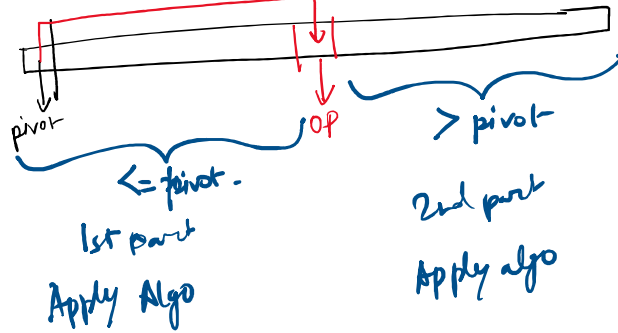


Quick sort -

Based on divide & conquer technique.



pivot element  
 ↓  
 3rd part choose  
 Karo.  
 ↓  
 arr ka first/  
 last element

• pivot = 35

35 50 15 25 80 20 90 45  
 (P) q % %

35 20 15 25 80 50 90 45  
 % % % (P) %

$n/2 - 1$   $n/2$   
 25 20 15 35 80 50 90 45  
 ↓ ↓  
 OP

$P \rightarrow \text{arr}[p] \leq \text{pivot}$   
 $q \leftarrow \text{arr}[q] > \text{pivot}$   
 ○ chk if p & q  
 have crossed  
 each other  
 → swap (pivot, arr[q])  
 ○ chk if p & q  
 havenot crossed  
 each other  
 → swap (arr[p], arr[q])

Left- 25 20 15  
 % % (P)  
 % (P)

15 20 25  
 ↓  
 OP  
 20 15 25  
 ↓  
 OP

pivot = 25  
 $\text{arr}[p] \leq \text{pivot}$   
 $\text{arr}[q] > \text{pivot}$

$15 \leq 25$   
 swap (arr[q], pivot)

80 50 90 45  
 % (P) (q)

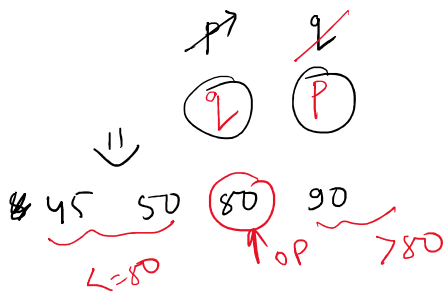
pivot = 80

$\text{arr}(p) \leq \text{pivot}$   
 $\text{arr}(q) > \text{pivot}$

swap (p, q)

80 50 45 90  
 % %  
 (P) (q)

swap (pivot, q)



Swap(pivot, q)

Merge Sort Recur<sup>n</sup>  $\Rightarrow T(n) = 2T(n/2) + n$

$TC = O(n \log n)$   
 $SC = O(n)$

Quick Sort SC -  $O(1)$   
Time Complexity -

① 1

① Best Case

② Pivot element  $\rightarrow$  original position  $\rightarrow$  middle pc

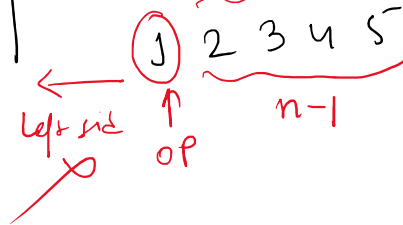
$$T(n) = 2T(n/2) + n$$

$$= O(n \log n)$$

② Worst Case - Sorted array as input

1 2 3 4 5

pivot = 1

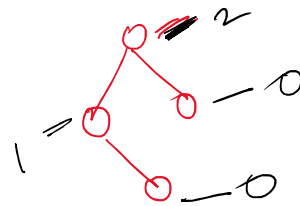


$$T(n) = T(n-1) + n$$

$$= O(n^2)$$

Trees -

Binary Trees - Nodes with 0, 1, 2 children.



Almost Complete Binary Tree

$\rightarrow$  Jab koi nodes insert karo, toh left direction se karo.

