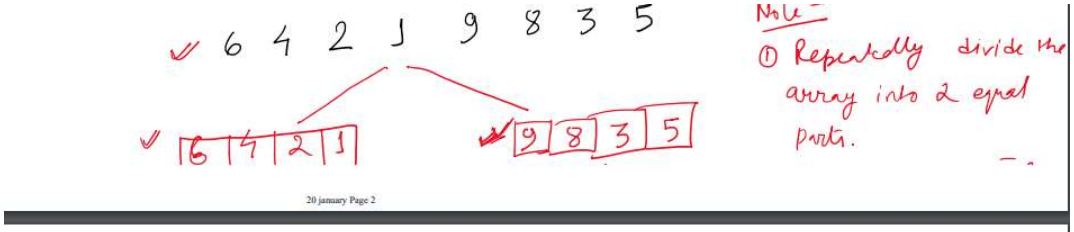
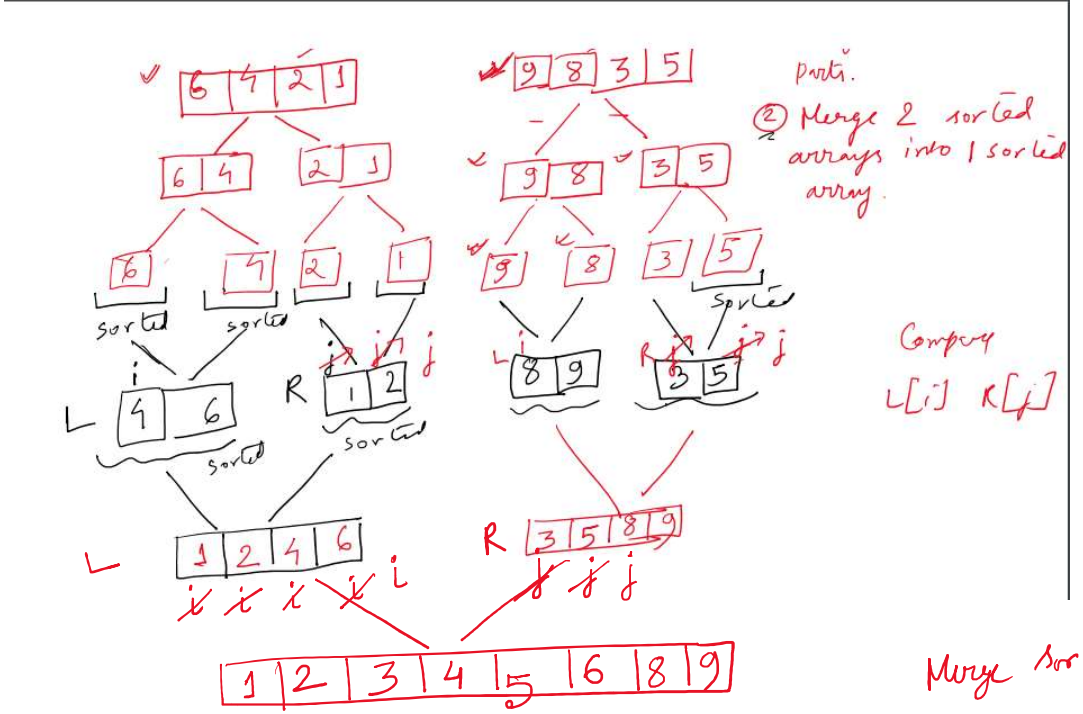
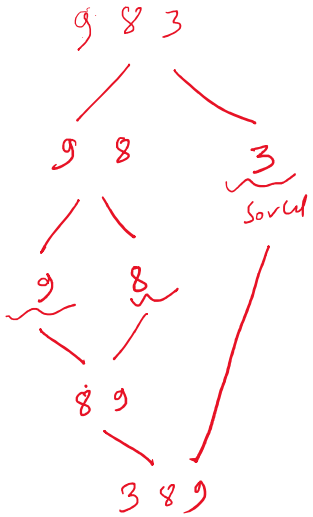






Merge sort is Recursive.

Recursive Code → Recurrence Relation (Solve) → Time Complexity

TC- T(n)

Suppose- Merge sort ka code T(n)

$T(n) = T(n/2) + T(n/2) + n$

Merge karne ka Time

Actual problem

2 subproblems

Master Th^m

$T(n) = 2T(n/2) + n$

$a=2, b=2, k=1, p=0$

$a \quad b^k$

$2 = 2^1$

2nd Condⁿ.

$k=1$

$p=0$

(2a) $\rightarrow$ $n^{\log_b a} \cdot \log^{p+1} n = n^{\log_2 2} \cdot \log^{0+1} n$

$= n \cdot \log n$

Quick Sort – Divide & Conquer Technique.

