Greedy (लालची) --> Greed (लालच)

Greedy algorithms are a class of algorithms that make **locally optimal** choices at each step with the hope of finding a **global optimum** solution. Examples of popular Greedy Algorithms are Fractional Knapsack, Dijkstra's algorithm, Kruskal's algorithm, Huffman coding and Prim's Algorithm.

Greedy: Think about present (**वर्तमान**)and not future (**भविष्य**).

Dynamic Programming: Think about future and not present. Think if in future it will give optimum solution (best possible solution) or not.

Eg: Cloth shop. During Festival time, we get offers on clothes in different shops. So what is our tendency, we move to that shop where:
1. Price is low ignoring the quality of the cloth. This is greedy.

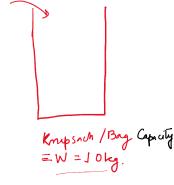DP= Low price along with quality of cloth matters.

Fractional Knapsack Problem

Given two arrays named **value** and **weight** of size n each, where **value[i]** represents the value and **weight[i]** represents the weight associated with the ith item. Also, we have been provided a knapsack with a maximum capacity of W.

The task is to pick some items (possibly all of them) such that the sum of values of all the items is maximized and the sum of weights of all the items is at most W.

As the name Fractional suggests, here we are allowed to take a fraction of an item.

For an item of weight 20 Kg. we can either take 1 Kg, 2 Kg, ..., or 19 Kg but we are not allowed to take 1.2 Kg or 5.7

Item	1	2	3	4	5
Value/Profit	10	15	10	20	8
Weight	3	3	2	5	1

Knapsach /Bag Capacity
= W = 10kg.

So, we can opt for a greedy algorithm. We can have several potentially correct strategies, some of the obvious ones are :

1. Picking the items with the largest values first.
2. Picking the items with the lowest weights first.
3. Picking the items based on some sort of ratio among their values and weights.

Need to consider both profit and weight at the same time. Value DIVIDE by weight gives the profit for 1 kg of item.

1.
Object |||| profit |||| weight |||| remaining_weight

1. Greedy about profit: Means place those items first in the bag/knapsack which has more profit. This will in turn maximise our profit. जिस object का सबसे ज्यादा मूल्य है, उसे बैग के अंदर डालूंगा।

Item	1	2	3	4	5
Value/Profit	10	15	10	20	8
Weight	3	3	2	5	1

Knapsach capacity = 10kg.

Item	Profit	Weight	Remaining weight
4	20	5	10 - 5 = 5.
2	15	3	5 - 3 = 2.
3	10	2	2 - 2 = 0.
	(45).		

Greedy about weight: Means place those items first in the bag/knapsack which has least weight. This will help us to put more items in our bag which will in turn maximise our profit. अगर मैं कम वजन के वस्तुओं को बैग के अंदर डालूंगा, तो मैं ज्यादा वस्तुएं बैग के अंदर डाल सकता हूँ।

Tuition Center - 1 student = Rs 500 fee.

Knapsack remaining capacity.

Item	Profit	Weight	Remaining weight
5	8	1	10−1=9
3	10	2	9−2=7
2	15	3	7−3=4
1	10	3	4−3=1
4	$\frac{20}{5}=4$	1	1−1=0

(47)

Item	1	2	3	4	5
Value/Profit	10	15	10	20	8
Weight	3	3	2	5	1

③ Greedy about profit ll weight (both).

① Find price of 1 kg for each item.

② Sort in descending order of
$\frac{Profit}{Weight}$.

Item	1	2	3	4	5
Value/Profit	10	15	10	20	8
Weight	3	3	2	5	1
1 kg price	3.33	5	5	4	8

Item	Profit	Weight	Remaining Wt
5	8	1	10−1=9
2	15	3	9−3=6
3	10	2	6−2=4
4	16	4	4−4=0

(59)

Item	2	3
Val	15	10
wt	3	2
1 kg	5	5

Knapsack = 64

Item   Profit   Weight   Remaining weight

Greedy Technique	DP (Dynamic Programming)
(i) Current step में Optimum Solution --hope for--> Future optimum Sol^n	(1) Current operations --> future optimum Solution.

Eg Dress.

Shop1   Shop2

50% Discount   20% Discount.

↓ Purchase.