ is a contractible Kan complex. The natural map $\mathrm{Hom}_{\mathsf{C}}(X, 0) \times \mathrm{Hom}_{\mathsf{C}}(0, Y) \to \mathrm{Hom}_{\mathsf{C}}(X, Y)$ then locates a unique element $0 \in \mathrm{Hom}_{\mathsf{hoC}}(X, Y)$. We also call this the *zero morphism*.

Definition 6.12. An ∞-category C is called *stable* if it has the following properties.

(1) C is pointed.

(2) Every morphism in C admits a fiber and cofiber.

(3) "Every morphism is the cokernel of its kernel and the kernel of its cokernel."

Condition (3) is imprecise, but that is to make obvious the analogy with abelian categories. OK now let me make it precise. This will also mean defining the things in (2).

Actually, triangulated categories – understand this?

Definition 6.13. Let C be a pointed ∞-category and $f : X \to Y$ a morphism. A *kernel* or *fiber* of f is a homotopy pullback of f with $0 \to Y$. We define a *cokernel* or *cofiber* of f dually. As in the following diagram.

$$\begin{tikzcd} \ker(f) \arrow[r, dashed] \arrow[d, dashed] & 0 \arrow[d] \\ X \arrow[r, "f"] & Y \end{tikzcd} \qquad \begin{tikzcd} X \arrow[r, "f"] \arrow[d] & Y \arrow[d, dashed] \\ 0 \arrow[r, dashed] & \mathrm{coker}(f) \end{tikzcd}$$

Definition 6.14. Let C be a pointed ∞-category. A *triangle* (g, f) in C is a commutative diagram

We write it (g, f) but a triangle consists of more data – write that in

$$\begin{tikzcd} X \arrow[r, "f"] \arrow[d] & Y \arrow[d, "g"] \\ 0 \arrow[r] & Z \end{tikzcd}$$

We say a triangle (g, f) is a *fiber sequence* (resp. *cofiber sequence*) if it is a pullback (resp. pushout).

Remark 6.15. If $f : X \to Y$ is a morphism in C, then a fiber of f is precisely a fiber sequence (f, g) for some g. Likewise, a cofiber of f is any cofiber sequence (g, f).

Definition 6.16. A category C is called *stable* if it has the following properties.

(1) C is pointed.
(2) C admits all fibers and cofibers.
(3) A triangle in C is a fiber sequence $\iff$ it is a cofiber sequence.

Remark 6.17. Condition (3) is like "cokernels of kernels are isomorhic to kernels of cokernels." Is there a precise restatement of (3) in this vein that isn't a mouthful?

Remark 6.18. Stability is a property, not a structure.

Remark 6.19. At least for me, "stability" here is used with an eye toward "stablizing the loop and suspension functors on Top*." And indeed, spectra will appear as precisely the stabilization of Top* with respect to those functors.

Definition 6.20. Let C denote a pointed ∞-category.

- If C admits cofibers, there is determined a *suspension functor* Σ: C $\to$ C which informally associates to X its homotopy pushout against two zero maps

$$\begin{tikzcd} X \arrow[r] \arrow[d] & 0 \arrow[d] \\ 0' \arrow[r] & \Sigma X \end{tikzcd}$$

- If C admits fibers, there is determined a *loop-space functor* Ω: C $\to$ C which informally associates to X its homotopy pullback along two basepoint-inclusions

$$\begin{tikzcd} \Omega X \arrow[r] \arrow[d] & 0 \arrow[d] \\ 0' \arrow[r] & X \end{tikzcd}$$

Remark 6.21. The formal definition of Σ, Ω in general takes some work, simply because it takes work to give any concrete description of functors between ∞-categories. Lurie's brief description of Σ, Ω begins in HA p. 23 at the bottom. A very nice unpacking of this construction is given by Alberto García-Raboso in this article.

The bottom line is that we get a "functorial construction of cofibers," which means a functor $\mathrm{Fun}(\Delta^1, \mathsf{C}) \to \mathrm{Fun}(\Delta^0, \mathsf{C}) \cong \mathsf{C}$ which on objects assigns a cofiber to each morphism in C. This is quite nice—there is really no good nice way to do this for triangulated categories. Grothendieck wrote a 1,976-page manuscript on "derivators" which are (a kind of?) "tool" to handle this problem that is just too hard and more concisely dealt with using ∞-categories. (Do people use derivators?) Grothendieck's manuscript affirms the belief that this is *definitely a problem worth solving*, in any case.

Learn triangulated categories and discuss the relation to stable ∞-categories; essentially, HA p. 24-

6.4 (6/20) Idempotents

Trying to read some algebraic K-theory papers, I come across *idempotent completions* (for ∞-categories). I've heard of these before but never really thought about them. I'll think about them today.

I'm looking at Kerodon Section 03Y9 and HTT 4.4.5 for this.

First we review some ordinary category theory.

Retracts should be maps fixing sub-sets/spaces/objects. Formally, denoting by C an ordinary category, we say *Y is a retract of X* if there exists some $r : X \to Y$ factoring $\mathrm{id}_Y : Y \to Y$, i.e. some diagram

$$\begin{tikzcd} & X \arrow[dr, "r"] & \\ Y \arrow[ur, "i"] \arrow[rr, "\mathrm{id}_Y"'] & & Y \end{tikzcd}$$

Remark 6.22. Let Ret or Idem^+ denote the category consisting of an "abstract retract diagram." (See Construction 03YB.) There is a tautological bijection $\{$functors $\mathsf{Ret} \to \mathsf{C}\} \cong \{$retract diagrams in $\mathsf{C}\}$.

Remark 6.23. See that $i \circ r : X \to X$ is an idempotent. In fact, this idempotent canonically determines X up to isomorphism: X is the equalizer of id_X and $i \circ r$. One may ask about the converse: if $\phi : X \to X$ is idempotent, does X have a retract Y such that $\phi = i \circ r$? (This Y is uniquely determined if it exists.) If this is the case, i.e. if the injection

$$\{\text{retracts of } X\}/\cong \quad \hookrightarrow \{\text{idempotent morphisms } X \to X\}$$

is a bijection for all X, we say that C is *idempotent complete*.

Definition 6.24. In general, we say an idempotent $\phi : X \to X$ is *split* if it arises from a retraction of X, i.e. if $\phi = i \circ r$ for some $r : X \to Y$ and $i : Y \to X$ satisfying $r \circ i = \mathrm{id}_Y$.

Remark 6.25. Thus, an ordinary category is idempotent complete $\iff$ every idempotent splits.

Proposition 6.26. *If an ordinary category C has equalizers or coequalizers, then it is idempotent complete.*

What's ado about retracts and idempotents in ∞-categories, then? Whatever they are, they should become ordinary retracts/idempotents upon passage to the homotopy category. Lurie explains two reasons that this is insufficient, though.

Explain the reasons? Not super important I think.

The ordinary story suggests that there should be a relationship (correspondence!) between ∞-categorical retracts and idempotents. Let's start with retracts.

Let C denote an ∞-category and X an object of C.

Definition 6.27. We say an object $Y \in \mathsf{C}$ is a *retract of X* if there exists $r : X \to Y$ and $i : Y \to X$ such that $r \circ i = \mathrm{id}_Y$, i.e. there exists a 2-cell witnessing that composition.

Proposition 6.28. *An object $Y \in \mathsf{C}$ is a retract of $X \iff$ there exists a functor $F : N\mathsf{Ret} \to \mathsf{C}$ taking the "abstract retract object" to Y and the other object to X. (Lurie would say the F "exhibits Y as a retract of X.")*

As in the ordinary case, we can classify retracts using (split)idempotents.

Definition 6.29. Define the category Idem to have one object $\tilde{X}$ and one non-identity morphism $e : \tilde{X} \to \tilde{X}$ with composition law $e \circ e = e$.

Definition 6.30. Let C be an ∞-category. An *idempotent* in C is a functor $F : N_\bullet\mathsf{Idem} \to \mathsf{C}$.

Remark 6.31 (We recover the ordinary case)**.** Suppose that C is an ordinary category. Since the nerve functor is fully faithful, functors $N_\bullet\mathsf{Idem} \to \mathrm{N}_\bullet\mathsf{C}$ are in bijection with functors $\mathsf{Idem} \to \mathsf{C}$. The former are precisely the idempotents in an ∞-category by definition, and the later are idempotents in C by inspection. Thus, ∞-categorical idempotents recover ordinary ones.

Definition 6.32. Let C be an ∞-category. An idempotent $I : N_\bullet\mathsf{Idem} \to \mathsf{C}$ is called *split* if there exists a functor $R : N_\bullet\mathsf{Ret} \to \mathsf{C}$ (a "retract diagram in C) extending I along the obvious functor $N_\bullet\mathsf{Idem} \hookrightarrow N_\bullet\mathsf{Ret}$.

Remark 6.33. These, too, recover ordinary idempotents.

Proposition 6.34. *Splittings of an idempotent are essentially unique. Precisely, we mean that restriction*

$$T : \mathrm{Fun}(N_\bullet \mathsf{Ret}, \mathsf{C}) \to \mathrm{Fun}(N_\bullet \mathsf{Idem}, \mathsf{C})$$

is fully faithful. (By definition, its essential image is the full subcategory of split idempotents in C*.)*

Proof. Idem sits naturally inside Ret. That's kind of the whole point. Clearly, Ret is generated by Idem under retracts; this implies that any $F :_\bullet \mathsf{Ret} \to \mathsf{C}$ is left and right Kan extended from its restriction $T(F) : N_\bullet \mathsf{Idem} \to \mathsf{C}$ (Proposition 03YQ). This lets us apply Corollary 030S (which I haven't really unpacked) to conclude that T is a trivial Kan fibration. Trivial fibrations of Kan complexes are fully faithful. (Consult Rune's notes.) □

Prove that given isomorphic idempotents $I \cong I'$, one splits iff the other splits. This isn't too hard—I just want to think about isofibrations?

Understand 7.3.6.15 one day?

Remark 6.35. Now let me rehash and give some ideas:

- If C is an ordinary category, we may speak of *idempotents* and *retracts*. By definition, given an object X, retracts of X (up to isomorphism) biject with split idempotents $X \to X$. But not every idempotent splits. Varying X, we get a fully-faithful functor $\mathrm{Fun}(\mathsf{Ret}, \mathsf{C}) \hookrightarrow \mathrm{Fun}(\mathsf{Idem}, \mathsf{C})$. The inverse problem amounts to taking a equalizer or coequalizer; if $\phi : X \to X$ is idempotent, then the (co)equalizer of ϕ with id_X realizes a retract of X (and retracts in general are uniquely determined).

- If C is an ∞-category, we define *retracts of* X as objects Y with a certain property analogous to ordinary retracts. If Y is a retract of X in this sense, there is determined a functor $N_\bullet \mathsf{Ret} \to \mathsf{C}$, unique up to isomorphism. We do a bit more work for idempotents: the most obvious "property" of idempotency actually is undesirabley ambiguous, so we *define* idempotents as functors $N_\bullet \mathsf{Idem} \to \mathsf{C}$. As in the ordinary case, we get a fully faithful (i.e. a trivial Kan fibration) functor

  Write this out?

  $$\mathrm{Fun}(N_\bullet \mathsf{Ret}, \mathsf{C}) \to \mathrm{Fun}(N_\bullet \mathsf{Idem}, \mathsf{C}).$$

  We can characterize its essential image as those idempotents with the *property* of splitness, defined analogously as in the ordinary setting. Once again, the inverse problem (i.e. determining whether an idempotent splits, i.e. extending some idempotent $N_\bullet \mathsf{Idem} \to \mathsf{C}$ to a retract diagram in a Kan-ny way?) amounts to finding the limit or colimit of $N_\bullet \mathsf{Idem} \to \mathsf{C}$.

- Also see Section 03Y9.

6.5 (6/21) I hate idempotents today. K-theory?

I physically cannot think about idempotents after yesterday. Luckily there is a different thing I also need to understand for modern algebraic K-theory, that being (modern?) algebraic K-theory. I already understand a little bit.

Let me collect some references.

- Rune Haugseng's 2010 notes The Q-construction for stable ∞-categories.
- The MO thread Motivation/interpretation for Quillen's Q-construction? and some of what's referenced therein.

Let's start somewhere classical. In *The Geometry of Iterated Loop Spaces*, May defined the little cubes operads $\mathcal{C}_k$, the prototypical E_k-operad. Given a space X, we can form the *free* $\mathcal{C}_k$*-algebra* $\mathcal{C}_k[X]$. Any k-fold loop space is a $\mathcal{C}_k$-algebra, in particular $\Omega^k \Sigma^k X$. We get a natural $\mathcal{C}_k$-algebra map $e : \mathcal{C}_k[X] \to \Omega^k \Sigma^k X$ from the natural map $X \to \Omega^k \Sigma^k X$ of spaces.[5] This e is a weak equivalence $\iff$ X is grouplike. This is the *approximation theorem*, which is used to prove the following.

What monoid structure?

Theorem 6.36 (May's recognition theorem). *X is a grouplike $\mathcal{C}_k$-algebra $\iff$ X is weakly equivalent to $\Omega^k Y$ for some Y.*

[5]Our spaces are based, of course. We *may* also want a based variant of the free $\mathcal{C}_k$-algebra on a space?

Let me sketch the hard direction of this (not going into detail about the hard parts). Suppose as given X a $\mathcal{C}_k$-algebra. There is a map of monads $\alpha_k : C_k X \to \Omega^k \Sigma^k X$ induced by the natural map $X \to \Omega^k \Sigma^k X$. The map $X \to \Omega^k \Sigma^k X$ then induces a map $X \to \Omega^k B(\Sigma^n, C_k, X)$, where that last space is a two-sided bar construction. We think of it as "like a k-fold delooping of X." Furthermore, α_k was a group completion,[6] so $X \to \Omega^k B(\Sigma^n, C_k, X)$ is a group completion also. (Here, the monoid structure on $\pi_0 X$ is induced by the $\mathcal{C}_k$-algebra structure.) Therefore, it is a weak equivalence if and only if X is grouplike, in which case we've realized X as a k-fold loop space.

What's the relation to K-theory? Take $k = 1$. The above says that an A_∞ structure on X lets us "deloop" X (that delooping being $B(\Sigma^n, C_k, X)$ above) and construct a group completion $X \to \Omega BX$. In other words, you get what you might expect: a "homotopy" monoid X has a "homotopy" group completion $X \to \Omega BX$.

Hey, that was kind of cool. Think more about this two-sided bar construction? All this operad-monad stuff fits together nicely.

Trying to sort my thoughts out about group completions is actually a bit of a pain.[7] I'm going to writhe in my stupidity and stop writing for today. Also see this MO post.

Write more about group completions?

6.6 (6/26) Structure of $\mathrm{Hom}_{\mathsf{C}}(-,-)$ for (various adjectives) categories

Recall that a *topological category* is one enriched over CG. In some form, Whitehead's theorem says that every space X is weakly equivalent to a CW-complex X', unique up to a unique weak equivalence. Therefore, $X \mapsto [X] := X'$ defines a functor $\theta : \mathsf{CG} \to \mathsf{hoCW}$. This functor exhibits hoCW as a localization $\mathsf{CG}[w^{-1}]$ at weak equivalences. It happens that θ preserves products. By general procedure, given such a nice functor θ, any CG-enriched category may now be canonically enriched (via θ) over hoCW. Now given a topological category C, we define hoC to have the same objects and we define $\mathrm{Hom}_{\mathsf{hoC}}(X,Y)$ to be the CW-approximation $[\mathrm{Hom}_C(X,Y)] \in \mathsf{hoCW}$.

Proposition 6.37. *The homotopy category of a topological category is canonically enriched (via Whitehead's theorem) over* hoCW*, i.e. the homotopy category of spaces.*

Now let C denote a quasicategory. Recall that a *morphism* $f : X \to Y$ between $X, Y \in \mathsf{C}$ is a 1-simplex such that $d_1(f) = X$ and $d_0(f) = Y$. Morphisms $f : X \to Y$ are in bijection with the vertices of the *mapping space* $\mathrm{Map}_{\mathsf{C}}(X,Y)$ defined as the fiber product

$$\begin{tikzcd} \mathrm{Map}_{\mathsf{C}}(X,Y) \arrow[r, dashed] \arrow[d, dashed] & \mathrm{Fun}(\Delta^1, \mathsf{C}) \arrow[d] \\ \Delta^0 \arrow[r, "*\mapsto(X,Y)"'] & \mathrm{Fun}(\partial\Delta^1, \mathsf{C}) \cong \mathsf{C} \times \mathsf{C} \end{tikzcd}$$

(One could also define this as a fiber product over $\mathrm{Fun}(\Delta^1, \mathsf{C})$.) This defines $\mathrm{Map}_{\mathsf{C}}(X,Y)$ as a simplicial set. In fact, $\mathrm{Map}_{\mathsf{C}}(X,Y)$ is a Kan complex. The shortest proof of this proceeds as follows.

