

# Chapter 2: Crypto Basics (密码基础)

MXDXBVTZWVMXNSPBQXLIMSCCSGXSCJXBOVQXCJZMOJZCVC  
TVWJCZAAXZBCSSCJXBQCJZCOJZCNSPOXBXSBTWVWJC  
JZDXGXXMOZQMSCSCJXBOVQXCJZMOJZCNSPJZHGXXMOSPLH  
JZDXZAAXZBXHCSCJXTCSGXSCJXBOVQX

— ciphertext

课后习题第10题

# Crypto

- **Cryptology** (密码学) — The art and science of making and breaking "secret codes"
- **Cryptography** (密码编码) — making "secret codes"
- **Cryptanalysis** (密码分析) — breaking "secret codes"
- **Crypto** — all of the above (and more)

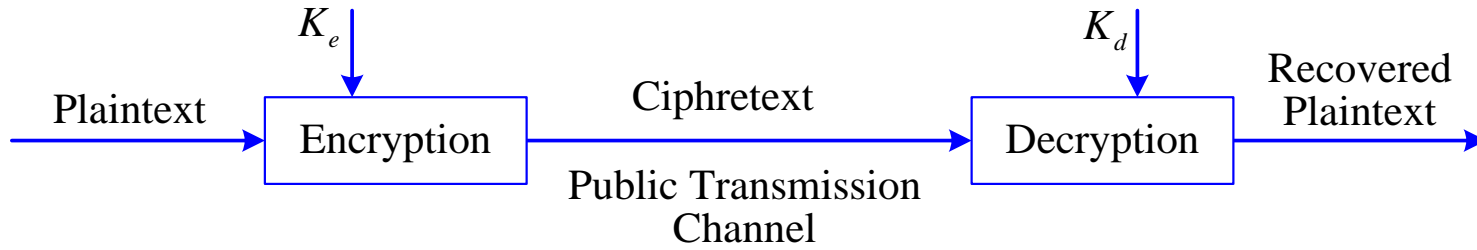
# How to Speak Crypto

- ❑ A *cipher* or *cryptosystem* is used to *encrypt* the *plaintext*
- ❑ The result of encryption is *ciphertext*
- ❑ We *decrypt* ciphertext to recover plaintext
- ❑ A *key* is used to configure a cryptosystem
- ❑ A *symmetric key* cryptosystem uses the same key to encrypt as to decrypt
- ❑ A *public key* cryptosystem uses a *public key* to encrypt and a *private key* to decrypt

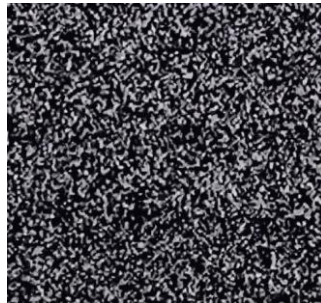
# Crypto

- ❑ Basic assumptions
  - The system is completely known to the attacker
  - Only the key is secret
  - That is, crypto algorithms are not secret
- ❑ This is known as **Kerckhoffs' Principle**
- ❑ Why do we make such an assumption?
  - Experience has shown that secret algorithms tend to be weak when exposed
  - Secret algorithms never remain secret
  - Better to find weaknesses beforehand

# 密码的一般性模型



0100100101101  
1100110011001  
10



Cryptology:

**Cryptography** studies how to design good encryption algorithms

**Cryptanalysis** tries to find security weaknesses of proposed algorithms and studies whether they are vulnerable to some attacks

# Kerckhoffs's Principle

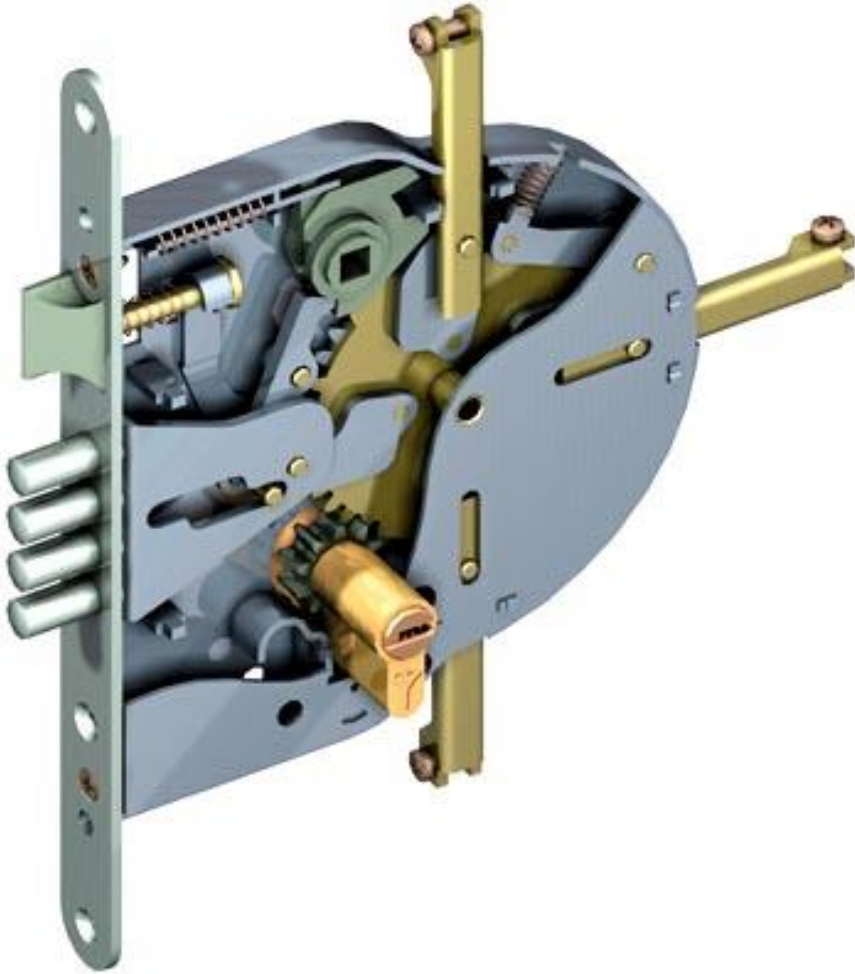


A cryptosystem should be secure even if everything about the system, except the **key**, is public knowledge.

