

Assignment - 1

Q. What is the time complexity of the following code?

```
int a = 0;  
for (i = 0; i < N; i++)
```

```
{  
    for (j = N; j > i; j--)
```

```
    {  
        a = a + i + j;
```

```
    }
```

```
}
```

sol: int a = 0; $\rightarrow 1$

for (i = 0; i < N; i++) $\rightarrow n$

{
 for (j = N; j > i; j--) $\rightarrow n(n-i)$

{

a = a + i + j; $\rightarrow n(n-i)$

}

}

$$T_n = 1 + n + n^2 - nk_1 + n^2 - nk_1$$

$$T_n = 2n^2 + (1 - 2k_1)n + 1$$

Drop the non-dominant terms.

$$T_n = 2n^2$$

Drop the constant coefficients.

$$T_n = n^2 = O(n^2)$$

② Let us say you need to double an array size whenever the array is full. What is the amortized cost of operations involved?

Sol: When array is not full, adding an element takes $O(1)$

When array is full, doubling its size involves

- i, Allocating a new array of double the size.
- ii, Copying all existing elements from the old array to the new array. This copying operation takes $O(n)$, where n is the size of the array at the time of resizing.

Amortized Cost:

Cost of each resizing operation is proportional to the size of the array being copied.

Consider n operations on the array.

Resizing occurs at array sizes

$$1, 2, 4, 8, \dots, 2^k \quad \text{where } 2^k \leq n$$

$$\text{Total cost (T)} = 1 + 2 + 4 + 8 + \dots + n = n$$

$$T = 2n - 1$$

$$T \approx 2n$$

$$\text{Amortized cost} = \frac{T}{n}$$

$$= \frac{2n}{n}$$

$$= 2$$

\therefore Amortized cost of appending an element to an array that

doubles in size when full is $O(1)$.

③ Show that if fewer than $\lceil n/2 \rceil$ unions are performed, then at least one set with a single element in it remains.

Sol: Consider n sets with each containing a distinct element.

Each union operation reduces the total number of sets by 1.

After k union operations, the no. of sets remaining will be $n - k$.

If fewer than $\lceil n/2 \rceil$ union operations are performed then

$$k < \lceil n/2 \rceil.$$

\therefore no. of sets remaining.

$$n - k > n - \lceil n/2 \rceil$$

$$n - \lceil n/2 \rceil$$

If n is even;

$$\lceil n/2 \rceil = n/2 \text{ so } n - n/2 = n/2$$

If n is odd,

$$\lceil n/2 \rceil = (n+1)/2$$

$$\begin{aligned} n - \lceil n/2 \rceil &= n - (n+1)/2 \\ &= (n-1)/2 \end{aligned}$$

In either case, $n - k > n - \lceil n/2 \rceil$ means more than one set remains. Since there are still multiple sets, at least one of these sets must still contain a single element.

\therefore If fewer than $\lceil n/2 \rceil$ union operations are performed, at least one set with a single element remains.

④ Show the resulting sequence of operations using Union by Rank. Union(1,3), Union(2,4), Union(3,5), Union(4,6), Union(7,8), Union(9,10), Union(5,7), Union(6,9), Union(8,10) with link and rank arrays.

Sol: link[]: Points to the parent of each element in the disjoint set forest

rank[]: Tracks the rank of the tree rooted at each element.

Initially, each element is its own parent and has a rank of 0.

link: [1, 2, 3, 4, 5, 6, 7, 8, 9, 10]

rank: [0, 0, 0, 0, 0, 0, 0, 0, 0, 0]

Operations:

1. Union(1,3):

→ Both are its own ~~parents~~ roots

→ Both ranks are 0.

→ Arbitrarily make 1 as root and link 3 to 1 and increase the rank of 1 by 1.

link: [1, 2, 1, 4, 5, 6, 7, 8, 9, 10]

rank: [1, 0, 0, 0, 0, 0, 0, 0, 0, 0]

2. Union(2,4):

→ Both 2 and 4 are their own ~~parents~~ roots.

→ Both ranks are 0.

→ Arbitrarily make 2 as root, link 4 to 2 and increase the rank of 2 by 1.

link: [1, 2, 1, 2, 5, 6, 7, 8, 9, 10]

rank: [1, 1, 0, 0, 0, 0, 0, 0, 0, 0]

3. Union (3, 5) :-

- Root of 3 is 1 and 5 is its own root.
- Rank of 1 is higher than rank of 5
- link 5 to 1.

link: [1, 2, 1, 2, 1, 6, 7, 8, 9, 10]

rank: [1, 1, 0, 0, 0, 0, 0, 0, 0, 0]

4. Union (4, 6) :-

- Root of 4 is 2 and 6 is its own root.
- Rank of 2 is 1 which is higher than rank of 6 is 0.
- link 6 to 2.

link: [1, 2, 1, 2, 1, 2, 7, 8, 9, 10]

rank: [1, 1, 0, 0, 0, 0, 0, 0, 0, 0]

5. Union (7, 8) :-

- 7 and 8 are their own roots
- Both ranks are 0.
- Arbitrarily make 7 as root, link 8 to 7 and then increase the rank of 7 by 1

link: [1, 2, 1, 2, 1, 2, 7, 7, 9, 10]

rank: [1, 1, 0, 0, 0, 0, 1, 0, 0, 0]

6. Union (9, 10) :-

- 9 and 10 are their own roots
- Both ranks are 0.
- Arbitrarily make 9 as root, link 10 to 9 and then increase the rank of 9 by 1.

link: [1, 2, 1, 2, 1, 2, 7, 7, 9, 9]

rank: [1, 1, 0, 0, 0, 0, 1, 0, 1, 0]

7. Union (5, 7):

→ Root of 5 is 1 and root of 7 is 7.

→ Both ranks of 1 and 7 are 1.

→ Arbitrarily make 1 as root and link 7 to 1, then increase the rank of 1 by 1.

link: [1, 2, 1, 2, 1, 2, 1, 7, 9, 9]

rank: [2, 1, 0, 0, 0, 0, 1, 0, 1, 0]

8. Union (6, 9):

→ Root of 6 is 2 and root of 9 is 9.

→ Both rank of 2 and 9 are 1.

→ Arbitrarily make 2 as root, link 9 to 2 and increase the rank of 2 by 1.

link: [1, 2, 1, 2, 1, 2, 1, 7, 2, 9]

rank: [2, 2, 0, 0, 0, 0, 1, 0, 1, 0]

9. Union (8, 10):

→ Root of 8 is 7 and Root of 10 is 9.

→ Both ranks of 7 and 9 are 1.

→ Arbitrarily make 7 as root and link 9 to 7 and increase the rank of 7 by 1.

link: [1, 2, 1, 2, 1, 2, 1, 7, 7, 9]

rank: [2, 2, 0, 0, 0, 0, 2, 0, 1, 0]