- The morphism $i : \partial\Delta^1 \to \Delta^1$ is a monomorphism and is bijective on vertices. This implies that the restriction map $i^* : \mathrm{Fun}(\Delta^1, \mathsf{C}) \to \mathrm{Fun}(\partial\Delta^1, \mathsf{C})$ is conservative (Charles' notes, 37.1; this works for map of simplicial sets.) Therefore the (trivial) map $\mathrm{Map}_{\mathsf{C}}(X,Y) \to \Delta^0$ is conservative. Since every edge in Δ^0 is an isomorphism, it must have been so in $\mathrm{Map}_{\mathsf{C}}(X,Y)$. Thus $\mathrm{Map}_{\mathsf{C}}(X,Y)$ is an ∞-groupoid, i.e. a Kan complex by Joyal's theorem.

Proposition 6.38. *A quasicategory is "enriched over spaces" in the sense that any mapping space between two objects is a Kan complex.*

This is sort of internal to a particular quasicategory. In the *Joyal model structure* on sSet, the fibrant objects are precisely the quasicategories.

There's a recurring theme of "infinity-categorical things have a *space* of morphisms between objects." Whatever "space" means to you. Now, let C denote a simplicially-enriched category. Simplicial sets are "space-ish," but we may ask about those simplicially-enriched categories whose hom-sets are all Kan complexes. That makes them *really* space-like.

[6] Is it just a group completion of spaces? Is there more to this statement?

[7] In the process, I found this nice article from Sanath Devalapurkar. Also these notes of Dylan Wilson's.

Proposition 6.39. *There is a model structure on the category of simplical categories, called the Bergner model structure whose fibrant objects are precisely the categories enriched in Kan complexes. Furthermore, there is a Quillen equivalence between this model category and* sSet *with the Joyal model structure.*

This is in HTT? I have heard it is not easy. I wonder if the proof uses anywhere that (by construction) the most natural definition of "the mapping space between objects of a quasicategory" form Kan complexes.

6.7 (6/28) Anima and other examples of ∞-categories via N^c

First, I just wanted to record a characterization of (ordinary) limits and colimits I learned. Let C denote a small category, J a small diagram category, and $\Delta : \mathsf{C} \to \mathrm{Fun}(\mathsf{J}, \mathsf{C})$ the "diagonal" functor $c \mapsto (\mathsf{J} \mapsto c)$. Furthermore, assume that C has all J-shaped colimits.

The assignment $(\phi : \mathsf{J} \to \mathsf{C}) \mapsto \varinjlim \phi$ defines a functor $\mathsf{C}^{\mathsf{J}} \to \mathsf{C}$. I want to show that this is left-adjoint to Δ. There are many ways to do this (some of which are kind of interesting to think about...) Maybe the quickest way uses the following:

- A functor $F : \mathsf{A} \to \mathsf{B}$ is a left adjoint iff one can specify, for each $b \in B$, an object $G_b \in \mathsf{A}$ and a "universal arrow" $\epsilon_b : b \to F(G_b)$.

That this is fulfilled for $F = \varinjlim : \mathsf{C}^{\mathsf{J}} \to \mathsf{C}$ is true, more-or-less by definition of colimits. The rest is vaguely "determined," maybe up to choice. One could also proceed via the following.

- Given functors $F : \mathsf{A} \to \mathsf{B} : G$, an adjunction $F \dashv G$ is determined by a universal natural transformation $\eta : \mathrm{id}_{\mathsf{A}} \Longrightarrow GF$.

Again, in our case, this is immediate by the definition of $\varinjlim$. Namely, given a diagram $\phi : \mathsf{J} \to \mathsf{c}$, its colimit $\varinjlim \phi$ is a universal cone under ϕ, which is precisely the data of a universal morphism (in c^{J}) $\phi \to \Delta_{\varinjlim \phi}$. Letting ϕ vary, we assemble η.

Proposition 6.40. *If* C *is small and* J *is a small diagram category, and furthermore* C *is* J*-cocomplete, then* $\varinjlim : \mathsf{C}^{\mathsf{J}}$ *is left-adjoint to the diagonal* $\Delta : \mathsf{C} \to \mathsf{J}$.

Analogously, the limit functor is right-adjoint to Δ. I think, in fact, that the existence of any left (resp. right) adjoint to Δ is *equivalent* to cocompleteness (resp. completeness) of C.

Anyway, that is not what I wanted to focus on. An *anima* is an ∞-category whose homotopy category is a groupoid. (In other words, anima are ∞-groupoids.)

Example 6.1. Joyal's theorem says that anima are precisely the Kan complexes.

Example 6.2 (Cores are anima)**.** The core of ∞-category, i.e. "the maximal ∞-subgroupoid," is an anima. Let's recall how to construct this. If C is an ordinary category, then C^{core} is the subcategory spanned by isomorphisms in C. For C an ∞-category, we define its core as the pullback

$$\begin{tikzcd}
\mathsf{C}^{core} \arrow[r, dashed] \arrow[d, dashed] \arrow[dr, phantom, "\lrcorner", very near start] & \mathsf{C} \arrow[d] \\
N(\pi \mathsf{C}^{core}) \arrow[r] & N(\pi \mathsf{C})
\end{tikzcd}$$

Example 6.3 (Hom-sets are anima)**.** As discussed previously, if C is a quasicategory, then $\mathrm{Hom}_{\mathsf{C}}(x, y)$ is an anima.

We want an ∞-*category of anima*. I think, at some point, that I defined the ∞-category of ∞-categories explicitly, based on Charles notes. This was *not* really the nerve of the ordinary category of ∞-categories. For some reason, the nerve isn't the "right" way to think about these kinds of constructions. This is intuitive: an ∞-category should express and organize some "homotopical phenomena," so if our input is an ordinary category C (which has no higher data) then whatever we extract from it (e.g. $N(C)$) probably won't have much good homotopical information. Write about N^c one day?

So there are two problems: we start with an insufficient amount of data (an ordinary category), and if we change that (by specifying more data), we need to "upgrade" the nerve $N_\bullet(-)$ to account for the additional data. The *homotopy-coherent nerve* resolves this; it is some technically-defined functor that sends simplicially-enriched categories to simplicial sets.

Definition 6.41. The *coherent/simplicial/homotopy-coherent nerve* is some functor

$$N^c(-) : \mathsf{Cat}^{\mathsf{sSet}} \to \mathsf{sSet}.$$

Proposition 6.42 (Fabian's notes, I.14)**.** *If* C *is a category enriched in Kan complexes, then* $N^c(\mathsf{C})$ *is an* ∞*-category. Moreover, there is a canonical homotopy equivalence of Kan complexes*

$$\mathrm{Hom}_{N^c(\mathsf{C})}(x,y) \simeq \mathrm{Hom}_{\mathsf{C}}(x,y).$$

Example 6.4. The ordinary full subcategories $\mathsf{Kan}, q\mathsf{Cat} \subseteq \mathsf{sSet}$ are Kan-enriched via

$$\mathrm{Hom}_{\mathsf{Kan}}(X,Y) = \mathrm{Fun}(X,Y) \quad \text{and} \quad \mathrm{Hom}_{q\mathsf{Cat}}(X,Y) = \mathrm{Hom}_{\mathsf{sSet}}(X,Y)^{core}.$$

The *∞-category of anima* is the coherent nerve $N^c(\mathsf{Kan})$. The *∞-category of quasicategories* is the coherent nerve $N^c(q\mathsf{Cat}')$, where $q\mathsf{Cat}'$ is the Kan-enriched category formed by replacing the hom-quasicategories in $q\mathsf{Cat}$ with their maximal ∞-groupoids.

Example 6.5 (Homotopy hypothesis, Fabian I.20)**.** Let X be a space. The singular complex $S_\bullet X$ is an ∞-groupoid. This functor $S_\bullet$ is right adjoint to geometric realization. A corollary to inspection of the skeletal filtration on simplicial sets says that a geometric realization has a canonical CW structure; thus $|-| : \mathsf{sSet} \to \mathsf{Top}$ lands in CW. Top, in particular CW, is Kan-enriched via $\mathrm{Hom}(X,Y) := \mathrm{Sing}_\bullet \mathrm{Hom}(X,Y)$. We can therefore pass to the simplicial nerve, and our adjunction becomes an *equivalence*

$$\mathsf{An} := N^c(\mathsf{Kan}) \underset{\mathrm{Sing}_\bullet}{\overset{|-|}{\rightleftarrows}} N^c(\mathsf{CW})$$

That this is an equivalence is Grothendieck's *homotopy hypothesis*.

Read example $N^c(Ch(R))$, i.e. nerve of category of chain complexes of R-modules. See Fabian, I.15(e).

7 July

7.1 (7/5) Derived ∞-categories I

There are a few ways to define, think about, characterize, etc. the *derived ∞-category* of an abelian category. I suppose it's the "correct" setting to do homological algebra with complexes modulo quasi-isomorphisms. In particular, they recover the triangulated structure we use in the ordinary setting for the purposes of homological algebra. Naturally, they are related to stable ∞-categories. I'll understand this eventually. Today I'll go over my takeaways from Achim Krause's talk about them.

To an abelian category $\mathcal{A}$, we want to associate an ∞-category $\mathcal{D}_\infty(\mathcal{A})$ that encodes the homotopy theory (where quasi's should be the weak equivalences) of $\mathrm{Ch}(\mathcal{A})$.

Definition 7.1. If $\mathcal{A}$ is an abelian category, we denote by $\mathrm{Ch}(\mathcal{A})$ the category of unbounded chain complexes in $\mathcal{A}$. Note that for $A, B \in \mathrm{Ch}(\mathcal{A})$, the differential on B gives rise to a dg structure on $\mathrm{Hom}_{\mathrm{Ch}(\mathcal{A})}(A, B)$, i.e. the structure of a complex of abelian groups. Thus, $\mathrm{Ch}(\mathcal{A})$ has a canonical enrichment over $\mathrm{Ch}(\mathbb{Z})$.

There's latent homotopical data in the structure of the $\mathrm{Ch}(\mathbb{Z})$-enrichment of $\mathrm{Ch}(\mathcal{A})$. (In particular, the $\mathrm{Ch}(\mathbb{Z})$-enrichment encodes the chain-homotopy equivalences.) We have a good way for translating this into our common language of ∞-categories: consider the composite

$$K : \mathrm{Ch}(\mathbb{Z}) \xrightarrow[\tau_{\geq 0}]{} \mathrm{Ch}_{\geq 0}(\mathbb{Z}) \xrightarrow[\Gamma]{\cong} s\mathsf{Ab} \xrightarrow[\text{forgor}]{} \mathsf{Kan}. \tag{II.2}$$

The forgetful functor lands in Kan because the underlying simplicial set of a simplicial abelian group is a Kan complex. The functor Γ is one direction of the *Dold-Kan correspondence*. The functor $\tau_{\geq 0}$ is the *canonical truncation*.[8]

Proposition 7.3. *Each category in Equation (II.2) is monoidal, with monoidal product given by (from left to right): the $\otimes$ of chain complexes, the same, pointwise $\otimes$ of abelian groups, the usual product. (In fact, they are symmetric monoidal.)*

Furthermore, each functor in Equation (II.2) is lax-monoidal. (Note that Γ is not lax-symmetric monoidal, although its inverse is.)

Thus, hitting hom-objects with K describes a functor

$$\{\mathrm{Ch}(\mathbb{Z})\text{-enriched categories}\} \to \{\mathsf{Kan}\text{-enriched categories}\}.$$

And we (vaguely, I have not yet really worked this out) know how to take the latter sort of category and reformulate it as an ∞-category.

Definition 7.4. If $\mathcal{A}$ is an abelian category, define its *homotopy category* as the ∞-category

$$\mathcal{K}_\infty(\mathcal{A}) := N^c(\mathrm{Ch}(\mathcal{A})_\Delta),$$

where $\mathrm{Ch}(\mathcal{A})_\Delta$ is the simplicial category obtained by applying K above to Hom-objects. Note that this is in ∞-category since K lands in Kan rather than just sSet.

The homotopy category $\mathcal{K}_\infty(\mathcal{A})$ is nice. For example, it is finitely bicomplete. Moreover, for chain complexes A and B we have by construction

$$\pi_n \mathrm{Hom}_{\mathcal{K}_\infty(\mathcal{A})}(A, B) \cong H_n \mathrm{Hom}_{\mathrm{Ch}(\mathcal{A})}(A, B).$$

In particular, $\pi_0 \mathrm{Hom}_{\mathcal{K}_\infty(\mathcal{A})}(A, B)$ consists of chain maps $A \to B$ modulo chain-homotopy equivalence.

Remark 7.5. There is a problem. Although chain-homotopic maps are identified upon passage to $\mathcal{K}_\infty(\mathcal{A})$, quasi-isomorphisms do not become isomorphisms. For example, the following chain map is a quasi-

[8] In negative degrees, $\tau_{\geq 0}$ naively truncates a chain complex. In degree 0, it sends A_0 to $\ker(d_1)$. It leaves positive degrees unchanged. This has the effect of preserving non-negative homology. This is contrast to the most obvious *stupid truncation*, which leaves A_0 unchanged. See [?, Section 0118].

isomorphism but is not invertible in $\mathcal{K}_\infty(\mathcal{A})$.

$$\begin{tikzcd}
\cdots \arrow[r] & 0 \arrow[r] & \mathbb{Z} \arrow[r, "2"] \arrow[d] & \mathbb{Z} \arrow[r] \arrow[d, "\pi"] & 0 \arrow[r] \arrow[d] & \cdots \\
& \cdots \arrow[r] & 0 \arrow[r] & \mathbb{Z}/2 \arrow[r] & 0 \arrow[r] & \cdots
\end{tikzcd}$$

We do not beat around the bush: to fix the problem, we just DK-localize it away. Let W be the set of quasi-isomorphisms of complexes in $\mathcal{K}_\infty(\mathcal{A})$. We define the *derived ∞-category* of the abelian category $\mathcal{A}$ to be the localization

$$\mathcal{D}_\infty(\mathcal{A}) := \mathcal{K}_\infty(\mathcal{A})[W^{-1}].$$

Write about DWyer-Kan localization, existence, examples?

7.2 (7/21) Ambidexterity I

A sheaf on X is called a *classical local system* if it is locally constant. For X locally connected, we have a categorical equivalence (the $\leftarrow$ direction of which is given by taking sections)

$$\{\text{classical local systems on } X\} \cong \{\text{covering maps } p : Y \to X\}.$$

Suppose that X is connected (and locally simply connected) and choose a basepoint x_0. Given a loop $g \in \pi_1(X, x_0)$, the pullback $g^*\mathcal{F}$ is constant, hence specifies an isomorphism $g^*\mathcal{F}_0 \cong g^*\mathcal{F}_1$, i.e. an automorphism of $\mathcal{F}_x$. This is suitably homotopy invariant and compatible with loop concatenation so as to define a group action $\pi_1(X, x_0) \to \mathrm{Aut}(\mathcal{F}_{x_0})$, the *monodromy representation* of $\mathcal{F}$ at x_0. If $\mathcal{F}$ is a sheaf of sets, the monodromy representation is a $\pi_1(X, x_0)$-set. If $\mathcal{F}$ is a sheaf of complex vector spaces, then the monodromy representation is a representation $\pi_1(X, x_0) \to \mathrm{GL}(\mathcal{F}_{x_0})$ and takes its sheaf of sections.

Proposition 7.6. *If X is connected (and locally simply connected),*[9] *then associating to a local system its monodromy representation defines one way of an equivalence of categories*

$$\{\text{classical local systems of } \mathbb{C}\text{-vector spaces on } X\} \xrightarrow{\sim} \{\text{complex representations } \pi_1(X, x_0)\}.$$

Remark 7.7. The inverse associates to a representation ρ a sheaf that tautologically has monodromy representation ρ. Some sources define this directly. I think it is exactly the following construction: given a representation $\rho : \pi_1(X, x_0) \to \mathrm{GL}(V)$, one takes the associated bundle $(\tilde{X} \times V)/\pi_1(X, x_0) \to X$

Remark 7.8. If $\mathcal{F}$ is a sheaf of sets, then one gets an equivalence between local systems on X and $\pi_1(X, x_0)$-sets.

Some good links for basics about the above stuff are Wikipedia, these short notes on "local systems and constructible sheaves" by P. Achar, and Szamuely's book.

Two observations: (1) if X is not connected, then the "local systems $\leftrightarrow \pi_1$ representations" picture gets awkward,[10] and (2) we can generalize our argument that a loop based at x_0 determines an automorphism $\mathcal{F}_{x_0} \xrightarrow{\sim} \mathcal{F}_{x_0}$ in a homotopy-invariant manner. By that I mean that the argument works to show *paths* from x_0 to x_1 induce isomorphisms $\mathcal{F}_{x_0} \xrightarrow{\sim} \mathcal{F}_{x_1}$ in a functorial, homotopy-invariant manner. The point (2) actually suggests a more general definition that addresses (1).

Definition 7.9. Let X be a topological space. A *local system* on X is a functor $\Pi_1 X \to \mathsf{D}$.[11]

This recovers the definition as classical local systems in the case that X is nice. (I don't actually know what the hypotheses for this are.) That is, for sufficiently nice X, we have an equivalence of categories

$$\{\text{classical local systems on } X\} \xrightarrow{\sim} \mathrm{Fun}(\Pi_1 X, \mathsf{Set}).$$

[9] Maybe you also need Hausdorff, locally path-connected, second countable...