**1883** La Cryptographie Militaire ("Military Cryptography"),  
le Journal des Sciences Militaires ("Journal of Military Science")

Dr. Jean Guillaume **Auguste** Victor François Hubert **Kerckhoffs**  
(1835-1903), Dutch (Hollander)

## Kerckhoffs's Principle (cont'd)



1. Key is reused
2. Structure is public

# Claude Shannon's maxim



Claude Shannon's maxim:  
"The enemy knows the system."

Claude Elwood Shannon (1916-2001), father of information theory and cryptography

Shannon Claude

**Communication Theory of Secrecy Systems**

*Bell System Technical Journal*, vol. 28 no. 4, pp. 656-715, 1949

Cited **13071** times in Scholar.Google

[www.bilibili.com/video/BV1YV411z7qo/?spm\\_id\\_from=333.337.search-card.all.click%20%0Bhttps://thebitplayer.com](https://www.bilibili.com/video/BV1YV411z7qo/?spm_id_from=333.337.search-card.all.click%20%0Bhttps://thebitplayer.com)

[Who meets Shannon or Shannon meets whom](#)



# Simple Substitution(简单替换)

□ Plaintext: **fourscoreandsevenyearsago**

□ Key:

Plaintext	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
Ciphertext	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C

□ Ciphertext:

**IRXUVFRUHDQGVHYHQBHDUVDJR**

□ Shift by 3 is "Caesar's cipher"

# Caesar's Cipher Decryption

- Suppose we know a Caesar's cipher is being used:

Plaintext	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
Ciphertext	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C

- Given ciphertext:

VSRQJHEREVTXDUHSDQWV

- Plaintext: spongebobsquarepants

# Caesar's Cipher

- Shift by  $n$  for some  $n \in \{0, 1, 2, \dots, 25\}$
- Then key is  $n$
- Example: key  $n = 7$

Plaintext

a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G

Ciphertext

# Cryptanalysis I: Try Them All

- ❑ We know Caesar's cipher and shift by  $n$  used
  - But the specific key is unknown
- ❑ Given ciphertext: **CSYEVIXIVQMREXIH**
- ❑ How to determine the key?
- ❑ Only 26 possible keys — try them all!
- ❑ **Exhaustive key search**(穷举搜索密钥)
- ❑ Solution: key is  $n = 4$

# Simple Substitution: General Case

- In general, simple substitution key can be any **permutation** of letters (**字母的任意置换**)
  - Not necessarily a Caesar's cipher (shift)
- For example

Plaintext	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
Ciphertext	J	I	C	A	X	S	E	Y	V	D	K	W	B	Q	T	Z	R	H	F	M	P	N	U	L	G	O

- In general,  $26! > 2^{88}$  possible keys

意味着Trudy用普通的穷举式密钥检索以每秒执行 $2^{40}$ 次（千亿次级别）密钥计算的超级计算机，约需花费890万年才能尝试完所有可能的密钥，平均工作量也是445万年。

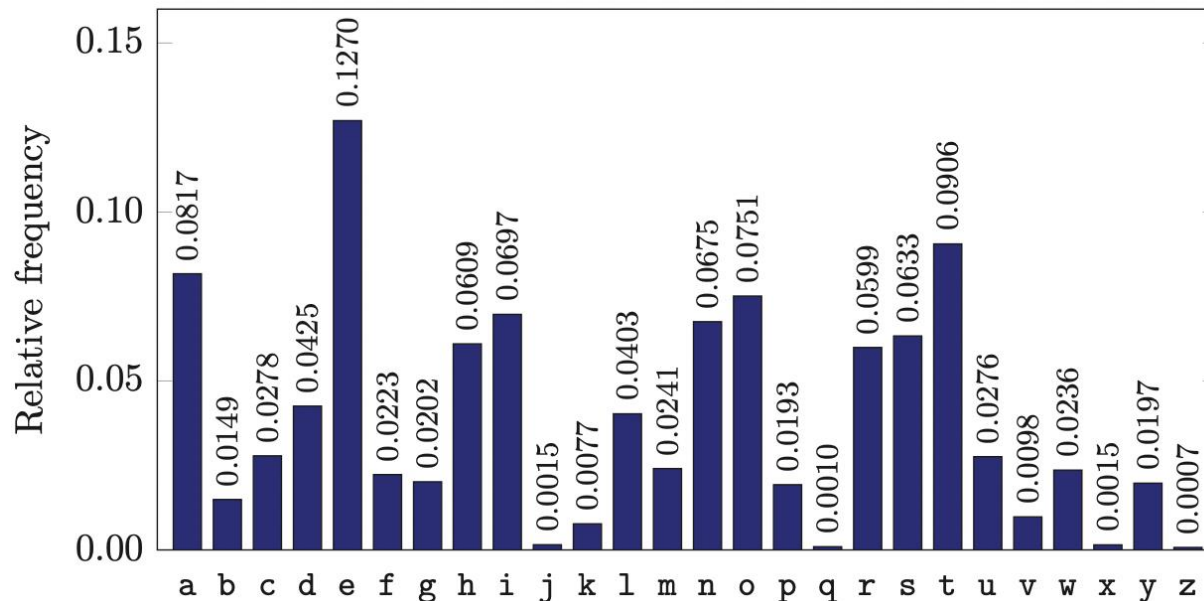
# Cryptanalysis II: Be Clever

- ❑ We know that a simple substitution used
- ❑ But *not* necessarily a Caesar's cipher (shift)
- ❑ Find the key given the ciphertext:

