Multi-threaded word count

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Free days remaining: 5 / Used: 0

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Introduction

The last assignment, we dealt with process management(fork(), execve()). In this assignment, we are going to implement a N:M producer and consumer multi-threading program. The goal of the program is to read a text file with a series of producer threads, and write each separated line into a shared object. Then with consumer threads, generate statistics of the read input file.

Concepts

Multi-threading is a technique that allows a single program to execute multiple threads concurrently. A thread is the smallest unit of a process. Unlike processes, which run in separate memory spaces, threads share the same memory space. But, to be more specific, there are types of memory regions (heap, static) that can be shared across multiple threads. This enables sharing data across threads more efficiently compared to that of processes. However, it is still required to coordinate (synchronize) each processes/threads for data consistency as well as avoiding deadlocks.

For implementation, we would be using the library pthread for the Linux environment. To be more specific, we will be using pthread_mutex as well as pthread_cond for sending conditions across multiple threads.

Implementation

Goal 1) Correct the given code to implement 1:1

First of all, I had to fix the primitive *prod_cons.c* code. When I executed the program, I used the LICENSE file as an input and setted up to create with 1 producer and 1 consumer. However, it didn't work because the program couldn't handle the input file properly.

```
choyoonseo@joyunseoui-MacBookAir 2024-os-hw2 % ./a.out ./LICENSE 1 1
main continuing
Cons: 0 lines
main: consumer_0 joined with 0
Prod_6aef7000: 21 lines
main: producer_0 joined with 21
```

<Figure 1 - Failure of 1 producer 1 consumer>

Due to a race condition between the single producer and consumer, I implemented a mutex lock to avoid it. In other words, I synchronized both threads.

Before entering the critical section, I used a mutex lock to guarantee mutual exclusion. pthread_mutex_lock() prevents other threads from entering the critical section. After locking the mutex, it is important to unlock it. pthread_mutex_unlock() allows other threads to enter the critical section. Without unlocking, other threads never get their turn, which can lead to deadlock.

Moreover, I used a condition variable because if I had only used a mutex, it would have been inevitable to experience busy waiting, which would waste CPU resources.

pthread_cond_wait() makes the thread remain blocked until another thread signals the condition variable using pthread_cond_signal() or pthread_cond_broadcast().

pthread_cond_signal() sends a signal which wakes up the waiting thread so it can continue processing.

It is important to initialize them before using those mutex and condition variables. I used pthread mutex init() and pthread cond init().

After using them, it is crucial to destroy them. pthread_mutex_destroy() and pthread_cond_destroy() inform the operating system that we've finished using the mutex or condition variable, and help clean up kernel resources.

```
choyoonseo@joyunseoui-MacBookAir 2024-os-hw2 % ./a.out ./LICENSE 1 1
main continuing
Cons_6b0c7000: [00:00] MIT License
Cons_6b0c7000: [01:01]
                                       Copyright (c) 2019 mobile-os-dku-cis-mse
Cons_6b0c7000:
                           [02:02]
Cons_6b0c7000:
                           [03:03]
Cons_6b0c7000:
                          [04:04] Permission is hereby granted, free of charge, to any person obtaining a copy [05:05] of this software and associated documentation files (the "Software"), to deal
                                                                                                                                                             '), to deal
Cons_6b0c7000:
                          [06:06] in the Software without restriction, including without limitation the rights [07:07] to use, copy, modify, merge, publish, distribute, sublicense, and/or sell [08:08] copies of the Software, and to permit persons to whom the Software is
Cons_6b0c7000:
Cons_6b0c7000:
Cons_6b0c7000:
Cons_6b0c7000:
                           [09:09]
                                         furnished to do so, subject to the following conditions:
Cons_6b0c7000:
                           [10:10]
                           [11:11]
[12:12]
[13:13]
Cons_6b0c7000:
                                        The above copyright notice and this permission notice shall be included in all
Cons_6b0c7000:
                                        copies or substantial portions of the Software.
Cons_6b0c7000:
Cons_6b0c7000:
                          [14:14] THE SOFTWARE IS PROVIDED "AS IS", WITHOUT WARRANTY OF ANY KIND, EXPRESS OR [15:15] IMPLIED, INCLUDING BUT NOT LIMITED TO THE WARRANTIES OF MERCHANTABILITY, [16:16] FITNESS FOR A PARTICULAR PURPOSE AND NONINFRINGEMENT. IN NO EVENT SHALL THE
Cons_6b0c7000:
Cons_6b0c7000:
Cons_6b0c7000: [17:17] AUTHORS OR COPYRIGHT HOLDERS BE LIABLE FOR ANY CLAIM, DAMAGES OR OTHER
Cons_6b0c7000: [18:18] LIABILITY, WHETHER IN AN ACTION OF CONTRACT, TORT OR OTHERWISE, ARISING FROM,
Cons_6b0c7000: [19:19] OUT OF OR IN CONNECTION WITH THE SOFTWARE OR THE USE OR OTHER DEALINGS IN THE
Cons_6b0c7000: [20:20] SOFTWARE.
Prod_6b03b000: 21 lines
Cons: 21 lines
main: consumer_0 joined with 21 main: producer_0 joined with 21
```

