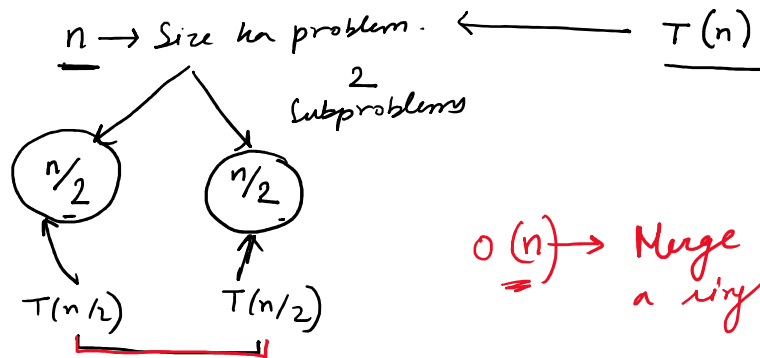


Merge Sort Analysis -

$$\rightarrow T(n) = \underbrace{2T(n/2)}_{\text{Divide}} + \underbrace{n}_{\text{Merge/Compare}}$$



$O(n) \rightarrow$ Merge 2 sorted arrays into a single sorted array.

$$T(n) = 2T(n/2) + n.$$

$$T(n) = aT(n/b) + n^k \log^p n$$

\Downarrow By Master Th^m \rightarrow

$$a=2, b=2, k=1, p=0$$

$$a < b^k$$

$$2 < 2^1$$

$$\text{Case 2a} \rightarrow T(n) = O\left(n^{\log_b a} \times \log^{p+1} n\right)$$

Property $\left(\log_a a\right) = 1$

$$= O\left(n^{\log_2 2} \cdot \log^{0+1} n\right)$$

$$= O(n^1 \cdot \log n) = n \log n.$$

Space Complexity - $O(n) \rightarrow$ Merge 2 sorted arrays ~~into~~ into 1 sorted array then it uses

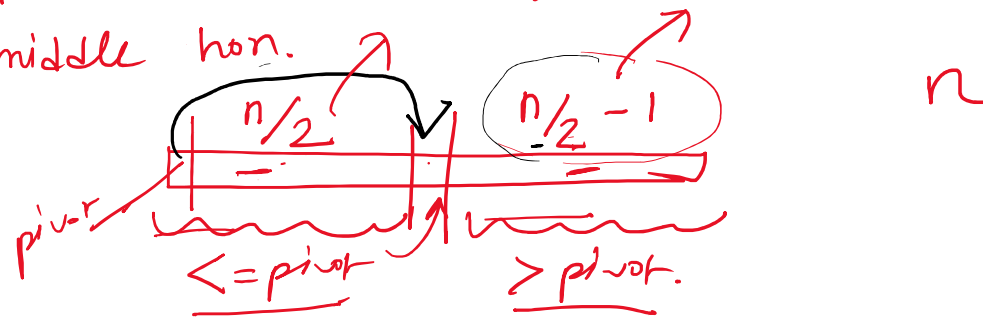
Space \rightarrow $O(1)$ \rightarrow Merge 2 sorted arrays \rightarrow \rightarrow into 1 sorted array then it uses a temporary array of size = left arr size + right arr size.

Temp arr ka max size = n

Quick Sort Analysis \rightarrow TC, SC.

SC $\rightarrow O(1) \rightarrow$ No XTRA space used.

TC \rightarrow Best Case = ① Pivot element ka original position Hamnecha middle hon.