[10] Since you would want to account for the varying π_1 between different connected components, but for an equivalence of some sort, you would need to pick a basepoint in each component, which you don't really do?

[11] Recall that the *fundamental groupoid* $\Pi_1 X$ is the category whose objects are points of X and whose morphisms are homotopy classes of maps.

Given a local system $\mathcal{F}$, the associated functor $\Pi_1 X \to \mathsf{Set}$ acts on objects by $x \mapsto \mathcal{F}_x$. Morphisms are sent to their associated monodromy representation.

Anyway, it's clear where we're going: we will replace $\Pi_1 X$ with its untruncated, derived version $\mathrm{Sing}(X)$ and our target category with some ∞-category.

Definition 7.10. Let X denote a Kan complex. An *∞-local system* valued in C is a functor $X \to \mathsf{C}$.

Remark 7.11. Although I'm not interested in it right now, the analogy with classical local systems and covering maps holds up in the ∞-categorical setting. The slogan is that spaces are ∞-topoi, and if $\mathcal{X}$ is an ∞-topos, then $\mathrm{Shv}(\mathcal{X})$ has a full subcategory of "locally constant sheaves" on $\mathcal{X}$. For $\mathcal{X}$ nice,[12] these locally constant sheaves turn out to be equivalent to some ∞-topos of the form $\mathsf{S}_{/K}$ for some Kan complex K. If $\mathcal{X} = \mathrm{Shv}(X)$ for some nice space X, then $K = \mathrm{Sing}(X)$. Wow! I'm not actually sure the precise relation of these locally constant sheaves to local systems in this setting.

The equivalence of classical local systems on X and $\pi_1 X$-representations (which holds for connected X) says that local systems "are" representation theory (you can recover the representation theory of G by taking $X = BG$). In an ∞-local system, one has a classical local system (given by truncating $\mathrm{Sing} X$ to $\Pi_1 X$), plus higher homotopy-coherent data coming from X. This is probably the easiest step down a road toward "higher representation theory."

Just as we often do representation theory with nice groups, e.g. G finite, in higher representation theory we should begin with some finiteness conditions. Our space X (really, its homotopy type) plays the role of G, so these should be conditions on X (really, on its homotopy groups).[13]

Definition 7.12 (Various finiteness conditions)**.** Let X denote a space. Let p be a prime and $m \in \mathbb{Z}_{\geq -2}$. Say X is...

(1) *m-finite* if...

- $m \geq -2$ and X is contractible; or
- $m \geq -1$, the set $\pi_0 X$ is finite, and the fibers of $\Delta : X \mapsto X \times X$ are all $(m-1)$-finite.

(2) *π-finite* if it is m-finite for some $m \geq -2$;

(3) A *p-space* if all its homotopy groups are p-groups; and

(4) *p-finite* if it is a π-finite p-space.

Example 7.1. If G is finite, then BG is π-finite. If G is a finite p-group, then BG is p-finite.

Let me wrap up by "doing something" with all this. Here's something we do in representation theory and the study of group actions.

Example 7.2 (Representation theory in characteristic zero)**.** Suppose as given a finite group G acting on an abelian group A. We define the *invariants* and *coinvariants* of the G-action as $A^G := \{a : ag = g \text{ for all } g \in G\}$ and $A_G := A/\{ga - a : a \in A, g \in G\}$, respectively. There is a natural *norm map*

$$N_G : A_G \to A^G$$

given by $\overline{m} \mapsto \sum_g gm$. (This is a kind of "averaging.") N_G is not an isomorphism in general, but it is if A is a rational vector space (i.e., multiplication by any n is invertible). In this case, the claim is that the composite $A^G \hookrightarrow A \twoheadrightarrow A_G$ is an inverse. For this, one shows that the composites both ways are multiplication by $|G|$, which is nonzero if A is a $\mathbb{Q}$-vector space. Thus, invariants and coinvariants coincide in representation theory in characteristic zero.

Remark 7.13. It is clear from the argument in the previous example that if p divides $|G|$, then N_G is *not* an isomorphism when A is an $\mathbb{F}_p$-vector space. In particular, if p divides $|G|$, then N_G is not an isomorphism for representation theory in characteristic p. In the edge case of $|G| = p^r$, we observe something called *unipotence* that is important but which I am not focused on right now.

[12]Here, "nice" is the higher version of "locally simply connected" or something. The precise term from HA is "locally of constant shape."

[13]Also, in doing chromatic and whatnot, finiteness conditions abound for other reasons too.

Suppose that a field k has characteristic zero. As we discussed, classical local systems of k-modules (i.e., k-vector spaces) over BG are "the same thing" as G-representations over k. Directly above, we said that given a G-representation $G \to \mathrm{GL}(V)$, the norm $N_G : V_G \to V^G$ is an isomorphism. We can ask what this statement translates to if we identify our representation with a local system. If we write $\mathcal{L}$ for that local system, I *think* the norm map is recovered as some canonical comparison map $\varinjlim_x \mathcal{L}_x \to \lim_x \mathcal{L}_x$, which is an isomorphism since $char(k)$ is zero.

If k has characteristic p, I also said that the norm is no longer an isomorphism. I'm not sure if the "classical local systems $\leftrightarrow$ G-representations over k" correspondence still works. Regardless, its our definitive model for "higher representation theory" (that is, I'm telling you that ∞-local systems *are* higher representation theory). And the definition of the comparison map as $\varinjlim_x \mathcal{L}_x \to \lim_x \mathcal{L}_x$ works for ∞-local systems.

So, now we can ask about how the norm map behaves for various ∞-local systems. We should have an eye toward situations with a notion of "characteristic," since in ordinary representation theory, the characteristic dictates useful phenomena via the norm. Chromatic homotopy theory strongly suggests certain examples, wherein the classical theory generalizes in structured but unexpected and interesting ways.

Consider a G-spectrum rather than a G-abelian group. Fix a prime p. We saw that for a G-abelian group, the norm $N_G : A_G \to A^G$ is an isomorphism in characteristic zero but not p. In chromatic land, to p we associate the sequence of Morava K-theories $K(n)$ which are "like primes" or which "intermediate characteristic zero and p."

Two good questions: if the $K(n)$ have "intermediate characteristic,"

(1) Can we do "representation theory of G over these intermediary-characteristic $K(n)$" and

(2) How does the norm behave in this representation theory?

Not taking this "higher representation theory" perspective, an answer to (2) was proven in 1996 (Clausen and Akhil have a short proof here):

Theorem 7.14. *Let G be a finite group and let X be a $K(n)$-local spectrum with a G-action (i.e. a functor $BG \to \mathrm{Sp}_{K(n)}$). Then the norm map $X_{hG} \to X^{hG}$ is an equivalence in $\mathrm{Sp}_{K(n)}$.*

This is surprising! In the case $n = 0$, we have $K(0) = H\mathbb{Q}$ and we work rationally, in which case the theorem reduces to knowing that composing N_G with $X^{hG} \to X \to X_{hG}$ is multiplication by $|G|$ which is invertible (we can divide by $|G|$). But for $n > 0$, we have that $K(n)_* \cong \mathbb{Z}/p[v_n, v_n^{-1}]$ in which $p = 0$. Thus, we cannot always "divide by $|G|$." Yet the theorem persists.

Hopkins-Lurie offer an insightful interpretation of all this. Remember that representation theory of G over k "is" local systems of k-modules over BG. One thing Hopkins-Lurie show is that we *can* do "higher representation theory" of G over $K(n)$, which we understand as ∞-local systems on BG of $K(n)$-modules. Then a norm map appears and is an isomorphism, as it is the map from Theorem 7.14 (I think). In fact, in the framework of HL one may take any space X in place of BG.

Theorem 7.15. *Let X be a π-finite space and let $\mathcal{L}$ be an ∞-local system of $K(n)$-module spectra over X, i.e. an object of the ∞-category $\mathrm{Fun}(X, \mathrm{Mod}_{K(n)})$. Then there is a canonical "norm" isomorphism*

$$N_X : C_*(X; \mathcal{L}) \xrightarrow{\sim} C^*(X; \mathcal{L}).$$

Example 7.3. Taking $\mathcal{L}$ =the trivial local system, we find that if X is a π-finite space, then $K(n)_* X \cong K(n)^* X$. We think of this as some general form of Poincaré duality. This generalizes work of Greenlees and Sadofsky, who proved the statement in the case that $X = BG$ in 1996.

HL develop a general framework to understand this. I'll get to that in a bit.

7.3 (7/25) Ambidexterity II

Continuing from last time. Let me recap for my own sake.

First we talked about some classical phenomena: given a finite group G acting on an abelian group A, we may form a "norm" map $N_G : A_G \to A^G$ whose composition with the canonical $A^G \hookrightarrow A \twoheadrightarrow A_G$ is multiplication by $|G|$, which is invertible if A is a rational vector space (in which case N_G exhibits $A_G \cong A^G$). Then I schizophrenically insisted that since $\mathbb{C}$-representation theory of G is "just" local systems of $\mathbb{C}$-modules ($\iff$ $\mathbb{C}$-vector spaces) over X (one recovers representations of G by taking $X = BG$), we should *define* "higher representation theory" to be the study of ∞-local systems.[14]

That established, I stated our first real result in this philosophy: *if X is a π-finite space, then given a local system $\mathcal{L}$ of $K(n)$-modules on X, there is a "norm" isomorphism*[15]

$$\varinjlim_x \mathcal{L}_x =: C_*(X;\mathcal{L}) \xrightarrow[N_X]{\sim} C^*(X;\mathcal{L}) := \lim_x \mathcal{L}_x$$

This is like the earlier result that $N_G : A_G \to A^G$ is an isomorphism if A is a $\mathbb{Q}$-vector space. *Ambidexterity* is supposed to give some framework to find and study canonical dualities between colimits and limits like this (e.g., as we had between G-orbits A_G and G-fixed points A^G.)

Let me make an educated guess as to how we will move forward and invent context for all this:

(1) We will formulate a "norm" map for any ∞-local system $X \to \mathsf{C}$, probably with reasonable stipulations on X and/or C. This will be a map

$$\varinjlim_x \mathcal{F}_x \to \lim_x \mathcal{F}_x.$$

(2) We will do so in such a way (or perhaps it will pop out from the formalism) that the composition with the canonical map $\lim_x \mathcal{F}_x \to \varinjlim_x \mathcal{F}_x$ is an endomorphism that we should call "multiplication by the cardinality $|X|$."

(3) We will study when this endomorphism is invertible.

(4) ??? Profit

Let's get to formalizing things.

Definition 7.16. Let X be a Kan complex and C an ∞-category admitting small limits and colimits. Denote by $\delta : \mathsf{C} \to \mathsf{C}^X$ the functor which maps an object C to the constant C-valued local system $\underline{C}_X$. Given a local system $\mathcal{L} \in \mathsf{C}^X$, define

$$C_*(X;\mathcal{L}) := \varinjlim_X \mathcal{L}_x \quad \text{and} \quad C^*(X;\mathcal{L}) := \lim_X \mathcal{L}_x.$$

I'll call these the *coinvariants* and *invariants* of $\mathcal{L}$. By construction, the functors $\mathcal{L} \mapsto C_*(X;\mathcal{L})$ and $\mathcal{L} \mapsto C^*(X;\mathcal{L})$ are left/right adjoint to the constant local system functor δ, respectively. Here's a diagram expressing this.

$$X \underset{\mathcal{L}\mapsto C^*(X;\mathcal{L})}{\overset{\mathcal{L}\mapsto C_*(X;\mathcal{L})}{\leftrightarrows}} \mathsf{C}^X$$

Definition 7.17. Let X be a Kan complex and C a category with small limits and colimits. Suppose as given a natural transformation $\mu : C^*(X;-) \to C_*(X;-)$ and a map of Kan complexes $f : X \to \mathrm{Hom}_{\mathsf{C}}(C,D)$, which we identify with its induced morphism $\underline{C}_X \to \underline{D}_X$. Consider the composite

$$C \to C^*(X;\underline{C}_X) \xrightarrow{f} C^*(X;\underline{D}_X) \xrightarrow{\mu} C_*(X;\underline{D}_X) \to D.$$

We call this the *integral of f with respect to μ* and denote it by $\int_X f d\mu$.

Remark 7.18. The first map $C \to C^*(X;\underline{C}_X)$ is the unit of the adjunction $\delta \dashv (\mathcal{L} \mapsto C^*(X;\mathcal{L}))$. The last map is the counit of the other adjunction. Maybe you, as I did, ask why we choose this combination of (co)units given our adjunctions. The simple answer is that this is the only way to get a map $C \to D$ given μ and f.

Remark 7.19. Since f is a family of things in $\mathrm{Hom}_{\mathsf{C}}(C,D)$, it makes sense that the "integral of f" should be a particular thing $\int_X f d\mu \in \mathrm{Hom}_{\mathsf{C}}(C,D)$.

[14] I honestly don't know if this is the right way to think about all this. Does "higher representation theory" already mean something definitive?

[15] This is(?) a more general form of a 1996 result of Greenlees-Sadofsky that exhibits an isomorphism $X_{hG} \xrightarrow{\sim} X^{hG}$ where X is a $K(n)$-local spectrum with finite G-action (which *maybe* reduces to showing that BG exhibits self-dual $K(n)$ (co)homology).

Recall that the norm $N_G : A_G \to A^G$ was given by $\overline{m} \mapsto \sum_g mg$, which we thought of as a kind of "averaging." Averaging is kind of like integrating. So, how can we use integrals (in the sense defined above) to find a "canonical" map

$$\varinjlim_X \mathcal{L}_x \to \lim_X \mathcal{L}_x?$$

To do so means to specify a map $\mathcal{L}_x \to \mathcal{L}_y$ for each path $x \to y$, in a manner functorial in x, y (and higher coherences?). The idea is to integrate—since we want a map $\mathcal{L}_x \to \mathcal{L}_y$, we should take $C = \mathcal{L}_x$ and $D = \mathcal{L}_y$. Denoting by $P_{x,y}$ the mapping space $\mathrm{Hom}_X(x, y)$, the system $\mathcal{L}$ determines a map $\phi_{x,y} : P_{x,y} \to \mathrm{Hom}_{\mathsf{C}}(\mathcal{L}_x, \mathcal{L}_y)$. Thus, given x, y and a local system $\mathcal{L}$, a natural transformation $\mu_{x,y} : C^*(P_{x,y}; -) \to C_*(P_{x,y}; -)$ specifies a map

$$\int_{P_{x,y}} \phi_{x,y} \, d\mu_{x,y} \in \mathrm{Hom}_{\mathsf{C}}(\mathcal{L}_x, \mathcal{L}_y).$$

If $\mu_{x,y}$ is functorial in x and y, then so is the above integral, whence we get a map $Nm_X(\mathcal{L}) : \varinjlim_X \mathcal{L}_x \to \lim_X \mathcal{L}_x$. This is also functorial in $\mathcal{L}$; we get a natural transformation

$$Nm_X : C_*(X; -) \to C^*(X; -).$$

7.4 (7/30) Ambidexterity III

Given a Kan complex X and an ∞-category C, I've described a procedure for "integrating maps" $f : X \to \mathrm{Hom}_{\mathsf{C}}(C, D)$ given some $\mu : C^*(X; -) \to C_*(X; -)$. Using this, I defined a "norm" $Nm_X : C_*(X; -) \to C^*(X; -)$ given (a functorial family of, for each $x, y \in X$) some $\mu : C^*(P_{x,y}, -) \to C_*(P_{x,y}; -)$. In the classical setting wherein a finite G acts on a rational vector space A, we found that the norm is an isomorphism $A_G \xrightarrow{\sim} A^G$ since $|G|$ was invertible. We also saw that if X is a π-finite space and $\mathsf{C} = \{K(n)\text{-module spectra}\}$, then $Nm_X : C_*(X; -) \to C^*(X; -)$ is an equivalence. THis gave us, for instance, a sort of $K(n)$-local Poincaré duality $K(n)_* X \cong K(n)^* X$. "Ambidexterity" means to describe daulity phenomena like this in general.

Definition 7.20. Suppose that X is a Kan complex and $\mathsf{C} \in \mathrm{Cat}_\infty$. We say that X is C-*ambidextrous* if

(1) X is n-truncated for some $n \geq -2$,

(2) For each pair $x, y \in X$, the path space $P_{x,y}$ is C-ambidextrous, and

(3) $Nm_X : C_*(X; -) \to C^*(X; -)$ is an equivalence.

If X is C-ambidextrous then we write μ_X: $C^*(X; -) \to C_*(X; -)$ for the inverse to Nm_X.

Remark 7.21. Note that if $n \geq -1$, then X is n-truncated $\implies P_{x,y}$ is $(n-1)$-truncated. This makes Definition 7.20 an inductive definition.

Prove this in my head.

Remark 7.22. We say X is (-2)-connected if it is contractible. Then $C \mapsto \underline{C}_X$ is an equivalence $\mathsf{C} \to \mathsf{C}^X$. In that case, it has naturally isomorphic left/right adjoints. So, if X is (-2)-connected, then X is automatically C-ambidextrous.

Beck-Chevalley fibrations?

Remark 7.23 (HL p. 91, right before §4.1)**.** C-ambidexterity of X imposes conditions on X and C. It is generally a finiteness condition on X, e.g. it often occurs that X has finite homotopy groups, analogous to asking that G be a finite group when we think about G-actions.