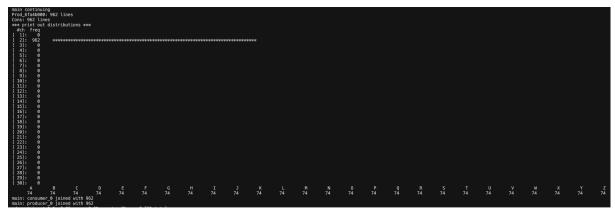
<Figure 2 - Success of 1 producer 1 consumer>

Finally, it works!

cons_6b037000 is the identifier of consumer thread, and Pros_6b03b000 is the identifier of producer thread. Cons_6b037000: [00:00] MIT License displays the work done by the consumer thread. As I have set up 1 consumer, this consumer reads every single line from the input file. Pros_6b03b000: 21 lines indicates that the producer thread has completed its task, which is reading from the input file. Cons: 21 lines shows that the consumer thread has finished its work. main: consumer_0 joined with 21 means that the main thread has completed synchronization with the consumer thread (consumer_0). main: producer_0 joined with 21 means that the main thread has completed synchronization with the producer thread (producer_0).

Goal 2) Add character statistics code

Now, let's count words! I added code to analyze and generate statistics for the string from *char_stat.c* to the consumer function.



<Figure 3 - Successful character statistics of 1 producer 1 consumer>

Yes it works!

To check the statistics, I created a simple input file that contains 74 instances of each letter of the alphabet.

Goal 3) Use FreeBSD9-orig.tar as an input

Until now, I checked that the program works well with a simple file like README.md or LICENSE. From now on, I'm going to use a 600MB size of FreeBSD9-orig.tar file as an input.

However, when I execute with the FreeBSD9-orig.tar file, the program stops due to the double free memory issue.

```
Cons_6d5b3000: [848553:848570] .\" 1. Redistributions of source code must retain the above copyright Cons_6d5b3000: [848554:848571] .\" notice, this list of conditions and the following disclaimer. Cons_6d5b3000: [848555:848572] .\" 2. Redistributions in binary form must reproduce the above copyright Cons_6d5b3000: [848556:848573] .\" notice, this list of conditions and the following disclaimer in the Cons_6d5b3000: [848557:848574] .\" documentation and/or other materials provided with the distribution. Cons_6d5b3000: [848558:848575] .\" Cons_6d5b3000: [848559:848576] .\" THIS SOFTWARE IS PROVIDED BY THE AUTHOR AND CONTRIBUTORS ``AS IS'' AND Cons_6d5b3000: [848560:848577] .\" ANY EXPRESS OR IMPLIED WARRANTIES, INCLUDING, BUT NOT LIMITED TO, THE Cons_6d5b3000: [848561:848578] .\" IMPLIED WARRANTIES OF MERCHANTABILITY AND FITNESS FOR A PARTICULAR PURPOSE Cons_6d5b3000: [848562:848579] .\" ARE DISCLAIMED. IN NO EVENT SHALL THE AUTHOR OR CONTRIBUTORS BE LIABLE Cons_6d5b3000: [848563:848580] .\" FOR ANY DIRECT, INDIRECT, INCIDENTAL, SPECIAL, EXEMPLARY, OR CONSEQUENTIAL Cons_6d5b3000: [848564:848581] .\" DAMAGES (INCLUDING, BUT NOT LIMITED TO, PROCUREMENT OF SUBSTITUTE GOODS a.out(38872,0x16d5b3000) malloc: **** set a breakpoint in malloc_error_break to debug Cons_6d5b3000: [848565:848581] zsh: abort ./a.out ./FreeBSD9-orig.tar 1 1
```