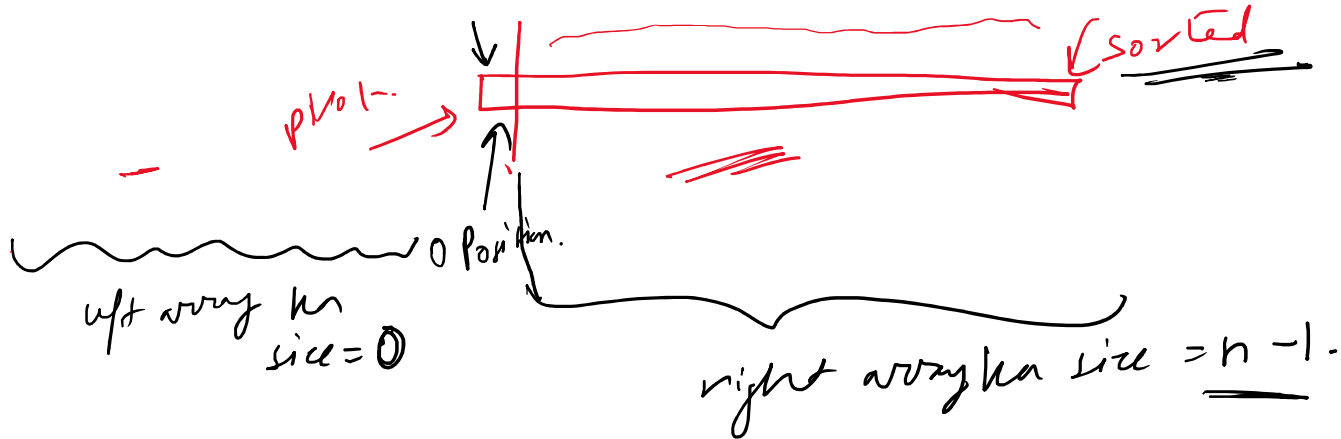
Recursive Relatⁿ = $T(n) = 2T(n/2) + n$ \rightarrow Swapping perform whole array traverse

~~Master Th^m~~

$T(n) = O(n \log n)$

$$T(n) = O(n \log n)$$

Worst case - ① Input array is already sorted and you r considering 1st element as pivot.



$$RR \Rightarrow T(n) = T(n-1) + n \quad \leftarrow \text{swap}$$

$$T(n) = aT(n/b) + (n^h \log^p n)$$

\Downarrow Substitution method.

$$T(n) = T(n-1) + n$$

$$T(n-1) = T(n-2) + (n-1)$$

$$T(n-2) = T(n-3) + (n-2)$$

⋮

$$T(2) = T(1) + 2.$$

$$T(1) = \cancel{T(0)} + 1.$$

$$T(n) + \cancel{T(n-1)} + \cancel{T(n-2)} + \dots + \cancel{T(2)} + \cancel{T(1)} = \cancel{T(n-1)} + n + \cancel{T(n-2)} + (n-1) + \cancel{T(n-3)} + (n-2) + \dots + \cancel{T(1)} + 2 + 1.$$

$$\Rightarrow T(n) = \underbrace{n + (n-1) + (n-2) + \dots + 2 + 1}_{\text{Sum of first } n \text{ natural numbers}}$$

$$= \frac{n \times (n+1)}{2} = \frac{n^2 + n}{2} \approx O(n^2).$$

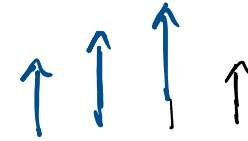
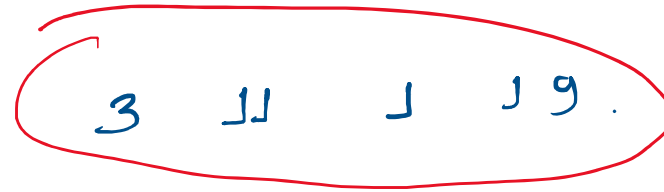
By Case - ~~TC~~ $TC \rightarrow O(n \log n).$

3, 3, 3, 3, 3

$SC \rightarrow O(1)$ — \because No extra space
is used.