More interestingly (to me), it is a general kind of additivity property of C, and results in a canonical "integration" or "summation" process for diagrams $X \to \mathsf{C}$.

At this point, Hopkins-Lurie discuss ambidexterity in the context of *Beck-Chevalley fibrations*, which I don't have the processing power to read right now. I'm going to see if I can just ignore it for a bit.

I'm going to the gym and will think a bit about how to proceed.

8 August

9 September

9.1 (9/15) Semiadditivity I

Some good references are this blog post and this paper, maybe these lecture notes. Parts of my notes here are just regurgitations of what I find, as usual.

Let's start somewhere familiar. A *preadditive category* is a category C with an Ab-enrichment. It is more or less standard that given any two objects A, B in a preadditive category, the properties of being their (co)product coincide. In fact, the proof only uses addition, so this is true for any CMon-enriched category. Maybe one way to understand this is to think that in any CMon-enriched category, there's a good way to "add" pairs of objects.

Nailing down what exactly it means for products and coproducts to coincide (i.e., for A, B to "have a biproduct") is subtle. More subtle than I had realized when I first learned about abelian categories and "biproducts" a long time ago. Previously, I thought that given $A, B \in \mathsf{C}$, their "biproduct" should just be the name we give to any object that is isomorphic to both their product and coproduct. But the resulting notion is not unique, not even up to isomorphism. Here's a definition.

Example?

Definition 9.1. Let C be a category with zero morphisms (e.g., a pointed category, or a CMon-enriched category). Given $A, B \in \mathsf{C}$, a *biproduct* for A and B is an object $A \oplus B$ together with maps

$$\begin{tikzcd} A \arrow[dr, "i_A"] & & B \arrow[dl, "i_B"'] \\ & A \oplus B \arrow[dl, "p_A"'] \arrow[dr, "p_B"] & \\ A & & B \end{tikzcd}$$

With the property that...

- We have $i_A p_A = \mathrm{id}_A, i_B p_B = \mathrm{id}_B, i_A p_B = 0$, and $i_B p_A = 0$,
- $(A \oplus B, p_A, p_B)$ is a product, and
- $(A \oplus B, i_A, i_B)$ is a coproduct.

This turns out to work. I point that we are crucially relying on a canonical choice of map between *any* two objects $x \to y$: choose the identity if $x = y$ and the zero map otherwise.

Theorem 9.2. *Biproducts of two objects are unique up to unique isomorphism.*

Theorem 9.3. *If* C *has finite biproducts, then the biproduct extends to a bifunctor* $\mathsf{C} \times \mathsf{C} \to \mathsf{C}$.

So in a category C with zero morphisms, a *biproduct* is an object which *coherently* satisfies two dual universal properties, and I've likened it to a "sum" of objects. Biproducts may not exist. I've told you that any CMon-enriched category has all finite biproducts. But if we're going to think of biproducts as a "sum," then maybe its natural to have a zero for this "sum," i.e. a zero object. This would, in particular, imply that zero morphisms exist. So, pointed categories seem like a good starting point to explore this notion of "a category whose objects we can add."

Definition 9.4. A category C is called *semiadditive* if it is pointed and admits all finite biproducts.

Remark 9.5. Here's a harrowing remark. We said a biproduct is a "coherent combination of a product and a coproduct." And now we're saying that pointed categories are a good place to study biproducts. But what is pointedness, i.e. what does it mean to have a zero object? A category C has a zero object precisely when the limit and colimit of the empty functor $\varnothing \to \mathsf{C}$ exist (i.e., when C has a terminal and initial object) and the two coincide (i.e., when the unique map $\varnothing \to *$ is an isomorphism, i.e. when $\mathrm{Hom}_{\mathsf{C}}(*, \varnothing) \neq \varnothing$). So to even start thinking about biproducts, we already have some sort of "double universal property" going on: namely, a coincidence of the limit and colimit of the empty functor. This is equivalent to the existence of left and right adjoints to the constant functor $\mathsf{C} \to *$ and a natural isomorphism between them.

That formulation of pointedness as a coincidence of left and right adjoints turns out to be a useful way to approach the notion of semiadditivity because the existence of biproducts can be formulated in the same fashion.

Proposition 9.6. *A category C is pointed if and only if the terminal functor $\mathsf{C} \to *$ admits naturally isomorphic left and right adjoints.*

Proposition 9.7. *Let $\Delta : \mathsf{C} \to \mathsf{C} \times \mathsf{C}$ denote the diagonal functor. Then C admits finite products (resp. coproducts) if and only if Δ admits a right (resp. left) adjoint.*

Proposition 9.8 (See this paper)**.** *Suppose that Δ admits naturally isomorphic left and right adjoints, which we refer to as a single functor $\oplus$. Consider the following condition.*

- *The unit $\mathrm{id}_{\mathsf{C}\times\mathsf{C}} \to \Delta\oplus$ of $\oplus \dashv \Delta$, which is given by the morphisms $(i_A, i_B) : (A, B) \to (A \oplus B, A \oplus B)$, is a section to the counit $\Delta\oplus \to \mathrm{id}_{\mathsf{C}\times\mathsf{C}}$ of $\Delta \dashv \oplus$, which is given by the morphisms $(p_A, p_B) : (A \oplus B, A \oplus B) \to (A, B)$.*

This condition holds if and only if C admits finite biproducts.

Ok, so we've interpreted pointedness and the admittance of biproducts as a sort of duality exhibited by quite natural functors: the terminal $\mathsf{C} \to *$ and the diagonal $\mathsf{C} \times \mathsf{C} \to \mathsf{C}$. This duality is succinctly expressed in terms of adjoints. Specifically, this duality amounts to the coincidence of left and right adjoints, plus some coherence between the evident pairs of (co)units (which degenerated in the case of $\mathsf{C} \to *$). Then semiadditivity translates into two similar duality conditions, one dependent on the other to make sense.

If you believe semiadditivity is important, then you may ask whether there are even "higher" duality conditions which C may satisfy, which we might hope are limit-colimit dualities with an interpretation using adjoints. And if C satisfies these "higher" conditions, then we should obtain some "higher" form of semiadditivity. The kicker is that semiadditivity is all about monoids, and our line of reasoning will reveal that "higher semiadditivity" means "*higher monoids.*"

Of course, I'm asking this with a lot of foresight. And of course, the full answer is ∞-categorical. But I hope the take I've sketched here provides a decent *low-level, 0-categorical* entry into the body of ideas I am trying to understand right now. I hope I can write a fuller such introduction eventually, but glad to blab it here for now.

9.2 (9/18) Semiadditivity II

Before we get weird—that is, before I talk about infinity categories—I should actually explain why semiadditive categories are interesting. I did not do that last time, I just basejumped off the hypothetical "If you believe semiadditivty is important..."

Definition 9.9. Let C be a category with products (resp. coproducts). Given an object $c \in \mathsf{C}$, its *diagonal* morphism is the arrow $\Delta_c : c \to c \times c$ (resp. its *codiagonal* is the arrow $\nabla_c : c \coprod c \to c$) induced in the left diagram (resp. right diagram).

$$\begin{tikzcd} & c \arrow[dl, "\mathrm{id}_c"'] \arrow[d, dashed, "\exists!\Delta_c"] \arrow[dr, "\mathrm{id}_c"] & \\ c & c \times c \arrow[l] \arrow[r] & c \end{tikzcd} \qquad \begin{tikzcd} c \arrow[r] \arrow[dr, "\mathrm{id}_c"'] & c \coprod c \arrow[d, dashed, "\exists!\nabla_c"] & c \arrow[l] \arrow[dl, "\mathrm{id}_c"] \\ & c & \end{tikzcd}$$

Proposition 9.10 (Possible reference)**.** *Let C be a category with coproducts (resp. products) regarded as a monoidal category under the coproduct (resp. product) bifunctor. Then each object $\in \mathsf{C}$ is a monoid when equipped with the codiagonal $\nabla_c : c \times c \to c$ and the initial map $\varnothing \to c$ (resp. a comonoid under the diagonal and terminal map). Furthermore, this monoid (resp. comonoid) structure is commutative (resp. cocommutative) and unique.*

Proof. We will work the coproduct / monoid case.

Uniqueness: given $c \in \mathsf{C}$, any unit map for a monoid structure on c must be a morhism $\varnothing \to c$, which is necessarily unique. Similarly, any map $m : c \coprod c \to c$ is induced by two maps $x, y : c \to c$ with id_c as a two-sided inverse, which necessitates $x = y = \mathrm{id}_c$. This forces $m = \nabla_c$.

-ativity: the monoidal product is the coproduct, i.e. we are working with a cocartesian symmetric monoidal category, so the braiding is the identity. This forces commutativity and associativity. □

Remark 9.11. It is easy that any morphism $f : c \to c'$ is also a morphism of monoids. Hence, $\mathsf{C} \cong \mathsf{CMon}(\mathsf{C}) \cong \mathsf{Mon}(\mathsf{C})$.

Remark 9.12. I believe, but do not know how to prove, the converse: if a category C admits finite coproducts, then it is semiadditive if and only if $\mathsf{CMon}(\mathsf{C}) \to \mathsf{C}$ is an equivalence. We will see this in much greater generality later.

Example 9.1. Take $\mathsf{C} = \mathsf{Mon}(\mathsf{Set})$, which has finite products and coproducts. The $\times$ is the cartesian product and the $\coprod$ is the free product. With respect to $\times$, the comonoidal product $\Delta_c : c \to c \times c$ is given by that in Set, namely the diagonal $x \mapsto (x, x)$. (To see this, hit Δ_c with $U : \mathsf{Mon} \to \mathsf{Set}$ and remember that it preserves limits.) With respect to $\coprod$, the monoidal product $\nabla_c : c \coprod c \to c$ is given by multiplication of elements with the monoid structure. (Note that although ∇_c is commutative by abstract nonsense, this does *not* imply that every monoid c is commutative because $c \coprod c$ is *not* $c \times c$.)

Example 9.2. But if $\mathsf{C} = \mathsf{CMon}(\mathsf{Set})$, then $\times$ and $\coprod$ coincide. Thus, each commutative monoid M is a categorical monoid and comonoid under the codiagonal and diagonal maps

$$\nabla_M : M \times M \to M \text{ and } \Delta_M : M \to M \times M.$$

Furthermore, these structures are unique, and are commutative and cocommutative, respectively.

Say actual things about semiadditive categories

9.3 (9/24) Just monoids

I want to think about monoids as functors out of span categories, because that seems entertaining and I haven't done it yet. Here are some references.

- Dan Freed's notes here.
- Schewede here, *Stable homotopical algebra and Γ-spaces*, 1999.
- Akhil's blog here.
- Did not really use this, but I did find stuff I want to get back to in Barwick's *On the algebraic K-theory of higher categories*, found here, in paritcular §3.
- This MO question.

Here's an ancient way to think about monoids. Regard FinSet_* as the category with objects $\{n_+ := \{0, 1, \ldots, n\} : n \geq 0\}$ and with morphisms based functions (those f with $f(0) = 0$).[16] For the purpose of accessing infinite loop spaces, i.e. $\mathbb{E}_\infty$-spaces, i.e. *homotopy coherent* monoids, Segal made systematic use of $\Gamma := \mathsf{FinSet}_*^{\mathrm{op}}$ which we now sometimes call *Segal's Γ category*. We can get a good demonstration of this machine by using it to encode just *ordinary* monoids, so we'll do that. The idea is that given a set X, a monoid structure on X is precisely the occurence of X as the image of 1_+ for a "nice" functor $\Gamma^{\mathrm{op}} \to \mathsf{Set}_*$.

Suppose as given a commutative monoid M. Define a functor $A_M : \Gamma^{\mathrm{op}} \to \mathsf{Set}_*$ as follows. On objects, $A_M(S) := \mathsf{Set}_*(S, M)$. Next, given a morphism of finite pointed sets $f : X \to Y$, we must define a "pushforward" $f_* := A_M(f) = \mathsf{Set}_*(X, M) \to \mathsf{Set}_*(Y, M)$. For this, given pointed $\theta : X \to M$ we define $f_*\theta : Y \to M$

$$f_*\theta(y) = \begin{cases} 0 & \text{if } y = 0, \\ \sum_{x \in f^{-1}(y)} \theta(x) & \text{otherwise.} \end{cases}$$

In words, for each $f : X \to Y$, the morphism $f_* = A_M(f)$ takes based functions $X \to M$ and uses the monoid structure of M to integrate along the fiber of f, resulting in functions $Y \to M$.

Think about this.

Remark 9.13. Is there a way to explain why we make the ad-hoc choice to define $f_*\theta(y) = 0$ if $y = 0$? It is obviously necessary that $0 \mapsto 0$, but it's ugly that in this one case we do *not* integrate over the fiber. I think there is a formulation of Γ-sets without basepoints and with partially defined maps, and we handle partial-definiteness by collapsing stuff to zero. Maybe in that formulation, the definition is more uniform...

Ask someone about this.

[16] So, we are replacing FinSet_* with a skeleton.

Proposition 9.14. *Given a commutative monoid M, the assignment $A_M : \Gamma^{\text{op}} \to \mathsf{Set}_*$ described above is functorial. Furthermore, it satisfies the following properties.*

(1) $A_M(*) = *$.

(2) $A_M(1_+) \cong M$ *canonically.*

(3) The natural maps[17] $X \vee Y \to X \times Y$ *induce isomorphisms* $A_M(X \vee Y) \xrightarrow{\sim} A_M(X) \times A_M(Y)$. *Thus,* A_M *takes coproducts to products.*

Proof. Associativity and commutativity are necessary for A_M to be covariant. Here are two cases to show this. For convenience, take $M = (\mathbb{Z}, +)$.

- Denote by $f : 3_+ \to 3_+$ and $g : 3_+ \to 2_+$ the maps $(0, 1, 1, 2)$ and $(0, 1, 2, 2)$, respectively. That $(1 + 2) + 3 = 1 + (2 + 3)$ exactly expresses $A_M(g \circ f) = A_M(g) \circ A_M(f)$.

- Denote by $t : 2_+ \to 2_+$ the twist map $(0, 2, 1)$. That $1 + 2 = 2 + 1$ exactly expresses $A_M(g \circ f) = A_M(g) \circ A_M(f)$.

This works in general. Probably, you would say that all maps in FinSet_* are composites of this form and work off that. Properties (1) and (2) are immediate. I think that (3) is equivalent to the fact that A_M takes coproducts to products.

□

Consider a functor $A : \Gamma^{\text{op}} \to \mathsf{Set}_*$. We call it a *$\Gamma$-set* if $A(*) = *$. We call it *special* if the natural map $A(S \vee T) \to A(S) \times A(T)$ is an isomorphism. Thus, the previous proposition says that if M is an abelian monoid, then A_M is a special Γ-set. Conversely, suppose that $F : \Gamma^{\text{op}} \to \mathsf{Set}_*$ is a special Γ-set. Consider the pair

$$(F_0 := F(1_+), F(k : 2_+ \to 1_+))$$

where k is the map $(0, 1, 1)$. Because F is special, $F(2_+) \cong F_0 \times F_0$, hence $F(k)$ is a map $F_0^2 \to F_0$. We think of F_0 as an "underlying set" and $F(k)$ as an "operation." Various commutative diagrams in FinSet_* imply that that this operation is associative, commutative, and unital. You use $F(*) = *$ to get the unit methinks. We may wrap this up into a condition which is sometimes called the *Segal condition*, which says that for all n we have $F(n_+) \xrightarrow{\sim} F(1_+)^{\times n}$. One allows $n = 0$, which we interpret as $F(*) \cong *$.

All this converges to the following.

Proposition 9.15. *The category* CMon *is equivalent to the category of special Γ-sets, i.e. the full subcategory of* $\mathrm{Fun}(\Gamma^{\text{op}}, \mathsf{Set}_*)$ *spanned by functors satisfying the Segal condition.*

So we have this categorical perspective on monoids: they are special Γ-sets. Let's get weirder. We will use *spans*. Given a category C, a *span* from x to y is a pair of arrows $x \leftarrow z \to y$. If C has pullbacks, then we define the *composite of spans* as the pullback along the "inner pair" of morphisms. We get a category $\mathsf{Span(C)}$ of spans in C, having the same objects as C but with spans as morphisms.

Now I will define a functor $\mathsf{CMon} \to \mathrm{Fun}(\mathsf{Span(FinSet)}, \mathsf{Set})$.

Definition 9.16. Let M denote an abelian monoid. Define a functor P_M: $\mathsf{Span(FinSet)} \to \mathsf{Set}$ as follows.

- On objects, we define $P_M([n]) := M^{\times n}$.

- On morphisms, say $f : [m] \to [n]$ given by $[m] \xleftarrow{g} Z \xrightarrow{h} [n]$, we define $P_M(f) : M^{\times m} \to M^{\times n}$ to be $(t_1, \ldots, t_n)$ where t_i is given by

$$t_i := \sum_{z_j \in h^{-1}(t_i)} m_{g(z_j)}.$$

Proposition 9.17. *The* $P_M : \mathsf{Span(FinSet)} \to \mathsf{Set}$ *just defined is a functor. Furthermore...*

(1) P_M *preserves products.*

(2) $P_M(\varnothing) = *$.