PBFPVYFBQXZTYFPBFEQJHDXXQVAPTPQJKTOYQWIPBVWLXTOX  
BTFXQWAXBVCXQWAXFQJWVLEQNTQZQGGQLFXQWAKVWLXQ  
WAEBIPBFXFQVXGTVJVWLBTPQWAEBFPBFHCVLXBQUFEVWLXGD  
PEQVPQGVPPBFTIXPFHXZHVFAGFOTHFEFBQUFTDHzBQPOTHXTY  
FTODXQHFTDPTOGHFQPBQWAQJJTODXQHFOQPWTBDHHIXQV  
APBFZQHCFWPFHPBFIPBQWKFABVYYDZBOTHBPBPQJTQOTOGHF  
QAPBFEQJHDXXQVAVXEBQPEFZBVFOJIWFFACFCCFHQWAUVWF  
LQHGFVAFXQHFUFHILTTAVWAFFAWTEVOITDHFHFQAITIXPFH  
XAFQHEFZQWGFLVWPTOFFA

# Cryptanalysis II

- ❑ Cannot try all  $2^{88}$  simple substitution keys
- ❑ Can we be more clever?
- ❑ English letter frequency counts... (英语字母的使用频率统计)



# Cryptanalysis II

## □ Ciphertext:

PBFPVYFBQXZTYFPBFEQJHDXXQVAPTPQJKTOYQWIPBVWLXTOXBTFXQ  
WAXBVCXQWAXFQJ VWLEQNTQZQGGQLFXQWAKVWLXQWAE BIPBFXFQ  
VXGTVJVWLBTPQWAE BFPBFHCVLXBQUFEVWLXGDPEQVPQGVPBPFTIXPFH  
XZHVFA GFOTHFEBQUFTD HZBQPOTHXTYFTODXQHFTDPTOGHFQPBQW  
AQJJTODXQHFOQPWTBDHHIXQVAPBFZQHCFWPFHPBFIPBQWK FABVYY  
DZBOTHBPQPJTQOTOGHFQAPBFEQJHDXXQVAVXEBQPEFZBVFOJIWFF  
ACFCCFHQWAUVWFLQHGF XVAFXQH FUFHILT TAVWAFFAWTEVOITDHFH  
FQAITIXPFH XAFQHEFZQWGFLVWPTOFFA

## □ Analyze this message using statistics below

Ciphertext frequency counts:

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
21	26	6	10	12	51	10	25	10	9	3	10	0	1	15	28	42	0	0	27	4	24	22	28	6	8



# Cryptanalysis II

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
21	26	6	10	12	51	10	25	10	9	3	10	0	1	15	28	42	0	0	27	4	24	22	28	6	8

$F \rightarrow e, Q \text{ 或 } P \text{ 或 } X \rightarrow t$

密文头几个字母PBFP...

1. 把F替换成e之后，PBeP....
2. 猜测P很可能是t
3. P替换成t之后，tBet....
4. 可以猜测B为h，因为the为一个常见的单词
5. 把密文中所有出现P, B, F的地方全部替换成明文再继续用类似的方法分析.....

(课后习题的第9题)

# Cryptanalysis: Terminology

- ❑ Cryptosystem is **secure** if best know attack is to try all keys
  - Exhaustive key search, that is
- ❑ Cryptosystem is **insecure** if *any* shortcut attack is known
- ❑ But then insecure cipher might be harder to break than a secure cipher!
  - What the ... ?

# Mono-Alphabetic Substitution Cipher

A mono-alphabetic cipher (simple substitution cipher) is a substitution cipher where each letter of the plain text is replaced with another letter of the alphabet. It uses a fixed key which consist of the 26 letters of a "shuffled alphabet".

You can generate your own encryption keys and encrypt your own messages using our online mono-alphabetic substitution engine:

### Mono-alphabetic Substitution Cipher

Plain text:

A long time ago, in a galaxy far, far away... It is a dark time for the Rebellion. Although the Death Star has been destroyed, Imperial troops have driven the Rebel forces from their hidden base and pursued them across the galaxy. Evading the dreaded Imperial Starfleet, a group of freedom fighters led by Luke Skywalker has established a new secret base on the remote ice world of Hoth. The evil lord Darth Vader, obsessed with finding young Skywalker, has dispatched thousands of remote probes into the far reaches of space...

1. Generate Key

Key: YOTDWSKCNZEPHGQARUMJILFXBV

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
Y	O	T	D	W	S	K	C	N	Z	E	P	H	G	Q	A	R	U	M	J	I	L	F	X	B	V

2. Start Substitution

Cipher text:

Y PQGK JNHW YKQ, NG Y KYPYXB SYU, SYU YFYB... NJ NM Y DYUE JNHW SQU JCW UWOWPENQG. YPJQCQIKC JCW DWYJC MJYU CYM OWWG DWMJUQBWD, NHAWUNYP JUQQAM CYLW DUNLWG JCW UWOWP SQUTWM SUQH JCWNU CNDDWG OYMW YGD AIUMIWD JCWH YTUQMM JCW KYPYXB. WLYDNGK JCW DUWYDWD NHAWUNYP MJYUSPWWJ, Y KUQIA QS SUWWDQH SNKCJWUM PWD OB PIEW MEBFYPEWU CYM WMJYOPNMCWD Y GWF MWTUWJ OYMW QG JCW UWHQJW NTW FQUPD QS CQJC. JCW WLNK PQUD DYUJC LYDWU, QOMMMWD FNJC SNGDNGK BQIGK MEBFYPEWU, CYM DNMAJYTCWD JCQIMYGDM QS UWHQJW AUQOWM NGJQ JCW SYU UWYTCWM QS MAYTW...