5

Insertion sort

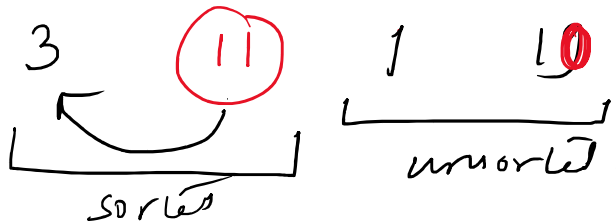


eg

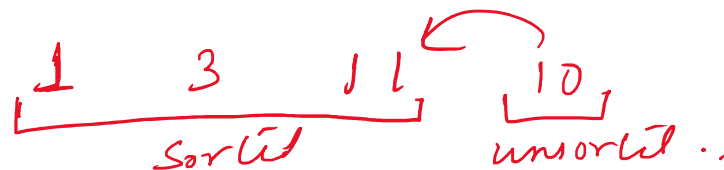
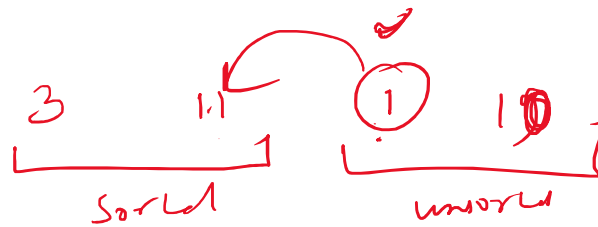
①

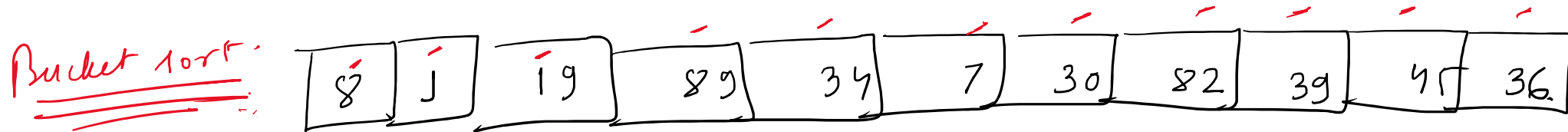


②

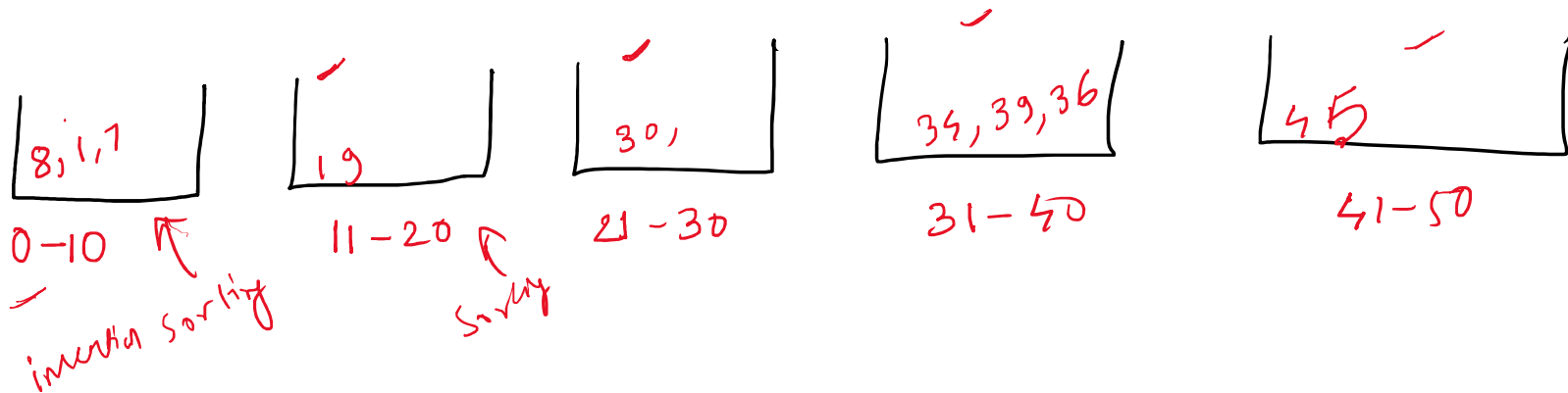


③



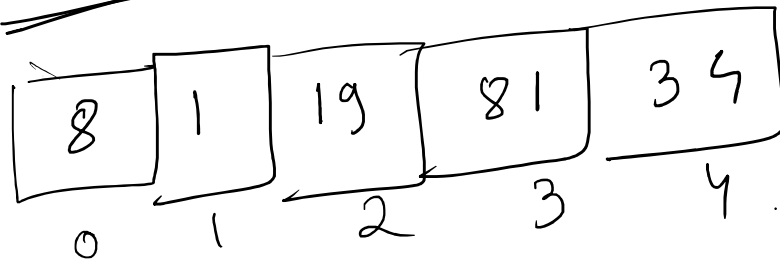


Q: ~~Bucket sort~~ No. of buckets = 10.