[17] The set-level definition of this map is clear. It is the one induced by the collapse maps $X \leftarrow X \vee Y \to Y$.

(3) $P(*) = M$.

(4) $P_M(\varnothing = \varnothing \to *)$ *is the morphism* $* \to M$ *which labels zero.*

Proof. Functoriality is probably elementary. Properties (2)-(4) are easy to check. Property (1) I insist is elementary once you figure out what products in $\mathsf{Span}(\mathsf{FinSet})$ are and work out an example. Let me do both:

- Products in $\mathsf{Span}(\mathsf{FinSet})$ are disjoint unions. The morphisms (spans) from a disjoint union to its constituents is the obvious one.

- Because the product of spans of sets is given by disjoint unions, we immediately have $P_M(f \times g) = P_M(f) \times P_M(g)$. write out example

□

Proposition 9.18. *The association* $M \mapsto P_M$ *extends to a fully-faithful functor* $\mathsf{CMon} \hookrightarrow \mathrm{Fun}(\mathsf{Span}(\mathsf{Fin}), \mathsf{Set})$ *whose essential image is the full subcategory* $\mathrm{Fun}^{\times}(\mathsf{Span}(\mathsf{Fin}), \mathsf{Set})$ *spanned by product-preserving functors.*

Proof. Extending $M \mapsto P_M$ to a functor is elementary. That this functor is fully-faithful is easy. In light of the previous proposition, to show that the image is the full subcategory of product-preserving functors, it suffices to show that any (product-preserving) functor F arises as P_M for some M. For this, take $M = F(*)$ and $\mu = F(\{1,2\} \xleftarrow{\mathrm{id}} \{1,2\} \to *)$. The functoriality of F implies that μ is an associative, unital, commutative operation on M, and it is clear that $P_M = F$. □

Corollary 9.19. *We have an equivalence* $\mathsf{CMon} \cong \mathrm{Fun}^{\times}(\mathsf{Span}(\mathsf{FinSet}), \mathsf{Set})$.

Remark 9.20. How to tie this back to Segal's special Γ-sets? A small hurdle is that one formulation uses pointed sets and the other does not. This is a nuisance!!! To start, convince yourself that pointed maps are essentially the same thing as partially-defined maps. More precisely, Set_* is equivalent to the category of sets and partially-defined maps. And we can relate partially-defined maps to spans: a partial map $(S_0 \subseteq S, f : S_0 \to T)$ is precisely a *restricted span*, i.e. a span where the left leg is an inclusion, namely the span $S \hookleftarrow S_0 \to T$. Now we can complete the picture: the inclusion $\mathsf{Span}^{rest}(\mathsf{FinSet}) \subseteq \mathsf{Span}(\mathsf{FinSet})$ induces

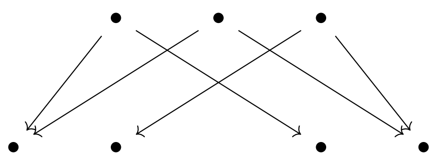
$$\mathrm{Fun}^{\times}(\mathsf{Span}(\mathsf{Fin}), \mathsf{Set}) \to \mathrm{Fun}(\mathsf{Span}^{rest}(\mathsf{Fin}), \mathsf{Set})$$

and we can identify the latter with $\mathrm{Fun}(\Gamma^{\mathrm{op}}, \mathsf{Set})$. The essential image of this restriction is the full subcategory spanned by those functors satisfying Segal's condition. What about the pointed-ness of Set?

Definition 9.21. If C admits finite products, then we define $\mathsf{CMon}(\mathsf{C}) := \mathrm{Fun}^{\times}(\mathsf{Span}(\mathsf{Fin}), \mathsf{C})$.

This general definition lets you do something fun.

Definition 9.22. Suppose that C admits finite products and let $M : \mathsf{Span}(\mathsf{FinSet}) \to \mathsf{C}$ denote a commutative monoid in C. We say that M is *grouplike* if the map $M^2 \to M^2$ given by $(a, b) \mapsto (a, a + b)$ is invertible i.e., the transformation given by the matrix $\left(\begin{smallmatrix} 1 & 1 \\ 0 & 1 \end{smallmatrix}\right)$, i.e., the map induced by the following span.



Example 9.3.

10 October

10.1 (10/8) The back and the shoulder joint

It is now properly Autumn (i.e., both temporally and meteorologically) and my client has back pain. Folks complain that higher categories and homotopy theory are too abstract and inapplicable, so let me include in my notes something seriously real-life: the anatomy and kinesiology I need to help my client.

The shoulder is a *ball-and-socket joint*, shallow in comparison to the hip joints. It is comprised of...

- The *humerus*, *clavicle*, and *scapula* bones, which are encasulated by
- Ligaments and capsule, which passively cushion and fasten the humerus' "ball" (its *head*) onto the scapula's "socket" (its *glenoid*), which are encapsulated by
- The *rotator cuff*, four small muscles that stabilizes / secures the humerus to the socket as it moves,
- The *prime movers*, by which I mean the larger muscles attached to the humerus and shoulder joint that move the arm. This includes your *lats*, *traps*, *deltoids*, and *pecs*.

As you move (or don't move), two important processes are carried out in your shoulder complex.

- Stabilization: things are pulling the humeral head into the center of the scapular socket. We classify these forces as either *static* or *dynamic*. Static forces arise from structure you cannot control, e.g. the pull of ligaments.[18] Dynamic forces are those you cause / control, e.g. the pull of your muscles.

- Scapular orientation: things are moving your scapulae so that your humeral joint is positioned effectively. Many muscles are involved. I'd classify them into two groups, based on their actions:

  - The upper back muscles attached to your scapulae (traps, rhomboids, teres major, serratus anterior) directly move your shoulder blades, and
  - The *erector spinae muscles* stabilize and extend the *thoracic spine*, positioning and orienting of the shoulder blades.

What may cause pain at the shoulder joint? If you visualize the shoulder joint as a ball in a socket, you see that the arm's range of motion is bounded by the scapular socket. *Impingement* is when the socket rubs against (impinges upon) the humeral head, its surrounding ligaments, or the rotator cuff muscles. This is the most common cause of shoulder pain. So, what causes impingement?

[18] In some sense, static forces describe your "passive" or "baseline" stability. Double-jointedness is partly a result of weak static forces on the joints in question.

- Anatomy: some people have more "hooked" sockets (specifically, the *acromion*, which is the "roof" of the socket) or less passive stability.
- Muscular imbalances: a weak or inactive rotator cuff may be overpowered by primary movers, which can move the arm beyond its appropriate ROM or excessively off-center the humeral head.
- Mobility limitations: stiff lats, pecs, or erector spinae / thoracic muscles may restrict or skew movement of the arm and shoulder blade. For example, tight or overdeveloped lats may cause excessive internal rotation, which may cause impingement when moving the arm overhead, since the humerus must externally rotate in order to not impinge upon the acromion.
- Coordination problems: movement is a muscular concert. Problems may occur if even one body part is mistimed or tone deaf. For instance, if we want to push overhead, then we may externally rotate the humerus, rotate the shoulder blade, and extend the thoracic spine to orient our shoulder blade "up." This helps us NOT impinge upon the "roof" of the socket. If we cannot coordinate all these muscle together (e.g., if they are too weak, or one lacks neuromuscular control) then this coordinated movement falls apart.

10.2 (10/9) Semiadditivity for ∞-categories

Ok, maybe now we can get weird. Suppose that C is pointed. We've discussed what it means for C to be semiadditive. One characterization (which I don't think I mentioned) is the following.

Definition 10.1. A pointed category C is semiadditive if and only if it admits finite products and coproducts and for every $A, B \in \mathsf{C}$, the "identity matrix" morphism $A \coprod B \to A \times B$ is an isomorphism.

And here is something nice that happens with semiadditive categories, which I somehow forgot to mention last time.

Proposition 10.2. *If C is semiadditive, then it is canonically enriched in CMon, with units the zero morphisms.*

Proof. The sum of two parallel arrows $f, g : A \to B$ is the composite $A \xrightarrow{\Delta} A \oplus A \xrightarrow{f+g} B \oplus B \xrightarrow{\nabla} B$. □

Remark 10.3. Recall that we *did* discuss the following fact: if C admits coproducts, and we consider C as a monoidal category under $\coprod$, then each object possesses a unique monoid structure given by $\nabla_X : X \coprod X \to X$ and $\varnothing \to X$. Furthermore, it is commutative. You might like to appeal to the general phenomenon that "the collection of maps into a CMon-ish thing form a CMon-ish thing."[19] And indeed, given parallel $f, g : X \to Y$, there is a composite $X \to X \coprod X \xrightarrow{f \coprod g} Y \coprod Y \xrightarrow{\nabla_Y} Y$. But a priori, we have the two "inclusions" $Y \to Y \coprod Y$ but *not* a "diagonal." Depending on which of the two inclusions you choose, this composite is more akin to $(f, g) \mapsto f + 0$ or $(f, g) \mapsto 0 + g$. We cannot "really" form $f + g$. Neither of those former maps function as a monoidal operation. There is manifestly no two-sided unit.

Think about this "category closed under $\mathrm{Hom}(-, X)$ for all X."

The definition of semiadditivity makes perfect sense for ∞-categories. In this case, the previous proposition translates to the following.

Proposition 10.4. *Let C be a semiadditive ∞-category. Then for every $X, Y \in \mathsf{C}$, the set $\pi_0 \mathsf{Map}_{\mathsf{C}}(X, Y)$ has a canonical commutative monoid structure.*

The homotopy theorists erases the π_0 symbol and asks what you get. The answer is $\mathbb{E}_\infty$-spaces.

10.3 (10/12) Higher semiadditivity: a first take

From last time: an ordinary semiadditive category has a canonical CMon-enrichment. This is true also for the π_0 of the mapping space of a semiadditive ∞-category. In fact, for a semiadditive ∞-category, the ordinary phenomenon lifts in the most satisfying way: its mapping spaces are *homotopy monoids*, i.e. $\mathbb{E}_\infty$-spaces.

As part of our ongoing campaign to subjugate chromatic homotopy theory, Hopkins-Lurie take very seriously the following observation:

[19] This is true in many cases: it is true for ordinary abelian monoids, it is true for the category of functors from any category to a semiadditive one, ...

Observation 10.5. A bicomplete ∞-category C is semiadditive if and only if (I) it is pointed and (II) the canonical $X \coprod Y \to X \times Y$ is an isomorphism. Notice that...

- Property (I) $\iff$ the limit and colimit of the empty functor $\varnothing \to \mathsf{C}$ exist (equivalently, C has an initial and terminal object) and there exists a map $* = \varinjlim \to \lim = \varnothing$, which is necessarily unique and an isomorphism if it exists.
- Property (II) is the same condition except we consider every functor $\{\bullet \quad \bullet\} \to \mathsf{C}$, and we ask that the “canonical" map $\phi : \varinjlim \to \lim$ is an isomorphism. Also, see that the definition of ϕ requires Property (I), for recall that it is the “identity matrix," which requires a zero map. By “identity matrix," we mean it is the induced map in the following diagram.

$$\begin{tikzcd}
X \arrow[d, "\mathrm{id}_X"'] \arrow[rrd, "0"] & X \coprod Y \arrow[l] \arrow[r] \arrow[d, dashed, "\phi"] & Y \arrow[d, "\mathrm{id}_Y"] \arrow[lld, "0"'] \\
X & X \times Y \arrow[l] \arrow[r] & Y
\end{tikzcd}
\qquad \approx \qquad
\begin{pmatrix} \mathrm{id}_X & 0 \\ 0 & \mathrm{id}_Y \end{pmatrix}$$

I encourage the reader to squint and see the silhouette of an inductive definition. If you can manage that, you are probably Jacob Lurie or Mike Hopkins. The rest of us have ground to cover. Let me state a informative "toy example" then outline my thought process as I explain the generalization to myself.

Example 10.1. If a finite group G acts on a rational vector space A, then the norm map $N_A : A_G \to A^G$ is an isomorphism. We can rephrase this categorically: our group action is equivalent to a functor $F : BG \to \mathsf{Vect}_{\mathbb{Q}}$ and our norm map is an isomorphism $N_A : \varinjlim F \to \lim F$. We can explicitly describe these things: F is the functor $* \mapsto A$ and $g \mapsto$ (the action of g on A) and N_A is the map $a \mapsto \sum_G ga$.

OK, that example did not accomplish what I wanted it to. (I was hoping N_A would have a nice matrix form, but I don't think it does in general.) Anyway, here's a stream of consciousness: Property (I) asks that the empty functor's limit and colimit are canonically isomorphic. Property (II) asks for a similar duality between every product and coproduct pair, i.e., it asks for binary biproducts, and thus finite biproducts. In which sense, (II) sort of asks for finite families of objects (0-types...) to have a "sum." And in the example, the norm N_A exhibits a duality between a BG-indexed family's limit and colimit, in which sense we could "integrate" maps $BG \to \mathsf{Vect}_{\mathbb{Q}}$. I point out that BG's are just (connected) 1-types, and in our case π_1 was finite since we assumed G was finite.

Definition 10.6. An ∞-category C with small colimits is called *m-semiadditive* if every finite m-type X is C-ambidextrous.

Remark 10.7. Thus, in the above observation, C would be 0-semiadditive. Note that the "definition" of semiadditive used in that observation is equivalent to the definition just given, see HL Prop. 4.4.9.

10.4 (10/21) Higher monoids via of finite spaces

I can continue writing now that I am out of the hell that is applying to the NSF GRFP. As of last time, we saw ordinary semiadditivity, it's direct generalization to ∞-categories which we called 0-semiadditivity, and then m-semiadditivity for all $m \geq -2$.

In the ordinary case, we saw that if C admits coproducts, then semiadditivity was equivalent to $\mathsf{CMon}(\mathsf{C}) \to \mathsf{C}$ being an equivalence. We also saw the nice property that hom-sets in C were canonically commutative monoids.

Now you consider 0-semiadditivity in the ∞-categorical case. You become pensive. Reflective. It is clear that if C is 0-semiadditive, then $\pi_0 \mathsf{Map}(X, Y)$ is canonically an abelian monoid. And I told you that $\mathsf{Map}(X, Y)$ is an $\mathbb{E}_\infty$-space. So, analogous to the ordinary case, is $\mathsf{CMon}(\mathsf{C}) \to \mathsf{C}$ an equivalence? Easy proof of this? Does this characterize the 0-semiadditivity of C? **Then what about m-semiadditivity, what replaces the $\mathbb{E}_\infty$-spaces?**

It turns out that we are seeing the first striation of a general phenomenon occuring across m. This is where Segal's infinite loop space machine comes in handy. Recall that $\mathsf{CMon}(\mathsf{C})$ consists of special

Γ-objects, which are equivalent to product-preserving functors $\mathsf{Span}(\mathsf{FinSet}) \to \mathsf{C}$. The point of Segal's machine was to be able to say that a homotopy commutative monoid, i.e. an $\mathbb{E}_\infty$-space, is the same thing as a product-preserving functor $\mathsf{Span}(\mathsf{FinSet}) \to \mathsf{Spaces}$. Using this perspective, we can ambitiously define an *m-commutative monoid* in C to be a nice functor $\mathsf{Span}(m\text{-finite spaces}) \to \mathsf{C}$. Here, "nice" should be some higher version of the Segal condition. These form a category we write $\mathsf{CMon}_m(\mathsf{C})$, and everything above generalizes: a category C is m-semiadditive $\iff$ the functor $\mathsf{CMon}_m(\mathsf{C}) \to \mathsf{C}$ is an equivalence, and an m-semiadditive category is canonically enriched in $\mathsf{CMon}_m(\mathsf{Spaces})$.

Definition 10.8. We denote by Spaces_n^m the ∞-category whose objects are the n-finite spaces and whose morphisms are spans such that the left leg is m-truncated.

Definition 10.9. Let C be an ∞-category admitting all K_m-indexed limits. A *m-commutative monoid* in C is a functor $F : \mathsf{Spaces}_m^{m-1} \to \mathsf{C}$ such that for every $X \in \mathsf{Spaces}_m^{m-1}$, the family of maps $\{(i_x)^* : F(X) \to F(*)\}_{x\in X}$ exhibits $F(X)$ as the limit of the constant X-indexed diagram to $F(*)$.

Example 10.2. Take $m = -1$. A space is (-1)-finite iff it empty or a singleton, and a map is always (-2)-connected, hence $\mathsf{Spaces}_{-1}^{-2} = \mathsf{Spaces}_{-1}$ is the ∞-category with two objects and a unique non-identity morphism $\varnothing \to *$. If C is an ∞-category admitting K_{-1}-limits (which amounts to admitting the empty limit and hence a final object $*$) then we may ask about (-1)-monoids $M : \mathsf{Spaces}_{-1} \to \mathsf{C}$. Such a thing amounts to a choice of morphism $M(\varnothing) \to M(*)$ with the property that $M(\varnothing) \to M(*)$ exhibits $M(\varnothing)$ as a limit over the constant functor $\mathsf{Spaces}_{-1} \mapsto M(*)$, which precisely says that $M(\varnothing)$ is a final object. Hence, a (-1)-commutative monoid M in C is an arrow $A \to B$ such that A is final. We can fix a final object and identify $\mathsf{CMon}_{-1}(\mathsf{C})$ with $\mathsf{C}_{*/}$, the ∞-category of pointed objects in C.

