

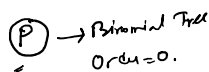
A **Binomial Tree** B_k is an ordered tree defined recursively, where k represents the order of the binomial tree.

- If the binomial tree is of order 0 (B_0), it consists of a single node.
- In general, a binomial tree of order k (B_k) consists of two binomial trees of order $k-1$, where one is linked as the left subtree of the other.

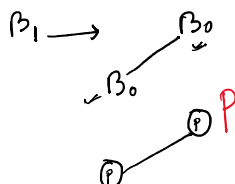
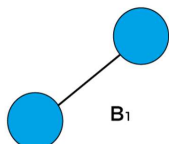
If B_0 , where k is 0, there would exist only one node in the tree.



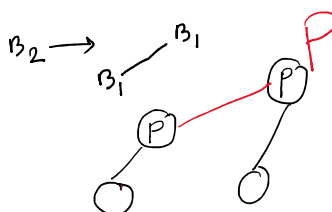
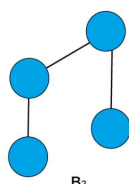
B_k .



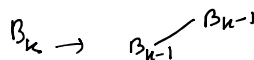
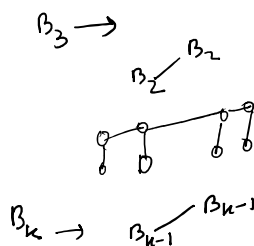
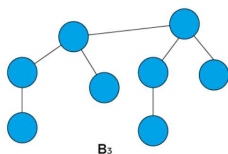
If B_1 , where k is 1. Therefore, there would be two binomial trees of B_1 in which one B_0 becomes the left subtree of another B_0 .



If B_2 , where k is 2. Therefore, there would be two binomial trees of B_2 in which one B_1 becomes the left subtree of another B_1 .

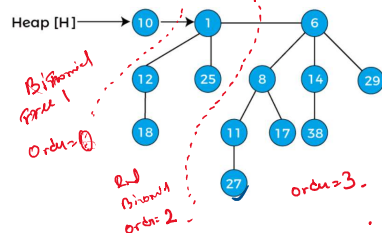


If B_3 , where k is 3. Therefore, there would be two binomial trees of B_3 in which one B_2 becomes the left subtree of another B_2 .



A **binomial heap** is a collection of binomial trees that satisfies the following binomial heap properties:

- No two binomial trees in the collection have the same order.
- Every binomial tree in the heap must follow the min-heap property, i.e., the value of a child node is greater than parent node.



Binomial Heap Union Operation

To perform the union of two binomial heaps, we have to consider the below cases -

Case 1: If $\text{degree}[x]$ is not equal to $\text{degree}[\text{next } x]$, then move pointer ahead.

Case 2: if $\text{degree}[x] = \text{degree}[\text{next } x] = \text{degree}[\text{sibling}(\text{next } x)]$ then,

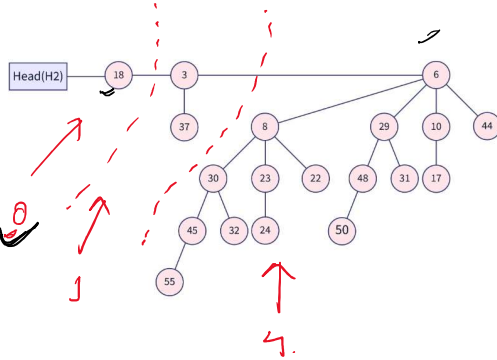
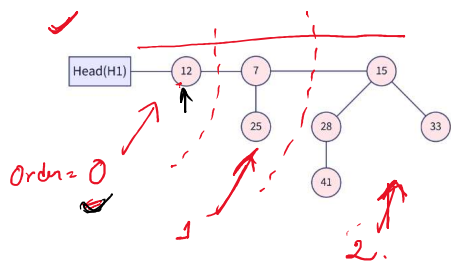
Move the pointer ahead.

Case 3: If $\text{degree}[x] = \text{degree}[\text{next } x]$ but not equal to $\text{degree}[\text{sibling}(\text{next } x)]$

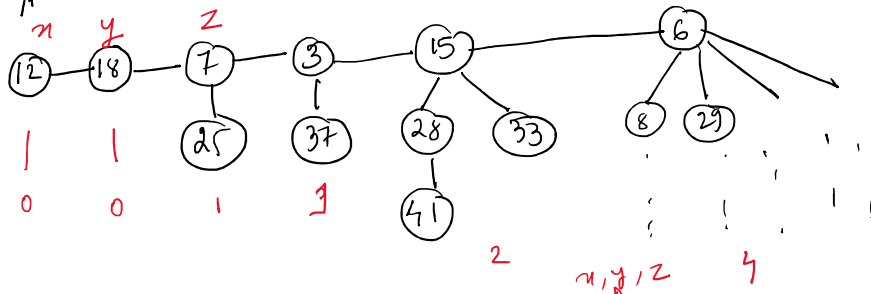
and $\text{key}[x] < \text{key}[\text{next } x]$ then remove $[\text{next } x]$ from root and attached to x .

Case 4: If $\text{degree}[x] = \text{degree}[\text{next } x]$ but not equal to $\text{degree}[\text{sibling}(\text{next } x)]$

and $\text{key}[x] > \text{key}[\text{next } x]$ then remove x from root and attached to $[\text{next } x]$.