# Vigenere Cipher

- Simple substitution is **monoalphabetic** (单字母替换)
- Vigenere cipher is simple example of a **polyalphabetic** substitution (多字母替换)
  - Caesars ciphers, based on a keyword **D** (shift by 3)
- For example, keyword **CAT** indicates shift by **2**, shift by **0**, shift by **19**
  - Then repeat as needed

# Vigenere Example

- ❑ Suppose that we want to encrypt attackatdawn

- ❑ Encryption:      keyword:    CATCATCATCAT  
                          plaintext:    attackatdawn  
                          ciphertext:    ctmccdctwcwg

- ❑ Ciphertext is ctmccdctwcwg

- ❑ How to decrypt? How to attack?

To attack a Vigenere ciphertext, we first need to **determine keyword length.**

Crack Vigenere Ciphertext using statistic method

<https://github.com/DavidLee528/VigenereCracker>

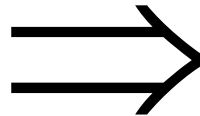
Part 1 — Cryptography

# Double Transposition

□ Plaintext: **attackxatxdawn**

	col 1	col 2	col 3
row 1	a	t	t
row 2	a	c	k
row 3	x	a	t
row 4	x	d	a
row 5	w	n	x

Permute rows  
and columns



	col 1	col 3	col 2
row 3	x	t	a
row 5	w	x	n
row 1	a	t	t
row 4	x	a	d
row 2	a	k	c

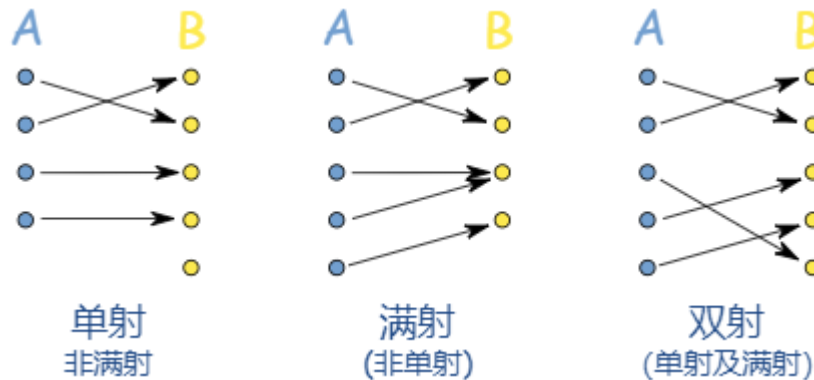
□ Ciphertext: **xtawxnattxadakc**

□ Key is matrix size (3, 5) and  
permutations: (1,4,2) and (1,3,2)

## Transposition Cipher(置换密码)

(置换) 有限集 $X$ 上的运算 $\sigma: X \rightarrow X$ ,  $\sigma$ 是一个双射函数。

➤  $\sigma$ 既是单射又是满射, 并且定义域和值域相同, 那么称 $\sigma$ 为一个置换。

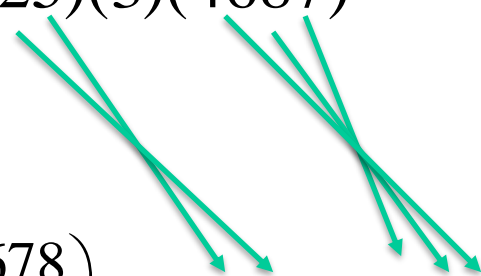


➤ 若 $\sigma$ 是一个置换,  $\forall x \in X$ , 存在唯一的 $x$ 使得 $\sigma(x) = x'$ 。

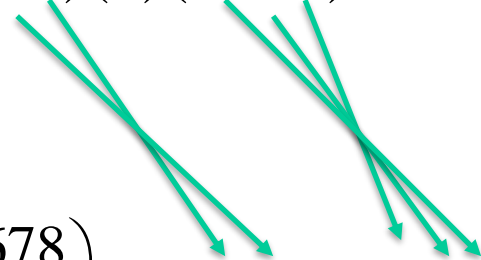
➤ 同理可以定义逆置换 $\sigma^{-1}: X \rightarrow X$ ,  $\sigma^{-1}$ 也是双射函数, 并且 $\sigma^{-1}$ 的定义域和值域相同, 即 $\forall x' \in X$ , 存在唯一的 $x \in X$ 使得 $\sigma^{-1}(x') = x$ 。

➤  $\sigma\sigma^{-1} = I$  (单位变换)

设有限集 $X=\{1,2,3,4,5,6,7,8\}$ ,  $\sigma$ 为 $X$ 上的一个置换, 并且满足 $\sigma(1)=2, \sigma(2)=5, \sigma(3)=3, \sigma(4)=6, \sigma(5)=1, \sigma(6)=8, \sigma(7)=4, \sigma(8)=7$ . 因为置换可以简单用对换表示, 所以上述置换 $\sigma$ 可以形式化表示为对换的乘积, 即

$$\sigma = \begin{pmatrix} 12345678 \\ 25361847 \end{pmatrix} = (125)(3)(4687)$$


则其逆置换 $\sigma^{-1}$ 可以表示为

$$\sigma^{-1} = \begin{pmatrix} 12345678 \\ 25361847 \end{pmatrix}^{-1} = \begin{pmatrix} 12345678 \\ 51372486 \end{pmatrix} = (152)(3)(4786)$$




# One-Time Pad: Encryption

e=000 h=001 i=010 k=011 l=100 r=101 s=110 t=111

Encryption: Plaintext  $\oplus$  Key = Ciphertext

	h	e	i	l	h	i	t	l	e	r
Plaintext:	001	000	010	100	001	010	111	100	000	101
Key:	111	101	110	101	111	100	000	101	110	000
Ciphertext:	110	101	100	001	110	110	111	001	110	101
	s	r	l	h	s	s	t	h	s	r