Example 10.3. Take $m = 0$. A space is 0-finite iff it is finite and discrete, and a map is (-1)-finite iff its fibers are all empty or contractible, hence iff it is an injection. Thus $\mathsf{Spaces}_0^{-1} = N\mathsf{Span}^{rest}(\mathsf{FinSet})$, the category of spans of finite sets with one leg an injection. A category C admits K_0-limits iff it admits finite products. A 0-monoid $M : \mathsf{Spaces}_0^{-1} \to \mathsf{C}$ is characterized by the property that for each finite set n, $M(n)$ is the limit of the constant functor $n \mapsto M(*)$, which I think of as saying $M(n) \cong M(*)^n$. We can identify $\mathsf{Span}^{rest}(\mathsf{FinSet})$ with Γ^{op}, and this property says that the resulting functor $M : \Gamma^{\mathrm{op}} \to \mathsf{C}$ satisfies Segal's condition. Leaving some details unchecked, we find that M is precisely an $\mathbb{E}_\infty$-monoid in C.

Example 10.4. I am still not sure how to say concrete things about higher monoids.

On the to-do list:

1. Figure out how to say something about a higher monoid
2. Figure out the basic, technical properties of Spaces_m^n that let us make basic proofs
3. Figure out what higher monoids are actually good for; modes?
4. Move on and think about semiadditivity and height; maybe higher monoids will reappear?

10.5 (10/24) How to feel a higher monoid

From the description of a monoid M as a product-preserving functor out of $\mathsf{Span}(\mathsf{FinSet})$, I can extract a "concrete description" of M: the zero object $* \to M$, the product operation $M^2 \to M$, and the various properties they satisfy. We want a similar concrete description of higher monoids. In what follows, let $F : \mathsf{Spaces}_m^{m-1} \to \mathsf{Spaces}$ denote an m-commutative monoid in Spaces, and write $M := F(*)$.

For this purpose, let me make some terminology. Suppose as given a morphism $h \in \mathsf{Spaces}_m^{m-1}$, thus a span $X \leftarrow Z \to Y$ of π-finite and m-truncated spaces whose left leg is $(m-1)$-truncated. Say h is a *lid* (resp. a *rid*) if its left (resp. right) leg is the identity.

First we think about rids. An $(m-1)$-truncated map $f : Y \to X$ is the same thing as a rid $\hat{f} : X \to Y$.

- For F to be an m-commutative monoid, it must be that given a space X, the family $\{\hat{i}_x : F(X) \to M\}_{x\in X}$ exhibits $F(X)$ as the limit of the constant functor $X \mapsto M$. Thus, $F(X) \cong \mathsf{Map}_{\mathsf{Spaces}}(X, M)$. Why?
- Do rids see more about $F(X)$? Suppose as given any old $(m-1)$-truncated map $f : X \to Y$. We get a map $F(\hat{f}) : \mathsf{Map}(Y, M) \to \mathsf{Map}(X, M)$. But note that in Spaces_m^{m-1}, we have $\hat{i}_x \circ \hat{f} = \hat{i}_{f(x)}$. Hitting this with F, we find that $F(\hat{f})$ is just restriction along f. Nothing to see here.

Next, we think about lids. Lids are simpler: a lid $g : X \to Y$ is just an m-truncated map $g : X \to Y$ of π-finite spaces. (Hence simpler, in the sense that there is no contravariance nor $(m-1)$-truncatedness.) Since F is an m-commutative monoid, we can identify $F(X) \cong \mathsf{Map}(X, M)$ and get a map $g_* := F(g) : \mathsf{Map}(X, M) \to \mathsf{Map}(Y, M)$. Now consider the following commutative diagram in Spaces_m^{m-1}.

$$\begin{tikzcd} X \arrow[r, "\hat{i}_{X_y}"'] \arrow[d, "g"] & X_y \arrow[d, "p_y"'] \\ Y \arrow[r, "\hat{i}_y"] & y \end{tikzcd}$$

Here, X_y is the homotopy fiber of g above y. Given $\phi \in \mathsf{Map}(X, M)$, this square tells us that

$$g_*\phi \in \mathsf{Map}(Y, M) \quad \text{is the map} \quad y \mapsto (p_y)_*(\phi|_{X_y}) \in M.$$

We should think of the restriction $\phi|_{X_y}$ as not varying too remarkably across y. (Maybe this is not correct. But we have some control over it, at least in the case of fibrations. Also, it is independent of F.) On the other hand, the action by projections (e.g., $(p_y)_*$) are totally native to the monoid F. Hence, the actions p_* by projection maps $p : X \to *$ constitute structure of F. Such an action is a map $p_* : \mathsf{Map}(X, M) \to M$, and we have an action for each m-truncated space X. You may say that the structure of actions by projections p_* is an "integration procedure" which associates to each m-truncated X-family of points $\phi : X \to M$ a new point $p_*(\phi) \in M$. Furthermore, consider a fibration $X \to Y$ and a map $f : X \to M$. We may form $\int_X f \in M$. But since $X \to Y$ is a fibration, we get a map $\mathsf{Map}(*, Y) \to \mathsf{Map}(X_y, X)$, and we can postcompose maps in the latter with f and integrate. Altogether this describes a point $\int_y \int_{X_y} f$. I think that the above commutative square relates these two integrals, which we christen a "Fubini-type relation." This relation is a path in M?

Do we?

10.6 (10/26) Categorical properties of spans of finite spaces

The property of m-finiteness is closed under pullbacks, so Spaces_m has pullbacks and these are computed in Spaces. In fact, the following is true.

Proposition 10.10 (Lemma 2.9). *Spaces_n admits K_n-colimits and those are preserved and detected by $\mathsf{Spaces}_n \hookrightarrow \mathsf{Spaces}$.*

Proof. In general, fully-faithful functors reflect limits and colimits, hence $\mathsf{Spaces}_n \hookrightarrow \mathsf{Spaces}$ does. That is, if a cone in Spaces_n is a (co)limit in Spaces, then it is a (co)limit in Spaces_n.

Now, we can show that Spaces_n admits K_n-colimits and $\mathsf{Spaces}_n \hookrightarrow \mathsf{Spaces}$ preserves them by proving that n-finiteness is closed under K_n-colimits in Spaces. (Reflectiveness is necessary here.) That is, it suffices to prove that if X is n-finite and $\phi : X \to \mathsf{Spaces}$ is such that $\phi(x)$ is always n-finite, then $\varinjlim(\phi)$ is n-finite. For this, consider that ϕ classifies some left Kan fibration $p : E_\phi \to X$. Since X is a Kan complex, so is E_ϕ, and p is a Kan fibration. The essential detail is that E_ϕ presents the colimit of ϕ, and by the long exact sequence of a Kan fibration, that X is n-finite $\implies E_\phi$ is n-finite. $\square$

Ask Charles about this.

We also need to know that the fibers of p are n-finite. Why is this the case?

Proposition 10.11 (Lemma 2.11). *Suppose that D admits K_n-colimits and let $F : \mathsf{Spaces}_n \to \mathsf{D}$ be any functor. Then F preserves K_n-colimits $\iff$ for every $X \in \mathsf{Spaces}_n$, the family $\{F(*) \to F(X)\}_{x\in X}$ exhibits $F(X)$ as the colimit of the constant $F(*)$-valued functor $\varinjlim(X \to \mathsf{D})$.*

The previous lemma says that $F : \mathsf{Spaces}_n \to \mathsf{D}$ preserves K_n-colimits $\iff$ it preserves the constant K_n-colimits with value $* \in \mathsf{Spaces}_n$. I would like to review the proof of this eventually. In particular, I would like to think about why this is true in the $n = 0$ case. Maybe it's intuitive in the ordinary case. In any case, this is a useful lemma.

Proposition 10.12. *The inclusion $\mathsf{Spaces}_n \hookrightarrow \mathsf{Spaces}_n^m$ preserves K_n-colimits.*

Corollary 10.13. *A functor $F : \mathsf{Spaces}_n^m \to \mathsf{D}$ preserves K_n-colimits $\iff$ its restriction $F : \mathsf{Spaces}_n \to \mathsf{D}$ preserves K_n-colimits.*

Proof. The $\implies$ direction is obvious. For the $\impliedby$ direction, suppose that the restriction preserves K_n-colimits. For some reason, for every Kan complex X, every diagram $X \to \mathsf{Spaces}_n^m$ arises from a diagram $X \to \mathsf{Spaces}_n$, whatever that means, and F preserves K_n-colimits of the latter kind. $\square$

Now we think about monoidal structure on Spaces_n^m. The Cartesian monoidal structure on Spaces_n induces a symmetric monoidal structure on Spaces_n^m. On objects, it is given by the Cartesian product of spaces. On morphisms, it is given by the Cartesian product of spans. This is not Cartesian (is there a Cartesian monoidal product on these spans?)

Proposition 10.14 (Cor. 2.17). *The symmetric monoidal product on* Spaces_n^m *preserves* K_n*-colimits in each variable seperately.*

Proof. Consider the following commutative diagram.

$$\begin{tikzcd}
S \times S \arrow[d, "\times"] & S_n \times S_n \arrow[l] \arrow[r] \arrow[d, "\times"] & S_n^m \times S_n^m \arrow[d, "\times"] \\
S & S_n \arrow[l] \arrow[r] & S_n^m
\end{tikzcd}$$

Where $\times$ denotes the Cartesian product on Spaces and the induced symmetric monoidal product on Spaces_n^m. The inclusion $\mathsf{Spaces}_n \hookrightarrow \mathsf{Spaces}$ preserves K_n-colimits, as does the Cartesian product of spaces in each variable; hence the left-down composite does too. Since $\mathsf{Spaces}_n \hookrightarrow \mathsf{Spaces}$ preserves K_n-colimits, we conclude that the Cartesian product on Spaces_n preserves K_n-colimits in each variable seperately. We just obtained that $\mathsf{Spaces}_n \hookrightarrow \mathsf{Spaces}_n^m$ preserves K_n-colimits, and furthermore any $X \to \mathsf{Spaces}_n^m$ factors through $\mathsf{Spaces}_n \hookrightarrow \mathsf{Spaces}_n^m$, so the colimit of $X \to S_n^m$ is computed in S_n. This implies that $\times : S_n^m \times S_n^m \to S_n^m$ preserves K_n-colimits. □

The ∞-category of small ∞-categories admitting K_n-colimits with morphisms the K_n-colimit preserving functors gets a symmetric monoidal product $\otimes_{K_n}$. As a tensor product, given $C, D \in \mathsf{Cat}_{K_n}$, there is a canonical map $C \times D \to C \otimes_{K_n} D$ with the universal property that restriction

$$\mathrm{Fun}_{K_n}(C \otimes_{K_n} D, E) \to \mathrm{Fun}(C \times D, E)$$

is fully-faithful with image the full subcategory spanned by functors $C \times D \to E$ preserving K_n-colimits in each variable seperately. We just saw that the Spaces_m^n has a symmetric monoidal product preserving K_n-colimits in each variable. Hence, the commutative algebras in Cat_{K_n} are symmetric monoidal ∞-categories admitting K_n-colimits and whose products preserve them in each variable separately. In particular, Spaces_n^m is a commutative algebra in Cat_{K_m}.

10.7 (10/28) Baby norms

The preprint for BHLS constructing counterexamples to the telescope conjecture came out Friday, and there was a conference about it Saturday (yesterday). I want to understand what they did, and (not coincidentally) ambidexterity and semiadditivity fit into the story, I think. Right now I think I'll try to understand *cyclotomic spectra*, starting with some discussion of norm maps and the Tate construction. Some of this will be a repeat of stuff from earlier. References include

- HA §6.1.6.
- Nikolaus-Scholze, lectures on TCH.

Definition 10.15. Let C be an ∞-category and consider a map $f : X \to Y$ of Kan complexes. Restriction defines a *pullback functor* f^*: $\mathrm{Hom}(Y, \mathsf{C}) \to \mathrm{Hom}(X, \mathsf{C})$. We define the *left and right Kan extensions* as the left and right adjoints $f_! \dashv f^* \dashv f_*$, respectively. (If they exist, these adjoints *are* Kan extensions by Corollary 030B.)

Example 10.5. Let C be an ∞-category and G a group. An *object with G-action* is a functor $F : BG \to \mathsf{C}$.[20] If C admits all BG-colimits, we define the *homotopy orbits* functor $(-)_{hG}$ by $F \mapsto \varinjlim_{BG} F$. If C admits all BG-limits, we define the *homotopy fixed points* functor $(-)^{hG}$ by $F \mapsto \lim_{BG} F$. The functors $(-)_{hG}$ and $(-)^{hG}$ occur as $f_!$ and f^* when f is the projection $BG \to *$.

[20] I want to also call these *G-objects*.

Construction 10.16 (The norm induced by f)**.** Let $f : X \to Y$ be a map of Kan complexes. The *diagonal map* $\delta : X \to X \times_Y X$ is the map induced by the identities id_X. Suppose we are given an equivalence $Nm_\delta : \delta_! \to \delta_*$ of the left and right adjoints to δ^*. We then may form the composite

$$p_0^* \xrightarrow{\text{unit of } \delta^* \dashv \delta_*} \delta_* \delta^* p_0^* \cong \delta_* \overset{\mathrm{Nm}_\delta^{-1}}{\cong} \delta_! \cong \delta_! \delta^* p_1^* \xrightarrow{\text{counit of } \delta_! \dashv \delta^*} p_1^*.$$

Where p_0, p_1 are the projections $X \times_Y X \rightrightarrows X$. By an adjunction, this is equivalent to a map $\mathrm{id}_{\mathsf{C}^X} \to p_{0*} p_1^*$. We may consider the composite

$$f^* f_* \xrightarrow{\text{counit of } f^* \dashv f_*} \mathrm{id}_{\mathsf{C}^X} \to p_{0*} p_1^*.$$

By abstract nonsense [HA, 6.1.6.3], this is an equivalence $f^* f_* \cong p_{0*} p_1^*$. Hence, we have a map $\mathrm{id}_{\mathsf{C}^X} \to f^* f_*$, and by an adjunction we get a map $f_! \to f_*$ in $\mathrm{Fun}(\mathsf{C}^X, \mathsf{C})$. We call this the *norm map induced by f*.

Here are some questions.

(I) What conditions on f and C beget the existence and equivalence of $\delta_!$ and δ_*? Furthermore, what begets the existence of $f_!$ and f_*, hence Nm_f?

(II) What conditions on f and C imply the norm induced by f is an equivalence?

Proposition 10.17 (Answer to I)**.** *The map f is (-1)-truncated $\iff$ δ is (-2)-truncated, i.e. is an equivalence, in which case $\delta_!$ and δ_* exist and are homotopy inverses to δ^*. This begets a canonical equivalence $Nm_\delta : \delta_! \cong \delta_*$. If in addition C has an initial and final object, then $f_!$ and f_* exist, hence Nm_f exists.*

Proposition 10.18 (Answer to II)**.** *If f is (-1)-truncated and C is pointed*[21]*, then $Nm_f : f_! \to f_*$ is an equivalence.*

Remark 10.19. The previous proposition actually arises from a close relationship between the pointedness of C and the fact that all its (-1)-truncated norms are equivalences. To be precise, suppose that C has an initial and terminal object. Then there exists a map $* \to \varnothing$ (i.e., C is pointed) $\iff$ the induced norm Nm_f is an equivalence for all (-1)-truncated f [HA, 6.1.6.7].

Now suppose that C is pointed. What if f is 0-truncated? In this case δ is (-1)-truncated, hence the previous propositions say that $\mathrm{Nm}_\delta : \delta_! \to \delta_*$ exists and is an equivalence. Now, in a fashion analogous to the above propositions, we wonder what is necessary for $f_!, f_*$ to exist, and for Nm_f to be an equivalence. When we assumed f was (-1)-truncated, we got $f_!$ and f_* by assuming that C had initial/terminal objects, and I think this arises from some construction taking the (co)limit over fibers, which by assumption were empty or contractible, hence the assumption. We can do something similar here, where this time fibers are not (-1)-truncated but 0-truncated, and so we should ask that C admits (co)products. But maybe this is a strong condition, so let's instead ask that f is 0-finite, so we only need C to admit *finite* (co)products. Then if our analogy holds water, $f_!$ and f_* should exist as soon as C admits finite (co)products, and the coincidence of (co)products should imply that Nm_f is an equivalence.

Figure this out.

Proposition 10.20 (HA, 6.1.6.12)**.** *Suppose that C is pointed and admits finite (co)products. If f is a 0-finite map, then f^* admits left and right adjoints $f_!, f_*$. Furthermore, TFAE.*

1. *For every 0-finite map $f : X \to Y$, the induced norm $Nm_f : f_! \to f_*$ is an equivalence.*
2. *C is 0-semiadditive. In other words, finite products and coproducts coincide in a canonical manner: for every finite set S and functor $F : S \to \mathsf{C}$, the map $\mathrm{colim}_S F \to \lim_S F$ determined by the "identity matrix"*[22] *is an equivalence.*

10.8 (10/31) Kan extensions

Last time, we got a headache thinking about surmounting hypotheticals. Specifically, given a map $f : X \to Y$ and a category C, we...

