

Master Theorem -

If the recurrence is of the form  $T(n) = aT(\frac{n}{b}) + \Theta(\tilde{n}^k \log^p n)$ , where  $a \geq 1, b > 1, k \geq 0$  and  $p$  is a real number, then:

- 1) If  $a > b^k$ , then  $T(n) = \Theta(n^{\log_b a})$
- 2) If  $a = b^k$ 
  - a. If  $p > -1$ , then  $T(n) = \Theta(n^{\log_b a} \log^{p+1} n)$
  - b. If  $p = -1$ , then  $T(n) = \Theta(n^{\log_b a} \log \log n)$
  - c. If  $p < -1$ , then  $T(n) = \Theta(n^{\log_b a})$
- 3) If  $a < b^k$ 
  - a. If  $p \geq 0$ , then  $T(n) = \Theta(n^k \log^p n)$
  - b. If  $p < 0$ , then  $T(n) = O(n^k)$

$$T(n) = T(n-1) + 2$$

$$T(n) = 3T(n/2) + n^2$$

$$n^k \log^p n$$

$$aT(n/b) \quad 3T(n/2)$$

$$a=3, b=2, k=2, p=0$$

$$a=3, b=2, k=2, b^k = 2^2 = 4$$

$$T(n) = \Theta(n^k \log^p n) \\ = \Theta(n^2 \log^0 n) = \Theta(n^2)$$

$$a < b^k \quad \text{Cond'n 3rd.}$$

$$p=0, \text{ Cond'n - } \underline{\underline{3a.}}$$

$$T(n) = 4T(n/2) + n^2$$

$$\textcircled{1} \quad a=4, b=2, k=2, p=0$$

$$\textcircled{2} \quad a=4, b^k = 2^2 = 4, a = b^k$$

If the recurrence is of the form  $T(n) = aT(\frac{n}{b}) + \Theta(\tilde{n}^k \log^p n)$ , where  $a \geq 1, b > 1, k \geq 0$  and  $p$  is a real number, then:

- 1) If  $a > b^k$ , then  $T(n) = \Theta(n^{\log_b a})$
- 2) If  $a = b^k$ 
  - a. If  $p > -1$ , then  $T(n) = \Theta(n^{\log_b a} \log^{p+1} n)$
  - b. If  $p = -1$ , then  $T(n) = \Theta(n^{\log_b a} \log \log n)$

②  $a=4, b^k = 2^2=4, a=b^k$

Cond<sup>n</sup> 2<sup>nd</sup>.

$p = 0 > -1$  (2<sup>a</sup>)

$$T(n) = \Theta \left( n^{\log_b a} \cdot \log^{p+1} n \right) = \Theta \left( n^{\log_2 4} \cdot \log^{0+1} n \right) = \Theta$$

