

# Network Project - Phase 2

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GSync v2 Protocol Mini-RFC

## Submitted By

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# GSync v2 Protocol Mini-RFC

**Protocol Name:** GSync v2 (Grid Synchronization Protocol Version 2) **Status:** Phase 2  
**Implementation Date:** December 2025

## 1. Introduction and Design Philosophy

GSync v2 is a **lightweight, UDP-based application-layer protocol** designed for **low-latency synchronization of game state** in multiplayer environments. It prioritizes real-time consistency and efficient bandwidth utilization.

The protocol addresses the unacceptable latency of existing TCP-based or heavyweight frameworks like ENet for fast-paced games.

### 1.1 Key Motivations & Constraints

The design goals reject reliability for **responsiveness**.

Feature	Specification/Target	Source
<b>Transport</b>	UDP only (connectionless, unreliable)	
<b>Update Rate</b>	20-60 Hz (Phase 2 uses 20 Hz)	
<b>Latency Target</b>	≤50 ms average end-to-end	
<b>Header Size</b>	Compact 28-byte header	
<b>Loss Tolerance</b>	Graceful degradation at 2–5% loss (Redundancy-based)	

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### 1.2 Core Design Principles

- **No Retransmission** and **No Acknowledgments** (reduces round-trip overhead).
- **Timestamp-based Reordering** (clients handle out-of-order packets).

- **Redundancy over Windowing** (simpler, effective for short timescales).

## 2. Protocol Architecture

### 2.1 Entities and Roles

- **Server (Authoritative):** Maintains ground truth grid state, broadcasts snapshots at fixed intervals, and resolves conflicts.
- **Clients (Followers):** Receive and apply snapshot updates, send input events (cell acquisitions), and render local game state.

### 2.2 Communication Flow

The communication flow shows the connection setup, continuous synchronization (20 Hz), and termination. > **Key Steps:** Client sends **INIT** → Server registers client → Server **Broadcasts SNAPSHOT+REDUNDANCY** (20 Hz) → Clients apply updates → Server sends **GAME\_OVER** (x2) on end condition.

## 3. Message Formats

### 3.1 Header Structure (28 bytes)

The header is a fixed 28-byte structure. The payload length is limited to 1200 bytes per packet.

Field	Offset	Size (bytes)	Type	Description
<b>Protocol ID</b>	0	4	ASCII	"GSYN" (0x4753594E)
<b>Version</b>	4	1	uint8	Protocol version = 1
<b>Message Type</b>	5	1	uint8	See section 3.2
<b>Snapshot ID</b>	6	4	uint32	Incremental counter

<b>Sequence Number</b>	10	4	uint32	Packet sequence for gap detection
<b>Server Timestamp</b>	14	8	uint64	ms since epoch (for latency calculation)
<b>Payload Length</b>	22	2	uint16	Bytes in payload (0-1200)
<b>CRC32 Checksum</b>	24	4	uint32	CRC32(header_zero + payload)

### 3.2 Message Types

Type	Value	Direction	Payload	Purpose
<b>INIT</b>	3	C → S	player_id (1 byte)	Register client
<b>SNAPSHOT</b>	0	S → C	Grid state (~303 bytes)	Broadcast state
<b>EVENT</b>	1	C → S	Acquire request (12 bytes)	Player action
<b>GAME_OVER</b>	4	S → C	Winner + scores	End game

## 4. Reliability and Performance Features

The protocol leverages redundancy instead of traditional retransmission to maintain low latency.

### 4.1 Redundancy-Based Reliability (K=3)

The SNAPSHOT payload includes **K=3 redundancy**, meaning the current snapshot is sent along with the two most recent previous snapshots (150ms history at 20 Hz).

- The Client extracts the first (newest) snapshot only.

**Redundancy provides insurance against burst loss and reordering, leading to better user experience than late corrections.**

### 4.2 Critical Event Reliability (Double-Send)

For critical but rare events, like the cell acquisition **EVENT**, the client performs a **double-send** with a 1ms spacing. This is effective simple redundancy to dodge the same loss pattern.

## 4.3 Checksum and Detection

- **CRC32 Checksum Validation:** The 4-byte CRC32 in the header detects bit-flip errors and is industry standard. Packets failing validation are discarded.  
  
**Duplicate Detection & Suppression:** Clients track the `last_seq_num` to skip duplicates and outdated snapshots.  
  
**Sequence Gap Detection:** Clients detect lost packets by checking if `seq_num > last_seq_num + 1`.

## 5. Experimental Evaluation

The protocol's performance is tested across four scenarios, including a **Baseline**, 2% and 5% **Packet Loss**, and **100ms Latency**.

### 5.1 Acceptance Criteria (Expected Results)

Scenario	Latency	Jitter	Loss Tolerance
<b>Baseline</b>	$\leq 50\text{ms}$	$< 10\text{ms}$	99% delivery
<b>Loss 5%</b>	$\leq 75\text{ms}$	$< 20\text{ms}$	95% delivery
<b>Delay 100ms</b>	100–150ms	$< 25\text{ms}$	95% delivery

### 5.2 Metrics Collected

Metrics include per-packet measurements logged to CSV files, such as **latency** (`recv_time - server_timestamp`), **jitter**, **duplicate flag**, and **gap count**. Aggregate metrics include mean, median, and 95th percentile for latency and jitter, duplicate rate, and server-side CPU utilization.

## **6. Evaluation**

Scenarios: Baseline, 2-5% loss, 100ms delay.

Metrics: Latency (50-75ms target), jitter, duplicates, delivery (95-99%).

Logging: CSV for snapshots/diagnostics; pcap captures.

## **7. Example Walkthrough**

Two-player game: INIT → SNAPSHOTs → EVENTs → GAMEOVER with scores.

## **8. Limitations & Future**

- Fixed 20 Hz; no encryption.
- Future: Adaptive rates, encryption, more players, use ack.