(I) May form the diagonal $\delta : X \to X \times_Y X$ and...

[21] I.e., (-1)-semiaddtive.

[22] That would be the collection of maps $\{\phi_{ij} : F(s_i) \to F(s_j)\}_{S \times S}$ given by $\phi_{ij} = \delta_{ij}$.

(II) Ask when $\delta^* : \mathsf{C}^{X\times_Y X} \to \mathsf{C}^X$ admits left and right adjoints $\delta_!, \delta_*$. When they exist, we may ask when there exists an equivalence $\mathrm{Nm}_\delta : \delta_! \cong \delta_*$. Then we...

(III) Ask when $f^* : \mathsf{C}^{X\times_Y X} \to \mathsf{C}^X$ admits left and right adjoints $f_!, f_*$. When they exist and the items in (II) exist, we may construct a comparison $\mathrm{Nm}_f : f_! \to f_*$ and ask when it is an equivalence.

Given that we are asking about adjoints to pullbacks, we expect Kan extensions to feature. Last time, I tried to be slightly more precise about this, and said something like "if f^* has left/right adjoints, then they are Kan extensions and may be computed by taking co/limits over the fibers of f." This suggests conditions on f and C necessary for the adjoints $\delta_!, \delta_*, f_!, f_*$ to exist and coincide: we need C to *have* those (co)limits in the first place, i.e. those indexed by the fibers of f. This gives existence, somehow by taking fiberwise (co)limits, which gives left/right Kan extensions. Then coincidence of these (co)limits implies coincidence of these left/right Kan extensions.

Asking for existence and coincidence of (co)limits over all fibers is a big ask. So, we first considered (-1)-truncated f. In this case, δ is (-2)-truncated (an equivalence) so that $\delta_! \cong \delta_*$ tautologically. Hence, we are only left to worry about the existence and equivalence of $f_!, f_*$. Our above spiel says that this concerns (co)limits over the fibers of f, which are $\varnothing$ or $*$. The limit/colimit of $\varnothing \to \mathsf{C}$ is an initial/terminal object, and the (co)limit of any $c : * \to \mathsf{C}$ is just c. Hence, all (co)limits indexed by fibers of f exist $\iff$ C has an initial and terminal object, in which case $f_!, f_*$ exist. Our above spiel says that in order for $\mathrm{Nm}_f : f_! \to f_*$ to be an equivalence, we need these (co)limits to coincide, thus C is pointed $\Longrightarrow \mathrm{Nm}_f$ is an equivalence.

Next, suppose as given a 0-*finite* map f. We must again check (I) and then (II). But note that the truncation order of δ is generally one less than that of f, so we can do something inductive: if we assume C is pointed, then the above $\Longrightarrow \delta_!, \delta_*$ exist and Nm_δ is an equivalence, so (I) is answered and we can ask (II). By the orange reasoning, $f_!, f_*$ should exist if C admits finite (co)products, since those are (co)limits over diagrams indexed by the spaces which can arise as fibers of f. Furthermore, Nm_f should be an equivalence once finite (co)products coincide in C. (This is why we asked f to be 0-finite rather than just truncated: if it were just truncated, we would want *countable* (co)products to coincide, but that is quite strong.) We saw all this was true.

OK, above was a rehash of last time. The explanation I tried to give is powered by the orange statement. I would like to flesh out that statement.

11 November

11.1 (11/2) Algebraic theories and their models

A central problem in homotopy theory is what I will call the *information crisis*: properly homotopy-invariant structures consist of too much data to handle "traditionally." Given a ring R, we naturally think of R as a set with distinguished elements $0, 1$ and compatible operations $+, \times$. But given an $\mathbb{E}_\infty$-ring R, you can surely write down the diagrams for the distinguished elements and operations, and you may understand how these pieces *should* fit together, but it takes an infinite amount of extra data (higher homotopies) to actually fit them together. For this reason, it is hard to work with R as we would an ordinary ring.

We want to express these structures in a more convenient[23] form. One helful idea is that we can think of an "algebraic structure" as some "instantiation" of fixed relations and rules, and we should try to rigorously (I) "isolate" those relations and rules and (II) concisely express and study how they are "instantiated." This idea is not precise and can be realized in different ways. *Operads* are one gadget which accomplishes this. I want to think about another (more general and concise) approach using *Lawvere theories*, which are also called *algebraic theories* I think. Before trying to go homotopy theory, we can think about algebraic theories in the context of ordinary algebra.

Lawvere theories are old, just a bit older than monads (Lawvere's thesis is dated 1963). They might be the first example of the idea that we can model things with structure and properties as structure-preserving functors $F : \mathsf{C} \to \mathsf{D}$. This philosophy is ubiquitous in category theory nowadays.

Definition 11.1. A category L is called a *Lawvere theory* if L admits finite products and there exists an object $x \in L$ such that every other object $y \in L$ is isomorphic to some cartesian power $y \cong x^n$.

Definition 11.2. Let L denote a Lawvere theory and C a category with finite products. A *model for L*, also called a *L-algebra*, is a product-preserving functor $L \to \mathsf{C}$.

Example 11.1 (Lawvere theory of groups)**.** Define $L_{grp} := \mathsf{fgFreeGp}^{\mathrm{op}}$. Each object is isomorphic to some $F(n)$. Furthermore, L_{grp} has coproducts since $\mathsf{fgFreeGp}$ has products. (In fact, we know that $F(m) \oplus F(n) \cong F(m+n)$.) Thus, L_{grp} is a Lawvere theory.

Groups are L_{grp}-algebras. To see this, fix a group G. Consider the functor $L_{grp} \hookrightarrow \mathsf{Gp}^{\mathrm{op}} \to \mathsf{Set}$ given by $F(n) \mapsto U\mathrm{Hom}_{\mathsf{Gp}}(F(n), G)$. This functor is product-preserving, since the inclusions of the full subcategory are, as well as the contravariant Hom. SO this defines a model for L_{grp}.

L_{grp}-algebras are groups. To see this, consider a product-preserving functor $T : L_{grp} \to \mathsf{Set}$. Here's how this works.[24]

- We get a set $G := T(F(1))$.
- Next, we want an operation $m : G^2 \to G$. Since T preserves products, this amounts to a morphism $F(1) \to F(2)$. For this, choose the morphism $x \mapsto ab$, where x and a, b are the generators of $F(1)$ and $F(2)$.
- The zero element $* \to G$ arises from the unique morphism $F(1) \to F(0) = *$. The inversion $G \to G$ arises from $x \mapsto -x^{-1}$.

I think the associativity, identity, and inverse equations hold because they hold for the $F(n)$'s. You can check that L_{grp}-algebra morphisms are precisely group homomorphisms, hence all the above describes an equivalence

$$\mathsf{Alg}_{L_{grp}} \cong \mathsf{Gp}.$$

Remark 11.3. It happens generally that if L is a Lawvere theory, then L is equivalent to the full subcategory spanned by the free L-algebras on finitely many generators. (Is this true precisely as stated? I think so. It's true for ∞-categories.)

Example 11.2. We can also ask for a Lawvere theory L_{set} whose models are sets. Sets have no operations. All a set S has is elements, which amount to a maps $* \to S$, one for each element. Writing $[n] := \{0, \dots, n\}$, you may feel that L_{set} should be the category generated by the projections $k : [n] \to [1]$ for each n and

[23] More concise, more conceptual, more calculationally friendly, etc.

[24] Andrej Bauer gives another, more elementary explanation of all this here.

$0 \leq k \leq n$. (These "pick out" the element k.) That was a good feeling, and a valid one, for indeed those morphisms generate $\mathsf{FinSet}^{\mathrm{op}}$ and we have

$$L_{set} \cong \mathsf{FinSet}^{\mathrm{op}}.$$

A *multisorted theory* is an annoying name for a natural object. Namely, it is like an algebraic theory, except for algebraic gadgets with ≥ 1 underlying sets.

Definition 11.4. A *multisorted algebraic theory* is a category C with finite products and a chosen set $S \subseteq Ob(\mathsf{C})$ such that every object is isomorphic to a finite product of objects in S.

Something something something about finitary vs. infinitary algebraic theories. I want to make this stuff homotopical. For this purpose, do we care about finitary vs. infinitary, since homotopical structures have a lot more pieces of data (infinitely more)? Think about this.

11.2 (11/4) Recap, why algebraic theories, rigs, spans?

Let me outline my thoughts the last few days and try to recall why I was thinking about algebraic theories, since I sort of forgot.

(1) Last month I started thinking about the relationship between semiadditivity and monoids, and whether we can make some analogy

$$\text{commutative monoids : semiadditivity :: commutative rigs : ???}$$

(2) A *rig* is a ring without inverses. To make sense of the above analogy, I wanted a good "perspective" of rigs. In the ordinary case, we may define $\mathsf{CRig} := \mathsf{CMon}(\mathsf{CMon}(\mathsf{Set}))$, keeping in mind that the first and second CMon are taken with respect to $\times$ and $\otimes_{\mathbb{Z}}$, respectively.[25]

(3) This suggests a detail toward answering (1). If C is pointed and admits small products, then it is semiadditive iff $\mathsf{CMon}(\mathsf{C}) \to \mathsf{C}$ is an equivalence.[26] In analogy, perhaps we should ask: given C with finite products and zero, when is $\mathsf{CRig}(\mathsf{C}) \to \mathsf{C}$ an equivalence?

(4) The above (3) generates some thoughts. *Should* we assume C is pointed with finite products? I think the assumption of finite products is necessary since it lets us form $\mathsf{CMon}(\mathsf{C})$. But pointedness I am not sure about. For CMon's, pointedness yields the additive identity, and biproducts yields the sum. But for CRig's, we want $+$ *and* $\times$, and separate units for them. So perhaps C should admit finite *products and coproducts*, encoding the sum and multiplication, as well as initial and terminal objects, encoding the additive and multiplicative identities. Does this work out?

(5) In any case, the stuff above is all some sort of abstract, meta investigation of algebraic structures (be it *structures on objects in a category*, such as commutative monoid objects, or *structures on categories*, for instance semiadditivity). I figure it may help to have a language for studying this stuff. That is why I am thinking about *Lawvere theories*. (I've also just heard of them before and wanted to know what they were.) Maybe modes... An example question I have is, how can we compare the monoidal structures on a given category? In our case, if we assume C admits finite products, we get a diagram

$$\begin{tikzcd}
\mathsf{CMon}(\mathsf{D},\times) \arrow[r, "\cong"] & \mathsf{D} := \mathsf{CMon}(\mathsf{C},\times) \arrow[r, "U"] & \mathsf{C} \\
\mathsf{CMon}(\mathsf{D},\otimes) \arrow[u] \arrow[ur, "U"] & &
\end{tikzcd}$$

---

[25]Note that $\mathsf{CMon}(\mathsf{Set})$ admits (co)products. Thus, we may also have formed $\mathsf{CMon}(\mathsf{CMon})$ with respect to the cartesian or cocartesian structures. But we know two things: (I) the finite products and coproducts coincide in CMon, and (II) in any cocartesian monoidal $(C, \cdot)$, each object has a unique monoid structure, and it is, in fact, commutative, and hence the forgetful $\mathsf{CMon}(C) \to C$ is an equivalence. Then in our case, if we take the (co)cartesian structure on CMon, then $\mathsf{CMon}(\mathsf{CMon})$ is not interesting: it is just CMon!

[26]In case you were wondering: once we assume C has products and zero, semiadditivity amounts to the canonical maps $X \to X \times Y \leftarrow Y$ exhibiting $X \times Y$ as a coproduct.

With the aim of studying this diagram, maybe we should treat $\mathsf{CMon}(-,-)$ as a functor? Or at least, we should systematically study the monoidal structures on a category this way. But this requires a formalism for doing so.

(6) And let's not forget about spans. We may, first of all, identify $\mathsf{CMon}(\mathsf{C}) \cong \mathrm{Fun}^{\times}(\mathsf{Span}(\mathsf{Fin}), \mathsf{C})$. Since CRig is formed as CMon's under the *tensor* product, we cannot iterate this "monoids are functors out of spans" construction, but I think there *is* something to be said. Not sure what yet. Is there a Span-ish description of commutative monoids with respect to tensor products? Can this, for instance, help us understand the above diagram?

(7) I've also had a lingering question which may explain or be explained by the above. That is, what's the deal with the Span construction? Why do spans *really* show up? I can (and did earlier in these notes) explain how to pivot from special Γ-objects to $\mathrm{Fun}^{\times}(\mathsf{Span}(\mathsf{Fin}), -)$, but I hope for a more natural appearance of spans. I have something in mind, but I can't write it out yet. An example question is, given a Lawvere theory T, what does the Lawvere theory $\mathsf{Span}(T^{\mathrm{op}})$ model? (That is a Lawvere theory, right? And the op should be there, right?)[27]

Some references I have not read yet are Gepner-Groth-Nikolaus, Haugseng, Elmanto-Haugseng, Cranch, Berman, ... Also, here's a quote by Lawvere in (an anniversary commentary on?) his thesis, which can be found here:

> *Algebras (and other structures, models, etc.) are actually functors to a background category from a category which abstractly concentrates the essence of a certain general concept of algebra, and indeed homomorphisms are nothing but natural transformations between such functors. Categories of algebras are very special, and explicit axiomatic characterizations of them can be found, thus providing a general guide to the special features of construction in algebra... The tools implicit [here] constitute a "universal algebra" which should not only be polished for its own sake but more importantly should be applied both to constructing more pedagogically effective unifications of ongoing developments of classical algebra, and to guiding of future mathematical research.*

11.3 (11/6) Spans

Last time, I did some review and asked the question "if we hit Lawvere theories with spans, what happens?" I was motivated to ask this question by the example of FinSet: recall that $\mathsf{FinSet}^{\mathrm{op}}$ classifies sets and $\mathsf{Span}(\mathsf{FinSet})$ classifies commutative monoids. There are many other examples. I want to think more about this, but first I want to think a bit harder about spans, and get some basic details sorted out.

An essential detail is that spans are naturally 2-categorical. I knew this before, but honestly I do not think about 2-categories, and I thought I could ignore this detail about spans. This was wrong, for reasons I will think about and explain later.

Definition 11.5. A *strict 2-category* is a category enriched in Cat. A *weak 2-category* is a category "weakly enriched" in Cat, in the sense that it has...

- Objects,
- Each pair of objects has a category of morphisms (whose 0-cells are *morphisms* and whose 1-cells are *2-morphisms*),
- For each pair of objects, *unitor* natural isomorphisms, and
- For each quadruple of objects, an *associator* natural isomorphism.

The unitors must satisfy unitality and the associators must satisfy the pentagon axiom. (Full def here.)

Proposition 11.6. *By some strictification procedure, weak 2-categories are equivalent to strict 2-categories.*

Definition 11.7. Let C denote a strict 2-category enriched in groupoids. The *nerve* of C is the ∞-category $N^{coh}(\mathsf{C}')$, whose objects are those of C and whose hom-spaces are defined by $\mathrm{Hom}_{\mathsf{C}'}(x,y) := N\mathrm{Hom}_{\mathsf{C}}(x,y)$.

[27]Think about this.

Proposition 11.8. *Given a strict 2-category* C *enriched in groupoids, its nerve* $N(\mathsf{C})$ *is a* $(2,1)$*-category in quasicategories, in the sense that all inner horns* $\Lambda^n_k \to N(\mathsf{C})$ *has a unique filler as soon as* $k \geq 3$ *(c.f. HTT 2.3.4.9).*

HTT 2.3.4 gives general information about (n-categories $\cap$ quasicategories). So does Cranch.

OK, that was what I wanted to know about 2-categories. We now consider *spans* in this context.

Definition 11.9 (Spans and their composites)**.** Fix a category C. A *span* in C is a diagram $X \leftarrow U \to Y$. Given fixed X, Y, a *morphism of spans* between two spans $(X \leftarrow U \to Y) \to (X \leftarrow U' \to Y)$ is an arrow $U \to U'$ such that the evident diagram commutes. Given fixed spans $s := (X \leftarrow U \xrightarrow{f} Y)$ and $t := (Y \xleftarrow{g} U' \to Z)$, we call a span $X \leftarrow U'' \to Z$ a *composite* of t with s if it is isomorphic to the outer span in the below diagram.

$$\begin{tikzcd}
 & & V \arrow[dl] \arrow[dr] \arrow[dd, phantom, "\lrcorner", very near start, rotate=-45] & & \\
 & U \arrow[dl] \arrow[dr, "f"] & & U' \arrow[dl, "g"'] \arrow[dr] & \\
X & & Y & & Z
\end{tikzcd}$$

If we have a functorial choice of pullbacks, we can define *the composite* $t \circ s$ as the outer span in that diagram. We will do this.

Definition 11.10 (*1-category of spans*)**.** Let C denote a category with pullbacks. We define $\mathsf{1Span(C)}$ to have the same objects as C, and for morphisms $X \to Y$ to be *isomorphisms classes of* spans $X \leftarrow U \to Y$. Composition is given by pullback, which is well-defined because we have taken *isomorphism classes* of spans for morphisms.

