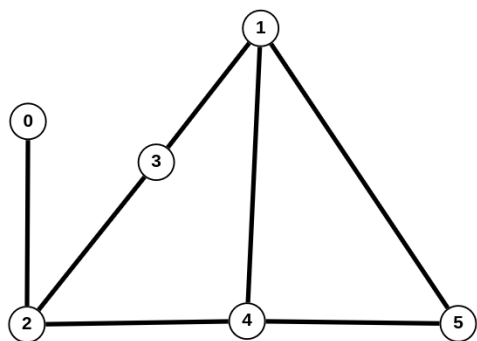


Problem Set 5

Name: Your Name

Collaborators: Name1, Name2

Problem 5-1.

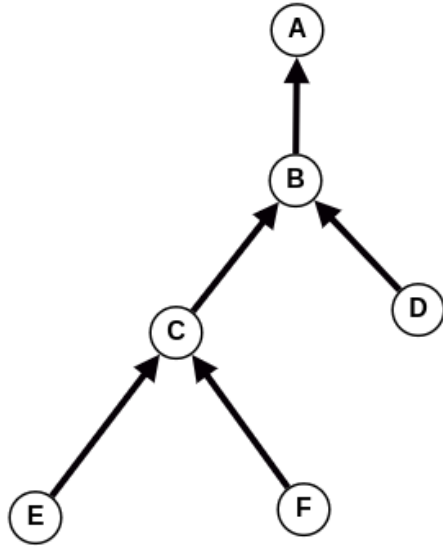


(a)

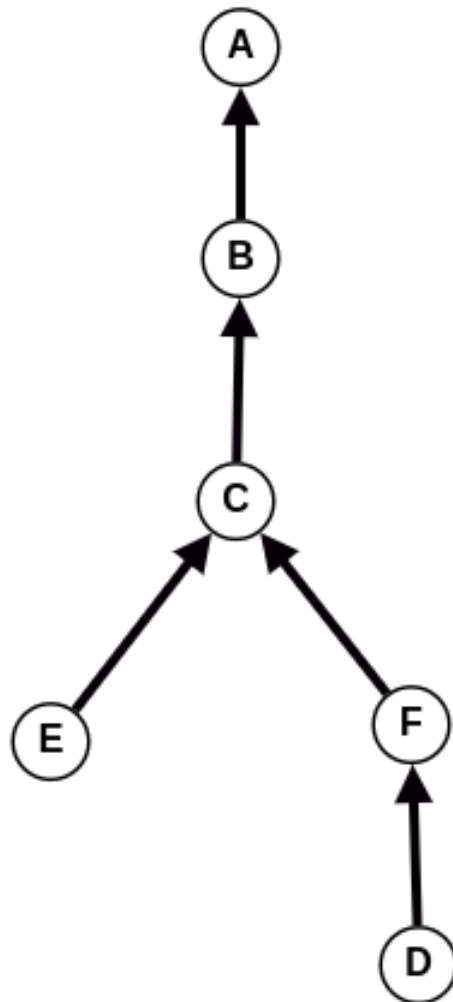
(b)

```
1 Adj = {
2     "A" : [ "B" ],
3     "B" : [ "C", "D" ],
4     "C" : [ "E", "F" ],
5     "D" : [ "E", "F" ],
6     "E" : [ ],
7     "F" : [ "E", "D" ],
8 }
```

(c) BFS : [A, B, C, D, E, F]



DFS : [A, B, C, E, F, D]



- (d) we can remove (F, D) which gives a **DAG** with a topological order of [A, B, C, D, F, E]
 and (D, F) which gives a **DAG** with a topological order of [A, B, C, F, D, E]

Problem 5-2. Construct a graph that contains the buildings and plants as its vertices and the edges are the wires in $O(n^4)$

Problem 5-3. Construct a graph $G = (V, E)$, V is the set of friends which every two short-circuiting friends are connected by a single an edge in E . to make this party possible we need G to be a bipartite graph. if G is a bipartite graph, then:

- G is 2-colorable.
- G does not contain any cycles of odd length.

The following algorithm will do the job:

- run **Full-BFS** on G and color every vertex at even level with the same color of the source vertex.
- check every edge in E to see if it connects two vertices with the same color, if so then G is not a bipartite graph and this party is not possible, else G is a bipartite graph.

Time complexity: $O(n)$

- Constructing G is $O(n)$, because we have n edges, and vertices are $O(n)$ because it is at most $2n$ vertices.
- Running Full-BFS is $O(n)$, since it is linear in the graph size which is $O(n)$.
- Checking every edge is $O(n)$.

Problem 5-4.

Problem 5-5.

Problem 5-6.

- (a)
- (b)
- (c)
- (d) Submit your implementation to `alg.mit.edu`.