<Figure 4 - Double free issue with FreeBSD9-orig.tar in 1 producer 1 consumer>

It is due to the dangling pointer. To deal with this issue, I added so-> line = NULL after free (so->line) to prevent attempts to free memory that had already been freed. In other words, I set the pointer to NULL so that it no longer points to an invalid memory address.

```
| Comp. Septimes | Dissistant |
```

<Figure 5 - Deadlock issue with FreeBSD9-orig.tar in 1 producer 1 consumer>

Right after solving the double free issue, a deadlock issue arose.

The program stopped after joining with the consumer thread. It seems like the producer thread cannot join with the main thread. It keeps waiting.

I thought it was interrupted by some other issue temporarily. So I forced the program to quit, and re-executed it 5 times. But every time I executed it, the total number of lines processed by the consumer thread was different.

So I asked ChatGPT why it is so.

It was because of a spurious wakeup.

From a software perspective, the code should execute normally, but the thread that should not wake up can wake up unexpectedly due to characteristics of the operating system or hardware. So I replaced the if statement with a while loop when I check whether the buffer is full or not in producer and consumer functions.

```
Com_GF791698: [21267911;21267911] Com_GF791698: [21267912;21267912] # $FreeSD: release/9.8.9/bin/cat/Makefile 198148 2889-10-15 18:17:29Z ru $

Com_GF791698: [21267915;21267918]

Com_GF791698: [21267915]

Com_GF791698: [21267918]

Com_GF791698: [21267915]

Com_GF791698: [21267916]

Com_GF791698: [21267916]

Com_GF791698: [21267916]

Com_GF791698: [21267916]

Com_GF791698: [21267916]

Com_GF791698: [21267916]

Com_GF791698: [21267916
```

<Figure 6 - Success with FreeBSD9-orig.tar in 1 producer 1 consumer>

Finally it works!

Although it took more than 5 minutes to execute, it worked!

```
./a.out FreeBSD9-orig.tar 1 1 70.16s user 198.04s system 66% cpu 6:43.85 total
```

<Figure 7 - Execution time with FreeBSD9-orig.tar in 1 producer 1 consumer>

I used the time command to measure the execution time.

It took 70.16 seconds in user mode, which represents the actual program code execution time, and 198.04 seconds in system mode, which represents time spent on system calls like

file I/O and memory management. The wall clock time, which is the total actual time taken from program start to finish was 6 minutes 43.85 seconds with a CPU utilization of 66%.

./a.out ./FreeBSD9-orig.tar 1 1 37.56s user 81_78s system 100% cpu 1:58.48 total <Figure 8 - Reduced time with FreeBSD9-orig.tar in 1 producer 1 consumer>

It took so much time, so I changed the order of pthread_mutex_unlock() and pthread_cond_signal() to reduce context switching and decrease execution time.
Also, I commented out unnecessary printf() statements.

As a result, the total execution time was reduced to % of the original time.

Goal 4) Implement 1:N from 1:1

Let's enhance the program so that it can execute with multiple consumers.

When handling 1 producer 1 consumer, the producer only reads a single line from the input file, and the shared object(buffer) can only hold a single line. However, when handling multiple consumers, the producer should read multiple lines from the input file and put them into the buffer.

So, I changed several things in the shared object structure.