# One-Time Pad: Decryption

e=000 h=001 i=010 k=011 l=100 r=101 s=110 t=111

**Decryption:** Ciphertext  $\oplus$  Key = Plaintext

	s	r	l	h	s	s	t	h	s	r
Ciphertext:	110	101	100	001	110	110	111	001	110	101
Key:	111	101	110	101	111	100	000	101	110	000
Plaintext:	001	000	010	100	001	010	111	100	000	101
	h	e	i	l	h	i	t	l	e	r

# One-Time Pad

Double agent claims following “**key**” was used:

	s	r	l	h	s	s	t	h	s	r
Ciphertext:	110	101	100	001	110	110	111	001	110	101
“ <b>key</b> ”:	101	111	000	101	111	100	000	101	110	000
“Plaintext”:	011	010	100	100	001	010	111	100	000	101
	k	i	l	l	h	i	t	l	e	r

e=000 h=001 i=010 k=011 l=100 r=101 s=110 t=111

# One-Time Pad

Or, might claim the key is...

	s	r	l	h	s	s	t	h	s	r
Ciphertext:	110	101	100	001	110	110	111	001	110	101
"key":	111	101	000	011	101	110	001	011	101	101
<hr/>										
"Plaintext":	001	000	100	010	011	000	110	010	011	000
	h	e	l	i	k	e	s	i	k	e

e=000 h=001 i=010 k=011 l=100 r=101 s=110 t=111

# One-Time Pad Summary

- ❑ **Provably** secure
  - Ciphertext gives **no** useful info about plaintext
  - All plaintexts are *equally likely*
- ❑ BUT, only when be used correctly
  - Pad must be random, used only once
  - Pad is known only to sender and receiver
- ❑ Note: pad (key) is same size as message
- ❑ So, why not distribute message itself, instead of the pad?

# Real-World One-Time Pad

- ❑ Project VENONA
  - Soviet spies encrypted messages from U.S. to Moscow in 30's, 40's, and 50's
  - Nuclear espionage, etc.
  - Thousands of messages
- ❑ Spy carried one-time pad into U.S.
- ❑ Spy used pad to encrypt secret messages
- ❑ Repeats within the "one-time" pads made cryptanalysis possible

# Venona project

The Venona project was a United States counterintelligence program initiated during World War II by the United States Army's Signal Intelligence Service and later absorbed by the National Security Agency (NSA), that ran from February 1, 1943, until October 1, 1980. It was intended to decrypt messages transmitted by the intelligence agencies of the Soviet Union (e.g. the KGB).

During the 37-year duration of the Venona project, the Signal Intelligence Service decrypted and translated approximately 3,000 messages. The signals intelligence yield included discovery of the Cambridge Five espionage ring in the United Kingdom and Soviet espionage of the Manhattan Project in the US.

Some of the espionage was undertaken to support the Soviet atomic bomb project. The Venona project remained secret for more than 15 years after it concluded. Some of the decoded Soviet messages were not declassified and published by the United States until 1995.

[https://en.wikipedia.org/wiki/Venona\\_project](https://en.wikipedia.org/wiki/Venona_project)

# VENONA Decrypt (1944)

From: NEW YORK

To: MOSCOW

No: 1340

21 September 1944

To VIKTOR[i].

Lately the development of new people [D% has been in progress]. LIBERAL[ii] recommended the wife of his wife's brother, Ruth GREENGLASS, with a safe flat in view. She is 21 years old, a TOWNSWOMAN [GOROZhANKA][iii], a GYMNAST [FIZKUL'TURNITSa][iv] since 1942. She lives on STANTON [STANTAUN] Street. LIBERAL and his wife recommend her as an intelligent and clever girl.

[15 groups unrecoverable]

[C% Ruth] learned that her husband[v] was called up by the army but he was not sent to the front. He is a mechanical engineer and is now working at the ENORMOUS [ENORMOZ][vi] plant in SANTA FE, New Mexico.

[45 groups unrecoverable]

detain VOLOK[vii] who is working in a plant on ENORMOUS. He is a FELLOWCOUNTRYMAN [ZEMLYaK][viii]. Yesterday he learned that they had dismissed him from his work. His active work in progressive organizations in the past was the cause of his dismissal.

In the FELLOWCOUNTRYMAN line LIBERAL is in touch with CHESTER[ix]. They meet once a month for the payment of dues. CHESTER is interested in whether we are satisfied with the collaboration and whether there are not any misunderstandings. He does not inquire about specific items of work [KONKRETNAYA RABOTA]. In as much as CHESTER knows about the role of LIBERAL's group we beg consent to ask C. through LIBERAL about leads from among people who are working on ENORMOUS and in other technical fields.

- ❑ "Ruth" == Ruth Greenglass
- ❑ "Liberal" == Julius Rosenberg
- ❑ "Enormous" == the atomic bomb



# Codebook Cipher

- ❑ Literally, a book filled with “codewords”
- ❑ Zimmerman Telegram encrypted via codebook

Februar	13605
fest	13732
finanzielle	13850
folgender	13918
Frieden	17142
Friedenschluss	17149
:	:

- ❑ Modern block ciphers are codebooks!
- ❑ More about this later...

# Codebook Cipher: Additive

- ❑ Codebooks also (usually) use **additive**
- ❑ Additive — book of “random” numbers
  - Encrypt message with codebook
  - Then choose position in additive book
  - Add additive sequence to get ciphertext
  - Send ciphertext and additive position (MI)
  - Recipient subtracts additives before decrypting
- ❑ Why use an additive sequence?