1, 7, 8, 19, 30, 34, 36, 39, 45, 82, 89 \rightarrow sorted

~~Linear~~ Searching - Array \rightarrow Element to Search -



$X = 81$ \leftarrow
return index;

```
ind = -1;  
for (i = 0; i <= n - 1; i++) {  
    if (arr[i] == X) {  
        ind = i;  
        print(ind);  
        break;  
    }  
}
```

Eg

8 1 5 20

$X = 1$

return index where X
occurs for 1st time.

return index where X

```
}  
if (ind == -1) {  
    print (Element not found)  
}  
}
```

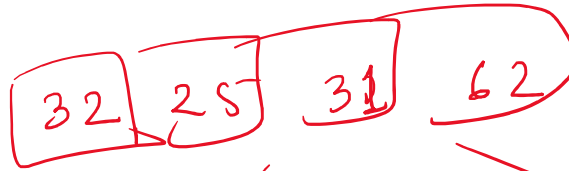
return index where X occurs for last time.

TC →

bcd bch sdbcs dbkcs

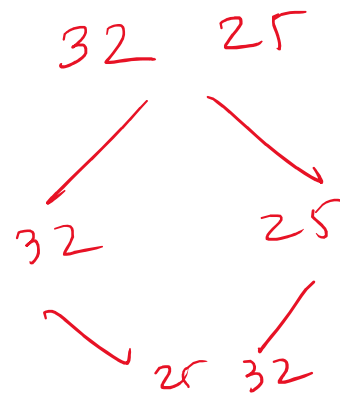
Algorithm

example →

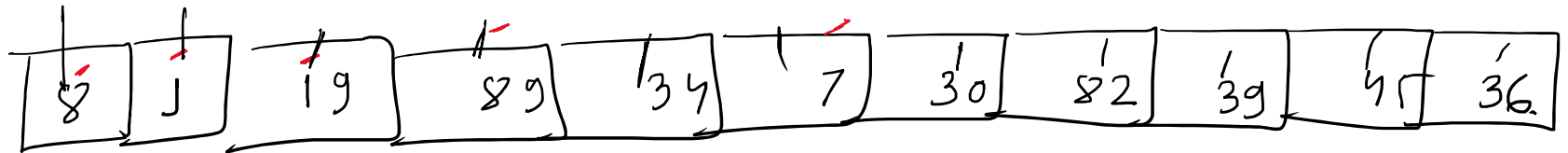


Apply Merge Sort.

⇒



Imp Arr =



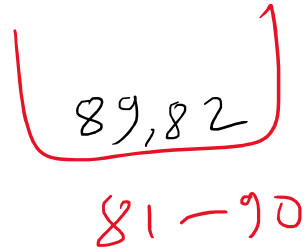
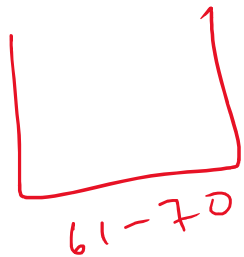
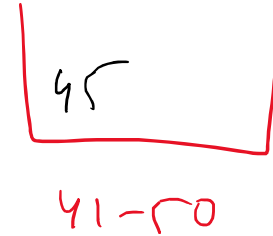
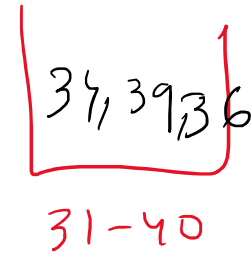
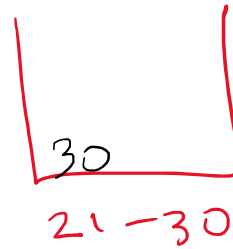
Prv =

No. of buckets = 10.

Min = 1

Max = 89.

Insertion Sort



1, 7, 8, 19, 30, 34, 36, 39, 45, 82, 89.