35 50 15 25 80 20 90 45

pivot = 35

p

q

50 > 35

Note –

① pivot element.

② pivot element = 1st element.

③ p → arr[p] <= pivot.

q → arr[q] > pivot.

= 50 / 35

35 20 15 25 80 50 90 45

p p p p q

q → arr[q] > pivot.

④ If p & q have not crossed each other, swap them

⑤ If p & q crossed each other, swap (pivot, arr[q])

25 20 15 | 35 | 80 50 90 45

<=35 | original position | >35

Same step apply | Same step apply here

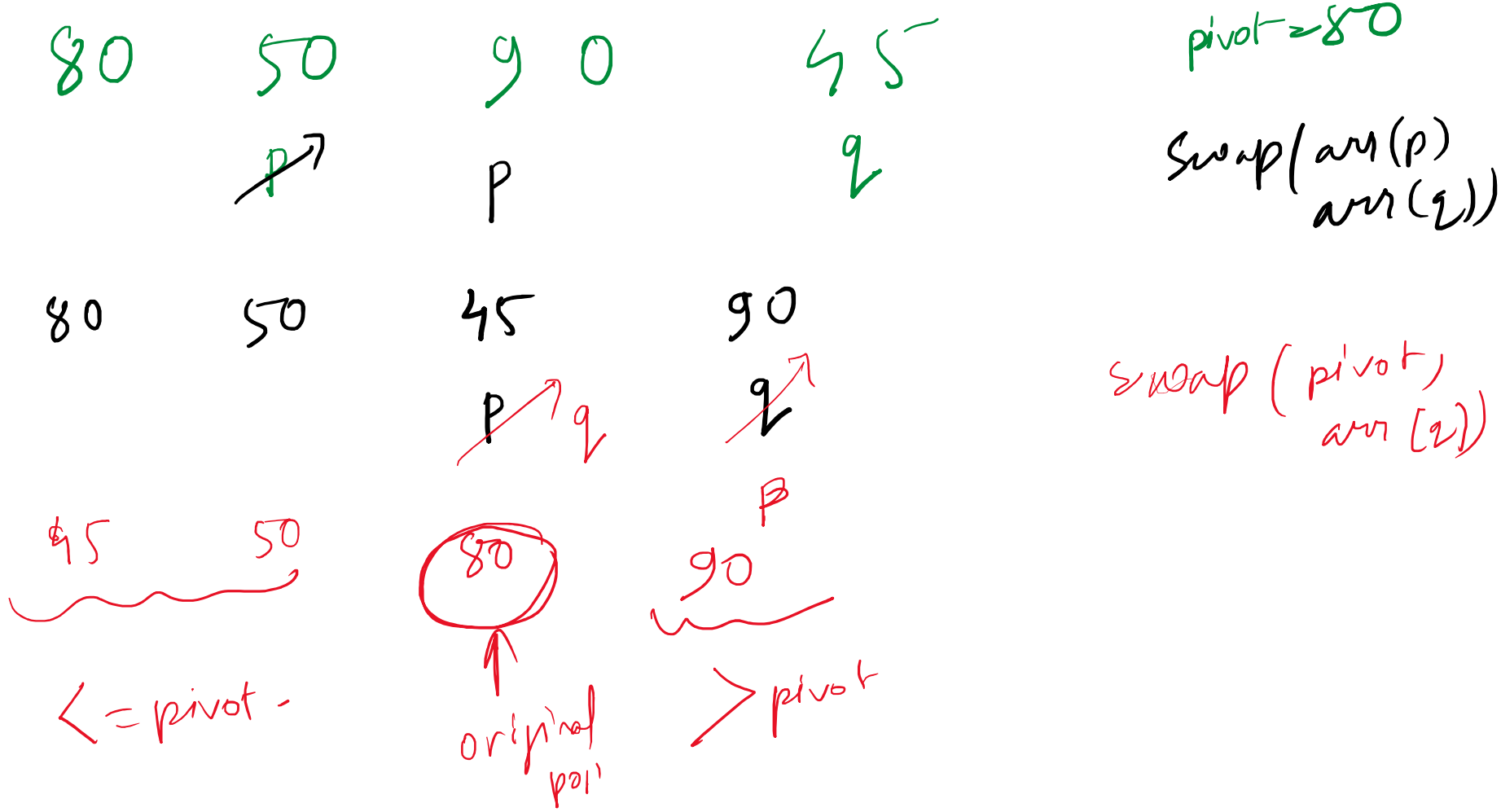
Left 25 20 15

p q

pivot = 25

swap (arr(q), pivot)

p p 15 20 25


80 50 90 45
p
p
q
pivot=80
swap(arr(p)
arr(q))
80 50 45 90
p
q
p
q
swap (pivot)
arr[q])
45 50
80
90
<=pivot
original pos
>pivot