# Zimmermann Telegram

- Perhaps most famous codebook ciphertext ever
- A major factor in U.S. entry into World War I

**CLASS OF SERVICE DESIRED**

Post Day Message	<input checked="" type="checkbox"/>
Day Letter	<input type="checkbox"/>
Night Message	<input type="checkbox"/>
Night Letter	<input type="checkbox"/>

Persons should mark as it appears on back hereof, which are hereby agreed to

**WESTERN UNION TELEGRAM**

NEW YORK, CARLTON, PRESIDENT

**CHARGE** *5387*

**via Galveston**

**GERMAN LEGATION MEXICO CITY**

**JAN 19 1917**

130	13042	13401	8501	115	3528	416	17214	8491	11310
18147	18222	21560	10247	11518	23677	13605	3494	14936	
98092	5905	11311	10392	10371	0302	21290	5161	39695	
23571	17504	11269	18276	18101	0317	0228	17694	4473	
22284	22200	19452	21589	87893	5569	13918	8958	12137	
1333	4725	4458	5905	17166	13851	4458	17149	14471	6708
13850	12224	6929	14991	7382	15857	67893	14218	36477	
5870	17553	67893	5870	5454	16102	15217	22801	17138	
21001	17388	7446	23638	18222	6719	14331	15021	23845	
3156	23552	22096	21604	4797	9497	22464	20855	4377	
23610	18140	22260	5905	13347	20420	39689	13732	20667	
6929	5275	18507	52262	1340	22049	13339	11265	22295	
10439	14814	4178	6992	8784	7032	7357	6926	52262	11267
21100	21272	9346	9559	22464	15874	18502	18500	15857	
2188	5376	7381	98092	16127	13486	9350	9220	76036	14219
5144	2831	17920	11347	17142	11264	7667	7762	15099	9110
10482	97556	3569	3670						

**BEPNSTORFF.**

**Charge German Embassy.**

# Zimmermann Telegram Decrypted

- ❑ British had recovered partial codebook
- ❑ Then able to fill in missing parts

MAILED  
October 1-8-58  
Washington, State Dept.  
By *Wm. A. Eckhoff*  
Date *Oct. 27, 1918*

TELEGRAM RECEIVED.

FROM 2nd from London # 5747.

"We intend to begin on the first of February unrestricted submarine warfare. We shall endeavor in spite of this to keep the United States of America neutral. In the event of this not succeeding, we make Mexico a proposal of alliance on the following basis: make war together, make peace together, generous financial support and an understanding on our part that Mexico is to reconquer the lost territory in Texas, New Mexico, and Arizona. The settlement in detail is left to you. You will inform the President of the above most secretly as soon as the outbreak of war with the United States of America is certain and add the suggestion that he should, on his own initiative, ~~invite~~ <sup>invite</sup> Japan to immediate adherence and at the same time mediate between Japan and ourselves. Please call the President's attention to the fact that the ruthless employment of our submarines now offers the prospect of compelling England in a few months to make peace." Signed, ZIMMERMANN.

# Random Historical Items(历史上密码事件)

- Crypto timeline
- Spartan Scytale — transposition cipher (大约公元前5世纪内, 置换密码)
- Caesar's cipher (公元前1世纪内, 替换密码)
- Edgar Allan Poe's short story: *The Gold Bug (1843)* (对字母的频率分析)
- Election of 1876 (置换与替换的结合)

# Election of 1876

- ❑ Rutherford Hayes vs Swindling Tilden
- ❑ Electoral college delegations for 4 states (including Florida) in dispute (选举团在4个州存在分歧)
- ❑ Commission gave all 4 states to Hayes
  - Vote on straight party lines
- ❑ Tilden accused Hayes of bribery
  - Was it true? (全部投给共和党被控告贿赂?)



# Election of 1876

- ❑ Encrypted messages by Tilden supporters later emerged
- ❑ Cipher: Partial codebook, plus transposition
- ❑ Codebook substitution for important words

## ciphertext

Copenhagen

Greece

Rochester

Russia

**Warsaw**

:

## plaintext

Greenbacks

Hayes

votes

Tilden

**telegram**

:

# Election of 1876

- ❑ Apply codebook to original message
- ❑ Pad message to multiple of 5 words (total length, 10,15,20,25 or 30 words)
- ❑ For each length, a fixed permutation applied to resulting message
- ❑ Permutations found by comparing several messages of same length
- ❑ Note that the **same key** is applied to all messages of a given length



# Election of 1876

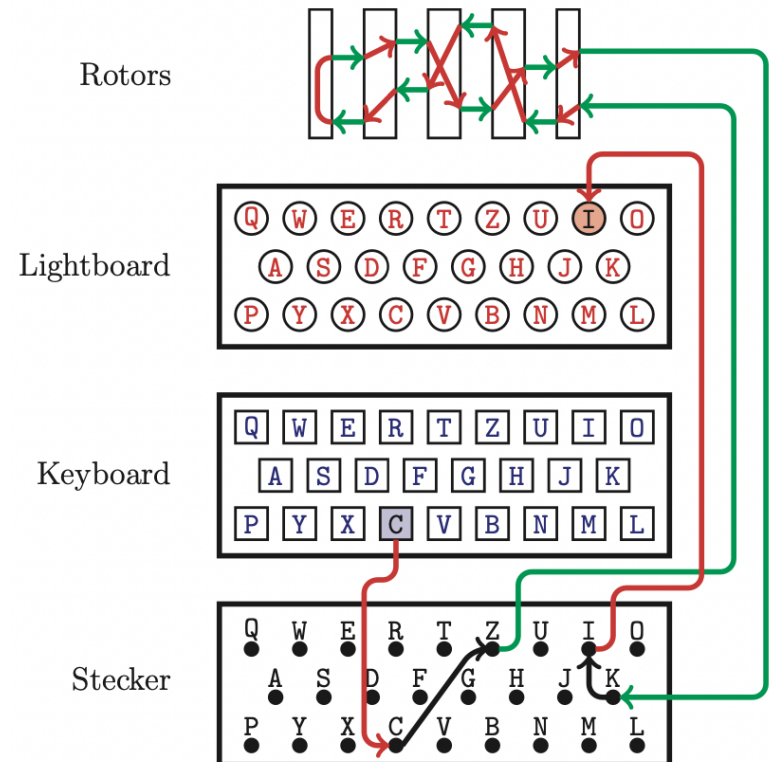
- ❑ Ciphertext: **Warsaw they read all unchanged last are idiots can't situation**
- ❑ Codebook: Warsaw == telegram
- ❑ Transposition: 9,3,6,1,10,5,2,7,4,8
- ❑ Plaintext: **Can't read last telegram. Situation unchanged. They are all idiots.**
- ❑ A weak cipher made worse by reuse of key
- ❑ Lesson? Don't overuse keys!