$$\log_2 4 = \log_2 2^2 = 2 \times \log_2 2 = 2$$

$$= \Theta \left( n^2 \cdot \log n \right)$$

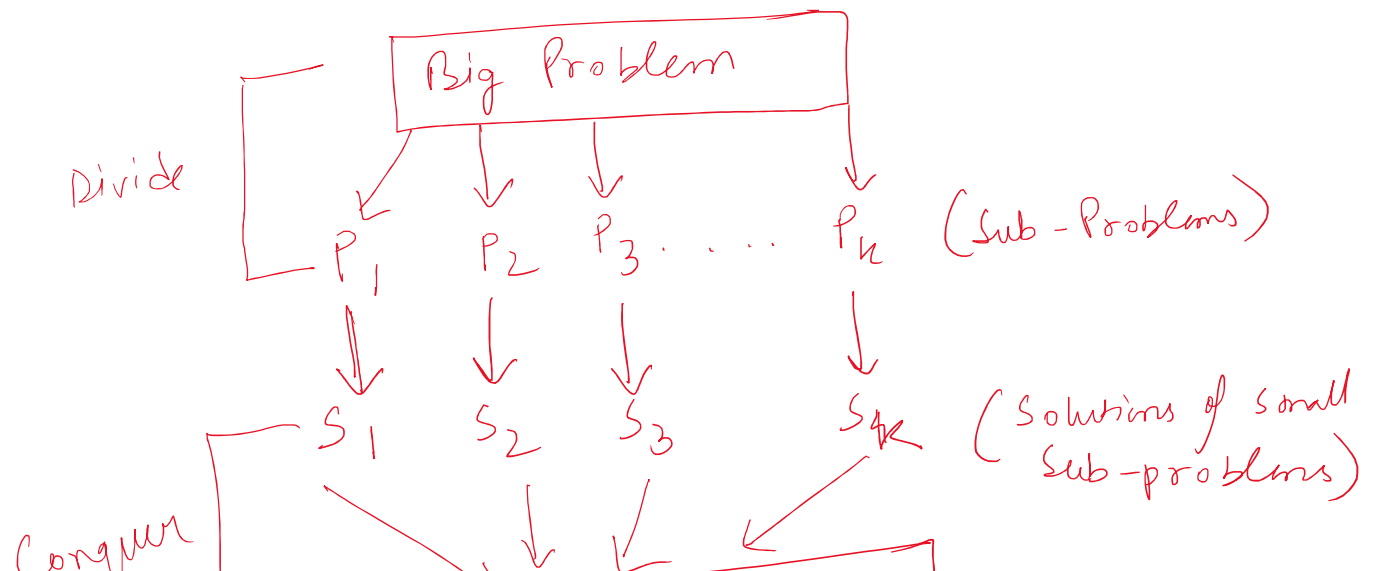
2) If  $a = b^k$

- a. If  $p > -1$ , then  $T(n) = \Theta(n^{\log_b a} \log^{p+1} n)$
- b. If  $p = -1$ , then  $T(n) = \Theta(n^{\log_b a} \log \log n)$
- c. If  $p < -1$ , then  $T(n) = \Theta(n^{\log_b a})$

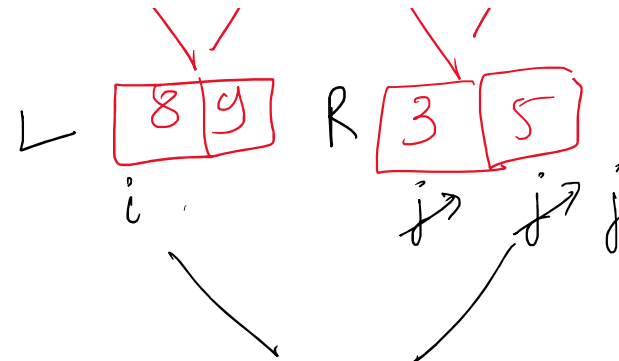
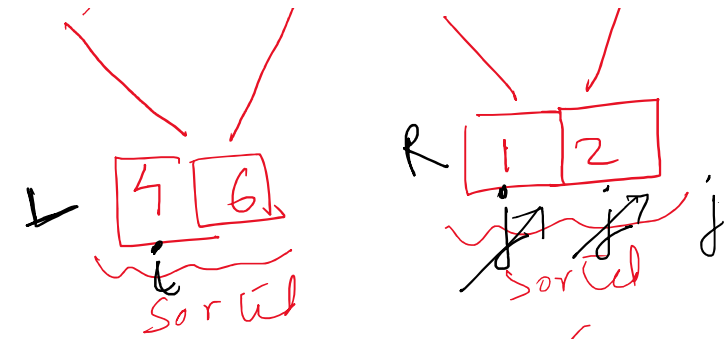
3) If  $a < b^k$

- a. If  $p \geq 0$ , then  $T(n) = \Theta(n^k \log^p n)$
- b. If  $p < 0$ , then  $T(n) = O(n^k)$

Divide & Conquer - Technique to ~~solve~~ design algorithms.







Compare  
 $L[i]$   $R[j]$