First, I changed a single pointer to a double pointer to store multiple strings. Second, I distinguished the condition variable like thread_cond_t has_data and pthread_cond_t has_space to better manage the buffer. Third, I used a done flag and a count variable instead of full, because the buffer contains multiple lines, not just a single line, and to clearly indicate whether the producer has finished reading the entire input file.

```
Cons_6f7fb000: [113:959] wW
Cons_6f4b3000: [98:959] vV
Cons_6f913000: [70:959] uU
Cons_6f5cb000: [107:959] tT
Cons_6f427000: [85:959] sS
Cons_6f427000: [86:959] rR
Cons_6f887000: [90:961] zZProd_6f39b000: 962 lines
Cons_6f657000: [83:961] yY
main: consumer_0 joined with 87
```

<Figure 9 - Deadlock issue with input file in 1 producer 10 consumer>

However, the deadlock issue occurred.

Some consumers might have been busy waiting because, before the producer finished, it only woke a single consumer using pthread cond signal().

So I changed it to pthread_cond_broadcast() to wake up every consumer thread. Since a single consumer changed to multiple consumers, the producer needs to take care of all the threads.

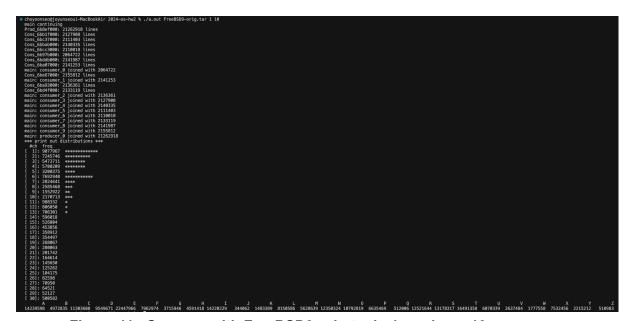
```
Cons_6b97f000:
Cons_6b7db000:
                         [109:954] sS
[101:961] zZProd_6b51f000: 962 lines
Cons_6b7db000:
Cons_6b5ab000:
                         [102:961] yY
[117:961] xX
Cons_6b867000:
Cons_6b867000:
Cons_6b6c3000:
Cons_6b74f000:
main: consumer_0
                            joined with 118
main: consumer_1
main: consumer_2
                            joined with 70
                            joined with 81
main: consumer_3
main: consumer_4
                            joined with 92
                            joined with
main: consumer_5
main: consumer_6
                            joined with 94
joined with 82
                           joined with 110
joined with 106
joined with 106
joined with 962
main: consumer_7
main: consumer_8
main: consumer_9
main: producer_0
./a.out ./input 1 10
                                   0.00s user 0.04s system 87% cpu 0.048 total
```

<Figure 10 - Race condition with input file in 1 producer 10 consumer>

Now, the multi-threaded execution works well, but another issue occurred.

The problem was that the character statistics code exists within the consumer function. Since multiple consumers are running, the character statistics should be calculated after all of the consumer threads have completed.

Therefore, I moved the character statistics code into a separate function.



<Figure 11 - Success with FreeBSD9-orig.tar in 1 producer 10 consumer>

Finally, it works!

As I used multiple consumers and commented out the print statements for the consumer threads' work, it only took around a minute. It is faster than 1 producer and 1 consumer.

But, there was a shocking thing for me to figure out.

<Figure 13 - Execution time in 1 producer 100 consumer>

How come 1:100 is much slower than 1:10? Since it uses more threads to process, I thought it would finish quicker than 1:10. But it was not.

So, I asked chatGPT about why it was so.

The main reason was the lock contention, which means that consumer threads have to wait to access the single buffer. As the number of consumers increases, parallel processing becomes possible, which is why 1:10 is much faster than 1:1. However, if the number of threads becomes too large, the consumer threads end up waiting for a long time.

Additionally, since I am using only one producer, it cannot provide sufficient data for multiple consumers, which could also be a reason.

Goal 5) Implement N:M with K shared objects from 1:N

Let's change a single producer to multiple producers to make it faster. I implemented multiple shared objects with a circular queue. In other words, I replaced the buffer with a circular queue.

Unlike 1:N, the shared object stores a single line from the input file at a time.

For example, with 3 producers and 5 consumers using 5 shared objects, each producer reads a single line from the input file and puts it in the shared object(so) in the following sequence.