# Early 20th Century

- ❑ WWI — Zimmermann Telegram
- ❑ “Gentlemen do not read each other's mail”
  - Henry L. Stimson, Secretary of State, 1929
- ❑ WWII — **golden** age of cryptanalysis
  - Midway/Coral Sea
  - Japanese **Purple** (codename **MAGIC**)
  - German **Enigma** (codename **ULTRA**)

# Enigma Cipher Machine

- ❑ Most famous cipher of WWII
  - Electro-mechanical device
  - Very rugged, used in the field



# Post-WWII History

- ❑ Claude Shannon — father of the science of information theory
- ❑ Computer revolution — lots of data to protect
- ❑ Data Encryption Standard (DES), 70's
- ❑ Public Key cryptography, 70's
- ❑ CRYPTO conferences, 80's
- ❑ Advanced Encryption Standard (AES), 90's
- ❑ The crypto genie is out of the bottle...

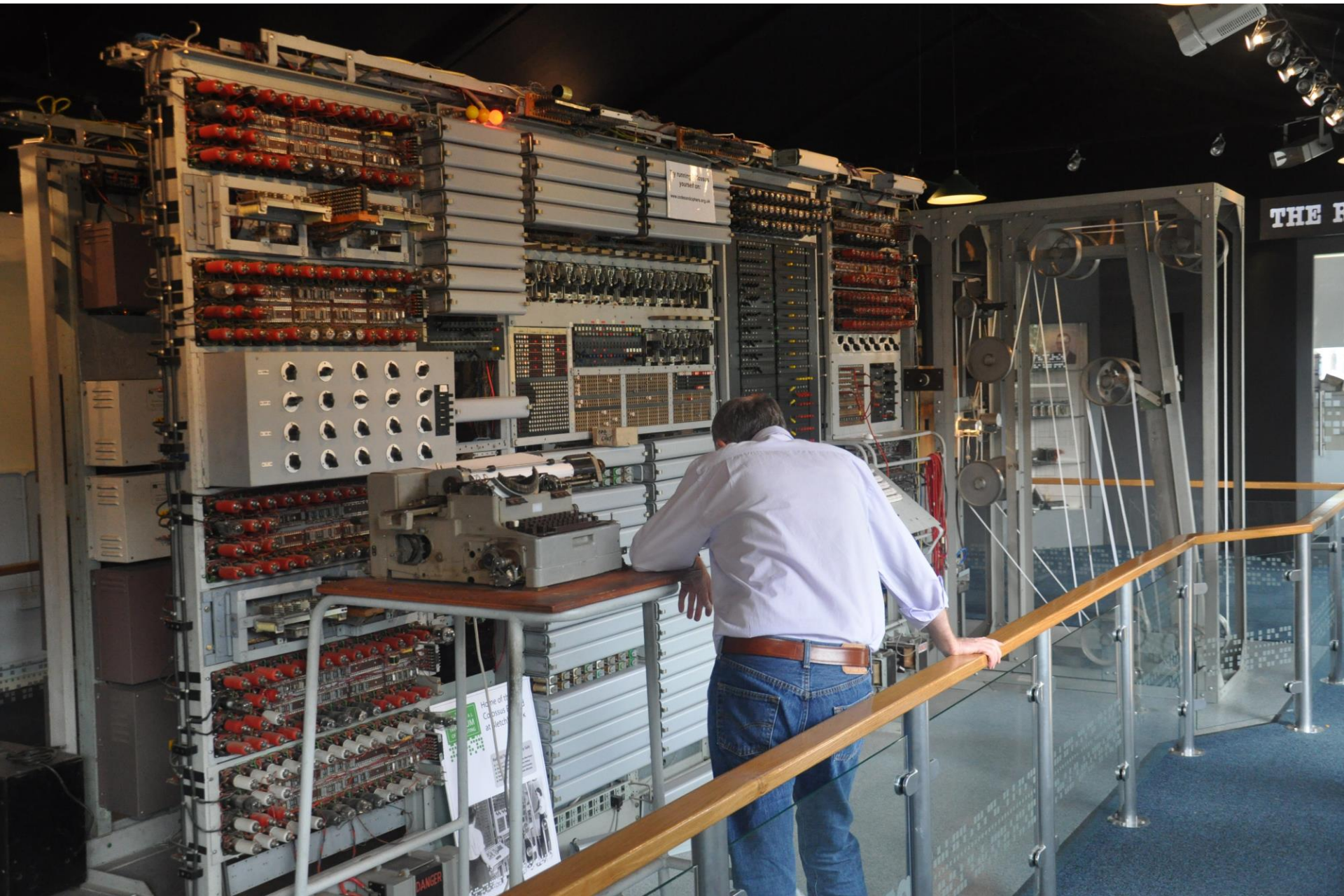


The bombe was an electro-mechanical device used by British cryptologists to help decipher German Enigma-machine-encrypted secret messages during World War II.

