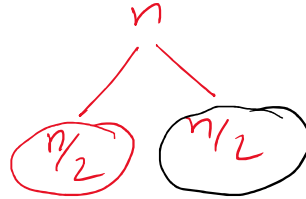
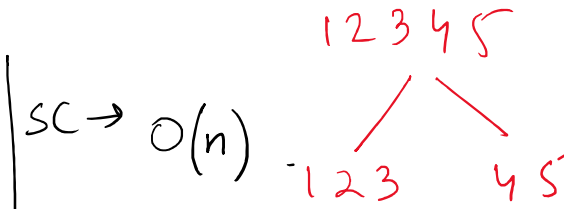


Merge Sort-

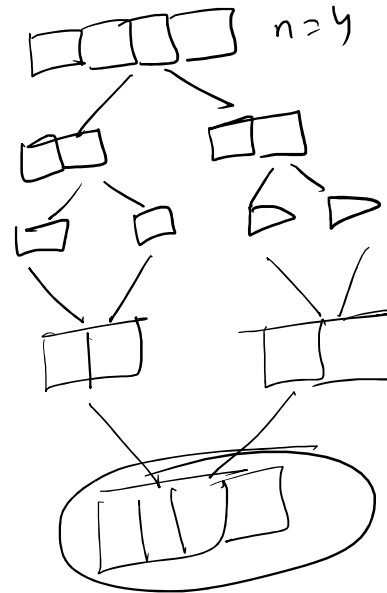


$$T(n) = T(n/2) + T(n/2) + n$$

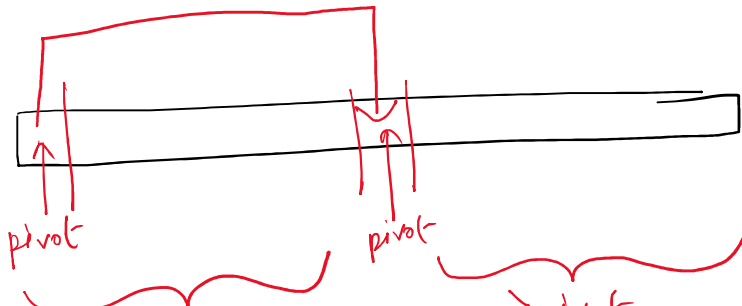
$$\boxed{T(n) = 2T(n/2) + n} \rightarrow O(n \log n)$$



Quick Sort- Based on divide & conquer technique



5 4 3 2 1
 \Downarrow
 4 3 4 (5)



pivot element.
 \downarrow
 3 4 4 5 choose
 4 as pivot

pivot $\underbrace{\hspace{1cm}}$ \leq pivot $\underbrace{\hspace{1cm}}$ $>$ pivot
 3rd element choose
 karrn.
 \downarrow
 1st / Last element

3

35 50 15 25 80 20 90 45
 P q

pivot = 35

$p \rightarrow arr[p] \leq pivot$

$q \leftarrow arr[q] > pivot$

35 20 15 25 80 50 90 45
~~P~~ ~~P~~ ~~P~~ P q

① Chk if p & q
 have crossed each
 other

Swap(arr[p], arr[q]);

$\overbrace{25 \ 20 \ 15}^{n/2 - 1}$ $\overbrace{80 \ 50 \ 90 \ 45}^{n/2}$
 35
 OP

② Chk if p & q have
 not crossed each
 other =

Swap(arr[p], arr[q]);

≤ 35

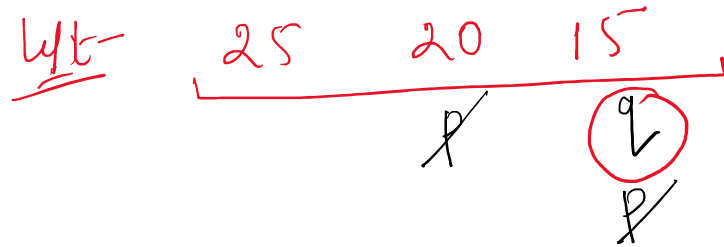
left

> 35

right

left

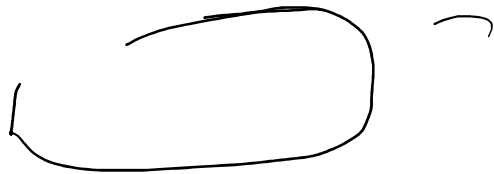
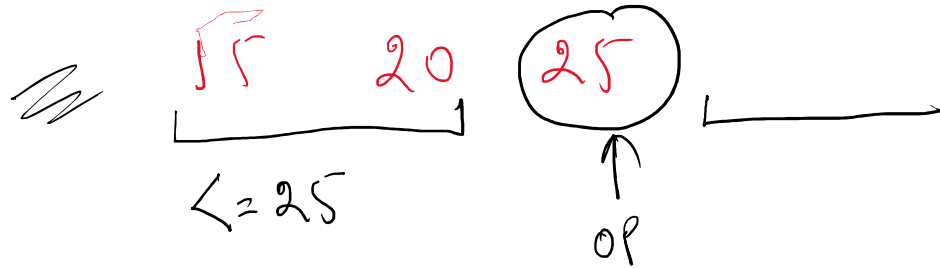
right



pivot = 25

$arr(p) \leq pivot$
 $arr(q) > pivot$

\Rightarrow
 $swap(pivot, arr(q))$



Right

80

50

90

45

pivot = 80

P

P

q

111



8

80

50

45

90

~~P~~

~~q~~

q

P



45

50

[45 50]

≤ 80

80

↑
op

[90]

> 80

SC - $O(1)$

TC - Best Case

pivot element → original position →

Worst Case

pivot element - - - 0
middle pc.

$$T(n) = T(n/2) + T(n/2) + n$$
$$= 2T(n/2) + n$$

$$T(n) = O(n \log n)$$