

If the recurrence is of the form $T(n) = aT(\frac{n}{b}) + \Theta(n^k \log^p n)$, where $a \geq 1, b > 1, k \geq 0$ and p is a real number, then:

- 1) If $a > b^k$, then $T(n) = \Theta(n^{\log_b a})$
- 2) If $a = b^k$
 - a. If $p > -1$, then $T(n) = \Theta(n^{\log_b a} \log^{p+1} n)$
 - b. If $p = -1$, then $T(n) = \Theta(n^{\log_b a} \log \log n)$
 - c. If $p < -1$, then $T(n) = \Theta(n^{\log_b a})$
- 3) If $a < b^k$
 - a. If $p \geq 0$, then $T(n) = \Theta(n^k \log^p n)$
 - b. If $p < 0$, then $T(n) = O(n^k)$

$$\left| \begin{array}{c} \Theta(n^k \log^p n) \\ \hline n^2 \end{array} \right|$$

$k=2, p=0$
 $(a)^0 = 1$

Theorem
~~Formal~~

$$T(n) = 3T(n/2) + n^2$$

$$a \times T(n/b) \quad 3T(n/2)$$

$a=3, b=2, k=2$

Step 1 - Chk if MT can be applied/not

Step 2 - $a, b, k, p \rightarrow ??$

Step 3 - $a < b^k$
 $3 < 2^2$

Step 4 - Condⁿ 3rd

Step 5 - $p = 0$ (condition 3a)
~~Step 4~~ $T(n) = \Theta(n^k \log^p n)$
 $= \Theta(n^2 (\log n)^0)$
 $= \Theta(n^2 \cdot 1)$
 $= \Theta(n^2)$

$$T(n) \leq T(n/2) + n^2.$$