## Admissible Strategies in Safety Games

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### AGENDA

INTRODUCTION AND PRELIMINARIES OS SAFETY GAMES AND WINNING STRATEGIE

WINNING STRATEGIES

RESULTS AND FUTURE WORK

# O1 INTRODUCTION AND PRELIMINARIES

## Introduction and Motivation

#### 1. Controller Synthesis

 System and Environment where system trying to achieve its goal.

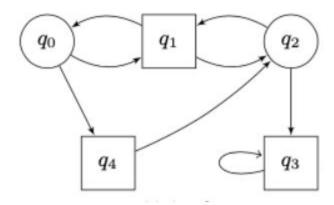
### 2. Graph Games and Safety Games

 Approach to attempt to solve controller synthesis problem.

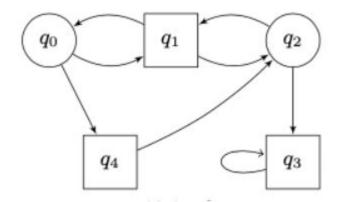
### 3. Winning Strategies and Admissible Strategies

Why do we need admissible strategies?

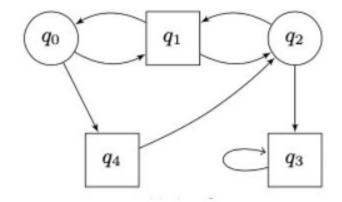
 Arena → Directed graph with finite vertices, vertices divided between players, edges as moves



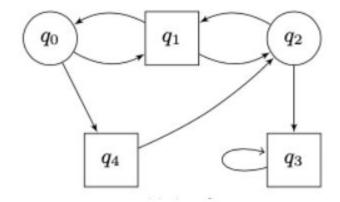
- Arena
- Winning Condition → Some given constraints which the players has to abide by to win the game. (ex. Fencing in a park)



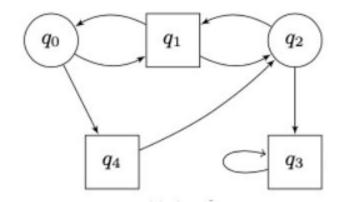
- Arena
- Winning Condition
- Multiplayer Game → A game played on a given arena with many players trying to achieve some given conditions.



- Arena
- Winning Condition
- Multiplayer Game
- Strategy → Given the past moves and the current state of a player, which state should it send the token next to.



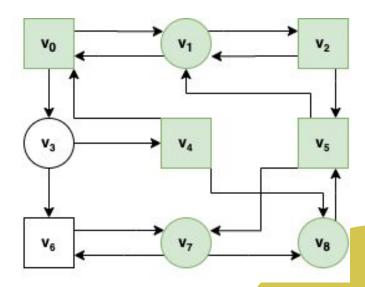
- Arena
- Winning Condition
- Multiplayer Game
- Strategy
- Winning Strategy → A strategy from a vertex which helps to satisfy the winning condition



# O2 SAFETY GAMES AND WINNING STRATEGIES

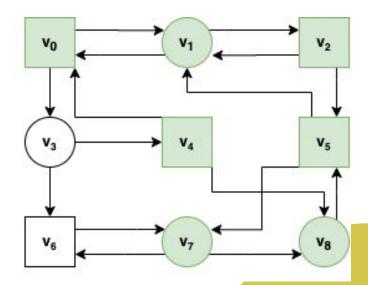
### What is a Safety Game?

- A game with a defined winning condition.
- We have a set of **safe vertices** S.
- Bad vertices → All vertices except S.
- Winning condition → Do not transition to any "bad" vertex.



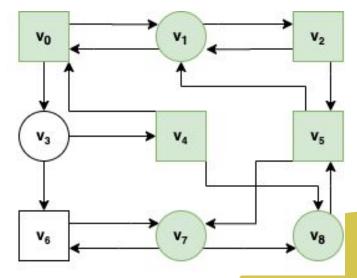
### **Brief Overview**

- Assume we are player 0 → if we start from v<sub>1</sub>, and we go to v<sub>0</sub>, then player 0 can send token to v<sub>3</sub> which is unsafe and we lose.
- But if we play strategically, and go to v<sub>2</sub>, no matter where player 1 sends the token, to v<sub>5</sub> and then to v<sub>1</sub> or v<sub>7</sub>, its safe.



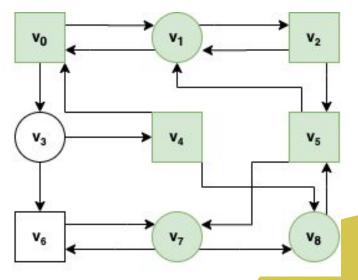
## Finding if Winning strategy exist in a Safety Game?

- Implemented the **algorithm** from scratch to find the winning strategy.
- Used concepts from Reachability Games to solve.
- Reachability Games → Transition to a "good" vertex belonging to a set R.
   (Opposite of a safety game)



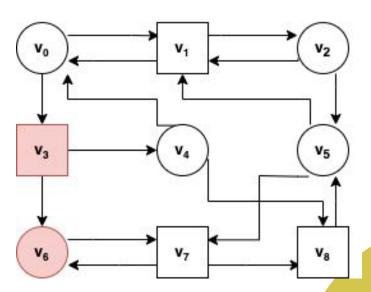
## Finding if Winning strategy exist in a Safety Game?

