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Problem Set 1

Problem Set 1

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Problem 1-1.

(a)

$$f_1 = \log n^n = O(n)$$

$$f_2 = (\log n)^n = O((\log n)^n)$$

$$f_3 = \log n^6 006 = O(\log n)$$

$$f_4 = (\log n)^6 006 = O((\log n)^c)$$

$$f_5 = \log \log 6006n = O(\log \log n)$$

solution:

$$(f_5, f_3, f_4, f_1, f_2)$$

(b) solution:

$$(f_1, f_2, f_5, f_4, f_3)$$

(c) solution:

$$(\{f_2, f_5\}, f_4, f_1, f_3)$$

(d) solution:

$$(f_5, f_2, f_1, f_3, f_4)$$

2 Problem Set 1

Problem 1-2.

(a) iterate over the given data structure starting from index i to (i + k - 1) // 2, swap i with (i + k - 1) using two variables x1 and x2 to store the returned values from delete_at(i) and delete_at(i + k - 1) then use insert_at(i, x2) and insert_at(i + k - 1, x1) to complete the swap. This recursive procedure makes no more than k/2 = O(k) recursive calls before reaching a base case, doing O(log n) work per call, so the algorithm runs in O(k log n) time.

```
1  def reverse(D, i, k):
2    if k < 2:
3        return
4
5        x1 = D.delete_at(i);
6        x2 = D.delete_at(i + k - 1);
7        D.insert_at(i, x2);
8        D.insert_at(i + k - 1, x1);
9        reverse(D, i + 1, k - 2)</pre>
```

(b) we can solve it recusively by setting the base case that is k == 0 then do nothing and recusively use $delete_at(i + k - 1)$ ans store the returned value in x1 then we check if j if bigger than i we do $insert_at(j - 1, x1)$ as deleting an element make a shift by one in the desired index to insert at else if not j > i we $insert_at(j, x1)$, then call move (D, i, k - 1, j) This recursive procedure makes no more than k = O(k) recursive calls before reaching a base case, doing O(log n) work per call, so the algorithm runs in O(k log n) time.

```
def move(D, i, k, j):
    if(k == 0):
        return

x1 = D.delete_at(i + k - 1);

if j > i:
        D.insert_at(j - 1, x1)
else:
        D.insert_at(j, x1)

move(D, i, k - 1, j)
```

Problem Set 1 3

Problem 1-3.

Problem 1-4.

(a) To support insert_first (x) in O(1) time we can use L. head pointer by make x.next equals the head pointer then make L. head points to the node x. This algorithm does occurrent operations in so it runs in O(1).

To support insert_last (x) in O(1) time we can use L.tail pointer by make L.tail.next points to the Node x then make L.tail points to the node x. This algorithm does occustant-time operations in so it runs in O(1).

To support delete_first () in O(1) time we make L.head points to L.head.next. (assuming that L.head exits) This algorithm does occurrent operations in so it runs in O(1).

To support delete_last() in O(1) time we make L.tail points to L.tail.prev. (assuming that L.tail exits) This algorithm does occurrent operations in so it runs in O(1).

- (b) make the head pointer of L2 points to x1.
 - make the tail pointer of L2 points to x2.
 - if x1 is the head
 - make x1.prev equals x2.next
 - if x1 is **NOT** the head
 - make x1.prev.next equals x2.next
 - if x2 is the tail
 - make the tail equals x1.prev
 - if x1 is **NOT** the tail
 - make x2.next.prev equals x1.prev
 - set x1.prev and x2.next to None
 - return L2
- (c) make L₂.head.prev equals to x.
 - save x.next in xn variable.
 - make x.next equals L2.head.
 - make L2.tail.next equals xn.
 - if x is the tail of L1
 - make L1.tail equals L2.tail.

4 Problem Set 1

• if x is **not** the tail of L1

- make xn.prev equals L2.tail.

```
(d)
   class Doubly_Linked_List_Node:
       def __init__(self, x):
           self.item = x
           self.prev = None
           self.next = None
      def later_node(self, i):
           if i == 0: return self
 9
           assert self.next
           return self.next.later_node(i - 1)
   class Doubly_Linked_List_Seq:
       def __init__(self):
           self.head = None
           self.tail = None
       def __iter__(self):
           node = self.head
19
           while node:
               yield node.item
               node = node.next
       def __str__(self):
24
           return '-'.join([('(%s)' % x) for x in self])
26
       def build(self, X):
           for a in X:
                self.insert_last(a)
       def get_at(self, i):
           node = self.head.later_node(i)
           return node.item
34
       def set_at(self, i, x):
           node = self.head.later_node(i)
36
           node.item = x
       def insert_first(self, x):
           node = Doubly_Linked_List_Node(x)
40
           if self.head is None:
41
               self.head = node
                self.tail = node
43
           else:
44
               node.next = self.head
45
               self.head.prev = node
46
               self.head = node
47
```

Problem Set 1 5

```
48
       def insert_last(self, x):
           node = Doubly_Linked_List_Node(x)
           if self.head is None:
               self.head = node
               self.tail = node
           else:
               self.tail.next = node
               node.prev = self.tail
               self.tail = node
62
       def delete_first(self):
63
           x = self.head.item
64
           self.head = self.head.next
           # if we the first ans last elements are the same node
66
           if self.head is None:
               self.tail = None
           else:
               self.head.prev = None
           return x
       def delete_last(self):
           x = self.tail.item
           self.tail = self.tail.prev
           # if we the first ans last elements are the same node
           if self.tail is None:
               self.head = None
           else:
               self.tail.next = None
           return x
83
       def remove(self, x1, x2):
85
           L2 = Doubly_Linked_List_Seq()
86
           L2.head = x1
           L2.tail = x2
           if x1 == self.head:
89
               x1.prev = x2.next
           else:
               x1.prev.next = x2.next
93
           if x2 == self.tail:
94
               self.tail = x1.prev
               x2.next.prev = x1.prev
98
```

6 Problem Set 1