<https://en.wikipedia.org/wiki/Bombe>

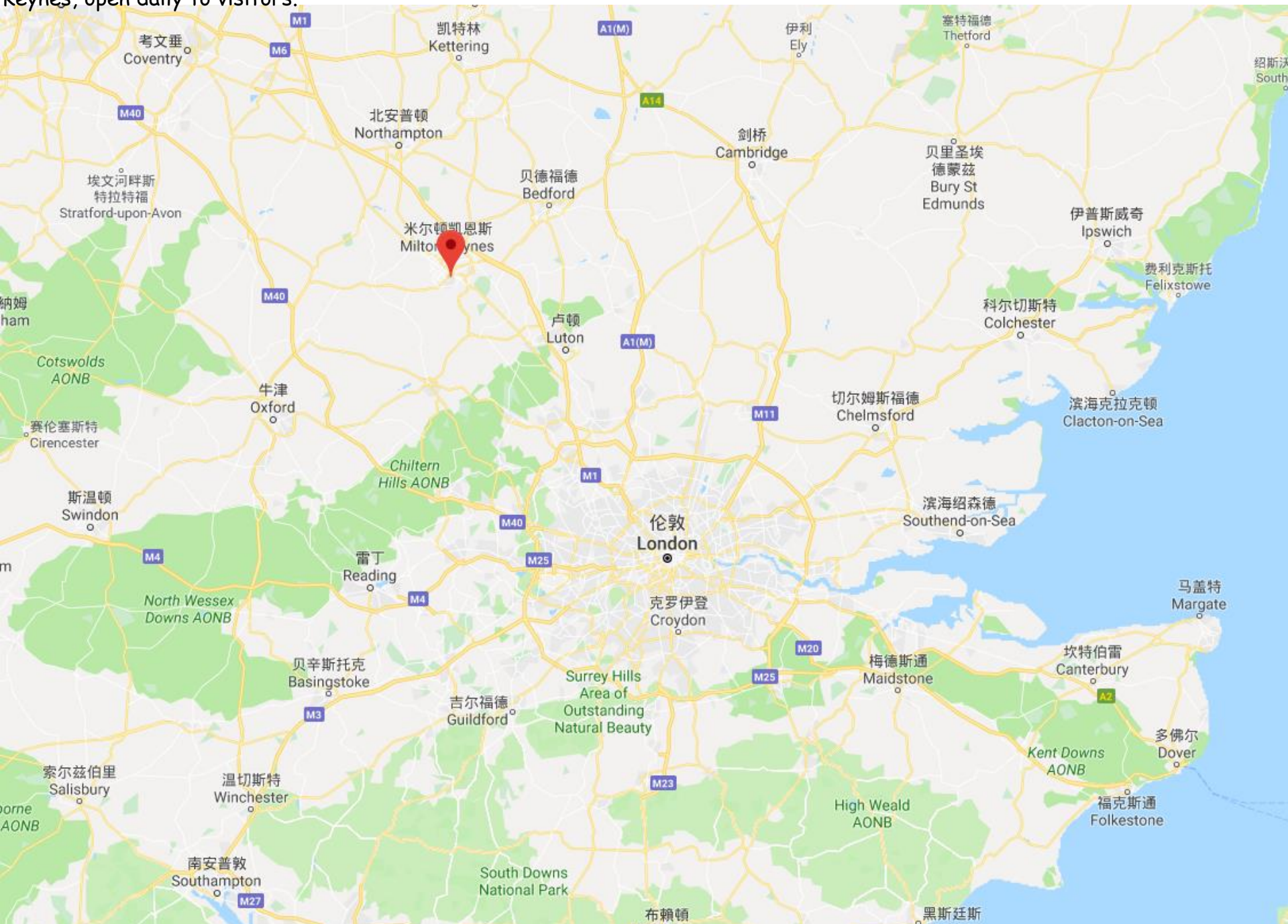








Bletchley Park, once the top-secret home of the World War Two Codebreakers, is now a vibrant heritage attraction in Milton Keynes, open daily to visitors.



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# Claude Shannon

- ❑ Founded field of information theory
- ❑ His 1949 paper: *Communication Theory of Secrecy Systems*
- ❑ Fundamental concepts
  - **Confusion** — obscure relationship between plaintext and ciphertext
  - **Diffusion** — spread plaintext statistics through the ciphertext
- ❑ Proved one-time pad is secure
- ❑ One-time pad is confusion-only, while double transposition is diffusion-only

# Taxonomy of Cryptography

## □ Symmetric Key

- Same key for encryption and decryption
- Modern types: Stream ciphers, Block ciphers

## □ Public Key (or “asymmetric” crypto)

- Two keys, one for encryption (public), and one for decryption (private)
- And digital signatures — nothing comparable in symmetric key crypto

## □ Hash algorithms

- Can be viewed as “one way” crypto

# Taxonomy of Cryptanalysis(密码分析的分类)

- From perspective of info available to Trudy
  - Ciphertext only( Trudy只知道密文)
  - Known plaintext(Trudy知道部分明文-密文对)
  - Chosen plaintext(Trudy可以选择部分明文来加密)
    - “Lunch time attack”(Alice吃午饭忘退出电脑, Trudy在她回来之前可以加密部分选定的消息)
    - Protocols might encrypt chosen data
  - Adaptively chosen plaintext(Trudy根据观测到的密文选择下一个明文)
  - Forward search (只适用于公钥密码, 用公钥加密假定的明文去与截获的密文匹配)
  - And others...

# 作业题

- 第2章的29个题目，每人所做题目为学号的后两位%29！学号末尾为29的同学做第29题。