 Duality Theorem → If player 0 wins in a safety game i.e. she is successful in keeping the token in S then player 1 wins the reachability game in the dual of the graph (with vertices swapped and R = complement of S)



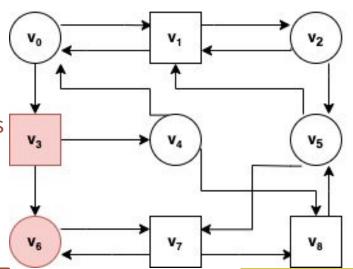
### **v**<sub>1</sub> $v_0$ V<sub>2</sub> $V_4$ $V_3$ $V_5$ V<sub>6</sub> V<sub>7</sub> **v**<sub>8</sub>

## Dual

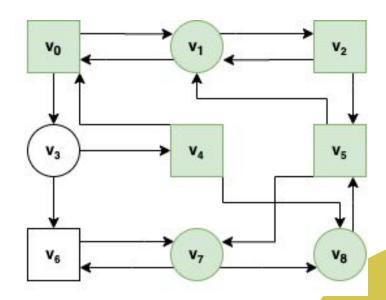


## Finding a Winning strategy in a Reachability Game?

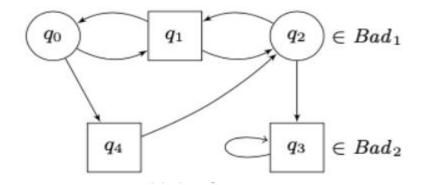
- **Intuition** → Start with set R and some player.
- Now expand R to include those states from where we are bound to reach R. (CPre.)
- Worst case → only one vertex added this time.
   Keep iterating and continue adding elements to R, eventually will converge since number of vertices finite. Will yield what is called the attractor for the region.



Code Demonstration to find if winning strategy exists



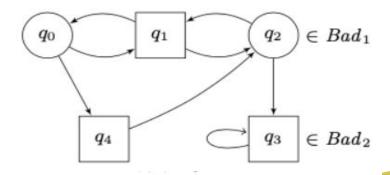
What if....?
No winning strategy for Player 0:(



# O3 ADMISSIBLE STRATEGIES

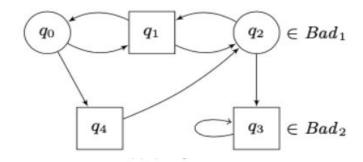
### Why Admissible strategies?

- Scenarios where Player 0 has no winning strategy
- What if Player 1 helps?
- Play rationally and not in an adversarial manner
- Both meet their objectives



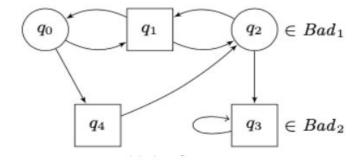
### Dominance of strategies

- Intuition → Consider q<sub>0</sub>. If we go to q<sub>4</sub>, we are bound to lose. How? q<sub>4</sub> will send it to q<sub>2</sub>.
- What if we go to q<sub>1</sub>? We may win in that case if Player 1 helps. As seen above, we look for a win-win situation.
- **Dominance** We can say  $q_0 \rightarrow q_1$  dominates  $q_0 \rightarrow q_4$ .

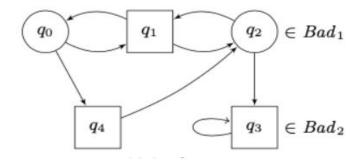


## Dominance of strategies

- Notion of value.
- **Use of subroutine** to find winning strategies.
- Implemented algorithm from scratch to iteratively find dominated strategies and remove them.
- And the strategies left are admissible in nature.



## Code Demonstration to find admissible strategies



# O4 RESULTS AND FUTURE WORK

### **RESULTS**

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#### **SAFETY GAMES**

Implemented algorithm to find winning strategies

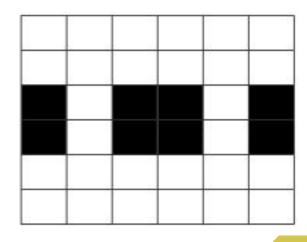
02

### **ADMISSIBLE STRATEGIES**

Implemented algorithm to find admissible strategies

## Problem in Robot Motion Planning

- Given an N x N grid with 2 robots with some obstacles, we have an uncontrolled robot which moves in an unconstrained way and a controlled robot which can sense the other robot in a 1 unit distance only (corners included).
- **Objective** → Avoid collision between the two and let them finish their tasks.



### WORK TILL NOW AND FUTURE WORK

#### PART 2

Algorithm to find strategies in **safety** games

#### PHASE 4

Describe Robot **Motion** Planning in LTL

#### PHASE 6

Merge everything to try to solve the robot motion problem.











Get acquainted with the domain

### PHASE 3

Algorithm to find **Admissible** strategies

#### PHASE 5

Convert LTL to a safety game (complex)

### **ACKNOWLEDGEMENT**

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We would like to express our gratitude to our mentor Prof. Purandar Bhaduri for giving us the opportunity to explore this new field of Infinite Games.

## THANK YOU

We are open to questions, if any.