The first producer:
$$so[0] \rightarrow so[1] \rightarrow so[2] \rightarrow so[3] \rightarrow so[4]$$

The second producer:
$$so[1] \rightarrow so[2] \rightarrow so[3] \rightarrow so[4] \rightarrow so[0]$$

The third producer:
$$so[2] \rightarrow so[3] \rightarrow so[4] \rightarrow so[0] \rightarrow so[1]$$

Also, each consumer reads a single line from the shared object in the following sequence.

The first consumer:
$$so[0] \rightarrow so[1] \rightarrow so[2] \rightarrow so[3] \rightarrow so[4]$$

The second consumer:
$$so[1] \rightarrow so[2] \rightarrow so[3] \rightarrow so[4] \rightarrow so[0]$$

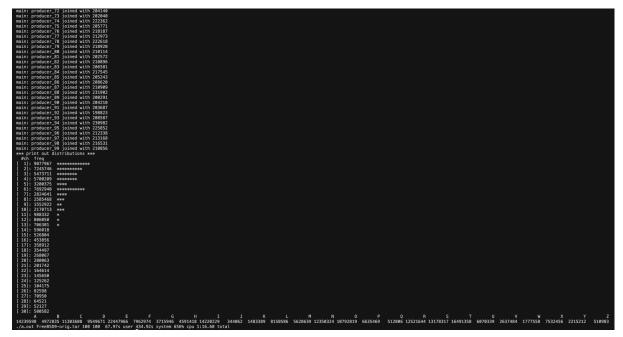
The third consumer:
$$so[2] \rightarrow so[3] \rightarrow so[4] \rightarrow so[0] \rightarrow so[1]$$

The fourth consumer:
$$so[3] \rightarrow so[4] \rightarrow so[0] \rightarrow so[1] \rightarrow so[2]$$

The fifth consumer:
$$so[4] \rightarrow so[0] \rightarrow so[1] \rightarrow so[2] \rightarrow so[3]$$

Also, as the number of producers reading the input file increases, I used a file_lock to avoid race condition. I also used a global_done variable to notify all threads(both producers and consumers) when the producer finish reading the input file. Additionally, I used a global_done_lock to ensure that only the producer who reads the last line of the

input file can change the flag. Finally, I used a so_index_lock to make sure that each producer and consumer uses a different shared object in each thread.



<Figure 14 - Success with FreeBSD9-orig.tar in 100 producer 100 consumer>

As you can see, it took a minute and 16.60 seconds to execute, which is shorter than 1:1 or 1:N.

The performance did increase, but not as much as I expected. I thought that the performance would increase dramatically due to having multiple producers and multiple shared objects, but it wasn't. The reason was that the number of mutex locks increased. Acquiring and releasing mutex locks has high performance impact caused by frequent context switching penalty and overhead.

Build

environment: visual studio code

programming language: C

compile: gcc Nprod_Mcons.c

execute: ./a.out FreeBSD9-orig.tar 100 100

Lesson

It was interesting to implement the concept of thread.

However, it was quite challenging because I've never used the pthread library before. While using functions from the library, ensuring proper synchronization between threads to avoid race conditions and deadlocks was one of the most difficult parts.

While implementing, it was surprising that the performance didn't dramatically improve from 1:1 to 1:N and from 1:N to N:M was surprising. Especially, the fact that acquiring and releasing mutex locks takes some time was one of the main reasons for only a slight, rather than dramatic, performance improvement from 1:N to N:M. According to Amdahl's Law, there is a fundamental limit to performance improvement, no matter how much the processors are parallelized.

When I began implementing N:M, I initially started with a single shared object. But I realized that this would be basically the same as 1:1. So, I tried implementing multiple shared objects using multiple buffers, with each producer mapped to a shared object, allowing each producer to read multiple lines from the input file independently. However, I discovered that this approach was still the same as 1:1. It's no difference from having multiple 1:1 pairs running independently. As a result, I changed multiple buffers to multiple circular queues. I realized that for N:M implementation to be effective, multiple shared objects are essential.

If time allows, I want to implement two types of N:M.

- 1) N:M with K shared objects using multiple buffers in a different algorithm from the one I used before.
- 2) N:M with K shared objects using a buffer and a circular queue.