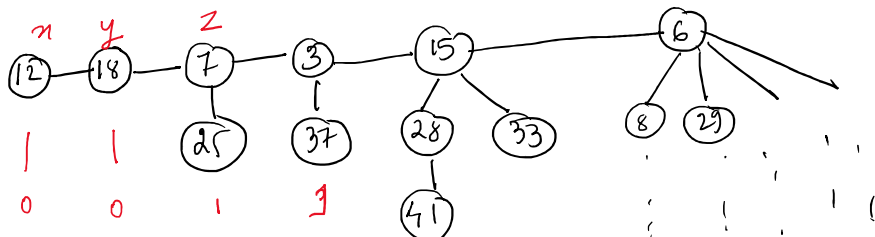
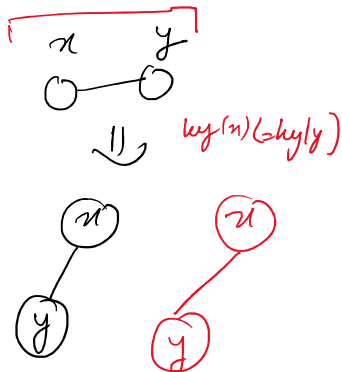
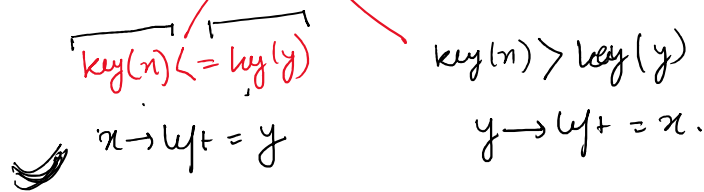


Order $\uparrow \rightarrow 0, 0, 1, 1, 2, 4$



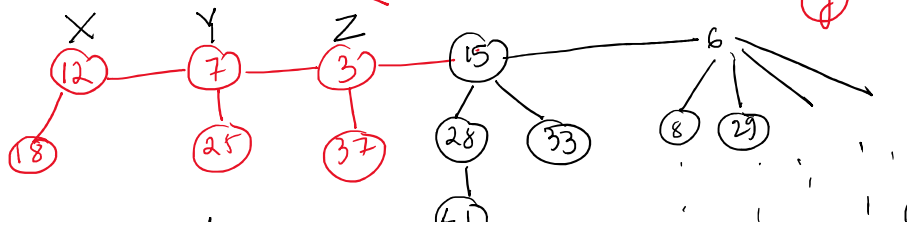
- Algo ① $\text{degree}(x) \neq \text{degree}(y) \rightarrow$ Move pointers in right direction by 1 place.
 ② $\text{degree}(x) = \text{degree}(y) = \text{degree}(z) \rightarrow$ Move pointers in right direction by 1 place.

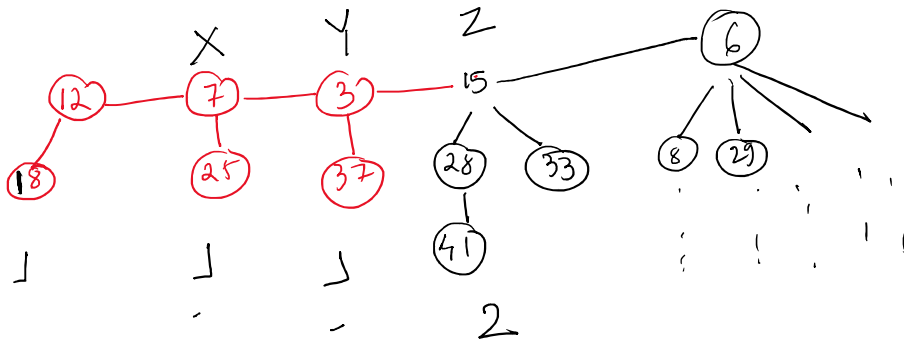
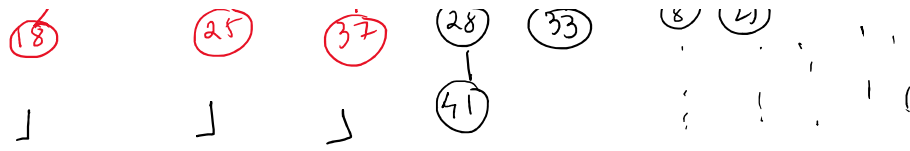
③ $\text{degree}(x) = \text{degree}(y) \neq \text{degree}(z)$.



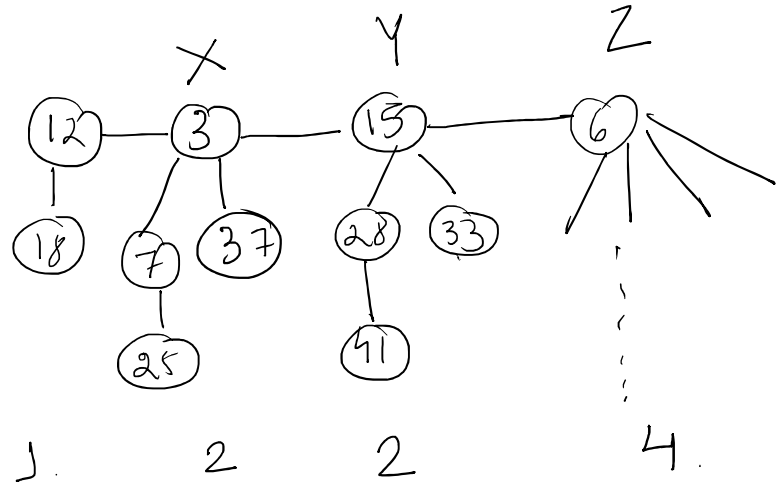
$\text{deg}(x) = 0, \text{deg}(y) = 0, \text{deg}(z) = 1$

Case 3 $\text{key}(x) < \text{key}(y)$





$\log(n)$ $\log(y)$
7 3



$\log(n)$ $\log(y)$
3 < 15.