Example 11.3. The category $\mathsf{1Span(FinSet)}$ is the Lawvere theory for commutative monoids in Set. That is, product-preserving functors $\mathsf{1Span(FinSet)} \to$ Set are precisely commutative monoids. We may replace Set with Top or sSet and ask what happens. Badzioch proved a "strictification" result: $\mathsf{1Span(FinSet)}$-models in Top are *equivalent* to "up to weak equivalence" $\mathsf{1Span(FinSet)}$-models. It is well-known (c.f. MSE) that the former (the topological abelian monoids) have the homotopy type of generalized Eilenberg-Maclane spaces. For the purposes of studying algebraic theories in homotopy theory, this means $\mathsf{1Span(FinSet)}$ is not good enough, since an algebraic theory for commutative monoids should give us $\mathbb{E}_\infty$-spaces.

Remark 11.11. Wait, but $\mathbb{E}_\infty$-spaces *are* modeled as the special Γ-spaces, i.e., as functors $\mathsf{FinSet}_* \to \mathsf{Top}$ weakly satisfying the Segal condition. I thought these could be identified with product-preserving functors out of $\mathsf{1Span(FinSet)}$? At least, I thought about this briefly when the target is Set, and Tomer spoke about this in the ∞-categorical case. But if this identification is possible 1-categorically, i.e., if

$$\mathrm{Fun}^\times(\mathsf{1Span(Fin)}, \mathsf{Top}) \cong \mathrm{Fun}^{segal}(\mathsf{Fin}_*, \mathsf{Top}),$$

then we find a contradiction with Badzioch's result, for it identifies the LHS with topological abelian monoids while the RHS is $\mathbb{E}_\infty$-spaces. I am probably mistaken, so that the $\cong$ does not exist. Let's assume I did *not* make a mistake identifying $\mathrm{Fun}^\times(\mathsf{1Span(Fin)}, \mathsf{Set}) \cong \mathrm{Fun}^{segal}(\mathsf{Fin}_*, \mathsf{Set})$. Then why does replacing Set with Top and weak-ifying the conditions on the functors destroy this equivalence? Somehow, when we do this, the Segal condition has remained "good enough" to recover $\mathbb{E}_\infty$-spaces, while the product-preserving condition on functors out of $\mathsf{1Span(Fin)}$ has not. Maybe this is because $\mathsf{1Span}$ lacks automorphism information (this is "stored 2-categorically," which we cannot access because we took isomorphism classes), but perhaps that automorphism information *does* exist in the 1-categorical structure of $\mathrm{Fun}^{segal}(\mathsf{Fin}_*, \mathsf{Top})$?

My remark kind of spoiled the story. Spans are naturally 2-categorical, and we are missing crucial information by flattening them to a 1-category. On the level of Set, this is inconsequential, but for studying homotopy-coherent structures in ∞-categories, this matters.

Definition 11.12 (*2-category of spans*)**.** Let C be a category with pullbacks (and a functorial choice of pullbacks). We define a weak 2-category $\mathsf{2Span(C)}$ to have the same objects as C; to have spans as morphisms; and to have as 2-morphisms (i.e., morphisms of spans between two fixed X, Y) the *isomorphisms* of spans. The rest of the structure is more-or-less immediate. (Maybe c.f. this.)

This extra isomorphism data is important. And if we want to turn $\mathsf{2Span}(\mathsf{C})$ into a quasicategory, we should take these isomorphisms into account. First, we can turn $\mathsf{2Span}(\mathsf{C})$ into a simplicial category: we may form $\mathsf{Span}_\Delta(\mathsf{C})$ to have the same objects as C and to have *mapping simplicial sets* $\mathsf{Map}_{\mathsf{Span}_\Delta(\mathsf{C})}(X, Y)$ given by the *groupoid* of spans $X \to Y$ and isomorphisms between them. Finally, we can define

$$\mathsf{Span}_\infty(\mathsf{C}) := N^{coh}(\mathsf{Span}_\Delta(\mathsf{C})).$$

Proposition 11.13 (Elementary description of $\mathsf{Span}_\infty(\mathsf{C})$)**.** *Let C be a category with pullbacks. We will construct a quasicategory Z equivalent to $\mathsf{Span}_\infty(\mathsf{C})$. Let C_n denote the poset of nonempty subintervals of $[n]$, with the reverse inclusion ordering. The mapping $[n] \mapsto C_n$ defines a cosimplicial object $\Delta \to \mathsf{Cat}$. Now define $Z_n :=$ the set of functors $C_n \to \mathsf{C}$ with the pullback property: if $I, J \subseteq [n]$ are such that $I \cap J \neq \varnothing$, then*

$$\begin{tikzcd} F(I \cap J) \arrow[r] \arrow[d] & F(I) \arrow[d] \\ F(J) \arrow[r] & F(I \cup J) \end{tikzcd}$$

is a pullback diagram. Since face and degeneracy maps respect pullbacks, the mapping $[n] \mapsto Z_n$ extends to a simplicial set $Z : \Delta^{\mathrm{op}} \to \mathsf{C}$. The claim is that Z is a quasicategory and is isomorphic to $\mathsf{Span}_\infty(\mathsf{C})$.

Proof. Cranch, Prop. 4.5. □

Remark 11.14. Given this description of $\mathsf{Span}_\infty(\mathsf{C})$, how does $\mathsf{Span}_\infty(\mathsf{C})$ compare to the ∞-category $\mathsf{Span}(\mathsf{D}, \mathsf{D}^\dagger)$ which Barwick associates to an ∞-category D with a coWaldhausen structure $\mathsf{D}^\dagger$?

Definition 11.15. Suppose that C admits pullbacks. We define functors

$$\mathsf{Lid} : \mathsf{C} \to \mathsf{Span}_\infty(\mathsf{C}) \quad \text{and} \quad \mathsf{Rid} : \mathsf{C}^{\mathrm{op}} \to \mathsf{Span}_\infty(\mathsf{C}).$$

On 0-cells, these are the identity. On 1-cells $f : X \to Y$, we define $\mathsf{Lid}(f)$ as the "left identity" span $X = X \to Y$. We define Rid dually.[28]

[28] By the previous proposition, $\mathsf{Span}_\infty(\mathsf{C})$ is the (homotopy coherent) nerve of a 2-category. A map $NC \to ND$ is precisely a map $C \to D$ of 1-categories; can we pull a similar trick here? Regardless, I do not think it is hard to directly define Lid and Rid as functors to $\mathsf{Span}_\infty(\mathsf{C})$. C.f. Cranch, Prop. 4.5.

12 December

12.1 (12/11) Ambidexterity and semiadditivity recollection

I have been applying to grad schools, am nearly done, I forgot a bit what I was doing before, and my interests are changing. I want to read "Ambidexterity and height." Let me remember and fix some ideas, notation, human spirit.

Example 12.1. G finite, A abelian, action $G \to \mathrm{Aut}(A)$. Tautological maps $A^G \hookrightarrow A \twoheadrightarrow A_G$. *Norm map* $Nm : A_G \to A^G$ given by $[a] \mapsto \sum_G ga$. (Pointedly, think of as "at each $a \in A/G$, take fiber of $A \to A/G$, sum over fiber.") The composite $A_G \to A^G \to A_G$ is multiplication by $|G| \implies$ if A is a char p vector space, then Nm cannot be an equivalence whenever p divides $|G|$. Turns out that if A is rational, then Nm is always an equivalence.

Example 12.2. Analogous situation for finite G acting on spectrum A (i.e., Tate vanishing if A is rational, not generally so if A is p-local). But now can also ask about "intermediary characteristic." That is, let G act on $K(n)$-local A. Turns out, still have Tate vanishing! This is a surprise. Unlike the rational case, have $p = 0$ in $K(n)_*$, yet we still have the isomorphism generally.

Other similar results, e.g., for $T(n)$-local spectra. Hopkins-Lurie: formal framework for capturing these sorts of phenomena? It's more likely than you think. Let's rephrase the previous example:

Example 12.3 (The same example, pointedly reworded)**.** Denote by $p : BG \to *$ the projection, and consider its pullback $p^* : \mathsf{Sp}_{K(n)} \to \mathrm{Fun}(BG, \mathsf{Sp}_{K(n)})$. The pullback p^* admits left and right adjoints, which take $E : BG \to \mathsf{Sp}_{K(n)}$ to E_{hG} and E^{hG}, respectively. If A is a $K(n)$-local spectrum with G-action, we get a functor $A : BG \to \mathsf{Sp}_{K(n)}$. From this perspective, the norm $N : A_G \to A^G$ occurs as part of a natural transformation

$$Nm : p_! \to p_*.$$

The previous example said the norm $Nm : A_G \to A^G$ was an isomorphism. In fact, this $Nm : p_! \to p_*$ is a natural isomorphism.

This suggests the general setup. Let C denote any category and "relativize" by replacing $p : BG \to *$ with any map $f : X \to Y$. We may still form the pullback $f^* : \mathsf{C}^Y \to \mathsf{C}^X$ and ask whether there exists an adjoint triple

$$f_! \dashv f^* \dashv f_*.$$

These adjoints existed in our motivating cases, wherein the norm map was a morphism $p_! \to p_*$. But for general C and f, the existence of $f_!$ and f_* is nontrivial, and even if they exist, *a priori* there is no canonical map $f_! \to f_*$. The is *the* fundamental problem. It is dauntingly general. But a simple observation cracks it wide open: if f is k-truncated, then its diagonal δ is $(k-1)$-truncated. This is a first step toward realizing a close relation between the theory of "abstract norms" and the higher algebra of C.

Let me explain this. Suppose that f is k-truncated. First, we want $\delta_!, \delta_*, f_!$, and f_* to exist. A standard detail is that these $\delta_!, f_!$ are left Kan extensions and δ_*, f_* are right Kan extensions. Thus, they are given by taking (co)limits over certain diagrams indexed by fibers. As f is k-truncated (resp. as δ is $(k-1)$-truncated), an easy way to make $f_!, f_*$ exist (resp. for $\delta_!, \delta_*$ to exist) is to assume C admits all (co)limits indexed over k-truncated diagrams (resp. $(k-1)$-truncated diagrams). In particular, if we suppose that f is k-finite, then all the functors exist as soon as C admits all (co)limits indexed by k-finite spaces. Analogously for δ.

We wanted conditions for $f_!, f_*$ to exist because how else could we conceive of a norm map $f_! \to f_*$. But why care about $\delta_!, \delta_*$? Because there is a natural and canonical association

$$\mathsf{Map}(\delta_*, \delta_!) \to \mathsf{Map}(f_!, f_*)$$

built from (co)units and certain obvious(?) identities.[29] In particular, if Nm_δ is an equivalence, the norm Nm_f should be defined as that associated to Nm_δ^{-1}. We still have no definition of Nm_δ, though.

...However, a map f is (-2)-truncated if and only if it is an equivalence, in which case $f_!, f_*$ are both homotopy inverses to f^* so we get a canonical $f_! \cong f_*$. In particular, if f is (-1)-truncated, then we define $\mathrm{Nm}_\delta : \delta_! \to \delta_*$ as this canonical isomorphism, and we define Nm_f as the transformation associated to Nm_δ^{-1}. All that is assuming the various functors left/right adjoints exist.

[29] Identities relating the diagonal $\delta : X \to X \times_Y X$ and the projections $X \times_Y X \to X$.

We have talked about the existence of $f_!$ and f_* and the construction of $\mathrm{Nm}_f : f_! \to f_*$. The map Nm_f is constructed from the same data for δ (whose truncation order is one lower) with the additional requirement that Nm_δ^{-1}. In the base case (truncation order -2), the situation degenerates, and Nm_δ is automatically defined and is an isomorphism. Let's collect all this (existence, duality, base case) into one inductive definition.

Definition 12.1. Let C be any category and $f : X \to Y$ a map of Kan complexes.

1. Say that f is *(-2)-ambidextrous* if and only if it is an equivalence. In this case, define Nm_f as the canonical isomorphism $f_! \cong f_*$.
2. For $n \geq -1$, say that f is *weakly n-ambidextrous* if and only if δ is $(n-1)$-ambidextrous.
3. For $n \geq -1$, say that f is *n-ambidextrous* if and only if
   (a) f is weakly n-ambidextrous (so that Nm_δ^{-1} exists),
   (b) C admits (co)limits indexed by the fibers of f (so that $f_!$ and f_* exist), and
   (c) The map $\mathrm{Nm}_f : f_! \to f_*$ associated to Nm_δ^{-1} is an equivalence.

Hence, ambidexterity is about norms—when they exist and are isomorphisms. We may say that f is *C-ambidextrous* if it is n-ambidextrous for some n. Let us investigate this notion for various f.

Example 12.4. Suppose that a map f is (-1)-truncated (i.e., its fibers are empty or contractible). Is it (-1)-ambidextrous? Since δ is (-2)-truncated it is an equivalence, so Nm_δ is an equivalence and f is weakly (-1)-ambidextrous. For f to be (-1)-ambidextrous, C must admit f-limits and f-colimits so that $f_!$ and f_* exist, which means admitting (co)limits over all empty and singleton diagrams. This amounts to C having initial and final objects. It turns out that if these coincide, i.e. if C is pointed, then Nm_f is an equivalence, in which case f is (-2)-ambidextrous.

Example 12.5. Suppose that a map f is 0-truncated (i.e., its fibers are disjoint unions of contractible spaces). Further assume that C is pointed, so that Nm_δ is automatically an equivalence by the previous example (since δ is (-1)-truncated). In order for $f_!$ and f_* to exist and thus for Nm_f to be defined, C must admit f-(co)limits, which amounts to C having countable (co)products. It turns out that if countable (co)products coincide, i.e. if C admits biproducts, then Nm_f is an isomorphism and so f is (-1)-ambidextrous.

Example 12.6. Suppose that f is 0-truncated and *the fibers of f are finite*. By the same argument as above, if C is pointed, then f is weakly 0-ambidextrous, and if C admits *finite* biproducts, then f is 0-ambidextrous. We generally like to impose such finiteness conditions on f because it is much easier for C to satisfy the conditions for Nm_f to be an equivalence. In this case, it is the difference between having countable versus finite biproducts—that is significant! For example, Ab has finite but not countable biproducts.

Example 12.7. Suppose that f is 1-truncated and that its fibers have finite homotopy groups. Thus, its fibers are of the form $\coprod_{i=1}^{N<\infty} BG_i$ where each G_i is finite. Since δ is 0-truncated and its fibers are finite, by the previous example Nm_δ is an equivalence once C admits finite biproducts, in which case f is weakly 1-ambidextrous. In order for $f_!$ and f_* to exist and thus for Nm_f to be defined, C must admit f-(co)limits, which amounts to admitting (co)limits indexed by spaces of the form $\coprod_{i=1}^{N<\infty} BG_i$ with G_i finite. If $\mathsf{C} = \mathsf{Sp}_{K(n)}$ for $n \neq \infty$, the classical theorem of [HS96] and [GS96]says that for G finite, the norm Nm_p associated to $p : BG \to *$ is an equivalence, in other words that $p : BG \to *$ is 1-ambidextrous. The power of our "abstract norm formalism" is that we can now say that *the only necessary hypotheses for this result are that $p : BG \to *$ is 1-truncated and has π_*-finite fibers.* Indeed, it turns out that if $\mathsf{C} = \mathsf{Sp}_{K(n)}$ and $f : X \to Y$ is *any* such map, then f is 1-ambidextrous. In fact, Hopkins-Lurie proves much more.

Theorem 12.2. *If $\mathsf{C} = \mathsf{Sp}_{K(n)}$ and $f : X \to Y$ is n-finite (to mean that it is n-truncated and its fibers have finite homotopy groups), then f is n-ambidextrous. In other words, for such f, the norm*

$$\mathrm{Nm}_f : f_! \to f_*$$

is an isomorphism.

Some remarks are in order. First, truncation order seems to be the right property by which to organize and study ambidexterity (more sophisticated results would confirm this). Second, as in Examples 12.6 and 12.7, it seems right to further restrict to the notion of *n-finite* maps, to mean any $f : X \to Y$ whose fibers are all n-truncated and have finite homotopy groups. For without finiteness, it is harder for C to admit all f-(co)limits and thus for Nm_f to exist, and *much* harder for f-(co)limits to coincide and thus for Nm_f to be an equivalence. These two points in mind and considering the theorem, it is natural to ask: given $n \geq -2$, which categories C are such that all n-finite maps are C-ambidextrous?

Definition 12.3. For $m \geq -2$, say that a category C is *m-semiadditive* if every m-finite map is C-ambidextrous.

This is a great definition. We know that for a chosen f, the existence of $f_!$ and f_* is about the existence of (co)limits indexed by its fibers, and that Nm_f is an isomorphism once these (co)limits canonically coincide. Hence, m-semiadditivity is about when a category C admits canonically isomorphic limits and colimits for every m-finite diagram $X \to$ C (the isomorphism being the norm map associated to $p : X \to *$, which is $\mathrm{Nm}_p : \varinjlim X = p_! \to p^* = \lim X$).

(Insert story: semiadditivity in ordinary situation is about commutative monoids, this "higher semiadditivity" leads us to realize "higher monoids," thus norms are related to semiadditivity are related to higher algebra of C.)

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